# Lecture 25: Conclusion 

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## Overview

- Administrivia
- Foundations of Symbolic Computation
- Computational Linear Algebra
- Modern Computational Algebra
- Computational Invariant Theory
- Topics I wish I had time to cover


## Rate this course!

## Please log in to

https://evaluate.uwaterloo.ca/

Today is the last day to provide us (and the school) with your evaluation and feedback on the course!

- This would really help me figuring out what worked and what didn't for the course
- And let the school know if I was a good boy this term!
- Teaching this course is also a learning experience for me:)
- Administrivia
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## Models of Computation

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(1) Dense representation

$$
p(x, y, z) \in \mathbb{F}[x, y, z] \quad \operatorname{deg}(P)=2
$$

ALL coaffienn (include $\left.\begin{array}{c}\text { ers coiffient) }\end{array}\right)$
(P000, $P_{100}, P_{010}, P_{001}, P_{110}, P_{101}, P_{011}, P_{200}$,

$$
\begin{aligned}
& \text { Pore, Poor) } \\
& P_{i j k}=\text { coefficient of } x^{i} y^{j} z^{h}
\end{aligned}
$$

$$
\prod_{i=1}^{n}(x ; 1)
$$

Models of Computation

- In addition to the standard bit representation of integers, we learned different models to represent algebraic objects:
(0) Dense representation
- Sparse representation
only tell non-zer coff, icienb
and their corresponding monomials

non-zers teams

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(9) Black-Box matrix representations

$$
\left[\begin{array}{lll}
a & b & c \\
d & a & b \\
e & d & a
\end{array}\right] \leftrightarrow(e, d, a, b, c)
$$

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- In addition to the standard bit representation of integers, we learned different models to represent algebraic objects:
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- Complexity of certain problems become vastly different depending on representation!
- Some open problems:
(1) factoring sparse (univariate or multivariate) polynomials fast
(2) factoring multivariate polynomials computed by algebraic circuits (without restriction on degree)
(3) testing whether two objects from the same model compute the same object

Given two straight-line programs, do they compute the same polynomial?
(9) more generally, the more succinct the representation, the harder it should be to efficiently solve problems

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- Polynomial multiplication quite often used in practice!
- Even more ubiquitous is the Discrete Fourier Transform!
- Used in audio and video compression [von zur Gathen, Gerhard 2013, Chapter 13] and references
- many more applications!


## Fundamental Operations - Euclidean Algorithm

- Learned how to compute the GCD between integers and two polynomials over $\mathscr{F}[x]$

Fundamental Operations - Euclidean Algorithm

$$
a f+b g=\operatorname{gcd}(f, g)
$$

- Learned how to compute the GCD between integers and two polynomials
- Extended Euclidean Algorithm fundamental for many other problems
(1) Compute inverses in modular computations
(2) Solving Pade Approximation problem in power series approximations (Much simpler to compute linear recurrence sequences! functions


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Can compute all these instances in parallel!
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- It can also be used to give fastest known algorithm for univariate polynomial factoring over finite fields! Kedlaya-Umans 2011


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$$
\mathbb{Z}[x] \longleftrightarrow
$$

$$
\frac{\mp_{p_{i}}[x]}{\substack{\text { Euclidean domains } \\ E F A}}
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- Allowed us to reduce the problem above to the Euclidean GCD algorithm
Euclidean algorithm only works for Euclidean Domains, and $\mathbb{Z}[x]$ is not an Euclidean domain
- Resultants also have nice theoretical properties
(1) Identifies the bad primes in modular algorithms
(2) Used as subroutine in factoring algorithms - when double roots appear
(3) Also used to prove upper bound in complexity of ideal membership problem!
(9) Many more applications!


## Polynomial Factoring

- Univariate polynomials over Finite Fields
- Cantor-Zassenhaus algorithm
- Berlekamp-Rabin algorithm

[^0]
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- Factoring bivariate polynomials
- Reduce to univariate factorization
- Use Hensel lifting to recover multivariate factorization

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- Applications in list decoding of Reed-Solomon codes!

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- Exponent of matrix multiplication ubiquitous in computational linear algebra!

One of the major open problems in computer science!

- Deep connections between matrix multiplication and ranks of tensors


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- Good thing about the recurrence we found is that parallel algorithms for matrix multiplication yield parallel algorithms for matrix inversion and determinant!


## Black-Box Linear Algebra

- An ubiquitous problem in scientific computing is to solve system of linear equations $A \mathbf{y}=\mathbf{b}$
(1) linear programming
(2) optimization
(3) polynomial multiplication
(3) factoring
(5) polynomial interpolation (DFT)
(0) computing GCD of polynomials (Resultants)
(3) many more


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- We have already done that many times!


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- Let $c(A)$ be the cost of multiplying $A$ by any vector $\mathbf{b}$, and $M(n)$ the cost of multiplying two degree $n$ polynomials

| Class of matrices | $c(A)$ |
| :--- | :---: |
| general | $2 n^{2}-2$ |
| Sylvester Matrix (Resulter $b)$ | $O(M(n))$ |
| DFT | $O(n \log n)$ |
| Vandermonde matrix | $O(M(n) \log n)$ |
| Berlekamp matrix over $\mathbb{F}_{q}$ | $O(M(n) \log q)$ |
| Sparse matrix with $s$ non-zero entries | $2 s$ |
| Toeplitz matrix | $O(M(n))$ |

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- In lecture 22 we devised much faster algorithms for inverting matrices with low $c(A)$ in black-box model
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## Connections Between Algebra \& Geometry

- Ideals in $\mathbb{C}\left[x_{1}, \ldots, x_{n}\right]$ are finitely generated

Hilbert's basis theorem.

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- (radical) Ideals in polynomial rings correspond to algebraic sets in finite-dimensional vector spaces

Hilbert's Nullstellensatz.

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Hilbert's Nullstellensatz.

- Central problems in modern commutative algebra:
(1) Ideal membership problem
(2) Solving System of Polynomial equations
(3) Extending partial solutions

Gröbner bases
Elimination Theory
Extension Theorem
(1) Implicitization Problem

## Applications of Symbolic Commutative Algebra

- Applications in mathematics
(1) Compute dimension of algebraic sets
(2) compute Hilbert Polynomials important numeric invariants
(3) Betti numbers
(c) Resolution of singularities
(6) Many more!


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- Robotics (robot motion and gemetric descriptions)
- Automatic Geometric Theorem Proving
- Methods to solve integer programming use Gröbner bases
- Bayesian Networks conditional dependencies define algebraic sets!
- Topological data analysis
- many more!
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- Topics I wish I had time to cover


## Finite Generation of Rings of Invariants

- We learned that Hilbert himself when he proved the Nullstellensatz and the basis theorem was after proving that

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- Invariants capture many interesting properties of our algebraic and geometric objects
(1) Whether a matrix is singular or not
(2) bipartite matching
(3) nilpotent matrices
(9) graph isomorphism
(3) word problem over free skew fields
(6) linear matroid intersection
(1) computation of optimal transport distances
(8) contingency tables
(9) Maximum Likelihood Estimation
(10) Symmetries in chemistry molecules
(1) many more


## Computational Aspects of Invariant Rings

- Algorithm (via Reynolds operator) to compute invariant polynomials of a certain degree
- Reynolds + Hilbert's argument gave us finite generation of ring of invariants


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## Computational Aspects of Invariant Rings

- Algorithm (via Reynolds operator) to compute invariant polynomials of a certain degree
- Reynolds + Hilbert's argument gave us finite generation of ring of invariants
- One major open question in computational invariant theory is to efficiently compute a generating set of invariants
- Depending on how efficient we can compute the invariants, it can have striking applications in computer science and other fields!

Topics I wish I had time to cover

- Solving Differential Equations
- Symbolic Integration
- Semialgebraic Systems of Equations
- Computing Radical of Ideal

$$
\begin{gathered}
\mathbb{R}[\bar{x}] \\
\psi \\
\left.P_{i}(\bar{x}) \geqslant 0\right\}_{i=1}^{t}
\end{gathered}
$$

- Checking Algebraic Independence
- Computing Primary Decompositions of Ideals
[- Complexity theory for algebraic computation
- Many more amazing topics in symbolic computation to explore!

Polynomial
finding
Testing

Thank you for taking the class!

## References I

von zur Gathen, J. and Gerhard, J. 2013.
Modern Computer Algebra
Cambridge University Press


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