Lecture 18: Multiplicative Weights Update

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Overview

- Multiplicative Weights Update
- Solving Linear Programs
- Conclusion
- Acknowledgements

Learning from Experts

- **Setup:** investing your co-op money on stock markets (or gambling).
- **Objective:** to get rich, but we don't know much about stock markets
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- Can we hope to do *as well as the best expert* in hindsight?

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  - Worst-case analysis.

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Total money we made: $\geq T - \log n$

Total money best expert made: T

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4. Each trading day, choose to trade based on *weighted majority* of the decisions of the experts

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Theorem (Multiplicative Weights Update)

Let M_t be the number of mistakes that our algorithm makes until time t, and let $M_t(i)$ be the number of mistakes that expert i made until time t. Then, for any expert $i \in [n]$, we have:

$$M_t \leq 2 \cdot (1+\varepsilon) M_t(i) + \frac{2 \log n}{\varepsilon}$$

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  - Initially $\Phi_1 = n$
  - $\Phi_t \geq 0$ for all t

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$$\Phi_{t+1} = \sum_{i \text{ right}} w_t(i) + (1 - \varepsilon) \cdot \sum_{j \text{ wrong}} w_t(j) \leq \left(1 - \frac{\varepsilon}{2}\right) \cdot \Phi_t$$

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- Thus,

$$\Phi_t \leq \Phi_1 \cdot \left(1 - \frac{\varepsilon}{2}\right)^{M_t} = n \cdot \left(1 - \frac{\varepsilon}{2}\right)^{M_t}$$

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4. Using inequality $-x - x^2 < \log(1-x) < -x$ for $x \in (0, 1/2)$, we get:

$$-\varepsilon/2 \cdot M_t + \log n > M_t(i) \cdot (-\varepsilon - \varepsilon^2)$$

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Which implies update rule is:

$$w_{t+1}(i) = \left(1 - \varepsilon \cdot \frac{m_t(i)}{w}\right) \cdot w_t(i)$$

$$p_{t+1}(i) = \frac{w_{t+1}(i)}{\Phi_{t+1}}$$

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Theorem (Multiplicative Weights Update)

With the setup above, after t rounds, for any expert $i \in [n]$, we have:

$$\sum_{t=1}^{T} \langle p_t, m_t \rangle \leq \sum_{t=1}^{T} m_t(i) + \varepsilon \cdot \sum_{t=1}^{T} |m_t(i)| + \frac{w \cdot \ln n}{\varepsilon}$$

Proof of the Theorem

The proof of the theorem in the previous slide simply follows from the same idea we had together with the following inequality:

$$(1 - \varepsilon x) \geq \begin{cases} (1-\varepsilon)^x, & \text{if } x \in [0,1] \\ (1+\varepsilon)^{-x}, & \text{if } x \in [-1,0] \end{cases}$$

when $\varepsilon \in (0, 1/2)$.

Also worth noting the inequalities (from Taylor expansion of ln) for $y \in (0, 1/2)$:

$$\ln(1+y) \geq y - y^2/2 \geq y - y^2$$
$$\ln(1-y) \geq -y - y^2$$

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3. More precisely: cost of i^{th} constraint

$$A_i x - b_i$$

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1. Think of each inequality $A_i x \geq b_i$ as an *expert* (A_i is i^{th} row of A)
2. Each constraint would like to be *the hardest constraint*, i.e. the one that is violated the most by the current proposed solution $x^{(t)}$
3. More precisely: cost of i^{th} constraint

$$A_i x - b_i$$

4. We would like to propose feasible solution (i.e. lower cost of *all constraints*). Hard to deal with all constraints at the same time.

Oracle

Definition (Oracle)

Let $A \in \mathbb{R}^{m \times n}$. We say that $\mathcal{O}$ is an oracle of *width w* for A if given a linear constraint

$$p^T A x \geq p^T b, \;\; x \geq 0$$

$\mathcal{O}(p)$ will return $y \geq 0$ such that

$$|A_i y - b_i| \leq w \quad \forall i \in [m]$$

Solving Linear Programs

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- MWU shows that over the long run:

The *total violation* of our *weighted constraints* will be close to the *total violation* of the *worst violated constraint*!

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    - Thus, we will assume that above never happens.

Theorem

Theorem (Multiplicative Weights Update)

Let $\delta > 0$ and suppose we are given an oracle with width w for A. The MWU algorithm either finds a solution $y \geq 0$ such that

$$A_i y \geq b_i - \delta \quad \forall i \in [m]$$

or concludes that the system is infeasible (and outputs a dual solution). Our algorithm makes $O(w^2 \log(m)/\delta^2)$ oracle calls.

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- Setting $\varepsilon = \delta / 2w$ and $T = \dfrac{4 \cdot w^2 \cdot \log m}{\delta^2}$ we get

$$\sum_{t=1}^{T} \frac{A_i x^{(t)} - b_i}{T} \geq -\delta$$

Conclusion

- Online Learning
  1. Experts are weak classifiers, want to choose hypothesis based on these experts
  2. Boosting (in learning theory)
- Solving linear programs! (today)
- Convex Optimization
- Computational Geometry
- many more

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- See Yaron's notes https://people.seas.harvard.edu/~yaron/AM221-S16/lecture_notes/AM221_lecture11.pdf
- See Elad's survey at https://arxiv.org/pdf/1909.05207.pdf
- See great survey on MWU ar https://www.cs.princeton.edu/~arora/pubs/MWsurvey.pdf