

CS 489 / 698: Software and Systems Security

Module 2: Program Security (Defenses)

fuzz testing

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Outline

- 1 Introduction
- 2 Evolution: from the rain-fuzzer to modern fuzzing
- 3 Program state coverage: “natural selection” in the fuzzing world
- 4 Loops: another trouble maker for branch coverage
- 5 Concolic execution: forced path exploration
- 6 Conclusion

History: why we call it “fuzzing”?

In 80's, someone remotely logged into a unix system over a dial-up network link **during a storm**.

The rain caused a lot of **random noise** on the dial-up link.

And these noise caused applications that were using data off the dial-up network line to **crash**.

Gist of the story? — The rain tests the program way better than human beings.

The goal of fuzzing

Q: What is fuzzing doing essentially? Try to describe it in a way that is as abstract/general as possible.

A: To **drive** the execution of a **system** into **desired states**.

Elaborating the definition

- What is special about the target **system**?
 - Do we know the source code?
 - Do we know the input format?
 - What are the challenges when executing the “system”?
- What do we mean by a **state**?
 - How can we tell that one state is different from another?
- What do we mean by **desired**?
 - New/unseen behavior?
 - Closeness to targeted execution points?
- What do we mean by **driving** the execution?
 - What can possibly be one mutation?
 - How do you select the next mutation?

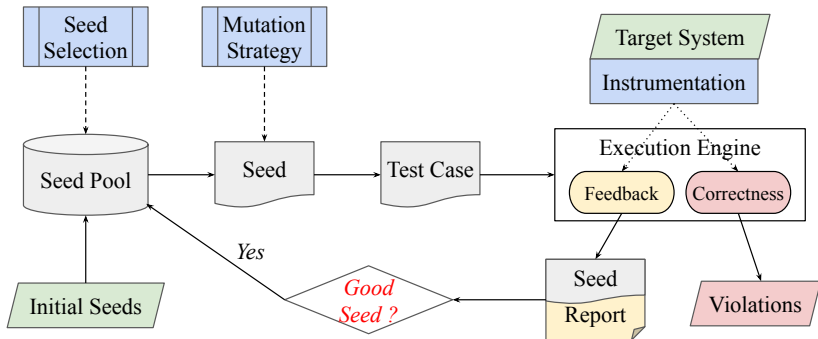
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Genetic algorithm

Training a program to play the snake game with genetic algorithm

Feedback-guided evolution process



Natural selection — survival of the fittest

Demo with AFL++

Acknowledgement: this demo is based on one of the examples used in the “[Fuzzing with AFL](#)” workshop by Michael Macnair.

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Intuition: what makes a high-quality seed?

```
1 pub fn foo(a: num, b: num) {
2   let c = if (a >= 0) {
3     1
4   } else {
5     2
6   };
7
8   // irrelevant operations
9
10  let d = if (b >= 0) {
11    2
12  } else {
13    3
14  };
15
16  // irrelevant operations
17
18  assert!(c != d);
19 }
```

Q: What is the testing plan?

- Cover every line?
- Cover every if-else branch?
- Cover every exit status?
- Cover every path?

⇒ if the fuzzer generates an input that **expands the coverage**, that input is a good seed.

Illustration of different coverage metrics

```
1 pub fn foo(a: num, b: num) {  
2   let c = if (a >= 0) {  
3     1  
4   } else {  
5     2  
6   };  
7  
8   // irrelevant operations  
9  
10  let d = if (b >= 0) {  
11    2  
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14  };  
15  
16  // irrelevant operations  
17  
18  assert!(c != d);  
19 }
```

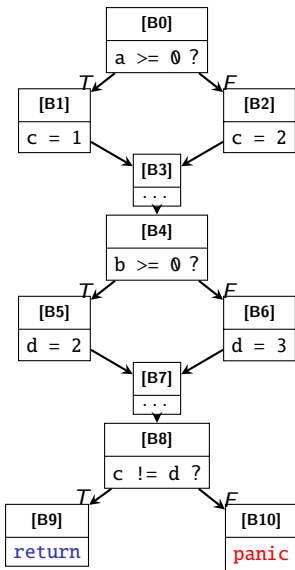
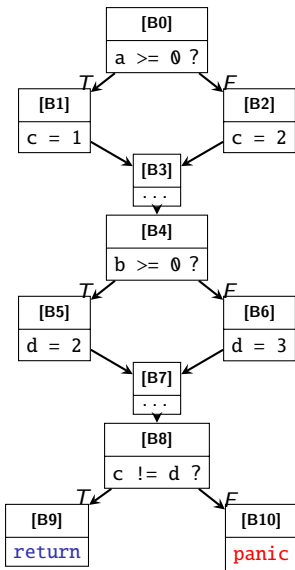


Illustration of different coverage metrics

- Cover every line?
 - Block coverage
- Cover every if-else branch?
 - Branch coverage
- Cover every exit status?
 - Return coverage
- Cover every path?
 - Path coverage



Path coverage: a theoretical optimum

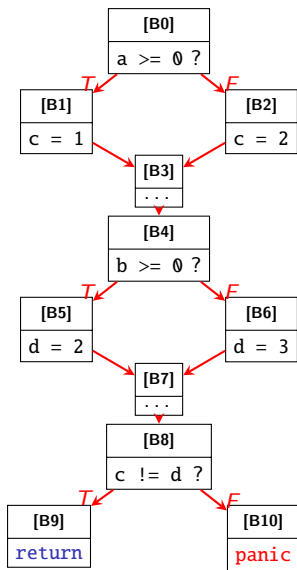
Claim: A program is **saturately tested** if we obtain a set of inputs that covers **every feasible path** of the program CFG.

NOTE: feasible paths include paths that leads to explicit and implicit panics.

Path coverage demo

- $a = 1, b = 1$
- $a = 1, b = -1$
- $a = -1, b = 1$
- $a = -1, b = -1$

No new program behaviors can be discovered \implies the program is saturately tested



Why not path coverage in practice?

Short answer: I don't know... AFL (American Fuzzy Lop) didn't adopt path coverage, so everyone follows suite...

Long answer:

- tracking block / branch coverage is **stateless** while tracking path coverage requires **stateful** instrumentations.
- different parts of the execution are not necessarily related, i.e., a new path does not necessarily mean interesting findings.
- it is hard to quantitatively measure the completeness of path coverage (because of infeasible paths). But by default, all branches should be somewhat feasible.

In practice, branch coverage hits a nice balance between effectiveness and easiness of instrumentation.

What's wrong with branch coverage?

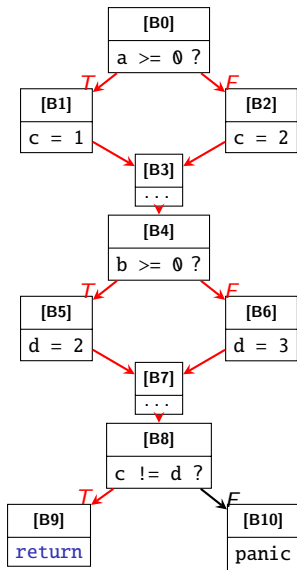
- $a = 1, b = 1$
- $a = -1, b = -1$

Two seeds already covered most of the branches.

- $a = 1, b = -1$

A seed that yields new path but is considered as a bad seed as it yields no new branch coverage.

⇒ fuzzer is not rewarded by mutating a and b , hence, lowering their priorities and the panic case may never be found, especially when fuzzing **complex** CFGs



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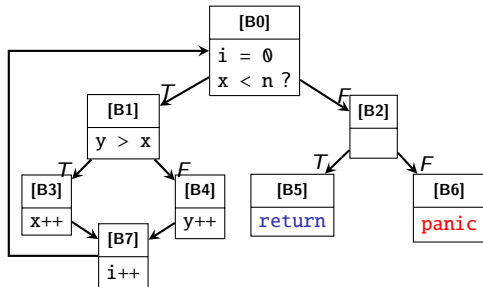
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Looping example

```

1  pub fn looping(
2    x: num,
3    y: num,
4    n: num
5  ) {
6    let i = 0;
7    while (x < n) {
8      if (y > x) {
9        x++;
10     }
11     else {
12       y++;
13     }
14     i++;
15  }
16  assert!(i != 7);
17  }

```



$y \leq x < n$

① $y++$ until $y == x$

② $y++$; $x++$ until $x == n$

$x < y \leq n$

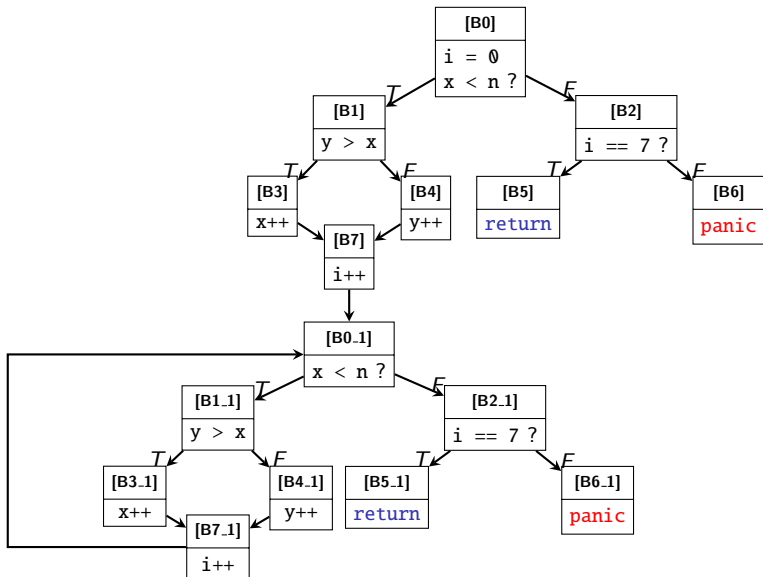
① $x++$ until $x == y$

② $y++$; $x++$ until $x == n$

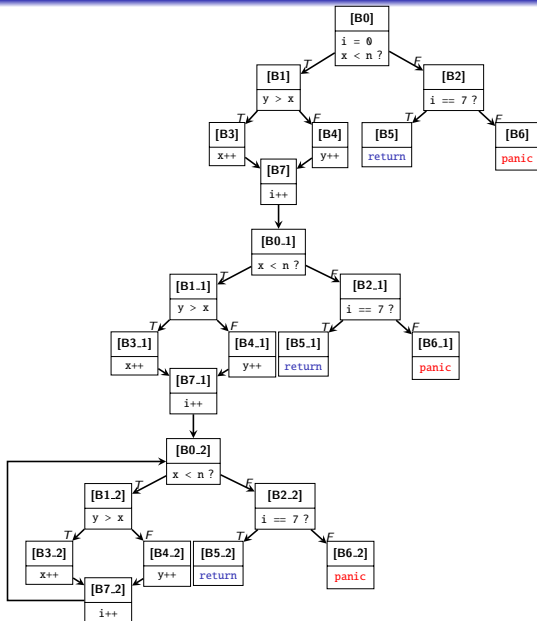
$x < n \leq y$

① $x++$ until $x == n$

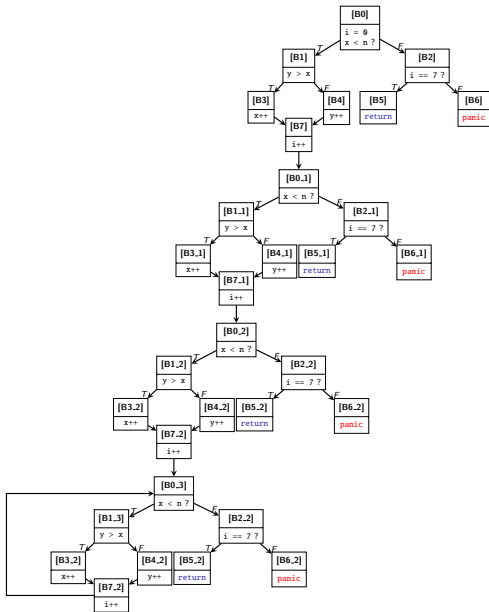
Solution: bounded loop unrolling



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Narrow-range constraints

Random input generation is not suitable for passing narrow-ranged constraints. For example:

```
1 fn foo(x: u64, y: u64) {  
2   if x + y = 42 {  
3     panic!();  
4   }  
5 }
```

If x and y are randomly generated $u64$, the chances that their sum equals 42 is extremely low.

On the other hand, this is much easier for SMT solvers to produce valid values for x and y that satisfies this constraint.

The general intuition behind concolic execution

Let fuzzing do most of the state exploration. If the coverage saturates, i.e., the fuzzer is not able to make progress on finding new coverage, invoke the symbolic reasoning engine to breakthrough.

Concolic execution demo

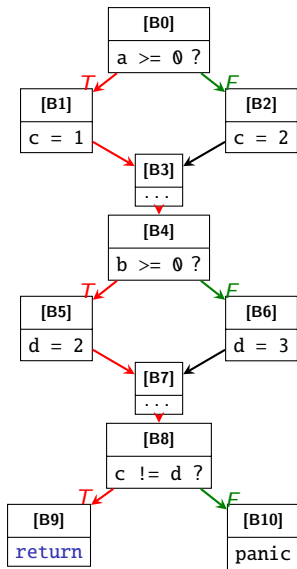
- $a = 1, b = 1$

We start with a sample input for the program, and execute the input concretely to obtain an execution trace.

Query 1: given constraint $a \geq 0 \wedge b \geq 0$ and the program, can we toggle $c \neq d$?
 \Rightarrow unsat, infeasible path

Query 2: given constraint $a \geq 0$ and the program, can we toggle $b \geq 0$?
 \Rightarrow sat, $a = 1, b = -1$

Query 3: given constraint *true* and the program, can we toggle $a \geq 0$?
 \Rightarrow sat, $a = -1$



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A comprehensive survey of current works

Fuzzing Family Tree

⟨ **End** ⟩