Efficient Crawling Through URL Ordering

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Outline

- Crawlers and Search Engines
- "Importance" Metrics
- Crawler Performance
- Crawler Models
- Experiments and Observations
- Conclusion
- Discussion

Search Engine Architecture



Crawling Web Pages

What pages should the crawler download?

- How should the crawler refresh pages?
- How should the load on the visited Web sites be minimized?
- How should the crawling process be parallelized?

What Pages Should the Crawler Visit First?

"important" pages

- Different importance metrics:
 - Similarity to a Driving query
 - Backlink Count
 - ◆ PageRank
 - Location
 - Forward Link Count

The Similarity Metric

Given a query q, the importance of a page p is defined to be the "textual similarity" between p and q [Arasu et al.] I(p) = IS(p, q)

The Backlink Count Metric

The importance of a page p is defined to be the number of links to p that appear over the entire Web

I(p) = IB(p)

IB'(p) denotes the estimated value of IB(p)

The PageRank Metric

The importance of a page p is defined to be the weighted sum of the importance of the pages that have backlinks to p I(p) = IR(p)

IR'(p) denotes the estimated value of IR(p)

 $\frac{IR(p)}{IR(p)} = (1-d) + d \left[\frac{IR(t_1)}{c_1} + \dots + \frac{IR(t_n)}{c_n} \right]$

Other Importance Metrics

Forward Link Count
Location Metric
Combinations of metrics

I(p) = k₁. IS(p, q) + k₂. IB(p)

The Question is

How to design a crawler that if possible visits high *I(p)* pages before lower ranked ones, for some definition of *I(p)*

Crawler performance

- Crawler models
 - Crawl and Stop
 - Crawl and Stop with Threshold
 - Limited Buffer Crawl

Crawl and Stop

 The crawler starts at an initial page and stops after visiting K pages

$$r_1, \dots, r_M, \dots, r_K$$

$$I(r_i) \ge I(rk)$$

P_{ST}(C) = M / K
P_{ST}(ideal crawler) = 1
P_{ST}(random crawler) = K / T, where T is the
total number of pages in the Web

Crawl and Stop with Threshold

The crawler starts at an initial page and stops after visiting K pages • Any page p with $I(p) \ge G$ (where G is a given importance target) is considered hot • $P_{ST}(C)$ = Total number of visited hot pages Total number of hot pages (H) P_{ST} (ideal crawler) = $\begin{cases} K/H, \text{ if } K < H\\ 1 & \text{otherwise} \end{cases}$ P_{ST} (random crawler) = $\begin{cases} K/T, \text{ if } K < T\\ 1 & \text{otherwise} \end{cases}$ CS 856

Limited Buffer Crawl

 The crawler can keep at most B pages in its buffer

Total number of hot pages in the buffer



Ordering Metrics

- A crawler keeps a queue of URLs it has seen during a crawl
- The crawler must select from this queue the next URL to visit
- An ordering metric O is used for this selection
- The chosen O metric should be suitable for the importance metric in mind

Experiments

 Define the entire Web to be a portion of the Stanford University Web pages (179,000 pages)

- Measure the performance of various ordering metrics for the importance metric *IB(p)*
- Measure the performance of various ordering metrics for the importance metric *IS(p, Q)*

Algorithm 5.1 Crawling algorithm (backlink based) Input: starting_url: seed URL Procedure:

```
[1] enqueue(url_queue, starting_url)
[2] while (not empty(url_queue))
[3]
     url = dequeue(url_queue)
[4]
     page = crawl_page(url)
[5]
     enqueue(crawled_pages, (url, page))
[6]
     url_list = extract_urls(page)
[7]
     foreach u in url list
[8]
        enqueue(links, (url, u))
        if (u \notin url_queue and (u, -) \notin crawled_pages)
[9]
           enqueue(url_queue, u)
[10]
     reorder_queue(url_queue)
[11]
```

Function description:

enqueue(queue, element):	append element at the end of queue
dequeue(queue) :	remove the element at the beginning
	of queue and return it
reorder_queue(queue) :	reorder queue using information in
	links (refer to Figure 2)

(1) breadth first
 do nothing (null operation)

(2) backlink count, IB'(p)
foreach u in url_queue
backlink_count[u] = number of terms (-,u) in links
sort url_queue by backlink_count[u]

(3) PageRank IR'(p)solve the following set of equations: $IR[u] = (1 - 0.9) + 0.9 \sum_{i} \frac{IR[v_i]}{c_i}$, where $(v_i, u) \in \texttt{links}$ and c_i is the number of links in the page v_i sort url_queue by IR(u)





Observations

Using an ordering metric outperforms random (or breadth) ordering
The ordering metric IR'(p) outperforms the IB'(p) one, even when the importance metric is IB(p)

IR'(p) vs IB'(p)



Algorithm 5.2 Crawling algorithm (modified similarity based) Input: starting_url: seed URL Procedure:

```
[1] enqueue(url_queue, starting_url)
[2] while (not empty(hot_queue) and not empty(url_queue))
[3]
     url = dequeue2(hot_queue, url_queue)
[4]
     page = crawl_page(url)
[5]
     enqueue(crawled_pages, (url, page))
[6]
     url_list = extract_urls(page)
[7]
     foreach u in url_list
[8]
        enqueue(links, (url, u))
[9]
        if (u \notin url_queue and u \notin hot_queue and (u,-) \notin crawled_pages)
           if (u contains computer in anchor or url)
[10]
[11]
             enqueue(hot_queue, u)
[12]
          else
[13]
             enqueue(url_queue, u)
[14]
     reorder_queue(url_queue)
[15]
     reorder_queue(hot_queue)
```

Function description:



```
Algorithm 5.3
                     Crawling algorithm (similarity based)
Input: starting_url: seed URL
Procedure:
  [1] enqueue(url_queue, starting_url)
  [2] while (not empty(hot_queue) and not empty(url_queue))
  [3]
        url = dequeue2(hot_queue, url_queue)
  [4]
       page = crawl_page(url)
  [5]
        if (page contains 10 or more computer in body
             or one computer in title)
  [6]
          hot[url] = True
        enqueue(crawled_pages, (url, page))
  [7]
  [8]
        url_list = extract_urls(page)
  [9]
        foreach u in url_list
  [10]
          enqueue(links, (url, u))
  [11]
          if (u \notin url_queue and u \notin hot_queue and (u,-) \notin crawled_pages)
  [12]
             if (u contains computer in anchor or url)
  [13]
                enqueue(hot_queue, u)
             else if (distance_from_hotpage(u) < 3)</pre>
  [14]
  [15]
                enqueue(hot_queue, u)
  [16]
             else
  [17]
                enqueue(url_queue, u)
        reorder_queue(url_queue)
  [18]
        reorder_queue(hot_queue)
  [19]
```



Observations

 When similarity is important, it is effective to use an ordering metric that considers 1) the contents of anchors and URLs and 2) the distance to the hot pages that have been discovered

Conclusion

 With a good ordering strategy, we can build crawlers that can obtain a significant portion of the hot pages relatively early

Limitations

- Experiments run only over a portion of the Stanford pages
- Is the definition of crawler performance adequate?
- What about other ordering metrics or combinations of metrics
- What about other types of crawlers e.g. focused crawlers

References

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