

Adversarial Search

CS 486 /686
May 19, 2005
University of Waterloo

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Introduction

- So far we have studied environments where there is only a single-agent
- Today we look at what happens if we are in a setting where there are multiple agents planning against each other
 - Game theory: zero sum games

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Outline

- Games
- Minimax search
- Evaluation functions
- Alpha-beta pruning
- Coping with chance
- Game programs

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Games

- Games are one of the oldest, most well-studied domains in AI
- Why?
 - They are fun
 - Games are usually easy to represent and the rules are clear
 - State spaces can be very large (so more challenging than "toy problems")
 - In chess the search tree has $\sim 10^{154}$ nodes
 - Like the "real world" in that decisions **have** to be made and time is vitally important
 - Easy to determine when a program is doing well
 - i.e. it wins

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Types of games

- Perfect vs imperfect information
 - Perfect info means that you can see the entire state of the game
 - Chess, checkers, othello, go,...
 - Imperfect info games include scrabble, poker, most card games
- Deterministic vs stochastic
 - Chess is deterministic
 - Backgammon is stochastic

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Games as search problems

- Consider a 2-player perfect information game
 - **State**: board configuration plus the player who's turn it is to move
 - **Successor function**: given a state returns a list of (move,state) pairs, indicating a legal move and the resulting board
 - **Terminal state**: states where there is a win/loss/draw
 - **Utility function**: assigns a numerical value to terminal states (e.g. In chess +1 for a win, -1 for a loss, 0 for a draw)
 - **Solution**: a strategy (way of picking moves) that wins the game

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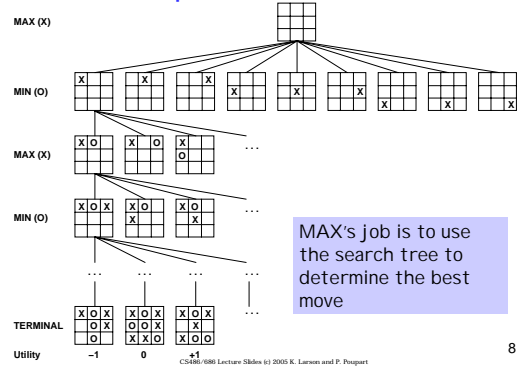
Game search challenge

- What makes game search challenging?
 - There is an opponent!
 - The opponent is malicious - it wants to win (i.e. it is trying to make you lose)
 - We need to take this into account when choosing moves
 - Simulate the opponent's behaviour in our search
- Notation: One player is called **MAX** (who wants to maximize its utility) and one player is called **MIN** (who wants to minimize its utility)

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Example: Tic-Tac-Toe



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Optimal strategies

- In standard search the optimal solution is a sequence of moves leading to a winning terminal state
- But MIN has something to say about this
- **Strategy** (from MAX's perspective):
 - Specify a move for the initial state, specify a move for all possible states arising from MIN's response, then all possible responses to all of MIN's responses to MAX's previous move....

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Optimal strategies

- Want to find the optimal strategy
 - One that leads to outcomes at least as good as any other strategy, given that MIN is playing optimally
 - Equilibrium (game theory)
 - Zero-sum games of perfect information are "easy games" from a game theoretic perspective

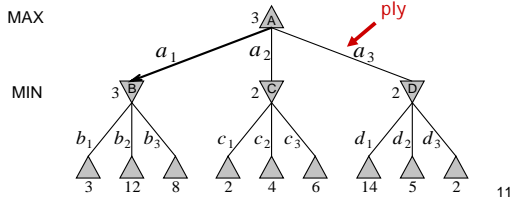
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Minimax Value

MINIMAX-VALUE(n) =

- $$\begin{cases} \text{Utility}(n) & \text{if } n \text{ is a terminal state} \\ \max_{s \in \text{Succ}(n)} \text{MINIMAX-VALUE}(s) & \text{if } n \text{ is a MAX node} \\ \min_{s \in \text{Succ}(n)} \text{MINIMAX-VALUE}(s) & \text{if } n \text{ is a MIN node} \end{cases}$$



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Minimax algorithm

function MINIMAX-DECISION($state$) *returns an action*

$v \leftarrow$ MAX-VALUE($state$)

return the action in SUCCESSIONS($state$) with value v

function MAX-VALUE($state$) *returns a utility value*

if TERMINAL-TEST($state$) **then** **return** UTILITY($state$)

$v \leftarrow -\infty$

for a, s in SUCCESSIONS($state$) **do**

$v \leftarrow$ MAX(v , MIN-VALUE(s))

return v

function MIN-VALUE($state$) *returns a utility value*

if TERMINAL-TEST($state$) **then** **return** UTILITY($state$)

$v \leftarrow \infty$

for a, s in SUCCESSIONS($state$) **do**

$v \leftarrow$ MIN(v , MAX-VALUE(s))

return v

Returns action corresponding to best possible move

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Properties of Minimax

- Complete if tree is finite
- Time complexity: $O(b^m)$
- Space complexity: $O(bm)$ (it is DFS)
- Optimal against an optimal opponent

m is depth of the tree

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Properties of Minimax

- Complete if tree is finite
- Time complexity: $O(b^m)$
- Space complexity: $O(bm)$ (it is DFS)
- Optimal against an optimal opponent
 - If MIN does not play optimally then we might be able to do better following a different strategy

m is depth of the tree

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Minimax and multi-player games

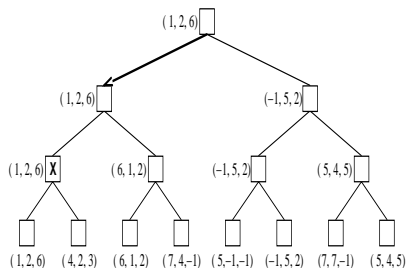
to move

A

B

C

A



Can not handle alliances, sidepayments....

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- Can we now write a program that will play chess reasonably well?

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- Can we now write a program that will play chess reasonably well?
 - For chess $b \sim 35$ and $m \sim 100$
 - Do we really need to look at all those nodes?

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Alpha-Beta Pruning

- No!
 - If we are smart (and lucky) we can do **pruning**
 - Eliminate large parts of the tree from consideration
- Alpha-Beta pruning applied to a minimax tree
 - Returns the same decision as minimax
 - Prunes branches that cannot influence final decision

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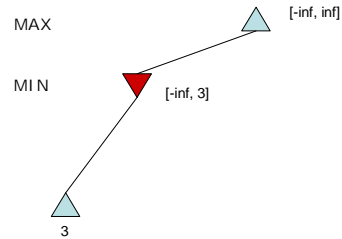
Alpha-Beta Pruning

- Alpha:
 - Value of best (highest value) choice we have found so far on the path for MAX
- Beta:
 - Value of best (lowest value) choice we have found so far on path for MIN
- Update alpha and beta as search continues
- Prune as soon as the value of the current node is known to be worse than current alpha or beta values for MAX or MIN

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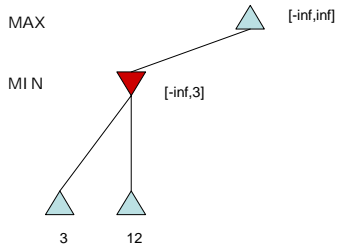
Alpha-Beta example



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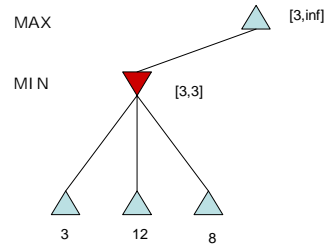
Alpha-Beta example



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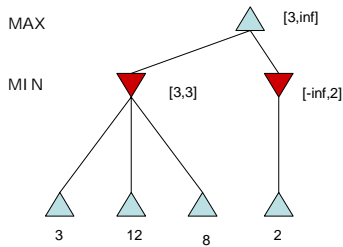
Alpha-Beta example



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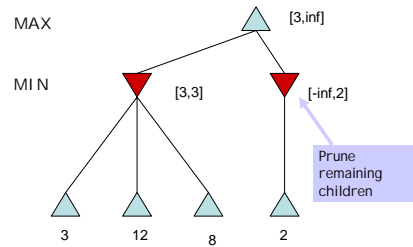
Alpha-Beta example



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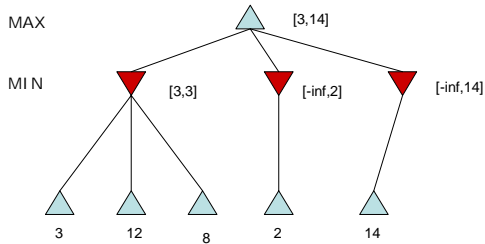
Alpha-Beta example



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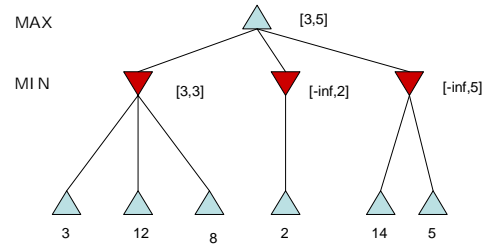
Alpha-Beta example



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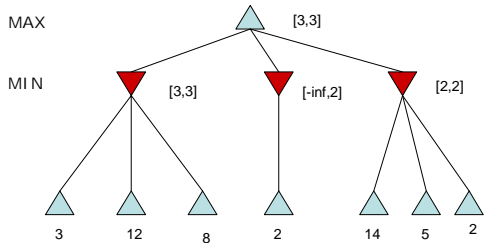
Alpha-Beta example



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Alpha-Beta example



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Properties of Alpha-Beta

- Pruning does not affect the final result
 - You prune parts of the tree that you would never reach in actual play
- The order in which moves are evaluated are important
 - With bad move ordering will prune nothing
 - With perfect node ordering can reduce time complexity to $O(b^{m/2})$

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Real-time decisions

- Alpha-beta can be a huge improvement over minimax
 - Still not good enough as we need to search all the way to terminal states for at least part of search space
 - Need to make a decision about a move quickly
- Heuristic **evaluation function** + **cutoff test**

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Evaluation functions

- Apply an evaluation function to a state
 - If terminal state, function returns actual utility
 - If non-terminal, function returns estimate of the expected utility (i.e. the chance of winning from that state)
- Function must be fast to compute

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Evaluation functions

- Evaluation functions can be given by the designer of the program (using expert knowledge) or learned from experience
- If features can be judged independently, a **weighted linear function** is good
 - $w_1f_1(s)+w_2f_2(s)+\dots+w_nf_n(s)$ with s as board state

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Cutting off search

- Instead of searching until we find a terminal state, we can cut search sooner and apply the evaluation function
- When?
 - Arbitrarily (but deeper is better)
 - Quiescent states
 - States that are "stable" - not going to change value (by a lot) in the near future
 - Singular extensions
 - Searching deeper when you have a move that is "clearly better" (i.e. moving the king out of check)
- Can be used to avoid the **horizon effect**

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Cutting off search

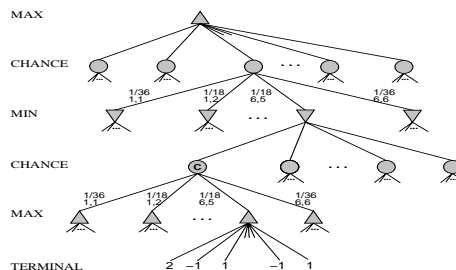
- How deep do we need to search?
 - Novice chess human player
 - 5-ply (minimax)
 - Master chess human player
 - 10-ply (alpha-beta)
 - Grandmaster chess human player
 - 14-ply + a fantastic evaluation function, opening and endgame databases, ..., special purpose hardware would be nice but is no longer really needed (Fritz)

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Stochastic games

- In games like Backgammon chance plays a roll



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Stochastic games

- Need to consider best/worst cases + probability they will occur
- Recall:** Expected value of a random variable x

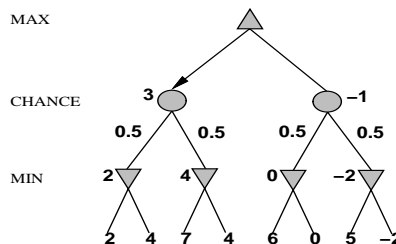
$$E[x]=\sum_{x \in X} P(x)x$$

- Expectiminimax** is like minimax but at chance nodes compute the **expected value**

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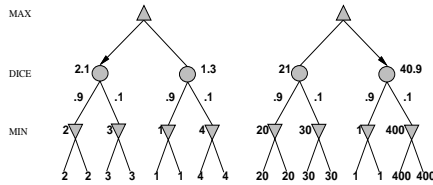
Expectiminimax



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Expectiminimax



WARNING: exact values do matter! Order-preserving transformations of the evaluation function can change the choice of moves. Must have **positive linear transformations** only

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Some Game Programs

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Checkers: Tinsley vs. Chinook



Mr. Tinsley suffered his 4th and 5th losses ever against Chinook

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Checkers

- Chinook: <http://www.cs.ualberta.ca/~chinook>
 - World Man-Machine Checkers Champion
 - Alpha-beta search
 - Opening database
 - Its secret weapon: **Endgame database**
 - Precomputed database of all 444 billion positions with 8 or fewer pieces, each with perfect win/loss/draw info
 - Perfect knowledge into the search
 - Checkers is now dominated by computers

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Chess: Kasparov vs. Deep Blue

Kasparov		Deep Blue
5'10"	Height	6' 5"
176 lbs	Weight	2,400 lbs
34 years	Age	4 years
50 billion neurons	Computers	32 RI SC processors + 256 VLSI chess engines
2 pos/sec	Speed	200,000,000 pos/sec
Extensive	Knowledge	Primitive
Electrical/chemical	Power Source	Electrical
Enormous	Ego	None

1997: Deep Blue wins by 3 wins, 1 loss, and 2 draws

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Chess

- Its secret:
 - Specialized chess processor + special-purpose memory optimization
 - Very sophisticated evaluation function
 - Expert features and hand-tuned weights
 - Opening and closing books
 - Alpha-beta + improvements (searching up to 40 ply deep!)
 - Search over 200 million positions per second (though lots of these possible moves are silly moves by human standards..)

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Chess

- There are now programs running on PCs that are on par with human champions
 - Deep Junior vs Kasparov in 2003: 3/3 tie
 - Deep Junior: 8 CPU, 8GB RAM, Windows 2000, 2000000 pos/second
- Is Chess still a human game or have computers conquered it?

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Backgammon

- TD-Gammon (Gerry Tesauro at IBM)
- One of the top players in the world
- But only searches two moves ahead!
- Its secret: One amazing evaluation function
 - Neural network trained with reinforcement learning during ~1million games played against itself
 - Humans play backgammon differently now, based on what TD-Gammon learned about the game
 - Very cool AI ☺



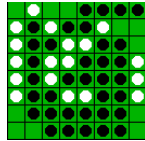
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Othello: Murakami vs. Logistello



Takeshi Murakami
World Othello Champion



1997: The Logistello software crushed Murakami by 6 games to 0

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Othello/Reversi

- Logistello (Michael Buro from U of Alberta)
- Human world champion crushed by the program
 - Humans no match for machine
- Its secret: Evaluation function
 - Automatically discovered and tuned knowledge
 - Samples patterns to see if its presence in a position can be correlated with success
 - Tuned 1.5 million parameters using self-play games with feedback

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Bridge

- GIB (Matt Ginsberg - U of Oregon)
 - World's first expert level bridge playing program (Finished 12th in human world championship in 1998)
 - Humans are still doing better, but the gap is narrowing quickly
- Its secrets:
 - Does simulations for each decision
 - Deals cards to opponents consistent with available information
 - Chooses action that maximizes expected return
 - Plus other tricks...

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Go: Goemate vs. ??



Name: Chen Zhixing
Profession: Retired
Computer skills:
self-taught programmer
Author of Goemate (one of the best
Go programs available today)



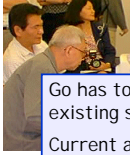
Gave Goemate a 9 stone
handicap and still easily
beat the program,
thereby winning \$15,000

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Go: Goemate vs. ??



Name: Chen Zhixing
Profession: Retired
Computer skills:

Go has too high a branching factor for existing search techniques ($b \sim 100$)

Current and future software must rely on huge databases and pattern-recognition techniques

Need to make strategic decisions – Which battle is worth fighting?

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Summary

- Games pose lots of fascinating challenges for AI researchers
- Minimax search allows us to play optimally against an optimal opponent
- Alpha-beta pruning allows use to reduce the search space
- A good evaluation function is key to doing well
- Games are fun

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Next class

- We will begin reasoning under uncertainty
– Chapter 13

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