

Periodicity, Morphisms, and Matrices

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Periodicity

Periodicity is an important property of words, with applications to

- ▶ string searching algorithms (e.g., Knuth-Morris-Pratt)
- ▶ formal languages (e.g., pumping lemmas)
- ▶ combinatorics on words (e.g., theorems of Thue, Lyndon-Schützenberger)

Periodicity

We say a sequence $(f_n)_{n \geq 0}$ is *periodic with period length* $h \geq 1$ if $f_n = f_{n+h}$ for all $n \geq 0$.

The following is a classical “folk theorem”:

Theorem. If $(f_n)_{n \geq 0}$ is a sequence that is periodic with period lengths h and k , then it is periodic with period length $\gcd(h, k)$.

Proof. By the extended Euclidean algorithm, there exist integers $r, s \geq 0$ such that $rh - sk = \gcd(h, k)$.

Then we have

$$f_n = f_{n+rh} = f_{n+rh-sk} = f_{n+\gcd(h,k)}$$

for all $n \geq 0$.

The 1965 Theorem of Fine-Wilf

Theorem. Let $(f_n)_{n \geq 0}$, $(g_n)_{n \geq 0}$ be two periodic sequences, of period lengths h and k respectively.

- (a) If $f_n = g_n$ for $0 \leq n < h + k - \gcd(h, k)$, then $f_n = g_n$ for all $n \geq 0$.
- (b) The conclusion in (a) would be false if $h + k - \gcd(h, k)$ were replaced by any smaller number.

Proof of the Fine-Wilf Theorem

(a) If $f_n = g_n$ for $0 \leq n < h + k - \gcd(h, k)$, then $f_n = g_n$ for all $n \geq 0$.

Proof of (a).

For the moment assume $\gcd(h, k) = 1$.

The proof is easy when $h = k = 1$, so assume WOLOG $h > k$.

Then we have

$$f_i = g_i = g_{i+k} = f_{i+k} = f_{(i+k) \bmod h}$$

for $0 \leq i < h - 1$.

Proof of the Fine-Wilf Theorem

We just proved that

$$f_i = f_{(i+k) \bmod h}$$

for $0 \leq i < h - 1$.

Start with f_{k-1} and apply this relation $h - 1$ times.

We get

$$f_{k-1} = f_{2k-1} = \cdots = f_{(h-1)k-1} = f_{hk-1},$$

where the indices are taken (mod h).

Since

$$\gcd(h, k) = 1,$$

it follows that all h indices (mod h) are represented in this equation.

Hence $f_i = f_0$ for all i , and the same result holds for g_i .

Proof of the Fine-Wilf Theorem

Now let us remove the restriction $\gcd(h, k) = 1$.

If $\gcd(h, k) = d$, group the symbols of f and g into groups of d symbols; call the result f' and g' .

If f and g agree on the first $h + k - \gcd(h, k)$ symbols, then f' and g' agree on the first $\frac{h}{d} + \frac{k}{d} - 1$ symbols.

Furthermore f' is periodic of period h/d and g' is periodic of period k/d .

From the results above $f' = g'$ and so $f = g$.

Proof of the Fine-Wilf Theorem

(b) The conclusion in (a) would be false if $h + k - \gcd(h, k)$ were replaced by any smaller number.

Proof of (b). Define strings $\sigma(h, k)$ as follows:

$$\sigma(h, k) = \begin{cases} 0, & \text{if } h = 0; \\ 0^{k-1}1, & \text{if } h \mid k; \\ \sigma(r, h)^q \sigma(r', r), & \text{if } h > 1 \text{ and} \\ & k = qh + r, \\ & h = q'r + r'. \end{cases}$$

Proof of the Fine-Wilf Theorem

Then it can be shown that if we construct periodic sequences f , g such that

- ▶ f is of period length k and has period $\sigma(h, k)$
- ▶ g is of period length h and has period $\sigma(k, h)$

then f and g agree on a prefix of a length

$$h + k - \gcd(h, k) - 1,$$

but disagree at the $h + k - \gcd(h, k)$ 'th term.

The Fine-Wilf Theorem

Remark. The maximal counter-examples in part (b) are the so-called **standard Sturmian words**, and play a role in the Knuth-Morris-Pratt string-matching algorithm.

For example, if $h = 5$ and $k = 8$ the maximal counter-examples are

f = 1011010110101101011010110...

g = 101101011011010110110101...

Variations on Fine-Wilf

Theorem. Let $\mathbf{f} = (f_n)_{n \geq 0}$, $\mathbf{g} = (g_n)_{n \geq 0}$ be two periodic sequences of real numbers, of period lengths h and k , respectively, such that

$$\sum_{0 \leq i < h} f_i \geq 0 \quad (1)$$

and

$$\sum_{0 \leq j < k} g_j \leq 0. \quad (2)$$

Let $d = \gcd(h, k)$.

(a) If

$$f_n \leq g_n \quad \text{for } 0 \leq n < h + k - d \quad (3)$$

then

- (i) $f_n = g_n$ for all $n \geq 0$; and
- (ii) $\sum_{j \leq i < j+d} f_i = \sum_{j \leq i < j+d} g_i = 0$ for all integers $j \geq 0$.

(b) The conclusion (a)(i) would be false if in the hypothesis $h + k - d$ were replaced by any smaller integer.

Sketch of Proof, Part (a)(i)

Define

$$P(z) = 1 + z + \cdots + z^{h-1} = (z^h - 1)/(z - 1);$$

$$Q(z) = 1 + z + \cdots + z^{k-1} = (z^k - 1)/(z - 1);$$

$$R(z) = (z^k - 1)/(z^d - 1); \quad d = \gcd(h, k)$$

$$S(z) = (z^h - 1)/(z^d - 1).$$

By hypothesis $P \circ \mathbf{f} \geq 0$, where by \circ we mean take the dot product of the coefficients of P to consecutive windows of \mathbf{f} . Then $R \circ (P \circ \mathbf{f}) \geq 0$. But then $RP \circ \mathbf{f} \geq 0$.

$$\begin{aligned}
P(z) &= 1 + z + \cdots + z^{h-1} = (z^h - 1)/(z - 1); \\
Q(z) &= 1 + z + \cdots + z^{k-1} = (z^k - 1)/(z - 1); \\
R(z) &= (z^k - 1)/(z^d - 1); \quad d = \gcd(h, k) \\
S(z) &= (z^h - 1)/(z^d - 1).
\end{aligned}$$

Similarly, by hypothesis $Q \circ (-\mathbf{g}) \geq 0$. Then $SQ \circ (-\mathbf{g}) \geq 0$. But $RP = SQ$, so

$$\sum_{0 \leq i < h+k-d} e_i (f_i - g_i) \geq 0. \tag{4}$$

where $R(z)P(z) = \sum_{0 \leq i < h+k-d} e_i z^i$.

Now if the coefficients e_i were strictly positive, then since $f_n \leq g_n$ for $0 \leq n < h + k - d$, we get $f_n = g_n$ for $0 \leq n < h + k - d$. By the Fine-Wilf theorem, $f_n = g_n$ for $n \geq 0$.

Positivity of the e_i

It remains to show that the coefficients e_i are positive.

To see this, note that

$$\begin{aligned}R(z)P(z) &= \frac{z^h - 1}{z^d - 1} \cdot \frac{z^k - 1}{z - 1} \\ &= (1 + z^d + z^{2d} + \cdots + z^{h-d})(1 + z + z^2 + \cdots + z^{k-1}).\end{aligned}$$

If $i < h$, write $i = qd + r$ where $0 \leq r < d$, and choose the term z^{qd} from the left factor and z^r from the right factor to see $e_i > 0$.

If $h \leq i < h + k - d$, choose z^{h-d} from the left factor and z^{i-h+d} from the right factor to see $e_i > 0$.

Maximal Counter-Examples

The maximal counter-examples in part (b) of the theorem turn out to be just the first differences of the maximal counter-examples to Fine-Wilf.

For example, for $h = 5$, $k = 8$ we have

n	0	1	2	3	4	5	6	7	8	9	10	11	12
f_n	-1	1	-1	0	1	-1	1	-1	0	1	-1	1	-1
g_n	0	1	-1	0	1	-1	1	-1	0	1	-1	0	1

Generalization to more than two periods

For integers $p \geq 1$ let ω_p denote a primitive p 'th root of unity, i.e.,
 $\omega_p := e^{2\pi\sqrt{-1}/p}$.

Define

$$R_p := \{\omega_p^i : 0 \leq i < p\} = \{\omega \in \mathbb{C} : \omega^p = 1\}.$$

Finally, for integers $h_1, h_2, \dots, h_r \geq 1$ define

$$\gamma(h_1, h_2, \dots, h_r) = |R_{h_1} \cup R_{h_2} \cup \dots \cup R_{h_r}|,$$

the number of distinct roots of unity among the h_1 'th, h_2 'th, etc., roots of unity.

By the principle of inclusion-exclusion, it follows that

$$\gamma(h_1, h_2, \dots, h_r) = \sum_{\substack{S \subseteq \{h_1, h_2, \dots, h_r\} \\ S \neq \emptyset}} \gcd(S) (-1)^{|S|+1},$$

where by $\gcd(S)$ for S a nonempty set we mean the greatest common divisor of all elements of S .

For example,

$$\gamma(6, 10, 15)$$

$$\begin{aligned} 6 + 10 + 15 - \gcd(6, 10) - \gcd(6, 15) - \gcd(10, 15) + \gcd(6, 10, 15) \\ = 22. \end{aligned}$$

Theorem

Let $(f_i(n))_{n \geq 0}$, $1 \leq i \leq r$, be r periodic complex-valued sequences with period lengths h_1, h_2, \dots, h_r , respectively. Suppose

$\sum_{1 \leq i \leq r} f_i(n) = 0$ for $0 \leq n < \gamma(h_1, h_2, \dots, h_r)$. Then

$\sum_{1 \leq i \leq r} f_i(n) = 0$ for all $n \geq 0$.

Proof. Any periodic complex-valued sequence $(f(n))_{n \geq 0}$ of period length p can be written in the form

$$f(n) = \sum_{0 \leq i < p} c_i \omega_p^{in}$$

for some coefficients c_0, c_1, \dots, c_{p-1} .

It follows that there exist coefficients $c_{i,j}$, $1 \leq i \leq r$ and $0 \leq j < h_i$ such that

$$f_i(n) = \sum_{0 \leq j < h_i} c_{i,j} \omega_{h_i}^{jn}.$$

Define

$$\begin{aligned} s &= [s_1, s_2, \dots, s_m] \\ &= [1, \omega_{h_1}, \omega_{h_1}^2, \dots, \omega_{h_1}^{h_1-1}, 1, \omega_{h_2}, \omega_{h_2}^2, \dots, \omega_{h_2}^{h_2-1}, \dots, \\ &\quad 1, \omega_{h_r}, \omega_{h_r}^2, \dots, \omega_{h_r}^{h_r-1}] \end{aligned}$$

where $m = h_1 + h_2 + \dots + h_r$.

Let $B := \gamma(h_1, h_2, \dots, h_r)$ and define the $B \times m$ matrix

$M = (t_{i,j})_{0 \leq i < B, 1 \leq j \leq m}$ by $t_{i,j} := s_j^i$.

Define the column vector

$$v := [c_{1,0}, c_{1,1}, \dots, c_{1,h_1-1}, c_{2,0}, c_{2,1}, \dots, c_{2,h_2-1}, \dots, \\ c_{r,0}, c_{r,1}, \dots, c_{r,h_r-1}]^T.$$

Then the hypothesis of the theorem is $Mv = 0$.

Some of the columns of M are identical because some of the entries in the vector s coincide.

We may delete the repeated columns of M and sum the corresponding entries of v to get $M'v' = 0$, where M' is a $B \times B$ matrix and v' is a column vector with B entries.

Now M' is a Vandermonde matrix and hence invertible, so $v' = 0$.

It follows that $\sum_{1 \leq i \leq r} f_i(n) = 0$ for all n .

Open Problem

Find a way to compute

$$\gamma(h_1, h_2, \dots, h_r) = |R_{h_1} \cup R_{h_2} \cup \dots \cup R_{h_r}|,$$

efficiently (in polynomial time in $\log h_1, \log h_2, \dots, \log h_r, r$).

The Matrix Associated with a Morphism

Given a morphism $\varphi : \Sigma^* \rightarrow \Sigma^*$ for some finite set $\Sigma = \{a_1, a_2, \dots, a_d\}$, we define the *incidence matrix* $M = M(\varphi)$ as follows:

$$M = (m_{i,j})_{1 \leq i,j \leq d}$$

where $m_{i,j}$ is the number of occurrences of a_i in $\varphi(a_j)$, i.e., $m_{i,j} = |\varphi(a_j)|_{a_i}$.

Example. Consider the morphism φ defined by

$$\varphi : a \rightarrow ab$$

$$b \rightarrow cc$$

$$c \rightarrow bb.$$

Then

$$M(\varphi) = \begin{pmatrix} 1 & 0 & 0 \\ 1 & 0 & 2 \\ 0 & 2 & 0 \end{pmatrix}$$

The Matrix Associated with a Morphism

The matrix $M(\varphi)$ is useful because of the following proposition.

Proposition. We have

$$\begin{bmatrix} |\varphi(w)|_{a_1} \\ |\varphi(w)|_{a_2} \\ \vdots \\ |\varphi(w)|_{a_d} \end{bmatrix} = M(\varphi) \begin{bmatrix} |w|_{a_1} \\ |w|_{a_2} \\ \vdots \\ |w|_{a_d} \end{bmatrix}.$$

Proof. We have

$$|\varphi(w)|_{a_i} = \sum_{1 \leq j \leq d} |\varphi(a_j)|_{a_i} |w|_{a_j}.$$

The Matrix Associated with a Morphism

Corollary.

$$\begin{bmatrix} |\varphi^n(w)|_{a_1} \\ |\varphi^n(w)|_{a_2} \\ \vdots \\ |\varphi^n(w)|_{a_d} \end{bmatrix} = (M(\varphi))^n \begin{bmatrix} |w|_{a_1} \\ |w|_{a_2} \\ \vdots \\ |w|_{a_d} \end{bmatrix}$$

The Matrix Associated with a Morphism

Hence we find

Corollary.

$$|\varphi^n(w)| = \begin{bmatrix} 1 & 1 & 1 & \cdots & 1 \end{bmatrix} M(\varphi)^n \begin{bmatrix} |W|_{a_1} \\ |W|_{a_2} \\ \vdots \\ |W|_{a_d} \end{bmatrix} .$$

The Length Sequence of an Iterated Morphism

We can now ask questions about the sequence of lengths

$$|x|, |h(x)|, |h^2(x)|, \dots$$

These questions were very popular in mathematical biology (L-systems) in the 1980's.

For example, here is a classical result:

Theorem. Suppose $h : \Sigma^* \rightarrow \Sigma^*$ is a morphism, and suppose there exist a word $w \in \Sigma^*$ and a constant c such that

$$|w| = |h(w)| = \dots = |h^n(w)| = c,$$

where $n = |\Sigma|$.

Then $c = |h^i(w)|$ for all $i \geq 0$.

Proof of the Theorem

It suffices to show $|h^{n+1}(w)| = c$, because then the theorem follows by induction on n .

Let M be the incidence matrix of h . By the Cayley-Hamilton theorem,

$$M^n = c_0 M^0 + \cdots + c_{n-1} M^{n-1}$$

for some constants c_0, c_1, \dots, c_{n-1} .

Define $f_i = |h^i(w)|$ and let

$$v = [|w|_{a_1} \ |w|_{a_2} \ \cdots \ |w|_{a_n}]^T.$$

Then for $0 \leq i < n$ we have

$$\begin{aligned} f_{i+1} - f_i &= [1 \ 1 \ \cdots \ 1](M^{i+1} - M^i)v \\ &= [1 \ 1 \ \cdots \ 1]M^i(M - I)v \\ &= [1 \ 1 \ \cdots \ 1]M^i v' \\ &= 0, \end{aligned}$$

where $v' := (M - I)v$.

Now

$$\begin{aligned}f_{n+1} - f_n &= [1 \ 1 \ \cdots \ 1]M^n v' \\&= [1 \ 1 \ \cdots \ 1](c_0 + \cdots + c_{n-1}M^{n-1})v' \\&= \sum_{0 \leq i < n} c_i [1 \ 1 \ \cdots \ 1]M^i v' \\&= 0,\end{aligned}$$

since each summand is 0.

Hence $f_{n+1} = f_n$.

Another Question

We might also ask, how long can the sequence of lengths

$$|x|, |h(x)|, |h^2(x)|, \dots$$

strictly decrease?

This question arose naturally in a paper on the two-sided infinite fixed points of morphisms, i.e., those two-sided infinite words \mathbf{w} such that $h(\mathbf{w}) = \mathbf{w}$.

The Length Sequence of an Iterated Morphism

If Σ has n elements, we can easily find a decreasing sequence of length n . For example, for $n = 5$, define h as follows:

$$h(a) = b$$

$$h(b) = c$$

$$h(c) = d$$

$$h(d) = e$$

$$h(e) = \epsilon$$

Then we have

$$h(abcde) = bcde$$

$$h^2(abcde) = cde$$

$$h^3(abcde) = de$$

$$h^4(abcde) = e$$

$$h^5(abcde) = \epsilon$$

So

$$\begin{aligned} |abcde| &> |h(abcde)| > |h^2(abcde)| > |h^3(abcde)| \\ &> |h^4(abcde)| > |h^5(abcde)| = 0. \end{aligned}$$

The Decreasing Length Conjecture

Conjecture. If $h : \Sigma^* \rightarrow \Sigma^*$, and Σ has n elements, then

$$|w| > |h(w)| > \cdots > |h^k(w)|$$

implies that $k \leq n$.

Another way to state the Decreasing Length Conjecture is the following:

Conjecture. Let M be an $n \times n$ matrix with non-negative integer entries. Let v be a column vector of non-negative integers, and let u be the row vector $[1 \ 1 \ 1 \ \cdots \ 1]$. If

$$uv > uMv > uM^2v > \cdots > uM^k v$$

then $k \leq n$.

Path Algebra

There is a nice correspondence between directed graphs and non-negative matrices, as follows:

If G is a directed graph on n vertices, we can construct a non-negative matrix

$$M(G) = (m_{i,j})_{1 \leq i,j \leq n}$$

as follows: let

$$m_{i,j} = \begin{cases} 1, & \text{if there is a directed edge from} \\ & \text{vertex } i \text{ to vertex } j \text{ in } G; \\ 0, & \text{otherwise.} \end{cases}$$

Then the number of distinct walks of length n from vertex i to vertex j in G is just the i,j 'th entry of M^n .

Similarly, given a non-negative $n \times n$ matrix $M = (m_{i,j})_{1 \leq i,j \leq n}$ we may form its associated graph $G(M)$ on n vertices, where we put a directed edge from vertex i to vertex j iff $m_{i,j} > 0$.

A Useful Lemma

Lemma. Let $r \geq 1$ be an integer, and suppose there exist r sequences of real numbers $\mathbf{b}_i = (b_i(n))_{n \geq 0}$, $1 \leq i \leq r$, and r positive integers h_1, h_2, \dots, h_r , such that the following conditions hold:

- (a) $b_i(n + h_i) \geq b_i(n)$ for $1 \leq i \leq r$ and $n \geq 0$;
- (b) There exists an integer $D \geq 1$ such that $\sum_{1 \leq i \leq r} b_i(n) > \sum_{1 \leq i \leq r} b_i(n + 1)$ for $0 \leq n < D$.

Then $D \leq h_1 + h_2 + \dots + h_r - r$.

Remark. When $r = 2$ and $\gcd(h_1, h_2) = 1$, then it can be shown that the bound in this Lemma is tight.

For example, for $h_1 = 5$, $h_2 = 8$ we find

n	0	1	2	3	4	5	6	7	8	9	10	11	12
$b_1(n)$	0	-16	8	-8	-24	0	-16	8	-8	-24	0	-16	8
$b_2(n)$	0	15	-10	5	20	-5	10	-15	0	15	-10	5	20
$b_1(n) + b_2(n)$	0	-1	-2	-3	-4	-5	-6	-7	-8	-9	-10	-11	28

Proof of the Decreasing Length Conjecture

Theorem. Suppose M is an $n \times n$ matrix with non-negative integer entries. If there exist a row vector u and a column vector v with non-negative integer entries such that

$$uv > uMv > uM^2v > \cdots > uM^k v,$$

then $k \leq n$. Also $k = n$ only if $M^n = 0$.

Proof.

- ▶ Let M be the matrix in the statement of the theorem and G its associated graph.
- ▶ Let $u = (u_1, u_2, \dots, u_n)$ and $v = (v_1, v_2, \dots, v_n)^T$.
- ▶ Let V be the set of vertices in G .
- ▶ Consider some maximal set of vertices forming disjoint cycles $\{C_1, C_2, \dots, C_r\}$ in G .
- ▶ Then V can be written as the disjoint union

$$V = C_1 \cup C_2 \cup \dots \cup C_r \cup W,$$

where W is the set of vertices which do not lie in any of the disjoint cycles.

- ▶ Any directed walk in G of length $|W|$ or greater must intersect some cycle C_i , for otherwise the walk would contain a cycle disjoint from C_1, C_2, \dots, C_r .
- ▶ Associate each walk of length $\geq |W|$ with the first cycle C_i it intersects.
- ▶ Define $P_{i,j,l}^s$ to be the number of directed walks of length s from vertex i to vertex j associated with cycle l .
- ▶ Also define

$$T_l^s := \sum_{1 \leq i, j \leq n} u_i v_j P_{ij,l}^s.$$

- ▶ Then for any $s \geq |W|$ we have

$$uM^s v = \sum_{1 \leq l \leq r} T_l^s. \quad (5)$$

- ▶ Then

$$T_i^s \leq T_i^{s+|C_i|},$$

since any walk of length s associated with cycle C_i can be extended to a walk of length $s + |C_i|$ by traversing the cycle C_i once.

- ▶ From the inequality $uM^s v > uM^{s+1} v$ for $0 \leq s \leq k - 1$ and Eq. (5) we have

$$\sum_{1 \leq i \leq r} T_i^s > \sum_{1 \leq i \leq r} T_i^{s+1}$$

for $|W| \leq s < k$.

- ▶ Now for $1 \leq i \leq r$ and $j \geq 0$ define $b_i(j) = T_i^{|W|+j}$ and $h_i = |C_i|$.
- ▶ Then the conditions of the previous Lemma are satisfied.

- ▶ We conclude that

$$k - |W| \leq |C_1| + |C_2| + \cdots + |C_r| - r.$$

- ▶ Moreover

$$|C_1| + |C_2| + \cdots + |C_r| + |W| = |V| = n$$

and so $k \leq n - r$.

- ▶ Finally $k = n$ implies that $r = 0$, so G is acyclic and $M^n = 0$.

For Further Reading

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