

# CS475/CM375

## Lecture 2: Sept 15, 2011

Cholesky factorization, tridiagonal,  
band matrices

Reading: [TB] Chapt. 23 p. 172-176

CS475/CM375 (c) 2011 P. Poupart & J. Wan

1

## Special Linear Systems

- Exploit special structures of linear systems
- More efficient  $LU$  factorization
- Symmetric systems
  - $LDM^T$  factorization (variant of  $LU$ )
- Symmetric positive definite systems
  - $GG^T$  factorization (a.k.a. Cholesky factorization)

CS475/CM375 (c) 2011 P. Poupart & J. Wan

2

## $LDM^T$ factorization

- Theorem: If all the leading principal submatrices of  $A$  are nonsingular, then there exist unique unit lower  $\Delta$  matrices  $L$  and  $M$ , and a unique diagonal matrix  $D$  such that  $A = LDM^T$ .
- Partial Proof:
  - Factor  $A = LU$
  - Define  $D = \text{diag}(d_1, \dots, d_n)$ ,  $d_i = u_{ii}$   $i = 1, \dots, n$
  - Let  $M^T = D^{-1}U = \text{unit upper } \Delta$  ( $M = \text{unit lower } \Delta$ )
  - Thus  $A = LU = LD(D^{-1}U) = LDM^T$
- Note:  $\text{flops}(LU) = \text{flops}(LDM^T)$

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

3

## Symmetric systems

- Theorem: If  $A$  is symmetric, then  $A = LDL^T$
- Proof:
  - By previous result,  $A = LDM^T$   
 $\Rightarrow M^{-1}AM^{-T} = M^{-1}LDM^TM^{-T} = M^{-1}LD$
  - Since  $M^{-1}AM^{-T}$  is symmetric, so is  $M^{-1}LD$
  - Also,  $M^{-1}L$  is lower  $\Delta \Rightarrow M^{-1}LD$  is lower  $\Delta$
  - So  $M^{-1}LD$  is both lower  $\Delta$  and symmetric  
 $\Rightarrow M^{-1}LD$  is diag  $\Rightarrow M^{-1}L$  is diag
  - Since  $M^{-1}L$  is also unit lower  $\Delta$ ,  
 then  $M^{-1}L = I \Rightarrow M = L$

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

4

## Symmetric systems

- Notes
  1. We can save about half the work by computing  $L$  and  $D$  only.
  2. One way is to compute the  $U$  factor only during the  $LU$  factorization.

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

5

## Positive definite systems

- Definition:  $A$  is positive definite iff  $x^T A x > 0$  for all  $x \neq 0$ .
- Properties of positive definite matrices:

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

6

## Positive definite systems

- Theorem: If  $A \in \mathbb{R}^{n \times n}$  is PD and  $X \in \mathbb{R}^{n \times k}$  has rank  $k$ , then  $B = X^T A X \in \mathbb{R}^{k \times k}$  is also PD
- Proof:
  - Let  $z \in \mathbb{R}^{k \times 1}$ . Then  $z^T B z = z^T X^T A X z = (Xz)^T A (Xz)$
  - If  $Xz = 0$ , then  $X$  is not rank  $k$ .
  - Hence  $z^T B z > 0$ .
- Corollary: If  $A$  is PD, then all its principal submatrices are PD. In particular, all diag entries are positive.

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

7

## Positive definite systems

- Corollary: If  $A$  is PD, then  $A = L D M^T$  and  $D$  has positive diag entries.
- Proof:
  - Let  $X = L^{-T}$ .  
Then  $X^T A X = L^{-1} (L D M^T) L^{-T} = D M^T L^{-T}$  is PD.
  - By previous corollary,  $\text{diag}(D M^T L^{-T})$  has positive entries.
  - Note that  $M^T$  and  $L^{-T}$  are unit upper  $\Delta$ .  
 $\Rightarrow M^T L^{-T}$  is also unit upper  $\Delta$   
 $\Rightarrow \text{diag}(D M^T L^{-T}) = D$

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

8

## Symmetric positive definite systems

- Theorem: If  $A$  is SPD, then there exists unique lower  $\Delta$   $G$  such that

$$A = GG^T$$

- Proof:

–  $A = LDL^T$  and  $D = \text{diag}(d_1, \dots, d_n)$ ,  $d_i > 0$ .

– Define  $D^{\frac{1}{2}} \equiv \text{diag}(\sqrt{d_1}, \dots, \sqrt{d_n})$

– Let  $G = LD^{\frac{1}{2}}$ . Then  $G$  is lower  $\Delta$

$$\Rightarrow GG^T = LD^{\frac{1}{2}}(LD^{\frac{1}{2}})^T = LD^{\frac{1}{2}}D^{\frac{1}{2}}L^T = LDL^T = A$$

## Symmetric positive definite systems

- Examples

## Cholesky factorization

- $A = GG^T$  is called the Cholesky factorization of  $A$  and the lower  $\Delta G$  is called the Cholesky factor.

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

11

## Cholesky factorization

- Algorithm big picture

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

12

## Cholesky factorization

```

For  $k = 1, 2, \dots, n$ 
   $a_{kk} = \sqrt{a_{kk}}$ 
  For  $j = k + 1, \dots, n$ 
     $a_{jk} = a_{jk} / a_{kk}$ 
  End
  For  $j = k + 1, \dots, n$ 
    for  $i = j, \dots, n$ 
       $a_{ij} = a_{ij} - a_{ik}a_{jk}$ 
    end
  End
End
 $flops(Cholesky) \approx \frac{n^3}{3}$ 

```

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

13

## Banded systems

- Definition:  $A$  has upper bandwidth  $q$  if  $a_{ij} = 0 \forall j > i + q$  and lower bandwidth  $p$  if  $a_{ij} = 0 \forall i > j + p$ .
- Picture

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

14

## Banded systems

- If  $A$  is banded, so are  $LU, GG^T, LDM^T$
- Theorem: Let  $A = LU$ . If  $A$  has upper bandwidth  $q$  and lower bandwidth  $p$ , then  $U$  has upper bandwidth  $q$  and  $L$  has lower bandwidth  $p$ .
- Picture

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

15

## Band Gaussian Elimination

```

For  $k = 1, 2, \dots, n - 1$ 
  For  $i = k + 1, \dots, \min(k + p, n)$ 
     $a_{ik} = a_{ik} / a_{kk}$ 
  end
  for  $i = k + 1, \dots, \min(k + p, n)$ 
    for  $j = k + 1, \dots, \min(k + q, n)$ 
       $a_{ij} = a_{ij} - a_{ik} a_{kj}$ 
    end
  end
End
If  $n \gg p$  and  $n \gg q$ , then  $\text{flops}(\text{band GE}) \approx 2npq$ 

```

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

16



## Tridiagonal systems

- Assume  $A$  is tridiagonal and symmetric
- Then
  - $L =$

$$- D =$$

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

17

## Tridiagonal System

- $A = LDL^T$  implies

$$\begin{aligned}
 a_{kk} &= (LDL^T)_{kk} \\
 &= \sum_i \sum_j l_{ki} d_{ij} l_{jk}^T \\
 &= \sum_i l_{ki} d_{ii} l_{ik}^T \\
 &= \sum_i l_{ki}^2 d_{ii} \\
 &= l_{k,k-1}^2 d_{k-1,k-1} + l_{kk}^2 d_{kk} \\
 &\quad (i = k-1) \\
 &= l_{k-1}^2 d_{k-1} + d_k
 \end{aligned}$$

$$\begin{aligned}
 a_{k,k-1} &= \sum_i \sum_j l_{ki} d_{ij} l_{j,k-1}^T \\
 &= \sum_i l_{ki} d_{ii} l_{i,k-1}^T \\
 &= \sum_i l_{ki} d_{ii} l_{k-1,i} \\
 &= l_{k,k-1} d_{k-1,k-1} l_{k-1,k-1} \\
 &\quad (i = k-1) \\
 &= l_{k-1} d_{k-1}
 \end{aligned}$$

CS475/CM375 (c) 2011 P. Poupart &amp; J. Wan

18

## Tridiagonal Factorization

- Algorithm

$$d_1 = a_{11}$$

for  $k = 2, \dots, n$

$$l_{k-1} = a_{k,k-1}/d_{k-1}$$

$$d_k = a_{kk} - l_{k-1}a_{k,k-1} \quad (l_{k-1}d_{k-1} = a_{k,k-1})$$

end

- $\text{flops}(\text{tridiag}) = O(n)$