



CS486/686 Lecture Slides (c) 2017 P.Poupart

Solving Problems by Searching

[RN2] Sec 3.1-3.5 [RN3] Sec 3.1-3.4

CS486/686 University of Waterloo Lecture 2: May 3, 2017

Outline

- Problem solving agents and search
- Examples
- Properties of search algorithms
- Uninformed search
 - Breadth first
 - Depth first
 - Iterative Deepening

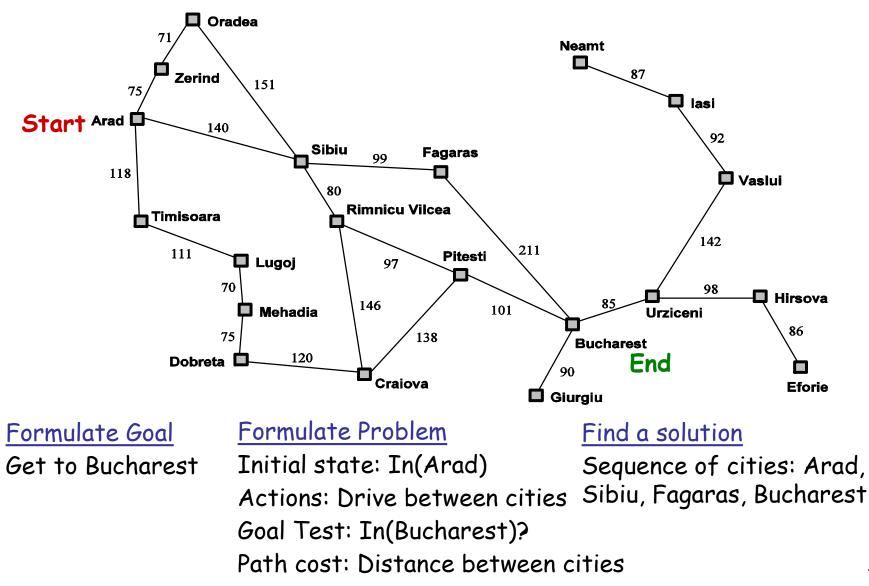
Introduction

- Search was one of the first topics studied in AI
 Newell and Simon (1961) *General Problem Solver*
- Central component to <u>many</u> AI systems
 - Automated reasoning, theorem proving, robot navigation, VLSI layout, scheduling, game playing,...

Problem-solving agents

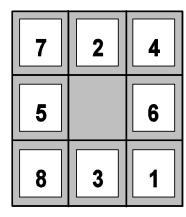
```
function SIMPLE-PROBLEM-SOLVING-AGENT( percept) returns an action
static: seq, an action sequence, initially empty
         state, some description of the current world state
         goal, a goal, initially null
         problem, a problem formulation
state \leftarrow UPDATE-STATE(state, percept)
if seq is empty then do
     goal \leftarrow FORMULATE-GOAL(state)
     problem \leftarrow FORMULATE-PROBLEM(state, goal)
     seq \leftarrow SEARCH(problem)
action \leftarrow FIRST(seq)
seq \leftarrow \text{Rest}(seq)
return action
```

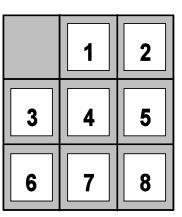
Example: Traveling in Romania



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Examples of Search Problems





Start State

Goal State

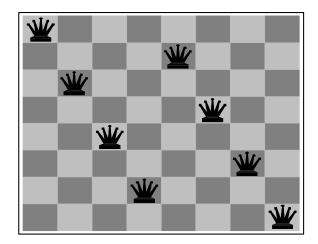
States: Locations of 8 tiles and blank

Initial State: Any state

Succ Func: Generates legal states that result from trying 4 actions (blank up, down, left, right)

Goal test: Does state match desired configuration

Path cost: Number of steps



States: Arrangement of 0 to 8 queens on the board

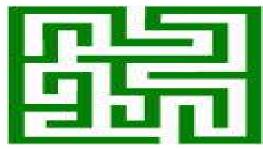
Initial State: No queens on the board Succ Func: Add a queen to an empty space

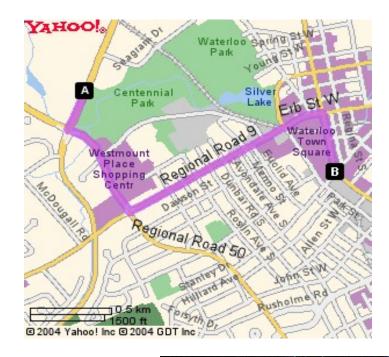
Goal test: 8 queens on board, none attacked

```
Path cost: none
```

More Examples











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Common Characteristics

- All of those examples are
 - Fully observable
 - Deterministic
 - Sequential
 - Static
 - Discrete
 - Single agent
- Can be tackled by **simple** search techniques

Cannot tackle these yet...

Chance

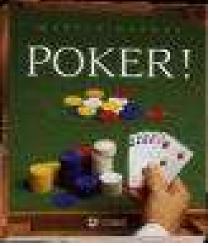
Infinite number of states





Games against an adversary





Hidden states

All of the above



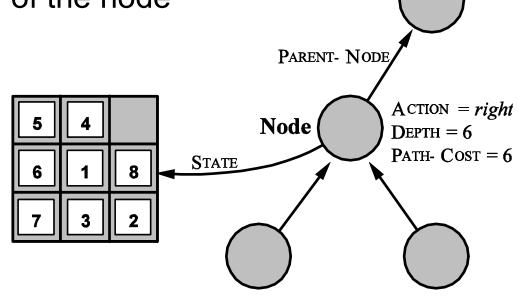
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Searching

- We can formulate a search problem
 - Now need to find the solution
- We can visualize a state space search in terms of trees or graphs
 - Nodes correspond to states
 - Edges correspond to taking actions
- We will be studying search trees
 - These trees are constructed "on the fly" by our algorithms

Data Structures for Search

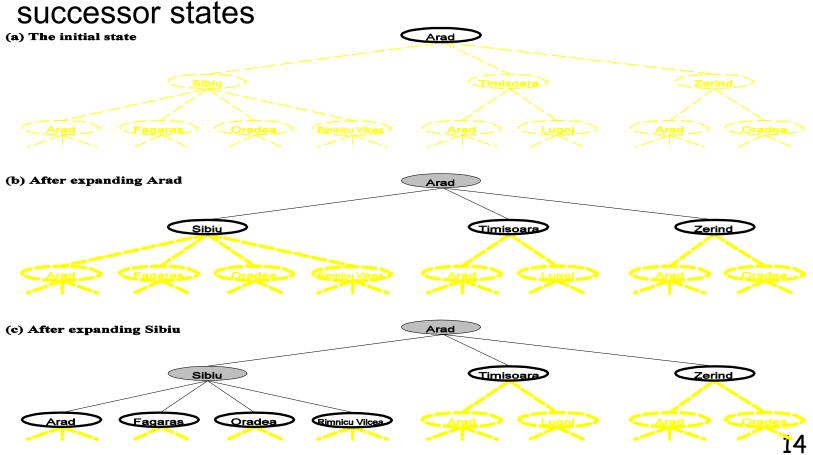
- Basic data structure: Search Node
 - State
 - Parent node and operator applied to parent to reach current node
 - Cost of the path so far
 - Depth of the node



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Expanding Nodes

- Expanding a node
 - Applying all legal operators to the state contained in the node and generating nodes for all corresponding



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Generic Search Algorithm

- 1. Initialize search algorithm with initial state of the problem
- 2. Repeat
 - 1. If no candidate nodes can be expanded, **return** failure
 - 2. Choose leaf node for expansion, according to search strategy
 - 3. If node contains a goal state, **return solution**
 - 4. Otherwise, expand the node, by applying legal operators to the state within the node. Add resulting nodes to the tree

Implementation Details

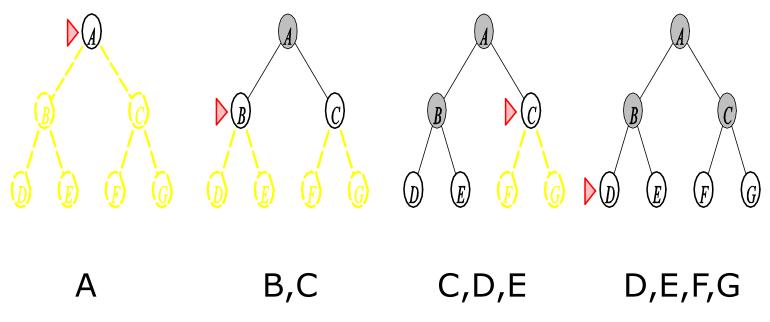
- We need to keep track only of nodes that need to be expanded (fringe)
 - Done by using a (prioritized) queue
- 1. Initialize queue by inserting the node corresponding to the initial state of the problem
- 2. Repeat
 - 1. If queue is empty, **return failure**
 - 2. Dequeue a node
 - 3. If the node contains a goal state, return solution
 - 4. Otherwise, expand node by applying legal operators to the state within. Insert resulting nodes into queue

Search algorithms differ in their queuing function!

Breadth-first search

All nodes on a given level are expanded before any nodes on the next level are expanded.

Implemented with a FIFO queue



Evaluating search algorithms

- Completeness: Is the algorithm guaranteed to find a solution if a solution exists?
- Optimality: Does the algorithm find the optimal solution (Lowest path cost of all solutions)
- Time complexity
- Space complexity

Variables

b	Branching factor			
d	Depth of shallowest goal node			
m	Maximum length of any path in the state space			

Judging BFS

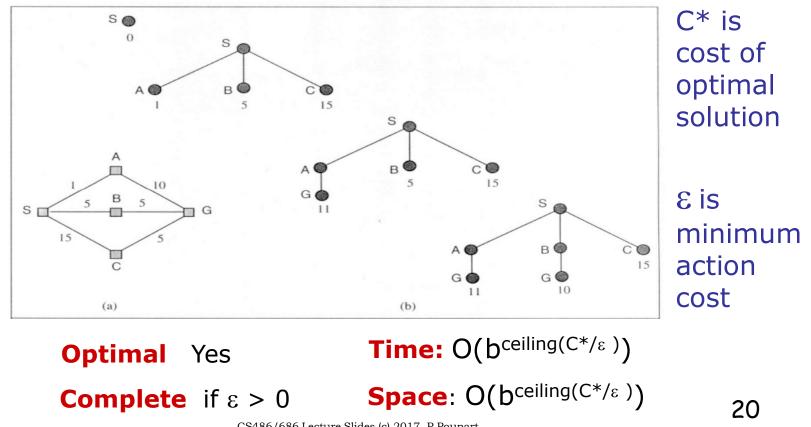
• Complete:

- Yes, if b is finite
- Optimal:
 - Yes, if all costs are the same
- Time:
 - $1+b+b^2+b^3+...+b^d = O(b^d)$
- Space:
 - O(b^d)

All uninformed search methods will have exponential time complexity \otimes

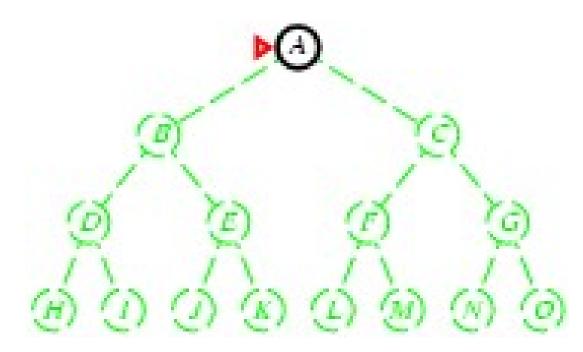
Uniform Cost Search

- A variation of breadth-first search
 - Instead of expanding shallowest node it expands the node with lowest path cost
 - Implemented using a priority queue



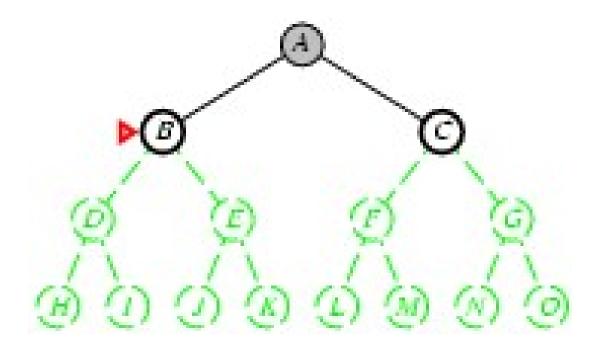
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The deepest node in the current fringe of the search tree is expanded first.



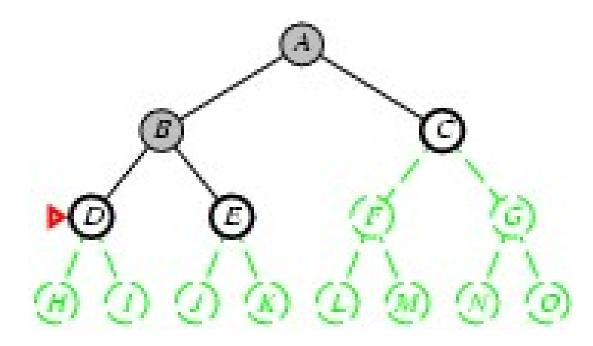
The deepest node in the current fringe of the search tree is expanded first.

Implemented with a stack (LIFO queue)

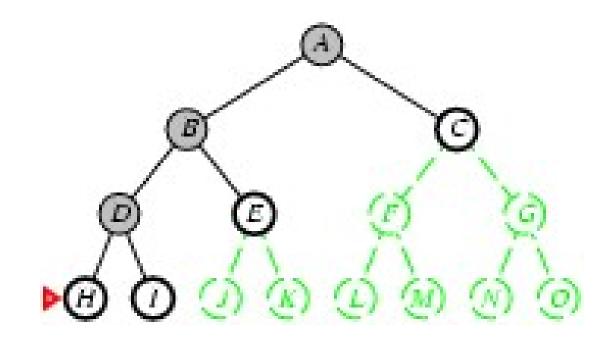


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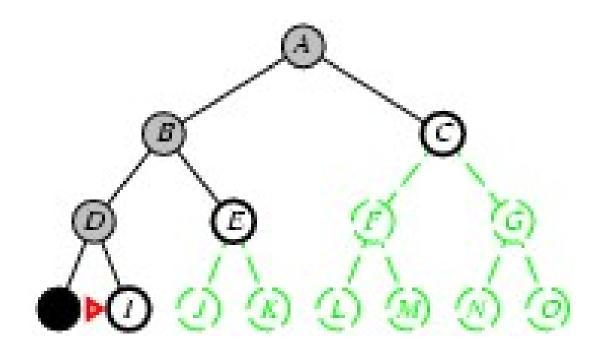
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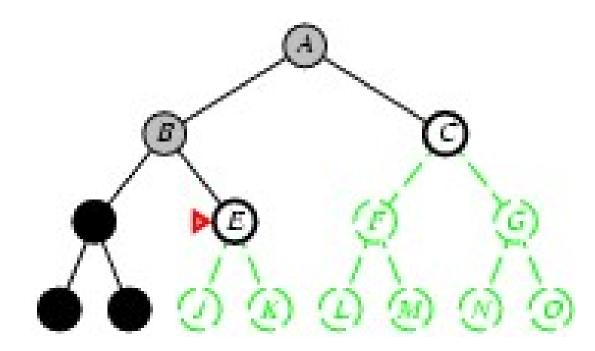


The deepest node in the current fringe of the search tree is expanded first.



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Implemented with a stack (LIFO queue)



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Judging DFS

• Complete:

- No, might get stuck going down a long path
- Optimal:
 - No, might return a solution which is deeper (i.e. more costly) than another solution
- Time:
 - O(b^m), m might be larger than d
- Space:
 - O(bm) 😳

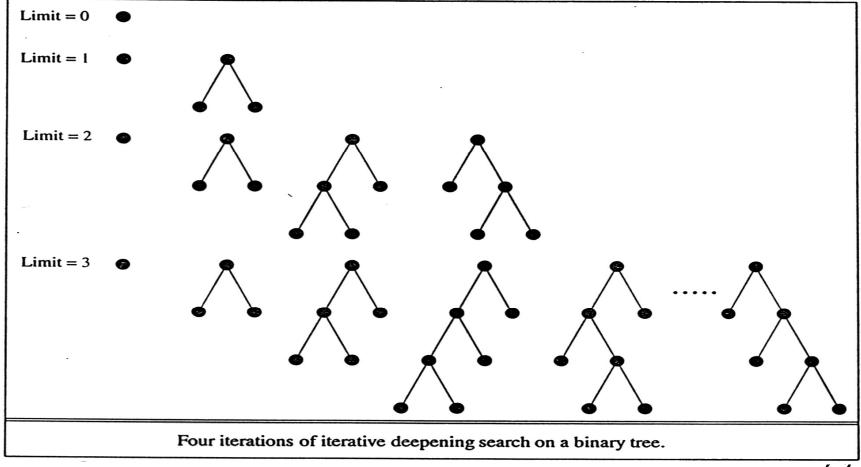
Do not use DFS if you suspect a large tree depth

Depth-limited search

- We can avoid the problem of unbounded trees by using a depth limit, l
 - All nodes at depth *l* are treated as though they have no successors
 - If possible, choose *l* based on knowledge of the problem
- **Time:** *O*(*b*^{*l*})
- Memory: O(bl)
- Complete?: No
- Optimal?: No

Iterative-deepening

• General strategy that repeatedly does depthlimited search, but increases the limit each time



Iterative-deepening

IDS is not as wasteful as one might think.

Note, most nodes in a tree are at the bottom level. It does not matter if nodes at a higher level are generated multiple times.

Breadth first search :

 $1 + b + b^2 + ... + b^{d-1} + b^d$ E.g. b=10, d=5: 1+10+100+1,000+10,000+100,000 = 111,111

Iterative deepening search :

 $(d+1)^*1 + (d)^*b + (d-1)^*b^2 + ... + 2b^{d-1} + 1b^d$ E.g. 6+50+400+3000+20,000+100,000 = 123,456

Complete, Optimal, O(b^d) time, O(bd) space



- Problem formulation usually requires abstracting away real-world details to define a state space that can feasibly be explored
- Variety of uninformed search strategies
 - Assume no knowledge about the problem (general but expensive)
 - Mainly differ in the order in which they consider the states

Criteria	BFS	Uniform	DFS	DLS	IDS
Complete	Yes	Yes	No	No	Yes
Time	O(b ^d)	O(b ^{ceiling(C*/ɛ)})	O(b ^m)	O(b ^I)	O(b ^d)
Space	O(b ^d)	O(b ^{ceiling(C*/ɛ)})	O(bm)	O(bl)	O(bd)
Optimal	Yes	Yes	No	No	Yes

Summary

- Iterative deepening uses only linear space and not much more time than other uninformed search algorithms
 - Use IDS when there is a large state space and the maximum depth of the solution is unknown
- Things to think about:
 - What about searching graphs?
 - Repeated states?