# Multiagent Systems: Intro to Game Theory 

CS 486/686: Introduction to Artificial Intelligence

## Introduction

- So far almost everything we have looked at has been in a single-agent setting
- Today - Multiagent Decision Making!
- For participants to act optimally, they must account for how others are going to act
- We want to
- Understand the ways in which agents interact and behave
- Design systems so that agents behave the way we would like them to

Hint for the final exam: MAS is my main research area. I like MAS problems. I even enjoy marking MAS questions. There will be a MAS question on the exam.

## Self-Interest

- We will focus on self-interested MAS
- Self-interested does not necessarily mean
- Agents want to harm others
- Agents only care about things that benefit themselves
- Self-interested means
- Agents have their own description of states of the world
- Agents take actions based on these descriptions


## What is Game Theory?

- The study of games!
- Bluffing in poker
- What move to make in chess
- How to play Rock-Paper-Scissors


But also

- auction design
- strategic deterrence
- election laws
- coaching decisions
- routing protocols


## What is Game Theory?

- Game theory is a formal way to analyze interactions among a group of rational agents that behave strategically
- Group: Must have more than 1 decision maker
- Otherwise, you have a decision problem, not a game



## What is Game Theory?

- Game theory is a formal way to analyze interactions among a group of rational agents that behave strategically
- Interaction: What one agent does directly affects at least one other
- Strategic: Agents take into account that their actions influence the game
- Rational: Agents chose their best actions


## Example



- Decision Problem
- Everyone pays their own bill
- Game
- Before the meal, everyone decides to split the bill evenly


## Strategic Game <br> (Matrix Game, Normal Form Game)

- Set of agents: $\mathrm{I}=\{1,2, .,,, \mathrm{N}\}$
- Set of actions: $A_{i}=\left\{a_{i}{ }^{1}, \ldots, a_{i}{ }^{m}\right\}$
- Outcome of a game is defined by a profile $a=\left(a_{1}, \ldots, a_{n}\right)$
- Agents have preferences over outcomes
- Utility functions $u_{i}: A->R$


## Examples

## Agent 2

One Two

|  | One | $2,-2$ |
| :---: | :---: | :---: |
|  | Agent 1 | $-3,3$ |
|  | Two | $-3,3$ |
|  |  |  |

## Zero-sum game. <br> $\sum_{i=1}{ }^{n} u_{i}(0)=0$

$\mathrm{I}=\{1,2\}$
$A_{i}=\{$ One,Two $\}$
$A_{n}$ outcome is (One, Two)
$\mathrm{U}_{1}(($ One,Two $))=-3$ and $\mathrm{U}_{2}(($ One,Two $))=3$

## Examples



Coordination Game

## Chicken

|  | T | C |
| :---: | :---: | :---: |
| T | -1,-1 | 10,0 |
| C | 0,10 | 5,5 |



Anti-Coordination Game

## Example: Prisoners' Dilemma



## Playing a Game

## - Agents are rational

- Let $p_{i}$ be agent $i$ 's belief about what its opponents will do
- Best response: $a_{i}=\operatorname{argmax} \sum_{a_{-i}} u_{i}\left(a_{i}, a_{-i}\right) p_{i}\left(a_{-i}\right)$


Notation Break: $a_{-i}=\left(a_{1}, \ldots, a_{i-1}, a_{i+1}, \ldots, a_{n}\right)$

## Dominated Strategies

- $a^{\prime}{ }_{i}$ strictly dominates strategy $a_{i}$ if

$$
u_{i}\left(a_{i}^{\prime}, a_{-i}\right)>u_{i}\left(a_{i}, a_{-i}\right) \forall a_{-i}
$$

- A rational agent will never play a dominated strategy!


## Example

## Confess Don't Confess

| Confess  <br> Don't <br> Confess $-5,-5$ <br>  $-10,0$ $0-1,-10$ |
| :--- | :--- | :--- |

## Strict Dominance Does Not Capture the Whole Picture

|  | A | B | c |
| :---: | :---: | :---: | :---: |
| A | 0,4 | 4,0 | 5,3 |
| B | 4,0 | 0,4 | 5,3 |
| $C$ | 3,5 | 3,5 | 6,6 |

## Nash Equilibrium

Key Insight: an agent's best-response depends on the actions of other agents

An action profile a* is a Nash equilibrium if no agent has incentive to change given that others do not change

$$
\forall i u_{i}\left(a_{i}^{*}, a_{-i}^{*}\right) \geq u_{i}\left(a_{i}^{\prime}, a_{-i}^{*}\right) \forall a_{i}^{\prime}
$$

## Nash Equilibrium

- Equivalently, a* is a N.E. iff

$$
\forall i a_{i}^{*}=\arg \max _{a_{i}} u_{i}\left(a_{i}, a_{-i}^{*}\right)
$$

|  | A | B | $C$ |
| :---: | :---: | :---: | :---: |
| A | 0,4 | 4,0 | 5,3 |
| B | 4,0 | 0,4 | 5,3 |
| $C$ | 3,5 | 3,5 | 6,6 |

(C,C) is a N.E. because

$$
\begin{aligned}
& u_{1}(C, C)=\max \left[\begin{array}{l}
u_{1}(A, C) \\
u_{1}(B, C) \\
u_{1}(C, C)
\end{array}\right] \\
& \text { AND } \\
& u_{2}(C, C)=\max \left[\begin{array}{l}
u_{2}(C, A) \\
u_{2}(C, B) \\
u_{2}(C, C)
\end{array}\right]
\end{aligned}
$$

## Nash Equilibrium

- If $\left(a_{1}{ }^{*}, a_{2}{ }^{*}\right)$ is a N.E. then player 1 won't want to change its action given player 2 is playing $a_{2} *$
- If $\left(a_{1}{ }^{*}, a_{2}{ }^{*}\right)$ is a N.E. then player 2 won't want to change its action given player 1 is playing $a_{1} *$

| $-5,-5$ | $0,-10$ |
| :--- | :--- |
| $-10,0$ | $-1,-1$ |


|  | A | B | c |
| :---: | :---: | :---: | :---: |
| A | 0,4 | 4,0 | 5,3 |
| B | 4,0 | 0,4 | 5,3 |
| $c$ | 3,5 | 3,5 | 6,6 |

## Another Example

B S


2 Nash Equilibria
Coordination Game

## Yet Another Example

Agent 2

|  | One | Two |
| :---: | :---: | :---: |
| One | $2,-2$ | $-3,3$ |
| Agent 1 1-2, |  |  |
| Two | $-3,3$ | 4,-4 |

## (Mixed) Nash Equilibria

- (Mixed) Strategy: $s_{i}$ is a probability distribution over $A_{i}$
- Strategy profile: $s=\left(s_{1}, \ldots, s_{n}\right)$
- Expected utility: $u_{i}(s)=\Sigma_{a} \Pi_{j} s\left(a_{j}\right) u_{i}(a)$
- Nash equilibrium: $s^{*}$ is a (mixed) Nash equilibrium if

$$
u_{i}\left(s_{i}^{*}, s_{-i}^{*}\right) \geq u_{i}\left(s_{i}^{\prime}, s_{-i}^{*}\right) \forall s_{i}^{\prime}
$$

## Yet Another Example

\[

\]

How do we determine $p$ and $q$ ?



## Yet Another Example

|  | q One |  | Two |
| :---: | :---: | :---: | :---: |
|  | One | $2,-2$ | $-3,3$ |
|  |  | $-3,3$ | $4,-4$ |
|  |  |  |  |

How do we determine p and q ?

## Exercise

|  | $B$ | S |
| :---: | :---: | :---: |
|  | 2,1 | 0,0 |
|  | 0,0 | 1,2 |
|  |  |  |

This game has 3 Nash Equilibrium (2 pure strategy NE and 1 mixed strategy NE).

## Mixed Nash Equilibrium

- Theorem (Nash 1950): Every game in which the action sets are finite, has a mixed strategy equilibrium.

John Nash

Nobel Prize in Economics (1994)


## Finding NE

- Existence proof is non-constructive
- Finding equilibria?
- 2 player zero-sum games can be represented as a linear program (polynomial)
- For arbitrary games, the problem is in PPAD
- $\quad$ Finding equilibria with certain properties is often NP-hard


## Repeated Games

Recall the Prisonner's Dilemma. What if the prisoners are habitual criminals?

| $-5,-5$ | $0,-10$ |
| :--- | :--- |
| $-10,0$ | $-1,-1$ |
| $-10,0$ | $-1,-1$ |
| $-5,-5$ | $0,-10$ |
| $-10,0$ | $-1,-1$ |

How do we define payoffs?
What is the strategy space?

## Repeated Games

Recall the Prisonner's Dilemma. What if the prisoners are habitual criminals?

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| :--- | :--- |
| $-10,0$ | $-1,-1$ |


| $-5,-5$ | $0,-10$ |
| :--- | :--- |
| $-10,0$ | $-1,-1$ |


| $-5,-5$ | $0,-10$ |
| :--- | :--- |
| $-10,0$ | $-1,-1$ |

How do we define payoffs?
Average reward
Discounted Awards

## Repeated Games

Recall the Prisonner's Dilemma. What if the prisoners are habitual criminals?

| $-5,-5$ | $0,-10$ |
| :--- | :--- |
| $-10,0$ | $-1,-1$ |$\quad$| $-5,-5$ | $0,-10$ |
| :--- | :--- | :--- |
| $-10,0$ | $-1,-1$ |$\quad$| $-5,-5$ | $0,-10$ |
| :--- | :--- | :--- |
| $-10,0$ | $-1,-1$ |

Strategy space becomes significantly larger!
$\mathrm{S}: \mathrm{H} \rightarrow \mathrm{A}$ where H is the history of play so far
Can now reward and punish past behaviour, worry about reputation, establish trust,...

## Repeated Games

Recall the Prisonner's Dilemma. What if the prisoners are habitual criminals?

| $-5,-5$ | $0,-10$ |
| :--- | :--- |
| $-10,0$ | $-1,-1$ |


| $-5,-5$ | $0,-10$ |
| :--- | :--- |
| $-10,0$ | $-1,-1$ |


| $-5,-5$ | $0,-10$ |
| :--- | :--- |
| $-10,0$ | $-1,-1$ |

Grim Strategy: In first step cooperate. If opponent defects at some point, then defect forever

Tit-for-Tat: In first step cooperate. Copy what ever opponent did in previous stage.

## Summary

# Definition of a Normal Form Game 

Dominant strategies
Nash Equilibria

