

Constraint Satisfaction

CS 486/686: Introduction to Artificial Intelligence

Outline

- What are Constraint Satisfaction Problems (CSPs)?
- Standard Search and CSPs
- Improvements
 - Backtracking
 - Backtracking + heuristics
 - Forward Checking

Introduction

- We have been solving problems by searching in a **space of states**
 - Treating states as black boxes
- Today, we study problems where **state structure** is important

Queens Problem

- **States:** All arrangements of 0, 1, ..., or 8 queens
- **Initial state:** 0 queens on the board
- **Successor function:** Add a queen to the board
- **Goal Test:** 8 queens on the board with no two attacking each other

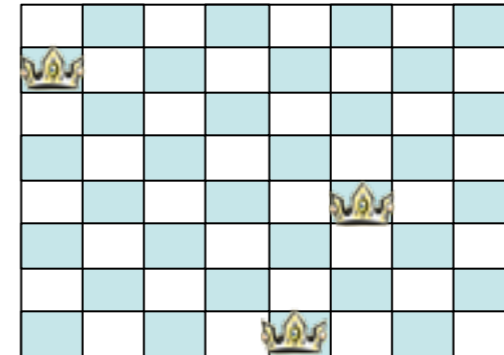
$64 \times 63 \times \dots \times 53 \sim 3 \times 10^{14}$ states



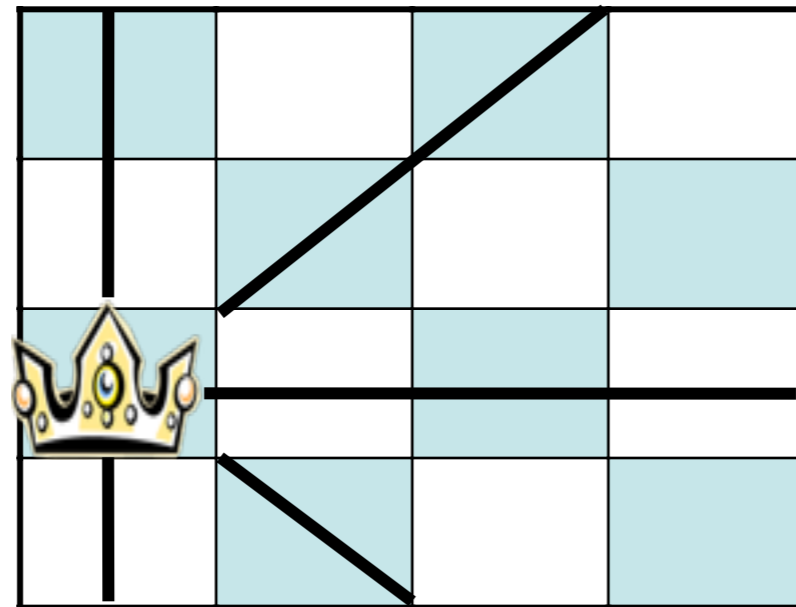
Queens Problem

- **States:** All arrangements of k queens ($0 \leq k \leq 8$), one per column in the leftmost k columns, with no queen attacking another
- **Initial state:** 0 queens on the board
- **Successor function:** Add queen leftmost empty column such that it is not attacked
- **Goal Test:** 8 queens on the board

2057 states



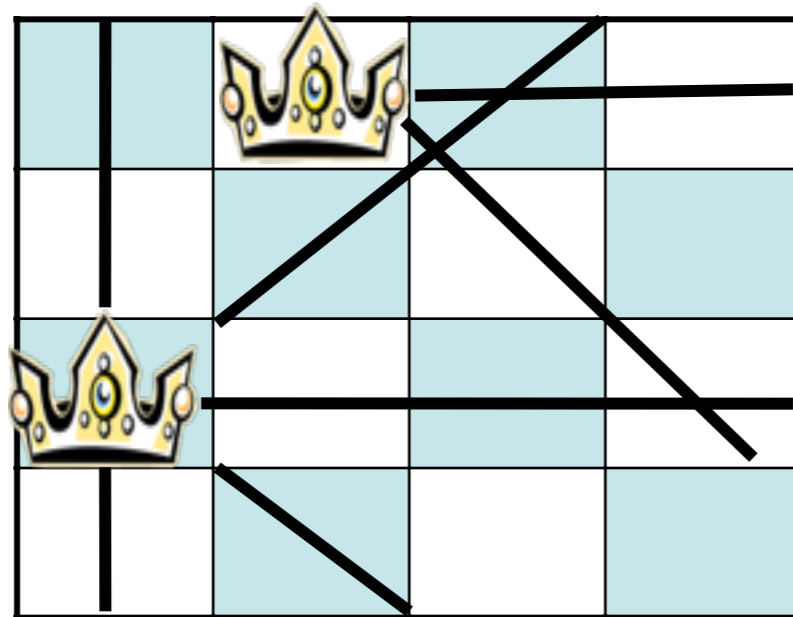
4 Queens: Constraint Propagation



Place a queen in a square

Remove conflicting squares from consideration

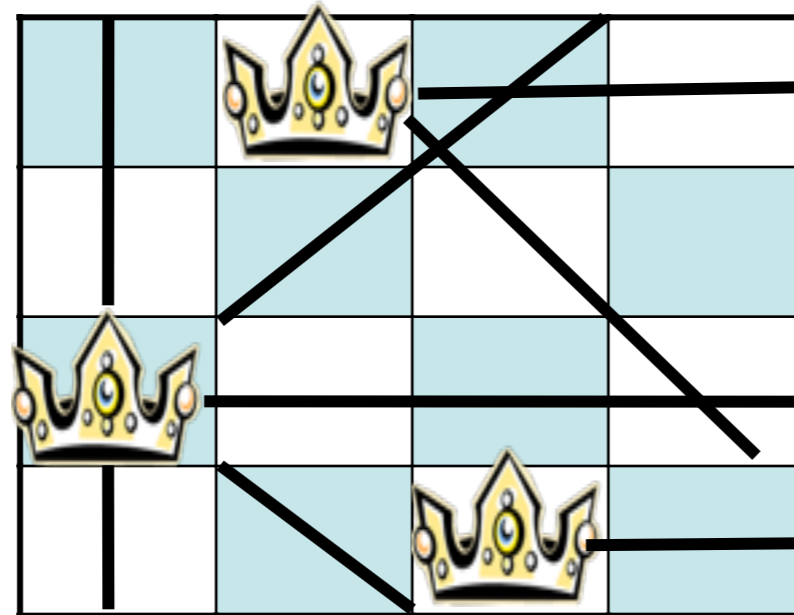
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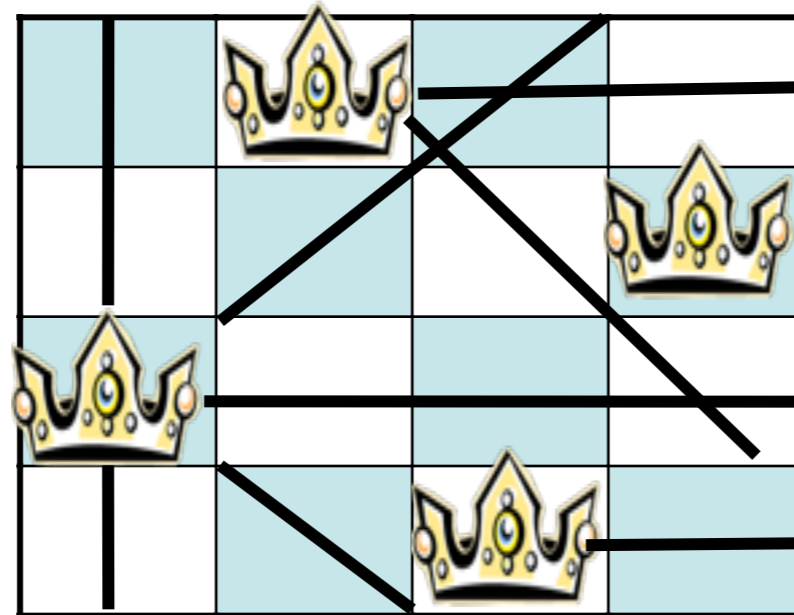
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4 Queens: Constraint Propagation



Place a queen in a square

Remove conflicting squares from consideration

CSP Definition

- A constraint satisfaction problem (CSP) is defined by $\{V, D, C\}$ where
 - $V = \{V_1, V_2, \dots, V_n\}$ is the set of **variables**
 - $D = \{D_1, D_2, \dots, D_n\}$ is the set of **domains**
 - $C = \{C_1, C_2, \dots, C_m\}$ is the set of **constraints**

CSP Definition

- A state is an **assignment** of values to some (or all) variables
 - $\{V_i=x_i, V_j=x_j, \dots\}$
- **Consistent** assignments
 - No violated constraints
- A **solution** is a complete, consistent assignment

Example: 8 Queens

- 64 variables V_{ij} , $i=1..8$, $j=1..8$
- Domain $D_{ij}=\{0,1\}$
- Constraints
 - $V_{ij}=1 \Rightarrow V_{ik}=0$ for $j \neq k$
 - $V_{ij}=1 \Rightarrow V_{kj}=0$ for $i \neq k$
 - Similar constraints for diagonals
 - $\sum_{ij} V_{ij}=8$

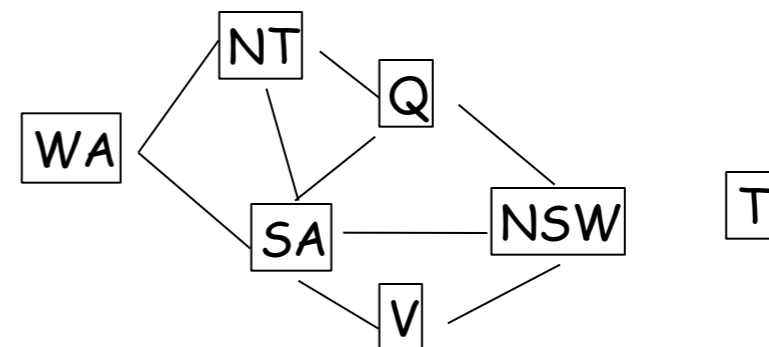
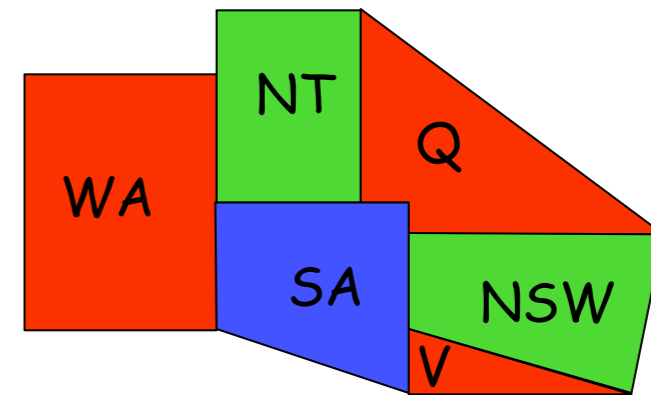
Binary Constraints:
Relate two variables

Example: 8 Queens

- 8 variables, V_i $i=1..8$
- Domains $D_i=\{1,2,\dots,8\}$
- Constraints
 - $V_i=k \Rightarrow V_j \neq k$ for all $j \neq i$
 - Similar constraints for diagonals

Example: Map Colouring

- 7 variables {WA, NT, SA, Q, NSW, V, T}
- Each variable has same domain {red, blue, green}
- No two adjacent variables have the same value



Constraint graph

Example: 3 Sat

- n Boolean variables V_1, \dots, V_n
- K constraints of the form $V_i^* \vee V_j^* \vee V_k^*$
where V_i^* is either V_i or $\sim V_i$
- NP-Complete

Variable Types of CSPs

- Discrete and finite
- Discrete variables with infinite domains
 - Constraint languages
- Continuous domains
 - Linear programming etc

Types of Constraints

- Unary
 - restricts a variable to a single value
- Binary
 - Constraint graph
- Higher order constraints involve three or more variables
 - Constraint hypergraphs

CSPs and Search

- n variables V_1, \dots, V_n
- **Valid assignment:** $\{V_1=x_1, \dots, V_k=x_k\}$ for $0 \leq k \leq n$ such that values satisfy constraints
- **States:** valid assignments
- **Initial state:** empty assignment
- **Successor:** $\{V_1=x_1, \dots, V_k=x_k\} \rightarrow \{V_1=x_1, \dots, V_k=x_k, V_{k+1}=x_{k+1}\}$
- **Goal test:** complete assignment
 - If all domains have size d then there are $O(d^n)$ complete assignments

Commutativity

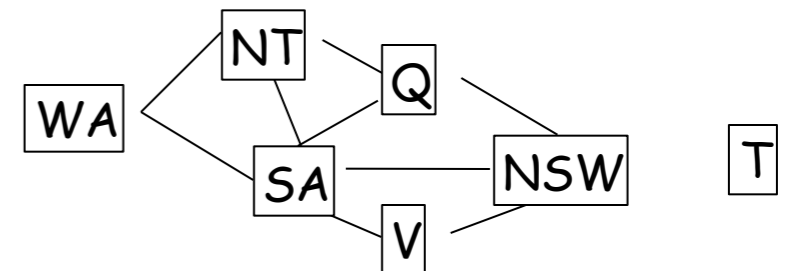
- CSPs are commutative
- Order of actions taken does not effect outcome
 - Can assign values to variables in any order
- CSP algorithms take advantage of this
 - Consider possible assignments for a **single variable at each node in search tree**

Backtracking Search

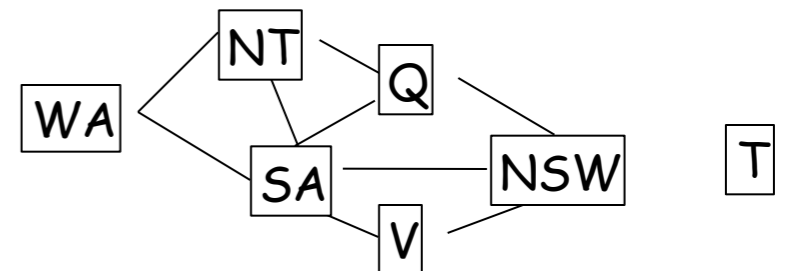
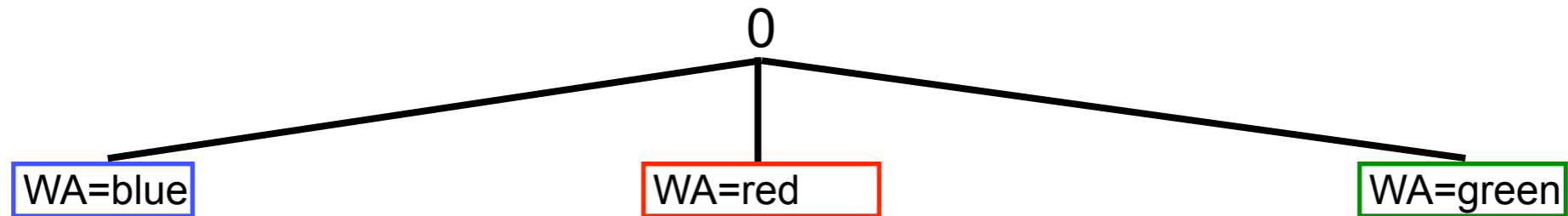
- Select unassigned variable X
- For each value $\{x_1, \dots, x_n\}$ in domain of X
 - If value satisfies constraints, assign $X=x_i$ and exit loop
- If an assignment is found
 - Move to next variable
- If no assignment found
 - Back up to preceding variable and try a different assignment for it

Backtracking Example

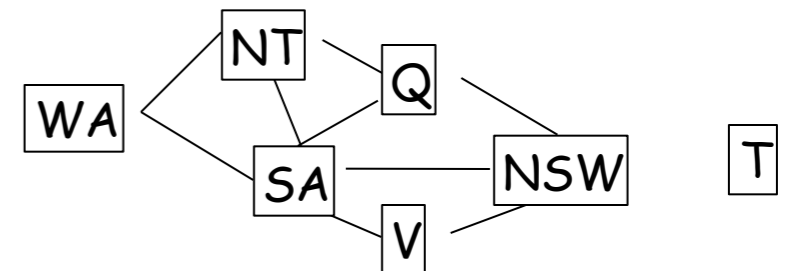
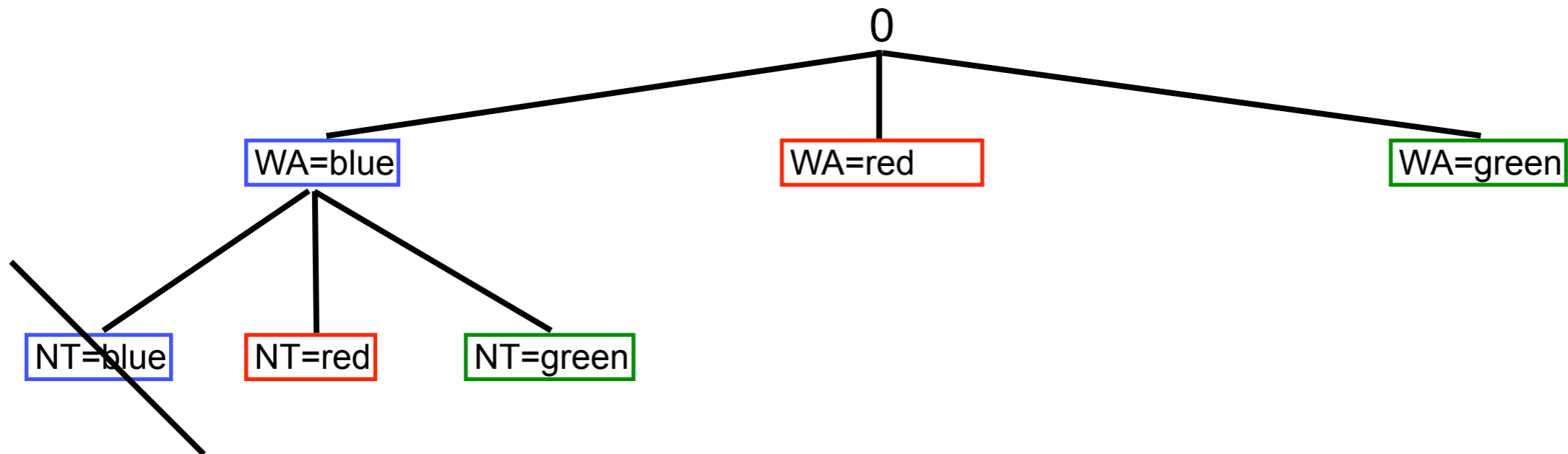
0



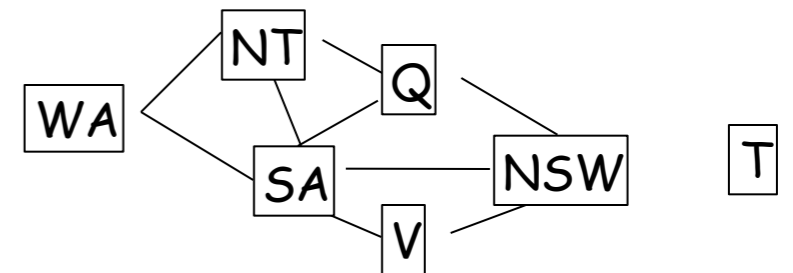
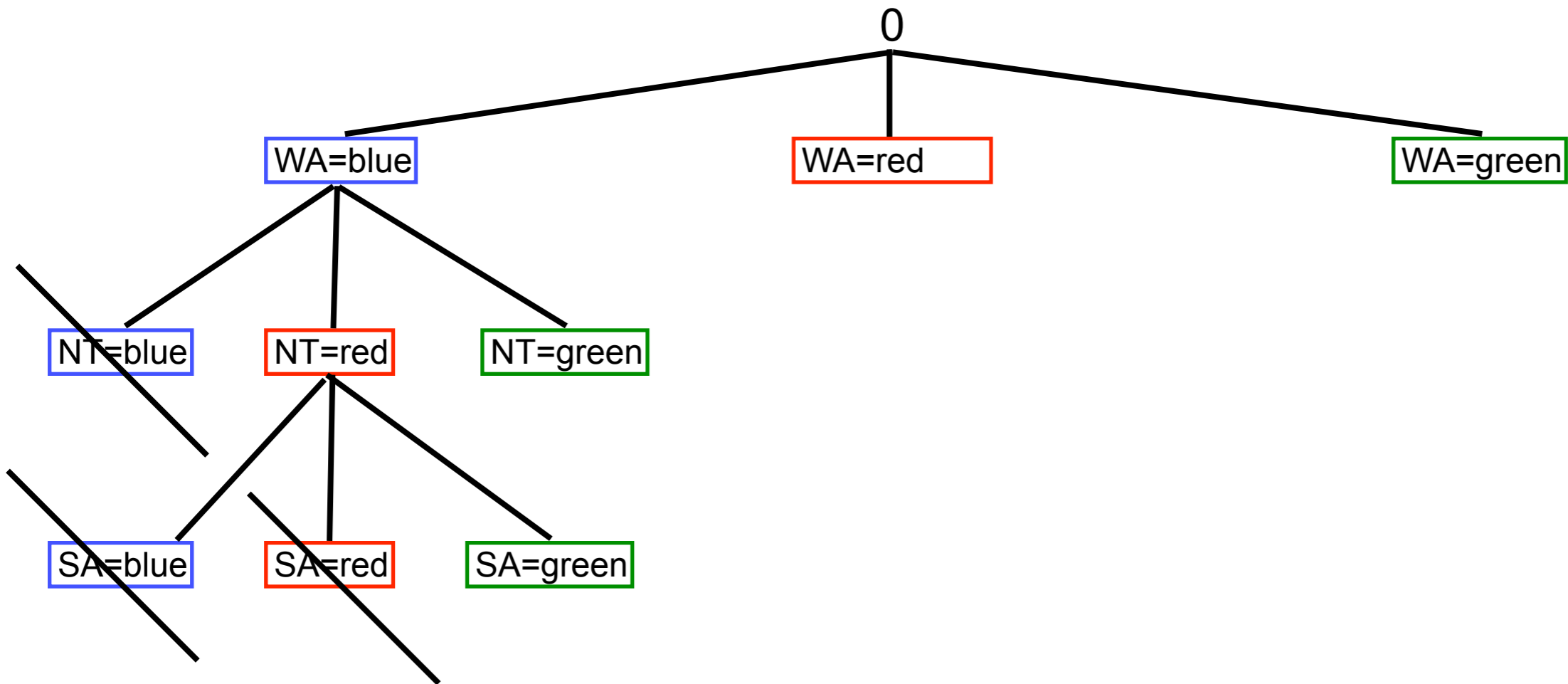
Backtracking Example



Backtracking Example



Backtracking Example

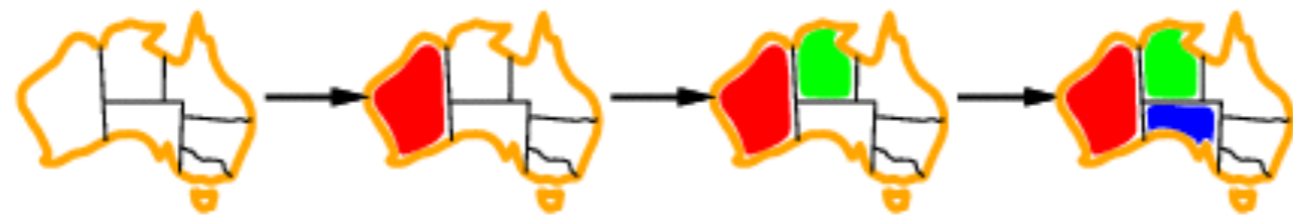


Backtracking and Efficiency

- Backtracking search = uninformed search method
 - Not very efficient
- Can we do better? Heuristics!
 - Which variable should be assigned next?
 - In which order should its values be tried?
 - Can we detect inevitable failure early?

Most Constrained Variable

- Choose the variable which has the fewest “legal” moves
 - **AKA minimum remaining values (MRV)**

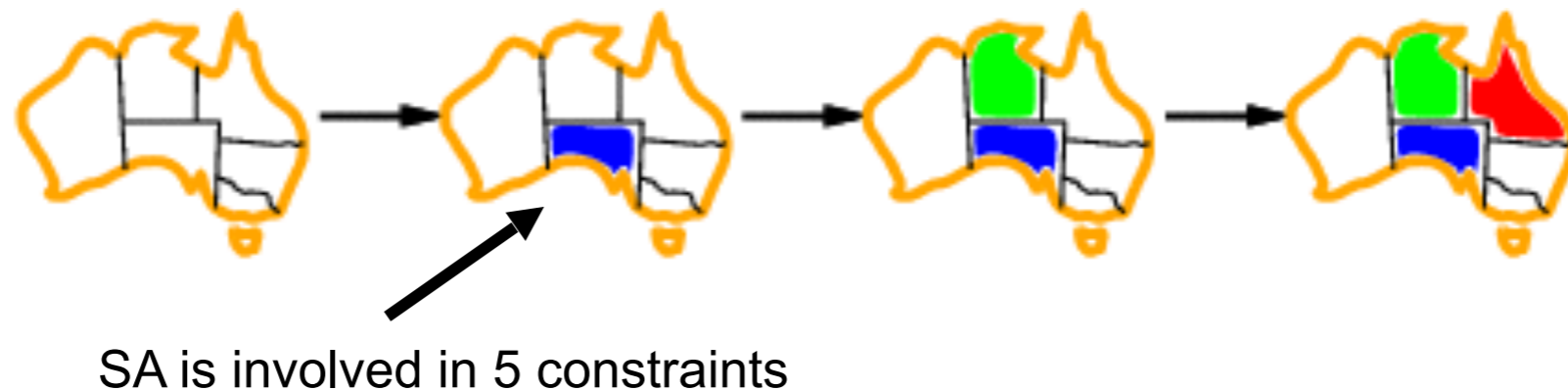


$D_{NT} = \{\text{green, blue}\}$
 $D_{SA} = \{\text{green, blue}\}$
 $D_{\text{others}} = \{\text{red, green, blue}\}$

$D_{SA} = \{\text{blue}\}$
 $D_Q = \{\text{blue, red}\}$
 $D_{\text{others}} = \{\text{red, green, blue}\}$

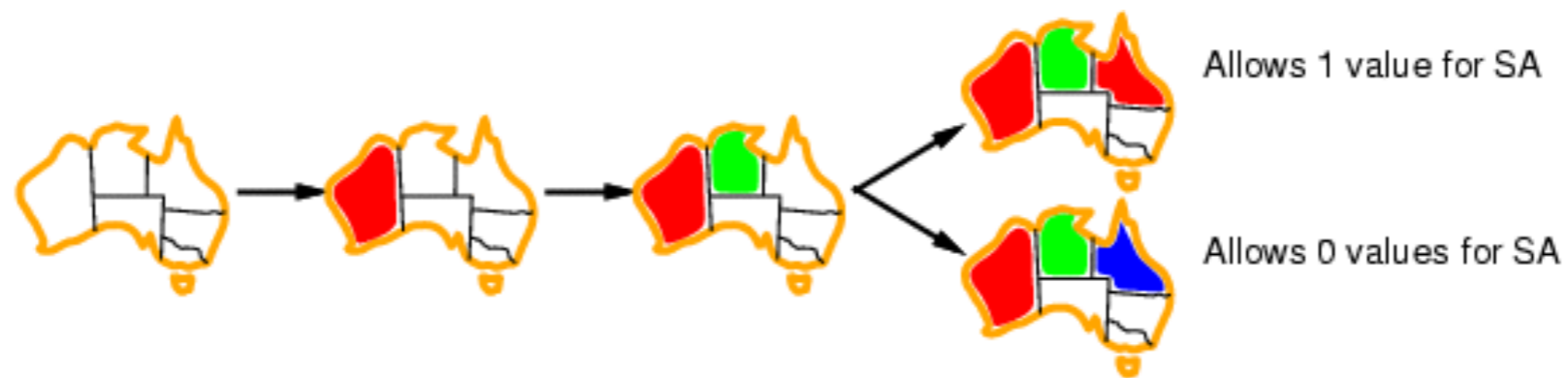
Most Constraining Variable

- Most constraining variable:
 - Choose variable with most constraints on remaining variables
- Tie-breaker among most constrained variables



Least-Constraining Value

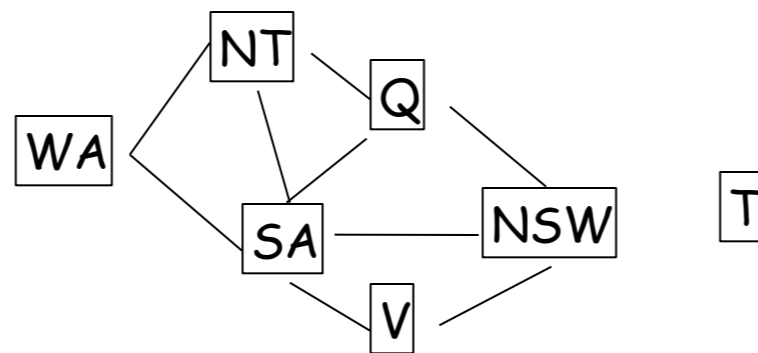
- Given a variable, choose the least constraining value:
 - The one that rules out the fewest values in the remaining variables



Forward Checking

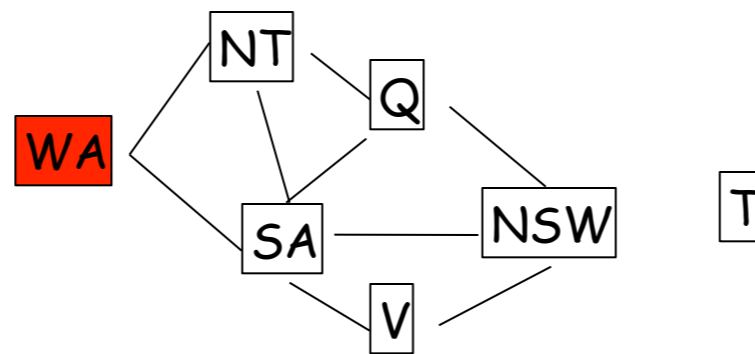
- Is there a way to detect failure early?
- Forward checking:
 - Keep track of remaining legal values for unassigned variables
 - Terminate search when any variable has no legal values

Example: Forward Checking



WA	NT	Q	NSW	V	SA	T
RGB	RGB	RGB	RGB	RGB	RGB	RGB

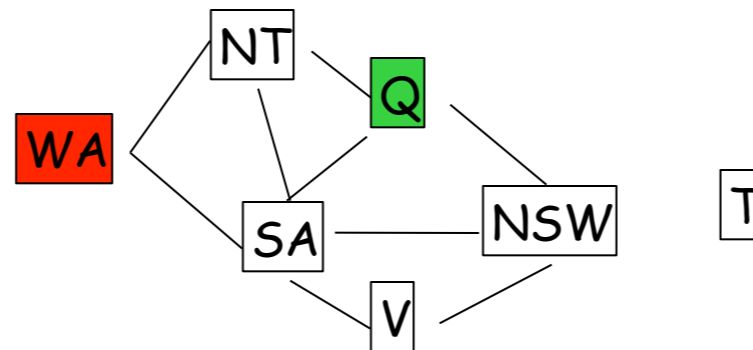
Example: Forward Checking



WA	NT	Q	NSW	V	SA	T
RGB	RGB	RGB	RGB	RGB	RGB	RGB
R	RGB	RGB	RGB	RGB	RGB	RGB

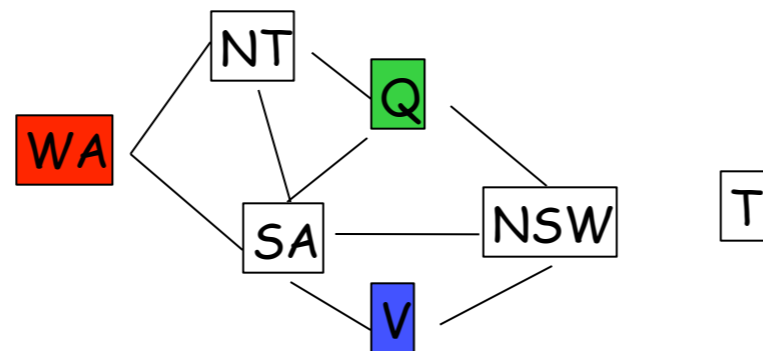
Forward checking removes the value Red of NT and of SA

Example: Forward Checking



WA	NT	Q	NSW	V	SA	T
RGB	RGB	RGB	RGB	RGB	RGB	RGB
R	GB	RGB	RGB	RGB	GB	RGB
R	GB	G	RGB	RGB	GB	RGB

Example: Forward Checking



WA	NT	Q	NSW	V	SA	T
RGB	RGB	RGB	RGB	RGB	RGB	RGB
R	GB	RGB	RGB	RGB	GB	RGB
R	B	G	RB	RGB	B	RGB
R	B	G	RB	B	B	RGB

Example: Forward Checking

Empty set: the current assignment
 $\{(WA \leftarrow R), (Q \leftarrow G), (V \leftarrow B)\}$
does not lead to a solution

WA	NT	Q	NSW	V	SA	T
RGB	RGB	RGB	RGB	RGB	RGB	RGB
R	GB	RGB	RGB	RGB	GB	RGB
R	B	G	RB	RGB	B	RGB
R	B	G	RB	B	B	RGB

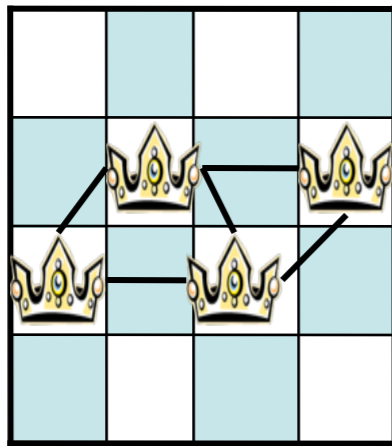
Arc Consistency

- Arc Consistency
 - Fast method of constraint propagation
 - Stronger than forward checking

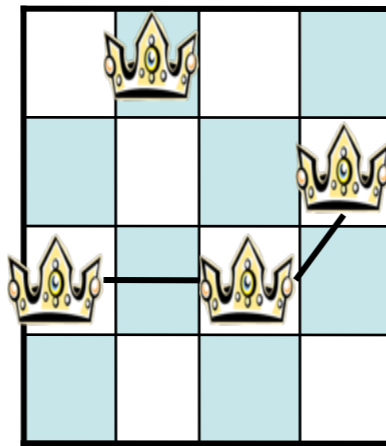
Iterative Improvement Algorithms

- Start with a broken but complete assignment of values to variables
 - Allow states to have variable assignments that do not satisfy constraints
- Randomly select conflicted variables
- Operators reassign values
 - **Min-conflict heuristic:** choose value that violates fewest constraints

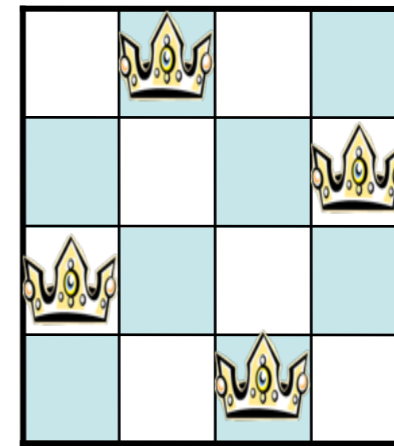
Example: 4 Queens



$h=5$



$h=2$



$h=0$

Given random initial state, can solve n-queens problem in almost constant time for arbitrary n with high probability (e.g. $n=10^7$)

Appears to be true for any randomly-generated CSP except in narrow range of ratio:

$$R = (\text{number constraints}) / (\text{number variables})$$

Summary

- How to formalize problems as CSPs
- Backtracking search
- Heuristics
 - variable ordering
 - value ordering
 - forward checking
 - arc consistency
- Iterative Improvement approaches