

# Optimal aggregation algorithms for middleware

CS856 Fall 2005 Presentation

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## About the paper

- Ronald Fagin, IBM Research
- Amnon Lotem, Maryland
- Moni Naor, Weizmann, Israel
  
- In *ACM Symp. Principles of Database Systems, 2001.*

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## Agenda

- Background
- Fagin's algorithm
- Threshold algorithm
- $\theta$ -approximation
- NRA algorithm
- Combined algorithm

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## Agenda

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### Background

-  Fagin's algorithm
  -  Threshold algorithm
  -   $\theta$ -approximation
  -  NRA algorithm
  -  Combined algorithm
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## Motivation: top- $k$ queries

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## Motivation: top- $k$ queries

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- Multimedia DB: "find 10 pictures that are funny and large in size"
  - Info. retrieval: "find 100 papers that are most relevant to my research areas"
  - Data stream: "find 5 users with the largest bandwidth usage"
  - Live examples:
    - QBIC: [www.qbic.almaden.ibm.com](http://www.qbic.almaden.ibm.com)
    - Flickr: [www.flickr.com](http://www.flickr.com)
    - WinFS for Windows Vista
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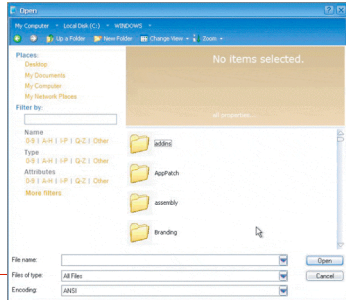
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## WinFS for Windows Vista




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## Data model

	$x_1$	$x_2$	$x_3$		$L_1$	$L_2$	$L_3$			
A	.1	.2	.7	→	E	.8	C	.9	D	.9
B	.5	.3	.3		B	.5	E	.4	A	.7
C	.1	.9	.4		D	.3	B	.3	E	.6
D	.3	.1	.9		C	.1	A	.2	C	.4
E	.8	.4	.6		A	.1	D	.1	B	.3

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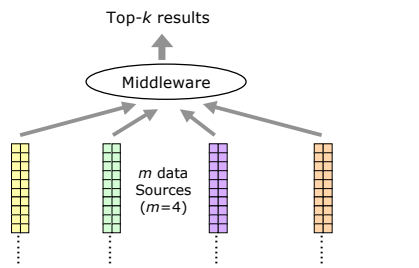
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## Data integration sys. (middleware)




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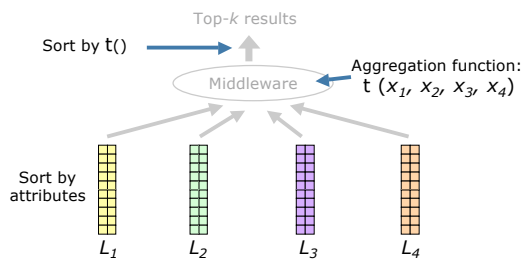
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## Problem definition



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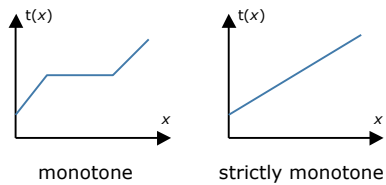
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## Aggregation functions

□ Can be  $\max()$ ,  $\min()$ ,  $\text{avg}()$ , ...



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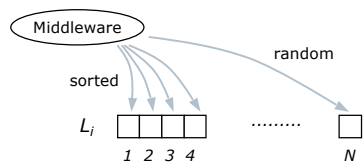
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## Sorted and random access



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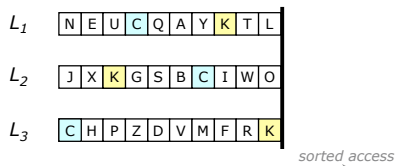
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## Fagin's algorithm (FA)

suppose  $m=3$ ,  $k=2$ ; objects are A, B, C, ..., Z



Stop when there are  $k$  objects, such that each of them has been seen in *each* list.

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## Fagin's algorithm (FA): step 2

- For each object  $R$  has been seen:
  - Do *random access* to get all of its attributes.
  - Calculate  $t(R)$ .
- Sort all these objects and output the first  $k$  objects.

FA is *correct*, but not always *optimal*.

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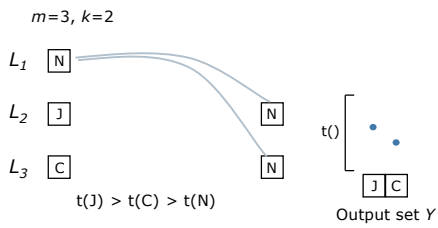
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## Threshold algorithm (TA)



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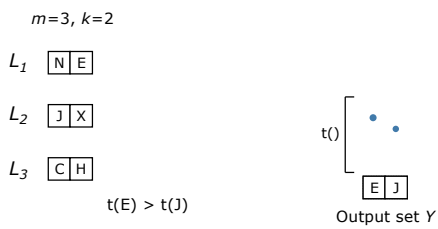
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## Threshold algorithm (TA)



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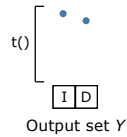
## Threshold algorithm (TA)

$m=3, k=2$

$L_1$  [N][E][U][C][Q][A][Y][K][T]

$L_2$  [J][X][K][G][S][B][C][I][W]

$L_3$  [C][H][P][Z][D][V][M][F][R]




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## Threshold algorithm (TA)

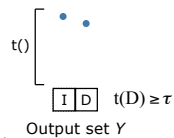
$m=3, k=2$

$L_1$  [N][E][U][C][Q][A][Y][K][T]  $x_1$

$L_2$  [J][X][K][G][S][B][C][I][W]  $x_2$

$L_3$  [C][H][P][Z][D][V][M][F][R]  $x_3$

Threshold value  
 $\tau = t(x_1, x_2, x_3)$



Stop when the grade of the last object in  $Y$  is equal or larger than the threshold value.

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## Middleware cost

- In this paper, we use middleware cost to measure optimality of an algorithm.
- To answer a query on database  $D$ , an algorithm  $A$  needs:
  - $s$  sorted accesses
  - $r$  random accesses
- The **middleware cost** of  $A$  on  $D$  is:

$$\text{cost}(A, D) = sC_S + rC_R$$

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### Instance optimality

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- A set of databases: **D**
- A set of (middleware) algorithms: **A**
- $B \in \mathbf{A}$  is instance optimality if:

$$\text{cost}(B,D) \leq c \cdot \text{cost}(A,D) + c'$$

for every  $A \in \mathbf{A}$  and  $D \in \mathbf{D}$

$c$ : optimality ratio

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### Instance optimality of TA

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- Assumptions
    - $t()$ : monotone
    - **D**: all
    - **A**: no wild guess
  - Optimality: TA is instance optimal, with optimality ratio  $m+m(m-1)c_R/c_S$
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### Instance optimality of TA (2)

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- Assumptions
    - $t()$ : strictly monotone
    - **D**: unique
    - **A**: all
  - Optimality: TA is instance optimal
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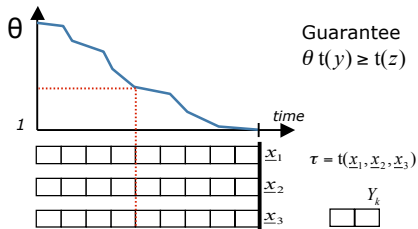
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## $\theta$ -approximation

$Y_k$  is the last object in  $Y$ ,  $\theta = \tau / t(Y_k)$



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## Instance optimality of $\theta$ -approximation

- Assumptions
  - $t(\cdot)$ : monotone
  - $D$ : all
  - $A$ : no wild guess
  - $\theta > 1$
- Optimality:  $\theta$ -approximation is instance optimal

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# Agenda

- 📁 Background
- 📄 Fagin's algorithm
- 📄 Threshold algorithm
- 📄  $\theta$ -approximation
- 📄 **No-Random-Access algorithm**
- ⌛ Combined algorithm

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# Lower/upper bound of an object

- Define Lower bound LB() as the value of  $t()$  when setting all unknown attributes to 0
- Define Upper bound UB() as the value of  $t()$  when setting all unknown attributes to  $\underline{x}_i$

$L_1$  [ ] [ ] [ ] [ ] [ ] [ ] 0.5  
 $L_2$  [ ] [ ] [ ] [ ] [ ] [ ] 0.2  
 $L_3$  [ ] [ ] [A] [ ] [ ] [ ] 0.3  
0.8

$LB(A) = t(0, 0, 0, 0.8)$

$UB(A) = t(0.5, 0.2, 0, 0.8)$

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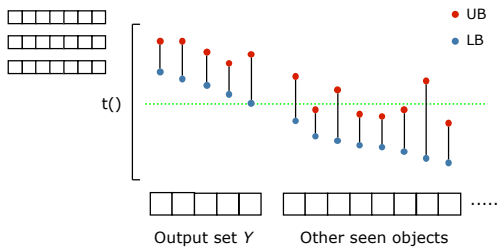


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# NRA algorithm




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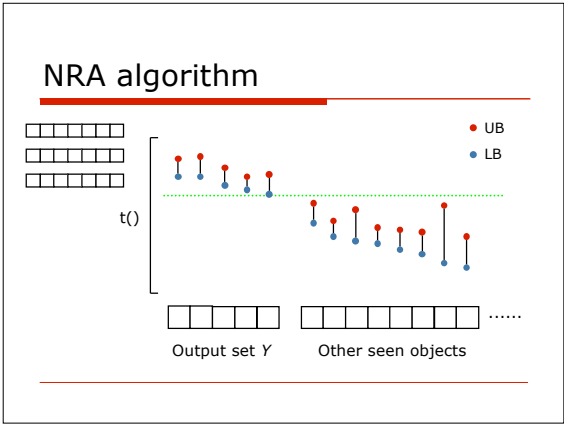
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- ### Instance optimality of NRA
- Assumptions
    - $t()$ : monotone
    - $\mathbf{D}$ : all
    - $\mathbf{A}$ : no random access
  - Optimality: NRA is instance optimal

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- ### Agenda
- 📁 Background
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  - 📖 Threshold algorithm
  - 📖  $\theta$ -approximation
  - 📖 NRA algorithm
  - 📖 **Combined algorithm**

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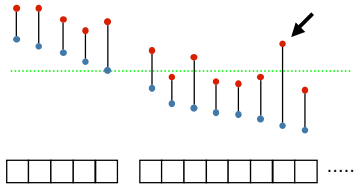
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## Combined algorithm (CA)



For every  $\lfloor c_R / c_S \rfloor$  step, obtain all unknown attr. of the object with the largest UB.

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## Instance optimality of CA

- Assumptions
  - $t()$ : strictly monotone in each argument
  - **D**: unique
  - **A**: all
- Optimality: CA is instance optimal

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## Agenda

- 📁 Background
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- 📖 Combined algorithm

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## Conclusion

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- TA is instance optimal in most cases
  - $\theta$ -approx: early stop
  - NRA: random access is not allowed
  - CA: random access is costly
  - Future work
    - *Tightly* instance optimal
    - More efficient structure of NRA
    - Compare CA vs. TA
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## Discussion

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- Object caching of TA
  - Grades output of NRA
  - Other metrics for algorithm optimality
  - Assumptions on databases
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## Backup slides

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### I.O. in other fields

- ❑ Competitive analysis
- ❑ Approximation algorithms
- ❑ The mean of Monte Carlo estimation (Dagum et al.)
- ❑ Operations on sorted sets (Demaine et al.)

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### Memory overhead

- ❑ FA: need to remember  $t()$  for all objects that have been seen.
- ❑ TA: only need to remember  $t()$  for objects in  $Y$ .
- ❑ NRA: similar to FA.

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### Wild guess example

$L_1$	$L_2$
(1, 1)	$(2n + 1, 1)$
(2, 1)	$(2n, 1)$
(3, 1)	$(2n - 1, 1)$
...	...
$(n + 1, 1)$	$(n + 1, 1)$
$(n + 2, 0)$	$(n, 0)$
$(n + 3, 0)$	$(n - 1, 0)$
...	...
$(2n + 1, 0)$	$(1, 0)$

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### Instance Optimality w/ wild guess

Assumptions

- $t()$ :  $\min(x_1, x_2)$
- **D**: all
- **A**: all

Optimality: no algorithms is instance optimal

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### Wild guess example w/ $\theta$

$L_1$	$L_2$
$(1, \cdot)$	$(2n + 1, \cdot)$
$(2, \cdot)$	$(2n, \cdot)$
...	...
$(n + 1, \frac{1}{\theta})$	$(n + 1, \frac{1}{\theta})$
$(n + 2, \frac{1}{2\theta^2})$	$(n, \frac{1}{2\theta^2})$
$(n + 3, \cdot)$	$(n - 1, \cdot)$
...	...
$(2n + 1, \cdot)$	$(1, \cdot)$

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### Instance Optimality w/ wild guess

Assumptions

- $t()$ :  $\min(x_1, x_2)$
- **D**: unique
- **A**: all
- $\theta > 1$

Optimality: no algorithms is instance optimal

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### Costs for top-1 & top-2 for NRA

$L_1$	$L_2$
$(R, 1)$	$(\cdot, \frac{1}{3})$
$(\cdot, \frac{1}{3})$	...
...	$(\cdot, \frac{1}{3})$
$(\cdot, \frac{1}{3})$	$(R, 0)$

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### Instance optimality of CA (2)

- Assumptions
  - $t()$ :  $\min()$
  - **D**: unique
  - **A**: all
- Optimality: CA is instance optimal

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### Instance optimality dependency

- Assumptions
  - $t()$ :  $\min()$
  - **D**: all
  - **A**: no wild guess
- Optimality: no algorithm has instance optimality independent of  $C_R/C_S$

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