

# Similarity Flooding: A versatile Graph Matching Algorithm and its Application to Schema Matching

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## Introduction to Schema Matching Problem

- The Goal is to map certain schema elements to other schema
- Applications:
  - Schema Integration / mediated schemas
  - Translate data between multiple databases
  - Databases Consolidation

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## Why schema matching is difficult?

- Imprecise wording, e.g. contact-info
- Different Ontology, e.g. 'Load' in electrical and mechanical contexts
- Schema and data maybe insufficient.
- Documentations and original schema designers are usually no available
- Matching decisions are highly subjective

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## Approaches

- Learning-based Approach
- Rules-based Approach
- Information Retrieval Approach

Most solutions requires user intervention either at the beginning (in case of learning) or after creating a mapping (correcting)

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## Similarity Flooding

- Based on Schema elements matching. (vs. schema + instances matching)
- Based on schema structure (vs. elements level matching)
- Each Schema is represented as a directed graph
- Based on the assumption :
  - Whenever any two elements in the graphs  $G_1$  and  $G_2$  are similar, their neighbors tend to be similar.

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## Similarity Flooding

```
CREATE TABLE Personnel(  
  □ Pno int,  
  □ Pname string,  
  □ Dept string,  
  □ Born date,  
  □ UNIQUE pkey(Pno)
```

Schema 1

```
CREATE TABLE Employee (  
  □ EmpNo int PRIMARY KEY,  
  □ EmpName varchar(50),  
  □ DeptNo int REFERENCES  
  □ Department,  
  □ Salary dec(15,2),  
  □ Birthdate date  
  □ )
```

```
CREATE TABLE Department (  
  □ DeptNo int PRIMARY KEY,  
  □ DeptName
```

Schema 2

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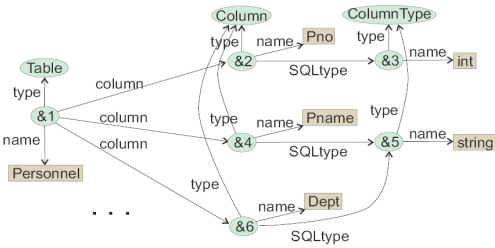
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## Equivalent Graph Representation




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## SF algorithm

1- Initial Mapping is made by simple string matching between graphs nodes.

Line#	Similarity	Node in $G_1$	Node in $G_2$
1.	1.0	Column	Column
2.	0.66	ColumnType	Column
3.	0.66	'Dept'	'DeptNo'
4.	0.66	'Dept'	'DeptName'
5.	0.5	UniqueKey	PrimaryKey
6.	0.26	'Pname'	'DeptName'
7.	0.26	'Pname'	'EmpName'
8.	0.22	'date'	'Birthdate'
9.	0.11	'Dept'	'Department'
10.	0.06	'int'	'Department'

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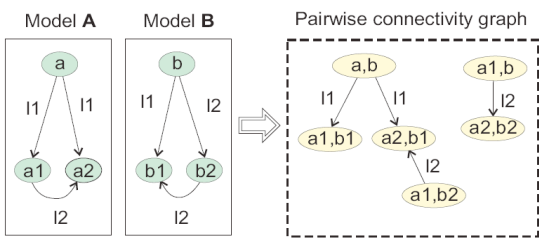
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## SF algorithm (cont'd)

2- Similarity Flooding




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## Assumptions

- Each edge type has the same contribution = 1.0
- Similarity contribution for edges with the same type outgoing from one node is evenly distributed

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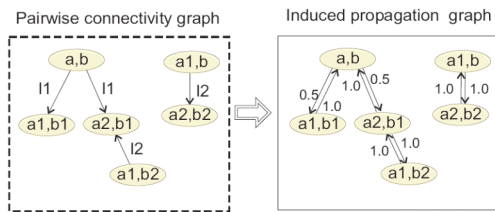
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## SF algorithm (cont'd)




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## SF algorithm (cont'd)

Propagation Rule :

$$\sigma^{i+1}(x, y) = \sigma^i(x, y) + \sum_{(a_u, p, x) \in A, (b_u, p, y) \in B} \sigma^i(a_u, b_u) \cdot w((a_u, b_u), (x, y)) + \sum_{(x, p, a_v) \in A, (y, p, b_v) \in B} \sigma^i(a_v, b_v) \cdot w((a_v, b_v), (x, y))$$

- Normalize similarity values after each iteration
- Stopping condition :  $\Delta(\sigma^i, \sigma^{i-1}) < \epsilon$
- Convergence can be guaranteed when all initial similarities for all pairs  $> 0$

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## SF algorithm (cont'd)

### 3- Filters

The goal of filtering is to choose the best match candidates from the output list

- ☞ Application-Specific Constraints filter : e.g. cardinality constraint
- ☞ Selection Metrics : e.g. stable marriage, maximal matching, assignment problem
- ☞ Selection Threshold ( $0 < t_{rel} \leq 1$ )

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## Matching Quality

- Accuracy : how much effort is needed to convert the output matching pairs to the intended one, i.e. the number of removing false positives and adding false negatives

$$Accuracy = 1 - \frac{(n - c) + (m - c)}{m}$$

n : number of returned matches

m : number of intended matches

c : number of returned intended matches

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## Intended Results Specification

Intended match result is categorized into 3 types:

- ☐ Sparse
- ☐ Expected
- ☐ Verbose

Sparse	Expected	Verbose	Node in $G_1$	Node in $G_2$
	+	+	[Table: Personnel]	[Table: Employee]
		+	[Table: Personnel]	[Table: Department]
	+	+	[UniqueKey: perskey]	[PrimaryKey: on EmpNo]
+	+	+	[Col: Personnel/Dept]	[Col: Department/DeptName]
		+	[Col: Personnel/Dept]	[Col: Department/DeptNo]
		+	[Col: Personnel/Dept]	[Col: Employee/DeptNo]
+	+	+	[Col: Personnel/Pno]	[Col: Employee/EmpNo]
+	+	+	[Col: Personnel/Pname]	[Col: Employee/EmpName]
+	+	+	[Col: Personnel/Born]	[Col: Employee/Birthdate]

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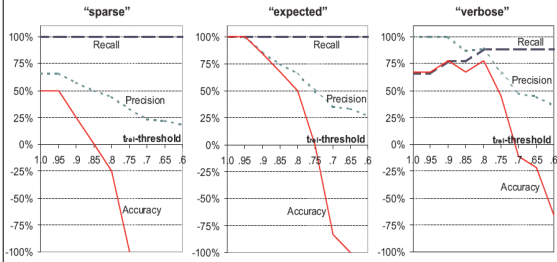
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## Results (cont'd)




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## Convergence Speed

Identifier	Fixpoint formula
Basic	$\sigma^{i+1} = \text{normalize}(\sigma^i + \varphi(\sigma^i))$
A	$\sigma^{i+1} = \text{normalize}(\sigma^0 + \varphi(\sigma^i))$
B	$\sigma^{i+1} = \text{normalize}(\varphi(\sigma^0 + \sigma^i))$
C	$\sigma^{i+1} = \text{normalize}(\sigma^0 + \sigma^i + \varphi(\sigma^0 + \sigma^i))$

Formula	$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$T_6$	$T_7$	$T_8$	$T_9$	Total
A (as is)	18	48	122	78	$\infty$	12	37	25	25	$\infty$
A (strongly connected)	15	56	89	81	1488	18	48	25	31	1851
B (as is)	8	428	17	39	8	13	10	24	21	568
B (strongly connected)	7	268	21	32	13	15	14	21	53	444
C (as is)	7	9	9	11	7	7	9	10	9	78
C (strongly connected)	7	9	8	11	7	5	9	7	9	72

Strongly connected :  $\sigma^0(x,y) > 0$ , for all  $x \in G_1, y \in G_2$

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## Pros

- Innovative method for quality matching.
- General Model, provided that there is a straightforward method to map the used schema to a graph.
- No learning phase is required before use
- Flexibility in filters to suits specific application constrains

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### Cons

- Weak basis for similarity propagation
- Estimation errors can also be propagated to neighbors
- Flooding techniques are usually slow. Not practical for large number of elements
- Initial similarities have huge impact on the output quality and the convergence speed. Which returns us to the first square : how to get good matching?
- Heterogeneous sources can be problematic when mapped to graphs
- Does not utilize data instances in the graph
- Unable to detect complex relations between elements

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### Recommendations

- Extension to N graph
- Adding more powerful more enhanced matcher to initialize similarities
- Adding sample of data to the graph and use instance matchers (e.g. format matcher) as initial similarity
- More useful in case of hierarchical schemas (e.g. XML)

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### Recommendations (cont'd)

- Restrict the graph to a hierarchical structure
  - The similarity will propagate in bottom-up and top-down fashions
  - Similarity propagation is much reasonable in case of parent/child relations
  - The performance and the convergence speed is improved due to limited propagation paths
  - Can fit most practical schemas, e.g., SQL /XML.

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## Open Questions

- How to improve convergence speed?
- Can a directed (ordered) flooding affect the convergence or matching quality?
- How to extend the model to N graph?
- Will using sample instances increase the matching quality? to what extent? at what cost?

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## Conclusion

- Similarity Flooding is a structural based approach that fits various data sources
- Generality is chosen over performance
- Useful for specific contexts (e.g. no complex relations)
- Further investigation is required to improve matching quality and convergence speed

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