

# CS 744 - Advanced Compiler Design

## Assignment 3

### Part 1: Partial Redundancy Elimination (20 marks)

In this part of the assignment, you will perform partial redundancy elimination (the variation due to Drechsler and Stadel [1] that was presented in class) on the following program. See the lecture slides from the course web page for details.

```
1 read(a, b, c, d, e, f);
2 c = 1;
3 if(a < 0) goto L1;
4 a = c + 2;
5 d = a + b;
6 b = d - 3;
7 if(d < 0) goto L2;
8 goto L3;
9 L1:
10 if(c >= 10) goto L2;
11 e = a + b;
12 Z[i] = e;
13 c = c + 1;
14 goto L1;
15 L2:
16 f = a + b;
17 write(f);
18 L3:
```

1. Identify the basic blocks, and draw a basic block graph.
2. Determine which basic blocks are **transparent** for the expression  $a + b$ , and in which basic blocks the expression is **locally available** and **locally anticipable**.
3. At the beginning and end of each basic block, determine whether the expression  $a + b$  is (globally) **available** and (globally) **anticipable**.
4. Perform the **earliest placement** computation. For each edge in the basic block graph, state whether the edge is the earliest place in which the expression  $a + b$  should be computed.
5. Perform the **latest placement** computation. For each edge in the basic block graph, state whether the expression  $a + b$  could be computed ?later? on the edge.
6. Determine on which edges a computation of  $a + b$  should be **inserted** by PRE.

- Determine from which basic blocks the first computation of  $a + b$  should be **deleted** by PRE.
- Write the code that results after partial redundancy elimination.

References [1] Karl-Heinz Drechsler and Manfred P. Stadel. A variation of Knoop, Ruthing, and Steffen's lazy code motion. SIGPLAN Not., 28(5):29-38, 1993.

## Part 2 :Register Allocation (15 marks)

In this part of the assignment, you will perform register allocation on the following program.

```
1 read(a, b, c, d, e, f, g);
2 L1:
3 c = a + b
4 d = c * b
5 e = c / d
6 f = e - d
7 a = e * f
8 b = a - f
9 g = g + b
10 if g < 10 goto L1;
11 write(g);
```

- Draw the control flow graph.
- Determine the set of variables that are live before and after each instruction.
- Draw the interference graph for the program.
- Trace two possible runs of using Briggs' register colouring algorithm to colour the interference graph with at most three registers, such that one run succeeds, and the other fails.
- Is it possible to colour the interference graph with only two registers? If so, give such a colouring. If not, show why not.