Informed Search

CS 486/686: Introduction to Artificial Intelligence

I

Outline

- Using knowledge
 - Heuristics
- Best-first search
 - Greedy best-first search
 - A* search
 - Variations of A*
- Back to heuristics

Last lecture

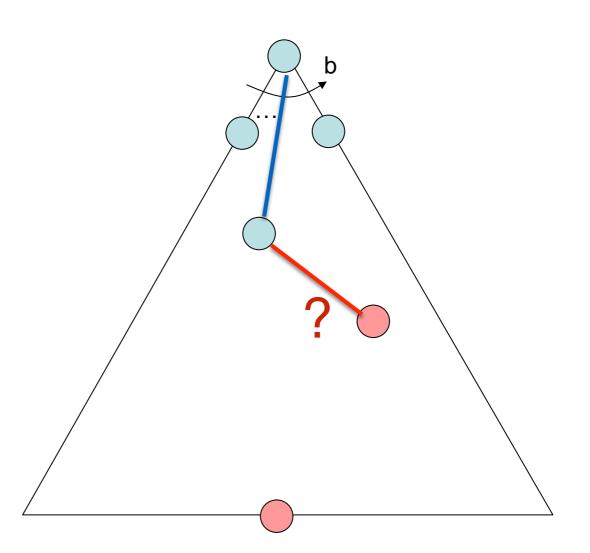
- Uninformed search uses no knowledge about the problem
 - Expands nodes based on "distance" from start node (never looks ahead to goal)
- Pros
 - Very general
- Cons
 - Very expensive
- Non-judgemental
 - Some are complete, some are not

Informed Search

- We often have additional <u>knowledge</u> about the problem
 - Knowledge is often merit of a node (value of a node)
 - Example: Romania travel problem?
- Different notions of merit
 - Cost of solution
 - Minimizing computation

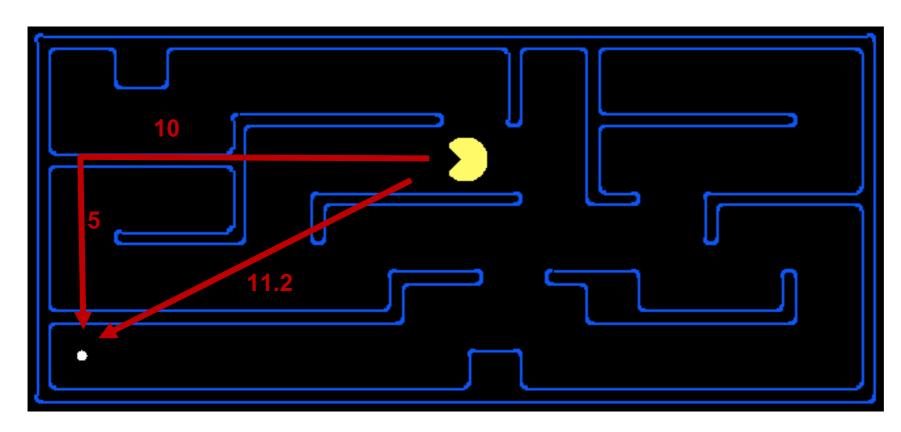
Uninformed vs Informed Search

- Uninformed search expands nodes based on distance from start node, d(n_{start}, n)
- Why not expand on distance to goal, d(n,ngoal)?



Heuristics

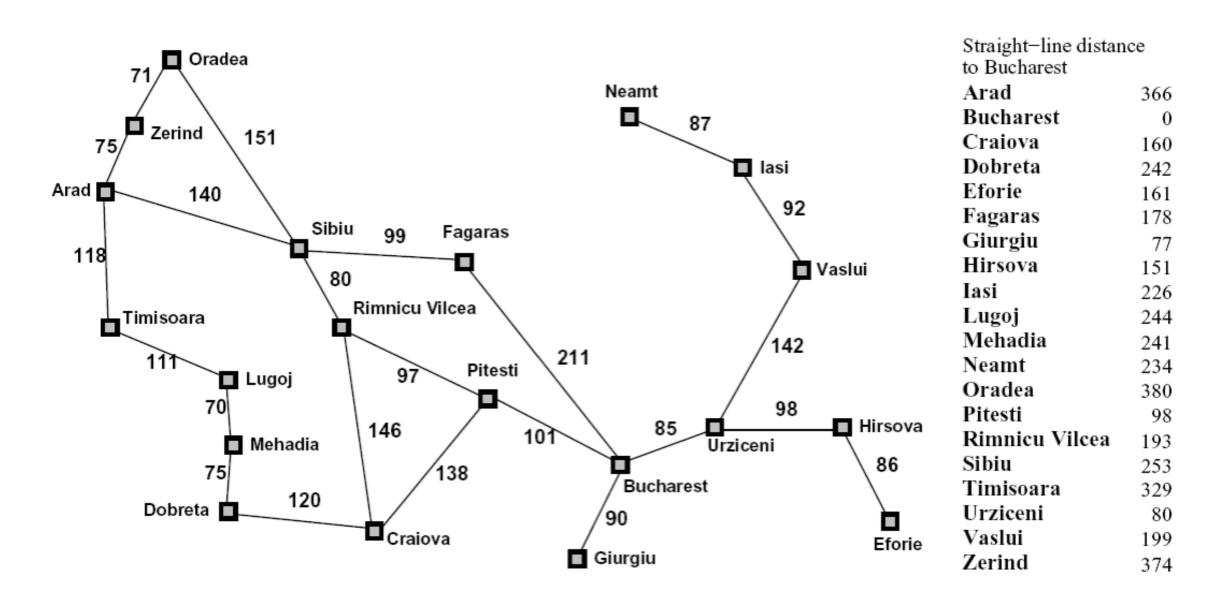
 A heuristic is a function that estimates the cost of reaching a goal from a given state



Examples:

- Euclidean distance
- Manhatten distance

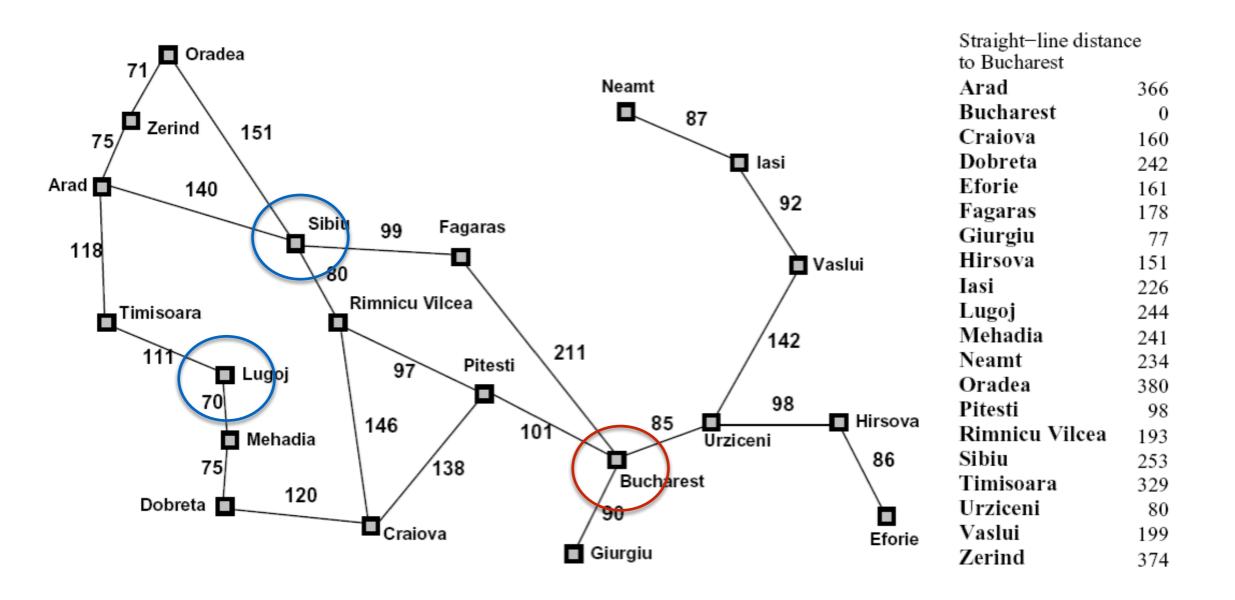
Example





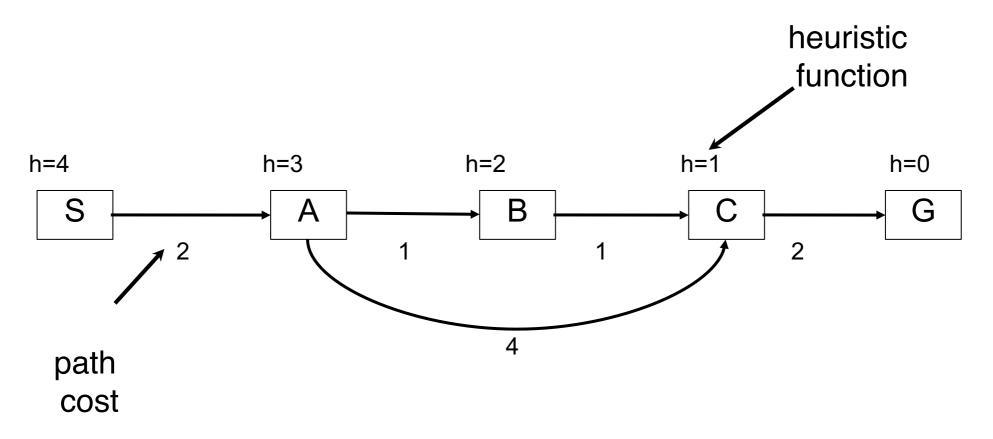
Heuristics: Structure

- If $h(n_1) < h(n_2)$ we guess it is cheaper to reach the goal from n_1 than n_2
- We require h(n_{goal})=0

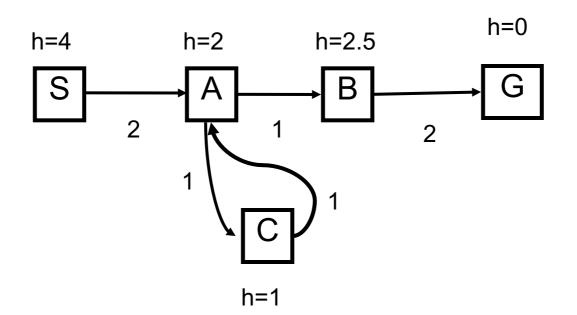


Example: Best First search

Search strategy: Expand the most promising node according to the heuristic



Example: Best First Search

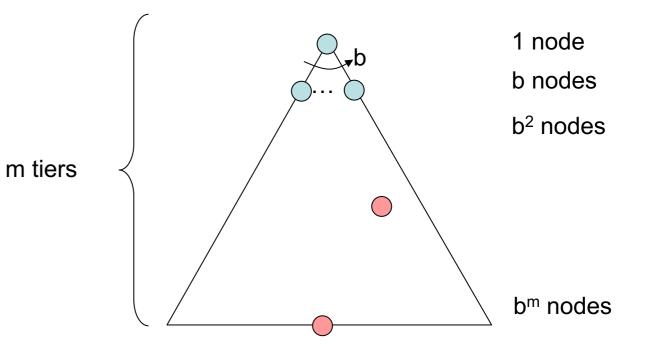




What can we do to make Best-first search simulate Breadth-first search?

Best First Search Properties

- Complete?
- Optimal?
- Time complexity
- Space complexity

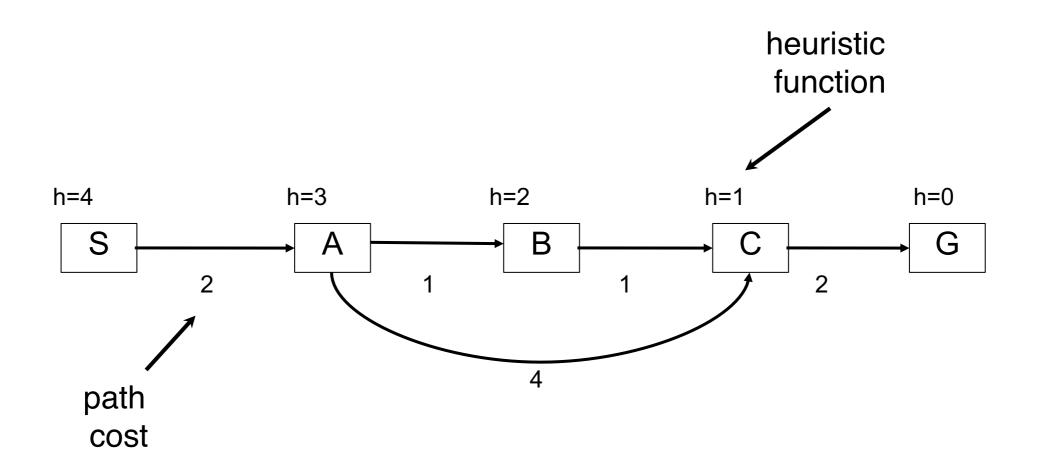


A* Search

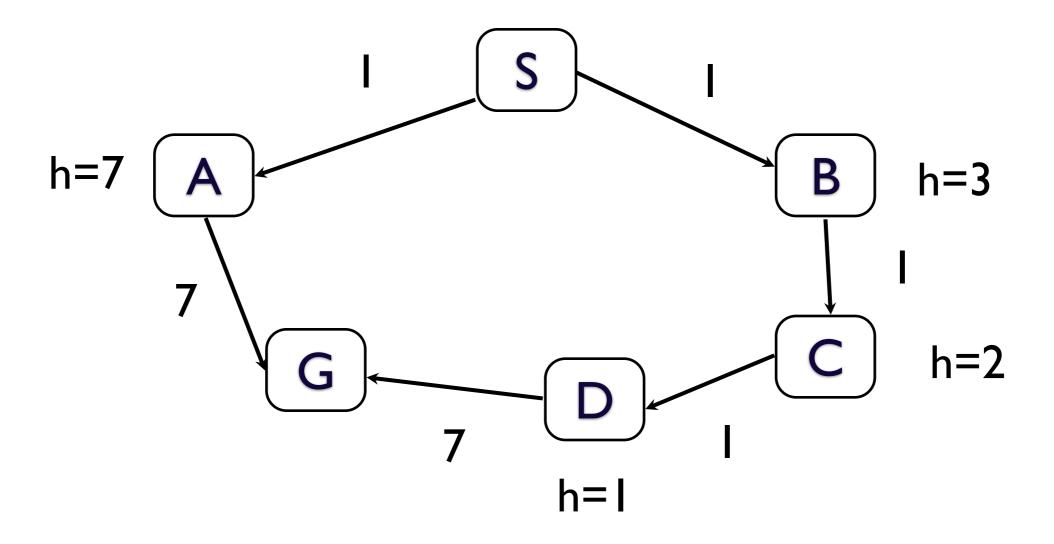
- Observations
 - Best first search ordered nodes by *forward cost to goal,* h(n)
 - Uniform cost search ordered nodes by backward cost of path so far, g(n)

- A* search
 - Expand nodes in order f(n)=g(n)+h(n)

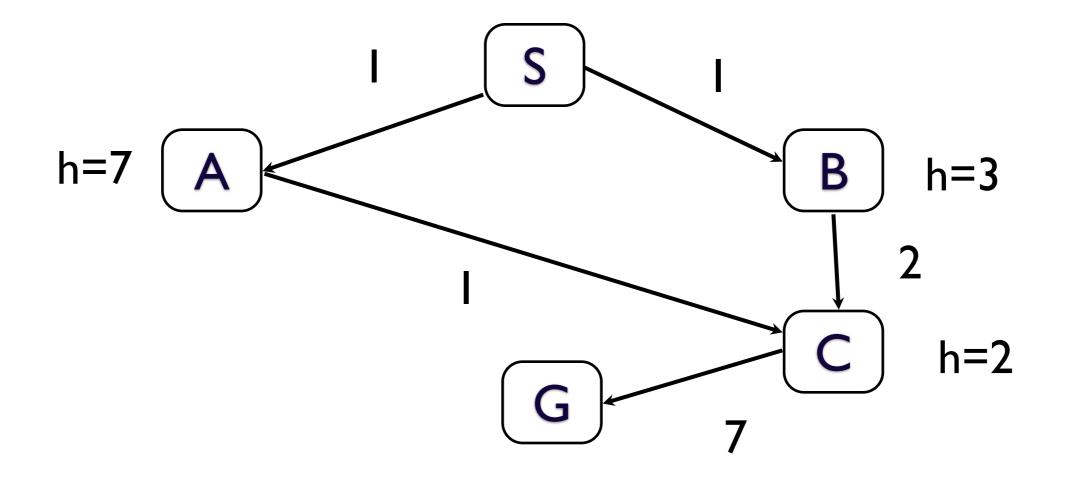
Example: A* search



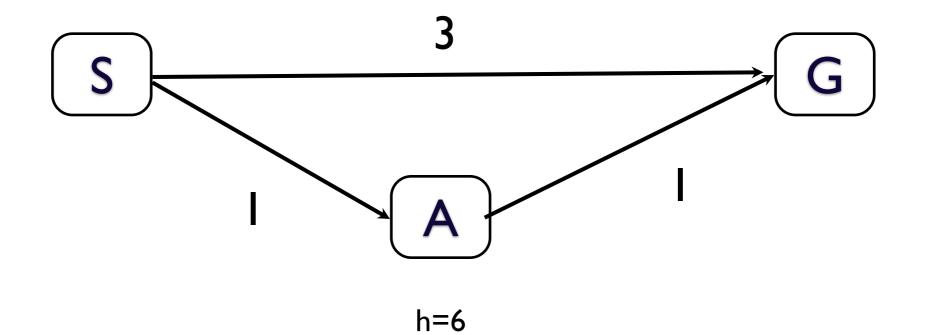
When Should A* Terminate?



A* and Revisiting States



Is A* Optimal?



Admissible Heuristics

A heuristic, h, is *admissible* if

 $0 \le h(n) \le h^*(n)$

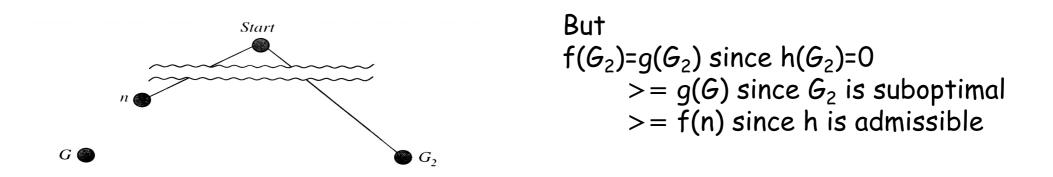
for all n, where h*(n) is the (true) shortest path from n to any goal state

Admissible heuristics are optimistic. Note that h(n)=0 is admissible.

Optimality of A*

 If the heuristic is admissible then A* with tree-search is optimal

Proof by contradiction Let goal G_2 be in the queue. Let n be an unexpanded node on the shortest path to optimal goal G. Assume that A* chose G_2 to expand. Thus, it must be that $f(n)>f(G_2)$



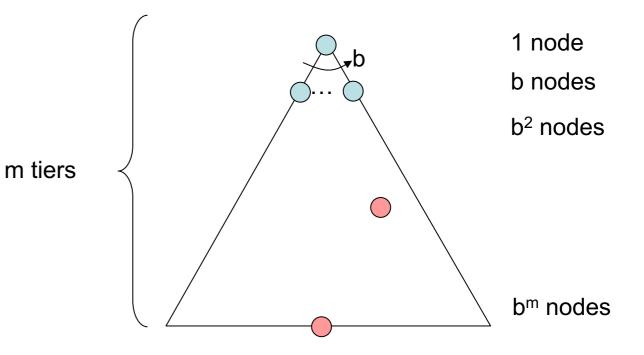
Contradiction. Therefore, A^* will never select G_2 for expansion.

Optimality of A*

- For graphs we require consistency
 - h(n)≤cost(n,n')+h(n')
 - Almost any admissible heuristic function will also be consistent
 - A* search on graphs with a consistent heuristic is optimal

A* Search Properties

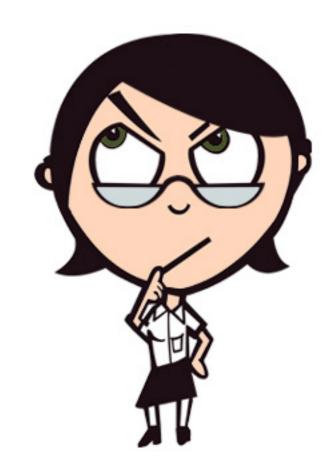
- Complete?
- Optimal?
- Time complexity
- Space complexity



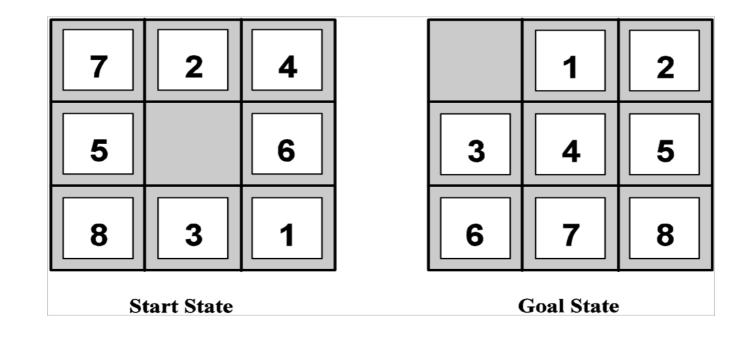
Heuristic Functions

• A good heuristic function can make all the difference!

How do we get heuristics



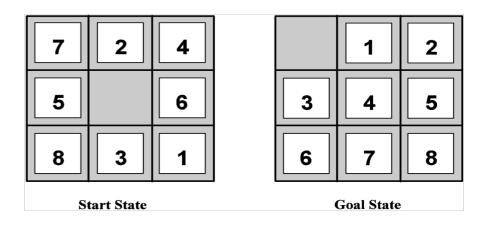
8 Puzzle



- Relax the game
 - 1. Can move from A to B is A is next to B
 - 2. Can move from A to B if B is blank
 - 3. Can move from A to B

8 Puzzle

- Can move from A to B: (*Misplaced Tile Heuristic*, h1)
 - Admissible?
- Can move from A to B if B is next to A:(*Manhatten Distance Heuristic*, h2)
 - Admissible?



Which is the better heuristic? (Which one dominates?)

8 Puzzle and Heuristics

Depth	IDS	A*(h ₁)	A*(h ₂)
2	10	6	6
4	112	13	12
8	6384	39	25
12	3644035	227	73
24	_	39135	1641

Designing Heuristics

• Relax the problem

- - -

- Precompute solution costs of subproblems and storing them in a pattern database
- Learning from experience with the problem class

Often there is a **tradeoff** between accuracy of your heuristic (and thus the amount of search) and the amount of computation you must do to compute it

Summary

- What you should know
 - Thoroughly understand A*
 - Be able to trace simple examples of A* execution
 - Understand admissibility of heuristics
 - Completeness, optimality

Memory-Bounded Heuristic Search

- Iterative Deepening A* (IDA*)
 - Basically depth-first search but using the f-value to decide which order to consider nodes
 - Use f-limit instead of depth limit
 - New f-limit is the smallest f-value of any node that exceeded cutoff on previous iteration
 - Additionally keep track of next limit to consider
 - IDA* has same properties as A* but uses less memory

Memory-Bounded Heuristic Search

- Simplified Memory-Bounded A* (SMA*)
 - Uses all available memory
 - Proceeds like A* but when it runs out of memory it drops the worst leaf node (one with highest fvalue)
 - If all leaf nodes have same f-value, drop oldest and expand newest
 - Optimal and complete if depth of shallowest goal node is less than memory size