Lecture 9: Random Walks, Markov Chains, Mixing Time

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Overview

Introduction

- Why Random Walks & Markov Chains?
- Basics on Theory of Finite Markov Chains

• Main Topics

- Stationary Distributions and Mixing Time
- Fundamental Theorem of Markov Chains

Conclusion

Acknowledgements

- Given a graph G(V, E)
 - **(**) random walk starts from a vertex v_0
 - at each time step it moves uniformly to a random neighbor of the <u>current vertex</u> in the graph

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Basic questions involving random walks:

• *Stationary distribution:* does the random walk converge to a "stable" distribution? If it does, what is this distribution?

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- *Cover time:* how long does it take to reach every vertex of the graph at least once?

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• Practice question: Compare question 2 to coupon collector problem!

What is a Markov Chain?

Random walk is a special kind of *stochastic process*:

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Process is "forgetful/memoryless"

Markov chain is characterized by this property.

Markov Chains and Random Walks are ubiquitous in randomized algorithms.

• Page Rank algorithm (next lecture)

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 Markov Chain *irreducible* if underlying directed graph is *strongly* connected (i.e. there is directed path from *i* to *j* for any pair *i*, *j* ∈ *V*)

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- $p_t \in \mathbb{R}^n$ probability vector: $p_t(i) := \Pr[\text{being at state } i \text{ at time } t]$
- Transition given by

$$p_{t+1} = p_t \cdot P$$

• *Period* of a state *i* is:

$$\mathsf{gcd}\{t \in \mathbb{N} \mid P_{i,i}^t > 0\}$$

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Lemma

For any finite, irreducible and aperiodic Markov Chain, there exists $T < \infty$ such that

$$P_{i,j}^t > 0$$
 for any $i, j \in V$ and $t \geq T$.

See proof in reference of [Häggström, Chapter 4].

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- Given two distributions $p, q \in \mathbb{R}^n$, their *total variational distance* is

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• p_t converges to q iff $\lim_{t \to \infty} \Delta_{TV}(p_t, q) = 0$

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The ε -mixing time of a Markov Chain is the smallest t such that

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• For complete graph, eigenvalues $\lambda_1 = 1, \lambda_2 = \cdots = \lambda_n = -1/(n-1)$, corresponding eigenvectors v_1, \ldots, v_n (orthonormal)

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Hitting Time

• Given states *i*, *j* in a Markov chain, the *hitting time* from state *i* to state *j* is defined as

$$T_{i,j} := \min\{t \ge 1 \mid X_t = j, X_0 = i\}$$

We say $T_{i,j} = \infty$ if the Markov chain never visits *j* starting from *i*.

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- The mean hitting time $\tau_{i,j} := \mathbb{E}[T_{i,j}]$
- *Hitting time lemma*: For any *finite*, *irreducible*, *aperiodic* Markov chain, and for any two states *i*, *j* (not necessarily distinct) we have that:

$$\Pr[T_{i,j} < \infty] = 1$$
 and $\mathbb{E}[T_{i,j}] < \infty$

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• Moreover, we can prove:

$$Pr[T_{i,j} > 2M] = Pr[T_{i,j} > M] \cdot Pr[T_{i,j} > 2M | T_{l,j} > M]$$

$$\leq (1 - \alpha) \cdot Pr[X_{2M} \neq j | T_{i,j} > M]$$

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- Thus, we have

$$\mathbb{E}[T_{i,j}] = \sum_{n \ge 1} \Pr[T_{i,j} \ge n] = \sum_{n \ge 0} \Pr[T_{i,j} > n] \le M/\alpha < \infty$$

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- The *return time* from state *i* to itself is *T_{i,i}*
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Theorem (Fundamental Theorem of Markov Chains)

Any finite, irreducible and aperiodic Markov Chain has the following properties:

- There exists a unique stationary distribution π, where π_i > 0 for all i ∈ [n]
- 2 The sequence of distributions {p_t}_{t≥0} will converge to π, no matter what the initial distribution is

$$\pi_i = \lim_{t \to \infty} P_{i,i}^t = \frac{1}{\tau_{i,i}}$$

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Intuition for proof of this theorem:

- two random walks are "indistinguishable" after they "<u>meet</u>" at the same vertex v at a particular time t
- By finiteness, irreducibility and aperiodicity, two walks will meet with positive probability (and thus by <u>Markov property</u>) become <u>same</u> <u>distribution</u>

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- Eigenvectors of *P* are $D^{-1/2}v_i$ where v_i are eigenvectors of *P'*. And v_i can be taken to form *orthonormal basis*.

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 - Graph strongly connected ⇒ *Perron-Frobenius* for irreducible non-negative matrices
 - unique eigenvector whose eigenvalue has max absolute value
 - eigenvector has all positive coordinates
 - eigenvalue is *positive*

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- Transition matrix: $P = D^{-1} \cdot A_G$
 - *P* similar to a symmetric matrix: $P' = D^{-1/2} A_G D^{-1/2}$
 - P and P' has same eigenvalues! And P' has only real eigenvalues!
 - Eigenvectors of *P* are $D^{-1/2}v_i$ where v_i are eigenvectors of *P'*. And v_i can be taken to form *orthonormal basis*.
 - Graph strongly connected ⇒ *Perron-Frobenius* for irreducible non-negative matrices
 - unique eigenvector whose eigenvalue has max absolute value
 - eigenvector has all positive coordinates
 - eigenvalue is *positive*
 - This eigenvector is $\pi!$

- Stationary distribution: $\pi_i = \frac{d_i}{2m}, \quad m = |E|$
- Transition matrix: $P = D^{-1} \cdot A_G$
 - *P* similar to a symmetric matrix: $P' = D^{-1/2} A_G D^{-1/2}$
 - P and P' has same eigenvalues! And P' has only real eigenvalues!
 - Eigenvectors of *P* are $D^{-1/2}v_i$ where v_i are eigenvectors of *P'*. And v_i can be taken to form *orthonormal basis*.
 - Graph strongly connected ⇒ *Perron-Frobenius* for irreducible non-negative matrices
 - unique eigenvector whose eigenvalue has max absolute value
 - eigenvector has all positive coordinates
 - eigenvalue is *positive*
 - This eigenvector is $\pi!$
 - All random walks converge to π , as we wanted to show.

Revisiting the complete graph

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Mixing time from eigenvalue gap

- write here that mixing time follows from eigenvalue gap
- in the next lecture, after Perron-Frobenius, revisit this point

Conclusion

Markov Chains and Random Walks are ubiquitous in randomized algorithms.

- Page Rank algorithm (next lecture)
- Approximation algorithms for counting problems [Karp, Luby & Madras]
 - Permanent of non-negative matrices [Jerrum, Vigoda & Sinclair]
- Sampling Problems
 - Gibbs sampling in statistical physics
 - many more places
- Probability amplification without too much randomness (efficient)
 - Random walks on expander graphs
- many more

Acknowledgement

- Lecture based largely on:
 - Lap Chi's notes
 - [Motwani & Raghavan 2007, Chapter 6]
 - [Häggström]
- See Lap Chi's notes at https://cs.uwaterloo.ca/~lapchi/cs466/notes/L11.pdf
- Also see Lap Chi's notes https://cs.uwaterloo.ca/~lapchi/cs466/notes/L14.pdf for a proof of fundamental theorem of Markov chains for undirected graphs.

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