Dominators

Definition

In a CFG, node a dominates b if every path from the start node to b passes through a. Node a is a dominator of b.

Property

The dominance relation is a partial order.

Definition

Node a strictly dominates b if $a \neq b$ and a dominates b.

Dominators

Theorem

IF a and b both dominate c,
THEN either a dominates b or b dominates a.

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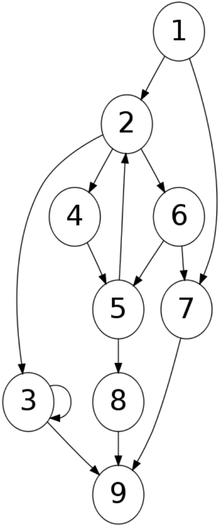
Corollary

Every node n has at most one immediate dominator idom(n) such that

- idom(n) $\neq n$
- idom(n) dominates n, and
- idom(n) does not dominate any other dominator of n.

Dominator Example



Computing Dominators

As a dataflow analysis

1. Forwards
2. Lattice is $(\mathcal{P}(Stmts), \supseteq)$
3. $\cap$
4. $\text{out}_\ell = \text{in}_\ell \cup \{\ell\}$
5. start node value is $\{\}$
6. $\bot = \{\text{all statements}\}$

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More efficient approaches

- Lengauer-Tarjan: see Appel book section 19.2
- Cooper, Harvey, Kennedy: `http://www.hipersoft.rice.edu/grads/publications/dom14.pdf`

Dominance Frontier

Definition

A node w is in the dominance frontier of x if:

- x does not strictly dominate w, and
- x dominates a predecessor of w.

Computing Dominance Frontier

$DF_{local}(x)$: the successors of x not strictly dominated by x.
$DF_{up}(y)$: nodes in $DF(y)$ not strictly dominated by $\mathrm{idom}(y)$.
$DF(x) = DF_{local}(x) \cup \bigcup_{\{y \mid \mathrm{idom}(y)=x\}} DF_{up}(y)$.

Computing Dominance Frontier

Algorithm DF(x):

$S = \{\}$
for all nodes $w \in \text{succ}(x)$ **do**
    if $\text{idom}(w) \neq x$ **then**
        $S \cup = \{w\}$
/* *S is now $DF_{local}(x)$* */
for all nodes y for which $\text{idom}(y) = x$ **do**
    /* *below we compute $DF_{up}(y)$* */
    for all nodes $w \in DF(y)$ **do**
        if x does not dominate w or $x = w$ **then**
            $S \cup = \{w\}$
return S

Computing Dominance Frontier (Alternative)

Restatement of definition of DF

$w \in DF(x)$ for every x that dominates a predecessor of w, but does not strictly dominate w.

Algorithm COMPUTE DFS():

for all nodes w **do**
    for all $p \in \text{preds}(w)$ **do**
        $x = p$
        while $x \neq \text{idom}(w)$ **do**
            $DF(x) \cup= \{w\}$
            $x = \text{idom}(x)$