# Introduction to Artificial Intelligence 

## Informed Search

CS 486/686
University of Waterloo

## Outline

- Using knowledge
- Heuristics
- Best-first search
- Greedy best-first search
- A* search
- Variations of A*
- Back to heuristics


## Last lecture

- Uninformed search uses no knowledge about the problem
- Expands nodes based on "distance" from start node (never looks ahead to goal)
- Pros
- Very general
- Cons
- Very expensive
- Non-judgmental
- Some are complete, some are not


## Informed Search

We often have additional knowledge about the problem

- Knowledge is often merit of a node (value of a node)
- Example: Romania travel problem?

Different notions of merit

- Cost of solution
- Minimizing computation


## Uninformed vs Informed Search

- Uninformed search expands nodes based on distance from start node, $\mathrm{d}($ nstart, n$)$
- Why not expand on distance to goal, $\mathrm{d}(\mathrm{n}$, goal)?



## Heuristics

A heuristic is a function that estimates the cost of reaching a goal from a given state


## Examples:

- Euclidean distance
- Manhatten distance


## Example



## Heuristics: Structure

- If $h(n 1)<h(n 2)$ we guess it is cheaper to reach the goal from $n 1$ than n2
- We require $h(n, g o a l)=0$


| Straight-line distance |  |
| :--- | ---: |
| to Bucharest |  |
| Arad | 366 |
| Bucharest | 0 |
| Craiova | 160 |
| Dobreta | 242 |
| Eforie | 161 |
| Fagaras | 178 |
| Giurgiu | 77 |
| Hirsova | 151 |
| Iasi | 226 |
| Lugoj | 244 |
| Mehadia | 241 |
| Neamt | 234 |
| Oradea | 380 |
| Pitesti | 98 |
| Rimnicu Vilcea | 193 |
| Sibiu | 253 |
| Timisoara | 329 |
| Urziceni | 80 |
| Vaslui | 199 |
| Zerind | 374 |

## Example: Best First search

Search strategy: Expand the most promising node according to the heuristic heuristic


## Example: Best First Search



## Best First Search Properties

- Complete?
- Optimal?
- Time complexity

- Space complexity


## A* Search

## Observations

- Best first search ordered nodes by forward cost to goal, h(n)
- Uniform cost search ordered nodes by backward cost of path so far, g(n)

A* search

- Expand nodes in order $f(n)=g(n)+h(n)$


## Example: A* search



## When Should A* Terminate?



## A* and Revisiting States



## Is A* Optimal?



## Admissible Heuristics

A heuristic is admissible if $0 \leq h(n) \leq h *(n)$ for all $n$, where $h^{*}(n)$ is the true shortest path cost from n to any goal state.

Admissible heuristics are optimistic. Note that (h)=0 is admissible.

## Optimality of A*

If the heuristic is admissible then $A^{*}$ with tree search is optimal.

If we have a graph, then we require a stronger property for the heuristic function.

A heuristic is consistent if $h(n) \leq \operatorname{cost}\left(n, n^{\prime}\right)+h^{\prime}(n)$
Almost any admissible heuristic will also be consistent.

## A* is Optimally Efficient

Among all optimal algorithms that start at the same start node and use the same heuristic, $A^{*}$ expands the fewest nodes

## A* Search Properties

- Complete?
- Optimal?
- Time complexity
- Space complexity



## Heuristic Functions

A good heuristic function can make all the difference!

How do we get heuristics?


## 8 Puzzle

- Relax the game
- Can move from $A$ to $B$ is $A$ is next to B
- Can move from $A$ to $B$ if $B$ is blank
- Can move from A to B


Start State


Goal State

## 8 Puzzle

> | Dominating heuristic: Given heuristics $\mathrm{h} 1(\mathrm{n})$ |
| :--- |
| and $\mathrm{h} 2(\mathrm{n}), \mathrm{h} 2(\mathrm{n})$ dominates $\mathrm{h} 1(\mathrm{n})$ if |
| $\forall n h 2(n) \geq h 1(n)$ and $\exists \mathrm{n} 2(n)>h 1(n)$ |

| 7 | 2 | 4 |
| :---: | :---: | :---: |
| 5 |  | 6 |
| 7 | 3 | 1 |

Start State


Goal State

Theorem: If $\mathrm{h} 2(\mathrm{n})$ dominates $\mathrm{h} 1(\mathrm{n}), \mathrm{A}^{*}$ using $h 2(n)$ will never expand more nodes that $A^{*}$ using h1(n).

- Can move from A to B: (Misplaced Tile Heuristic, h1)
- Can move from $A$ to $B$ if $B$ is next to $A:(M a n h a t t e n ~ D i s t a n c e ~ H e u r i s t i c, ~$ h2)


## 8 Puzzle and Heuristics

| Depth | IDS | $A^{*}\left(\mathbf{h}_{\mathbf{1}}\right)$ | $A^{*}\left(\mathbf{h}_{\mathbf{2}}\right)$ |
| :--- | :--- | :--- | :--- |
| $\mathbf{2}$ | 10 | 6 | 6 |
| $\mathbf{4}$ | 112 | 13 | 12 |
| $\mathbf{8}$ | 6384 | 39 | 25 |
| $\mathbf{1 2}$ | 3644035 | 227 | 73 |
| $\mathbf{2 4}$ | - | 39135 | 1641 |

## Designing Heuristics

- Relax the problem
- Precompute solution costs of subproblems and storing them in a pattern database
- Learning from experience with the problem class
- ...

Often there is a tradeoff between accuracy of your heuristic (and thus the amount of search) and the amount of computation you must do to compute it

## Summary

- What you should know
- Thoroughly understand A*
- Be able to trace simple examples of A* execution $^{*}$
- Understand admissibility of heuristics
- Completeness, optimality


## Some Things to Think About

- What is the relationship between A* search and Dijkstra's algorithm?
- A* search can be very memory intensive. Can you think of some variants of $A^{*}$ search that might reduce the memory overhead?

