

# Solving Problems by Searching

CS486/686  
University of Waterloo  
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## Outline

- Problem solving agents and search
- Examples
- Properties of search algorithms
- Uninformed search
  - Breadth first
  - Depth first
  - Iterative Deepening

## Introduction

- Search was one of the first topics studied in AI
  - Newell and Simon (1961) *General Problem Solver*
- Central component to many AI systems
  - Automated reasoning, theorem proving, robot navigation, VLSI layout, scheduling, game playing,...

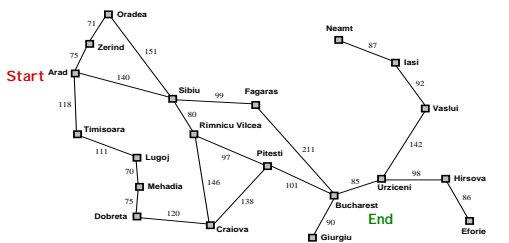
## Problem-solving agents

```

function SIMPLE-PROBLEM-SOLVING-AGENT(percept) returns an action
  static: seq, an action sequence, initially empty
         state, some description of the current world state
         goal, a goal, initially null
         problem, a problem formulation

  state ← UPDATE-STATE(state, percept)
  if seq is empty then do
    goal ← FORMULATE-GOAL(state)
    problem ← FORMULATE-PROBLEM(state, goal)
    seq ← SEARCH(problem)
  action ← FIRST(seq)
  seq ← REST(seq)
  return action
    
```

## Example: Traveling in Romania



**Formulate Goal**  
Get to Bucharest

**Formulate Problem**  
Initial state: In(Arad)

**Find a solution**  
Sequence of cities: Arad, Sibiu, Fagaras, Bucharest  
Goal Test: In(Bucharest)?  
Path cost: Distance between cities

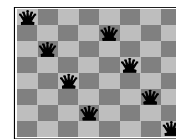
## Examples of Search Problems



Start State



Goal State



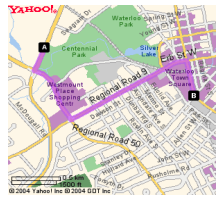
**States:** Locations of 8 tiles and blank  
**Initial State:** Any state  
**Succ Func:** Generates legal states that result from trying 4 actions (blank up, down, left, right)

**Goal test:** Does state match desired configuration  
**Path cost:** Number of steps

**States:** Arrangement of 0 to 8 queens on the board  
**Initial State:** No queens on the board  
**Succ Func:** Add a queen to an empty space

**Goal test:** 8 queens on board, none attacked  
**Path cost:**

## More Examples



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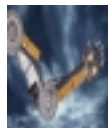
## Common Characteristics

- All of those examples are
  - Fully observable
  - Deterministic
  - Sequential
  - Static
  - Discrete
  - Single agent
- Can be tackled by **simple** search techniques

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## Cannot tackle these yet...



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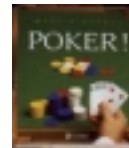
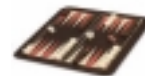
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## Cannot tackle these yet...

Chance

Infinite number of states

Games against an adversary



Hidden states



All of the above

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## Searching

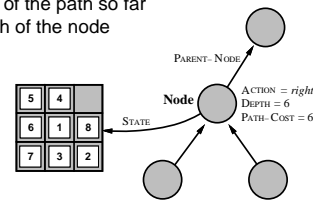
- We can formulate a search problem
  - Now need to find the solution
- We can visualize a state space search in terms of trees or graphs
  - Nodes correspond to states
  - Edges correspond to taking actions
- We will be studying **search trees**
  - These trees are constructed “on the fly” by our algorithms

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## Data Structures for Search

- Basic data structure: **Search Node**
  - State
  - Parent node and operator applied to parent to reach current node
  - Cost of the path so far
  - Depth of the node

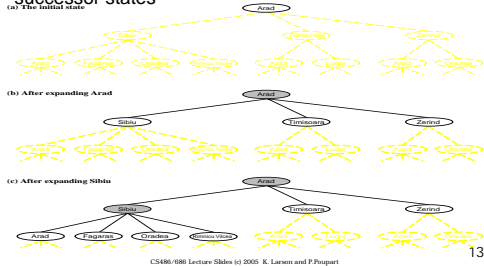


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## Expanding Nodes

- Expanding a node
  - Applying all legal operators to the state contained in the node and generating nodes for all corresponding successor states



## Generic Search Algorithm

- Initialize search algorithm with initial state of the problem
- Repeat**
  - If no candidate nodes can be expanded, **return failure**
  - Choose leaf node for expansion, according to **search strategy**
  - If node contains a goal state, **return solution**
  - Otherwise, expand the node, by applying legal operators to the state within the node. Add resulting nodes to the tree

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## Implementation Details

- We need to only keep track of nodes that need to be expanded (**fringe**)
    - Done by using a (prioritized) **queue**
- Initialize queue by inserting the node corresponding to the initial state of the problem
  - Repeat**
    - If queue is empty, **return failure**
    - Dequeue a node
    - If the node contains a goal state, **return solution**
    - Otherwise, expand node by applying legal operators to the state within. Insert resulting nodes into queue

**Search algorithms differ in their queuing function!**

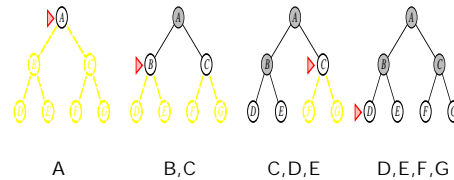
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## Breadth-first search

All nodes on a given depth are expanded before any nodes on the next level are expanded.

Implemented with a FIFO queue



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## Evaluating search algorithms

- Completeness:** Is the algorithm guaranteed to find a solution if a solution exists?
- Optimality:** Does the algorithm find the optimal solution (Lowest path cost of all solutions)
- Time complexity**
- Space complexity**

Variables

b	Branching factor
d	Depth of shallowest goal node
m	Maximum length of any path in the state space

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## Judging BFS

- Complete:** Yes, if b is finite
- Optimal:** Yes, if all costs are the same
- Time:**  $1+b+b^2+b^3+\dots+b^d = O(b^d)$
- Space:**  $O(b^d)$

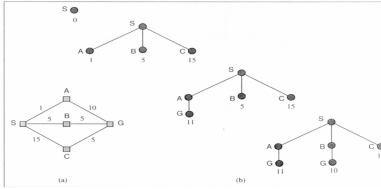
All uninformed search methods will have exponential time complexity ☹

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## Uniform Cost Search

- A variation of breadth-first search
  - Instead of expanding shallowest node it expands the node with lowest path cost
  - Implemented using a priority queue



$C^*$  is cost of optimal solution

$\epsilon$  is minimum action cost

**Optimal**  
**Complete** if  $\epsilon > 0$

**Time:**  $O(b^{\lceil C^*/\epsilon \rceil})$   
**Space:**  $O(b^{\lceil C^*/\epsilon \rceil})$

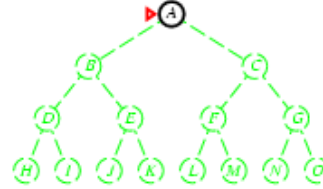
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## Depth-first search

The deepest node in the current fringe of the search tree is expanded first.

Implemented with a stack (LIFO queue)



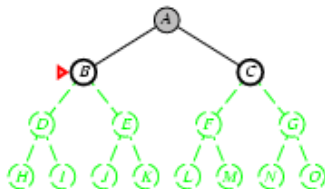
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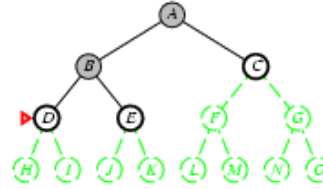
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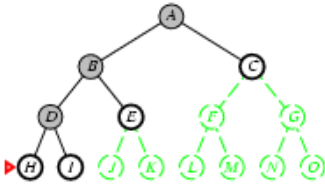
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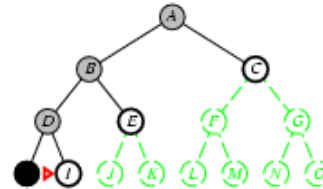
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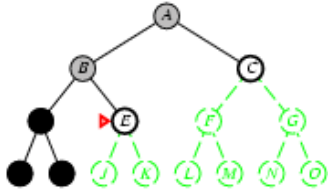
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## Depth-first search

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Implemented with a stack (LIFO queue)



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## Judging DFS

- **Complete:** No
  - Might get stuck going down a long path
- **Optimal:** No
  - Might return a solution which is at greater depth (i.e. more costly) than another solution
- **Time:**  $O(b^m)$ 
  - $m$  might be larger than  $d$
- **Space:**  $O(bm)$  😊

Do not use DFS if you suspect a large tree depth

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## Depth-limited search

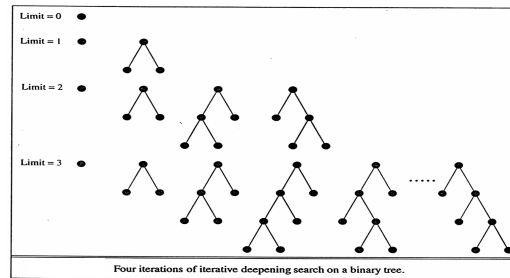
- We can avoid the problem of unbounded trees by using a **depth limit,  $l$** 
  - All nodes at depth  $l$  are treated as though they have no successors
  - If possible, choose  $l$  based on knowledge of the problem
- **Time:**  $O(b^l)$
- **Memory:**  $O(b^l)$
- **Complete?:** No
- **Optimal?:** No

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## Iterative-deepening

- General strategy that repeatedly does depth-limited search, but increases the limit each time



Four iterations of iterative deepening search on a binary tree.

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## Iterative-deepening

IDS is not as wasteful as one might think.

Note, most nodes in a tree are at the bottom level. It does not matter if nodes at a higher level are generated multiple times.

**Breadth first search :**

$$1 + b + b^2 + \dots + b^{d-1} + b^d$$

E.g.  $b=10, d=5: 1+10+100+1,000+10,000+100,000 = 111,111$

**Iterative deepening search :**

$$(d+1)*1 + (d)*b + (d-1)*b^2 + \dots + 2b^{d-1} + 1b^d$$

E.g.  $6+50+400+3000+20,000+100,000 = 123,456$

Complete, Optimal,  $O(b^d)$  time,  $O(bd)$  space

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## Summary

- Problem formulation usually requires abstracting away real-world details to define a state space that can feasibly be explored
- Variety of uninformed search strategies
  - Assume no knowledge about the problem ) general but expensive
  - Mainly differ in the order in which they consider the states

Criteria	BFS	Uniform	DFS	DLS	IDS
<b>Complete</b>	Yes	Yes	No	No	Yes
<b>Time</b>	$O(b^d)$	$O(b^{\lceil C^2/k \rceil})$	$O(b^m)$	$O(b^l)$	$O(b^d)$
<b>Space</b>	$O(b^d)$	$O(b^{\lceil C^2/k \rceil})$	$O(bm)$	$O(bl)$	$O(bd)$
<b>Optimal</b>	Yes	Yes	No	No	Yes

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## Summary

- Iterative deepening uses only **linear space** and not much more time than other uninformed search algorithms
  - Use IDS when there is a large state space and the maximum depth of the solution is unknown
- Things to think about:
  - What about searching graphs?
  - Repeated states?

## Next Class

- Informed search techniques
- Russell & Norvig: Sections 4.1-4.2