Lecture 25: Conclusion

Rafael Oliveira

University of Waterloo Cheriton School of Computer Science

rafael.oliveira.teaching@gmail.com

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Overview

- Administrivia
- Foundations of Symbolic Computation
- Computational Linear Algebra
- Modern Computational Algebra
- Computational Invariant Theory
- Topics I wish I had time to cover

Please log in to

https://evaluate.uwaterloo.ca/

Today is the **last day** to provide us (and the school) with your evaluation and feedback on the course!

- This would really help me figuring out what worked and what didn't for the course
- And let the school know if I was a good boy this term!
- Teaching this course is also a learning experience for me :)

Administrivia

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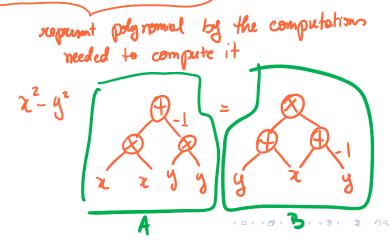
Dense representation

(x,y,2) E [[x,y,2] deg(P = 2 ALL coefficients (include (Poo, Pio, Poio, Poor, Pllo, Pior, Poir, Pzoo, TI (x,11) P020, Por2) Pise = coefficient of xigizh

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 - Operation Sparse representation

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 - Black-Box matrix representations

$$\begin{bmatrix} a & b \\ d & a \\ e & d \\ e$$

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 - Algebraic circuits (straight-line programs)
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- Complexity of certain problems become vastly different depending on representation!
- Some open problems:
 - **1** factoring sparse (univariate or multivariate) polynomials fast
 - factoring multivariate polynomials computed by algebraic circuits (without restriction on degree)
 - testing whether two objects from the same model compute the same object

Given two straight-line programs, do they compute the same polynomial?

Improve generally, the more succinct the representation, the harder it should be to efficiently solve problems

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- Learned how to multiply polynomials much faster!
 - I Highly non-trivial algorithms!
 - Ø Karatsuba: reduce number of multiplications needed!

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- Polynomial multiplication quite often used in practice!
- Even more ubiquitous is the Discrete Fourier Transform!
 - Used in audio and video compression [von zur Gathen, Gerhard 2013, Chapter 13] and references
 - many more applications!

Fundamental Operations - Euclidean Algorithm

• Learned how to compute the GCD between integers and two polynomials over **F**(x)

Fundamental Operations - Euclidean Algorithm

$$af + bg = gcd(f_1g)$$

- Learned how to compute the GCD between integers and two polynomials
- Extended Euclidean Algorithm fundamental for many other problems
 - Compute inverses in modular computations
 - Solving Pade Approximation problem in power series approximations

Much simpler to compute linear recurrence sequences! approximate pouver series by rational functions

Modular Computation and Chinese Remandering Theorem

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Modular Computation and Chinese Remandering Theorem

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- Chinese Remaindering Theorem allows us to
 - Parallelize" the problem: compute many instances of the problem modulo small primes

Can compute all these instances in parallel!

2 Control intermediate coefficients

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2 Control intermediate coefficients

• It can also be used to give fastest known algorithm for univariate polynomial factoring over finite fields! Kedlaya-Umans 2011

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• Use resultant to compute a modular GCD algorithm for two polynomials over $\mathbb{Z}[x]$

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Euclidean algorithm only works for Euclidean Domains, and $\mathbb{Z}[x]$ is not an Euclidean domain

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Euclidean domains

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- Resultants also have nice theoretical properties
 - Identifies the bad primes in modular algorithms
 - **2** Used as subroutine in factoring algorithms when double roots appear
 - Also used to prove upper bound in complexity of ideal membership problem!
 - Many more applications!

- Univariate polynomials over *Finite Fields*
 - Cantor-Zassenhaus algorithm
 - Berlekamp-Rabin algorithm

 $^{^1}$ Lovasz won the Abel prize this year - the Nobel prize for mathematics. In their acknowledgments, they mentioned this algorithm as one of his (many) remarkable worksl $_\odot$

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 - Hensel lifting
 - Bounds on coefficients of factors

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 - Reduce to univariate factorization
 - Use Hensel lifting to recover multivariate factorization

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Berlekamp's decoding algorithm of Reed-Solomon Codes!

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- Applications in list decoding of Reed-Solomon codes!

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One of the major open problems in computer science!

• Deep connections between matrix multiplication and ranks of tensors

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Known as *backpropagation* in Machine Learning.

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• Good thing about the recurrence we found is that *parallel algorithms* for matrix multiplication yield *parallel algorithms* for matrix inversion and determinant!

 An ubiquitous problem in scientific computing is to solve system of linear equations Ay = b

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- linear programming
- optimization
- olynomial multiplication
- 4 factoring
- oplynomial interpolation (DFT)
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- We have already done that many times!

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- Let c(A) be the cost of multiplying A by any vector **b**, and M(n) the cost of multiplying two degree n polynomials

Class of matrices	c(A)
general	$2n^2 - 2$
Sylvester Matrix (Resultion b)	O(M(n))
DFT	$O(n \log n)$
Vandermonde matrix	$O(M(n) \log n)$ $O(M(n) \log q)$
Berlekamp matrix over $\mathbb{F}_{m{q}}$	$O(M(n)\log q)$
Sparse matrix with s non-zero entries	2 <i>s</i>
Toeplitz matrix	O(M(n))

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- In lecture 22 we devised much faster algorithms for inverting matrices with low c(A) in black-box model

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Connections Between Algebra & Geometry

• Ideals in $\mathbb{C}[x_1, \dots, x_n]$ are *finitely generated* Hilbert's basis theorem. Connections Between Algebra & Geometry

 \bullet Ideals in $\mathbb{C}[x_1,\ldots,x_n]$ are finitely generated

Hilbert's basis theorem.

• (radical) Ideals in polynomial rings correspond to algebraic sets in finite-dimensional vector spaces

Hilbert's Nullstellensatz.

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Connections Between Algebra & Geometry

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Hilbert's Nullstellensatz.

- Central problems in modern commutative algebra:
 - Ideal membership problem
 - Solving System of Polynomial equations
 - Section 2 Sec
 - Implicitization Problem

Gröbner bases Elimination Theory Extension Theorem

Applications of Symbolic Commutative Algebra

- Applications in mathematics
 - Compute dimension of algebraic sets
 - 2 compute Hilbert Polynomials
 - 8 Betti numbers
 - Resolution of singularities
 - Many more!

important numeric invariants

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- Methods to solve integer programming use Gröbner bases
- Bayesian Networks conditional dependencies define algebraic sets!
- Topological data analysis
- many more!

important numeric invariants

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Finite Generation of Rings of Invariants

• We learned that Hilbert himself when he proved the Nullstellensatz and the basis theorem was after proving that

Ring of invariant polynomials are finitely generated

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- Invariants capture many interesting properties of our algebraic and geometric objects
 - Whether a matrix is singular or not
 - Ø bipartite matching
 - Inilpotent matrices
 - graph isomorphism
 - S word problem over free skew fields
 - Iinear matroid intersection
 - computation of optimal transport distances
 - Contingency tables
 - Maximum Likelihood Estimation
 - Symmetries in chemistry molecules
 - 🗓 many more

Computational Aspects of Invariant Rings

- Algorithm (via Reynolds operator) to compute invariant polynomials of a certain degree
- Reynolds + Hilbert's argument gave us finite generation of ring of invariants

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Computational Aspects of Invariant Rings

- Algorithm (via Reynolds operator) to compute invariant polynomials of a certain degree
- Reynolds + Hilbert's argument gave us finite generation of ring of invariants
- One major open question in computational invariant theory is to *efficiently compute* a generating set of invariants
- Depending on how efficient we can compute the invariants, it can have striking applications in computer science and other fields!

Topics I wish I had time to cover

- Solving Differential Equations
- Symbolic Integration
- Semialgebraic Systems of Equations
- Computing Radical of Ideal
- Checking Algebraic Independence
- Computing Primary Decompositions of Ideals
- Complexity theory for algebraic computation
- Many more amazing topics in symbolic computation to explore!

1 p(x̄) ≥ 0

Thank you for taking the class!

References I



von zur Gathen, J. and Gerhard, J. 2013.

Modern Computer Algebra

Cambridge University Press

