

Solving Problems by Searching

CS486/686
University of Waterloo
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Outline

- Problem solving agents and search
- Examples
- Properties of search algorithms
- Uninformed search
 - Breadth first
 - Depth first
 - Iterative Deepening

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Introduction

- Search was one of the first topics studied in AI
 - Newell and Simon (1961) *General Problem Solver*
- Central component to many AI systems
 - Automated reasoning, theorem proving, robot navigation, VLSI layout, scheduling, game playing,...

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Problem-solving agents

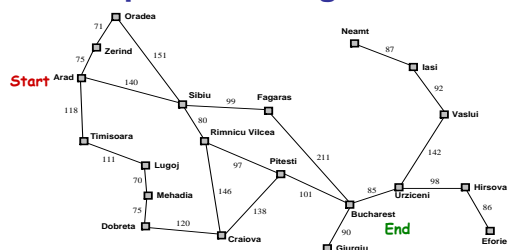
```
function SIMPLE-PROBLEM-SOLVING-AGENT( percept) returns an action
static: seq, an action sequence, initially empty
        state, some description of the current world state
        goal, a goal, initially null
        problem, a problem formulation

state ← UPDATE-STATE(state, percept)
if seq is empty then do
  goal ← FORMULATE-GOAL(state)
  problem ← FORMULATE-PROBLEM(state, goal)
  seq ← SEARCH( problem)
action ← FIRST(seq)
seq ← REST(seq)
return action
```

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Example: Traveling in Romania



Formulate Goal
Get to Bucharest

Formulate Problem
Initial state: In(Arad)
Actions: Drive between cities
Goal Test: In(Bucharest)?
Path cost: Distance between cities

Find a solution
Sequence of cities: Arad, Sibiu, Fagaras, Bucharest

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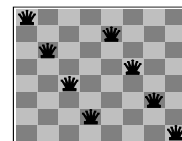
Examples of Search Problems



Start State



Goal State



States: Locations of 8 tiles and blank on the board

Initial State: Any state

Succ Func: Generates legal states that result from trying 4 actions (blank up, down, left, right)

Goal test: Does state match desired configuration

Path cost: Number of steps

States: Arrangement of 0 to 8 queens on the board

Initial State: No queens on the board

Succ Func: Add a queen to an empty space

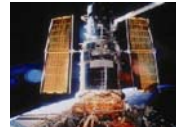
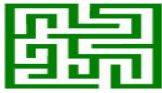
Goal test: 8 queens on board, none attacked

Path cost:

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More Examples



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Common Characteristics

- All of those examples are
 - Fully observable
 - Deterministic
 - Sequential
 - Static
 - Discrete
 - Single agent
- Can be tackled by **simple** search techniques

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Cannot tackle these yet...



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Cannot tackle these yet...

Chance



Infinite number of states



Games against an adversary



Hidden states

All of the above



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Searching

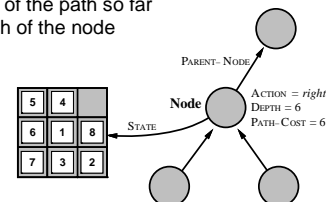
- We can formulate a search problem
 - Now need to find the solution
- We can visualize a state space search in terms of trees or graphs
 - Nodes correspond to states
 - Edges correspond to taking actions
- We will be studying **search trees**
 - These trees are constructed "on the fly" by our algorithms

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Data Structures for Search

- Basic data structure: **Search Node**
 - State
 - Parent node and operator applied to parent to reach current node
 - Cost of the path so far
 - Depth of the node

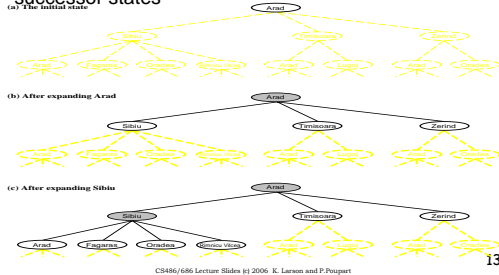


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Expanding Nodes

- Expanding a node
 - Applying all legal operators to the state contained in the node and generating nodes for all corresponding successor states



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Generic Search Algorithm

1. Initialize search algorithm with initial state of the problem
2. **Repeat**
 1. If no candidate nodes can be expanded, **return failure**
 2. Choose leaf node for expansion, according to **search strategy**
 3. If node contains a goal state, **return solution**
 4. Otherwise, expand the node, by applying legal operators to the state within the node. Add resulting nodes to the tree

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Implementation Details

- We need to only keep track of nodes that need to be expanded (**fringe**)
 - Done by using a (prioritized) **queue**
1. Initialize queue by inserting the node corresponding to the initial state of the problem
 2. **Repeat**
 1. If queue is empty, **return failure**
 2. Dequeue a node
 3. If the node contains a goal state, **return solution**
 4. Otherwise, expand node by applying legal operators to the state within. Insert resulting nodes into queue

Search algorithms differ in their queuing function!

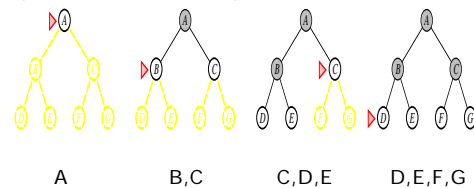
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Breadth-first search

All nodes on a given level are expanded before any nodes on the next level are expanded.

Implemented with a FIFO queue



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Evaluating search algorithms

- Completeness:** Is the algorithm guaranteed to find a solution if a solution exists?
- Optimality:** Does the algorithm find the optimal solution (Lowest path cost of all solutions)
- Time complexity**
- Space complexity**

Variables

b	Branching factor
d	Depth of shallowest goal node
m	Maximum length of any path in the state space

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Judging BFS

- Complete:** Yes, if b is finite
- Optimal:** Yes, if all costs are the same
- Time:** $1 + b + b^2 + b^3 + \dots + b^d = O(b^d)$
- Space:** $O(b^d)$

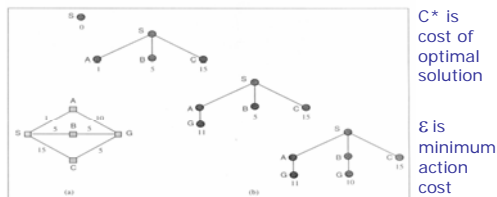
All uninformed search methods will have exponential time complexity ☹

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Uniform Cost Search

- A variation of breadth-first search
 - Instead of expanding shallowest node it expands the node with lowest path cost
 - Implemented using a priority queue



Optimal
Complete if $\epsilon > 0$

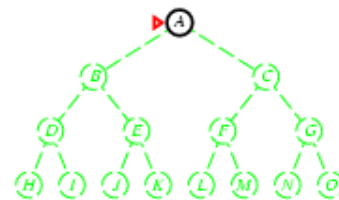
Time: $O(b^{\lceil C^*/\epsilon \rceil})$
Space: $O(b^{\lceil C^*/\epsilon \rceil})$

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Depth-first search

The deepest node in the current fringe of the search tree is expanded first.

Implemented with a stack (LIFO queue)



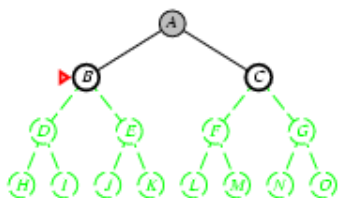
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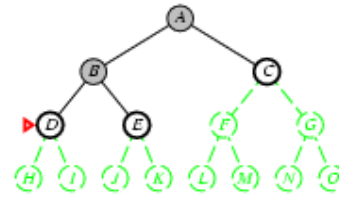
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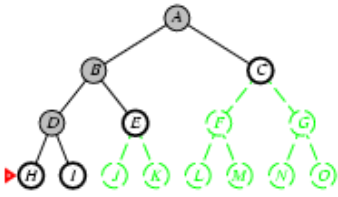
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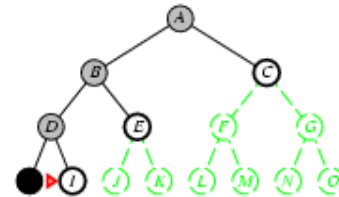
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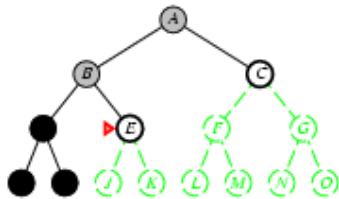
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Depth-first search

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Implemented with a stack (LIFO queue)



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Judging DFS

- **Complete:** No
 - Might get stuck going down a long path
- **Optimal:** No
 - Might return a solution which is at greater depth (i.e. more costly) than another solution
- **Time:** $O(b^m)$
 - m might be larger than d
- **Space:** $O(bm)$ 😊

Do not use DFS if you suspect a large tree depth

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Depth-limited search

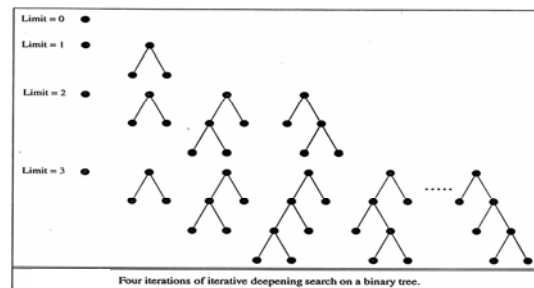
- We can avoid the problem of unbounded trees by using a **depth limit, l**
 - All nodes at depth l are treated as though they have no successors
 - If possible, choose l based on knowledge of the problem
- **Time:** $O(b^l)$
- **Memory:** $O(bl)$
- **Complete?:** No
- **Optimal?:** No

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Iterative-deepening

- General strategy that repeatedly does depth-limited search, but increases the limit each time



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Iterative-deepening

IDS is not as wasteful as one might think.

Note, most nodes in a tree are at the bottom level. It does not matter if nodes at a higher level are generated multiple times.

Breadth first search :

$$1 + b + b^2 + \dots + b^{d-1} + b^d$$

E.g. $b=10, d=5: 1+10+100+1,000+10,000+100,000 = 111,111$

Iterative deepening search :

$$(d+1)*1 + (d)*b + (d-1)*b^2 + \dots + 2b^{d-1} + 1b^d$$

E.g. $6+50+400+3000+20,000+100,000 = 123,456$

Complete, Optimal, $O(b^d)$ time, $O(bd)$ space

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Summary

- Problem formulation usually requires abstracting away real-world details to define a state space that can feasibly be explored
- Variety of uninformed search strategies
 - Assume no knowledge about the problem) general but expensive
 - Mainly differ in the order in which they consider the states

Criteria	BFS	Uniform	DFS	DLS	IDS
Complete	Yes	Yes	No	No	Yes
Time	$O(b^d)$	$O(b^{\text{celling}(C^*/a)})$	$O(b^m)$	$O(b^l)$	$O(b^d)$
Space	$O(b^d)$	$O(b^{\text{celling}(C^*/a)})$	$O(bm)$	$O(bl)$	$O(bd)$
Optimal	Yes	Yes	No	No	Yes

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Summary

- Iterative deepening uses only **linear space** and not much more time than other uninformed search algorithms
 - Use IDS when there is a large state space and the maximum depth of the solution is unknown
- Things to think about:
 - What about searching graphs?
 - Repeated states?

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Next Class

- Informed search techniques
- Russell & Norvig: Sections 4.1-4.2

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