

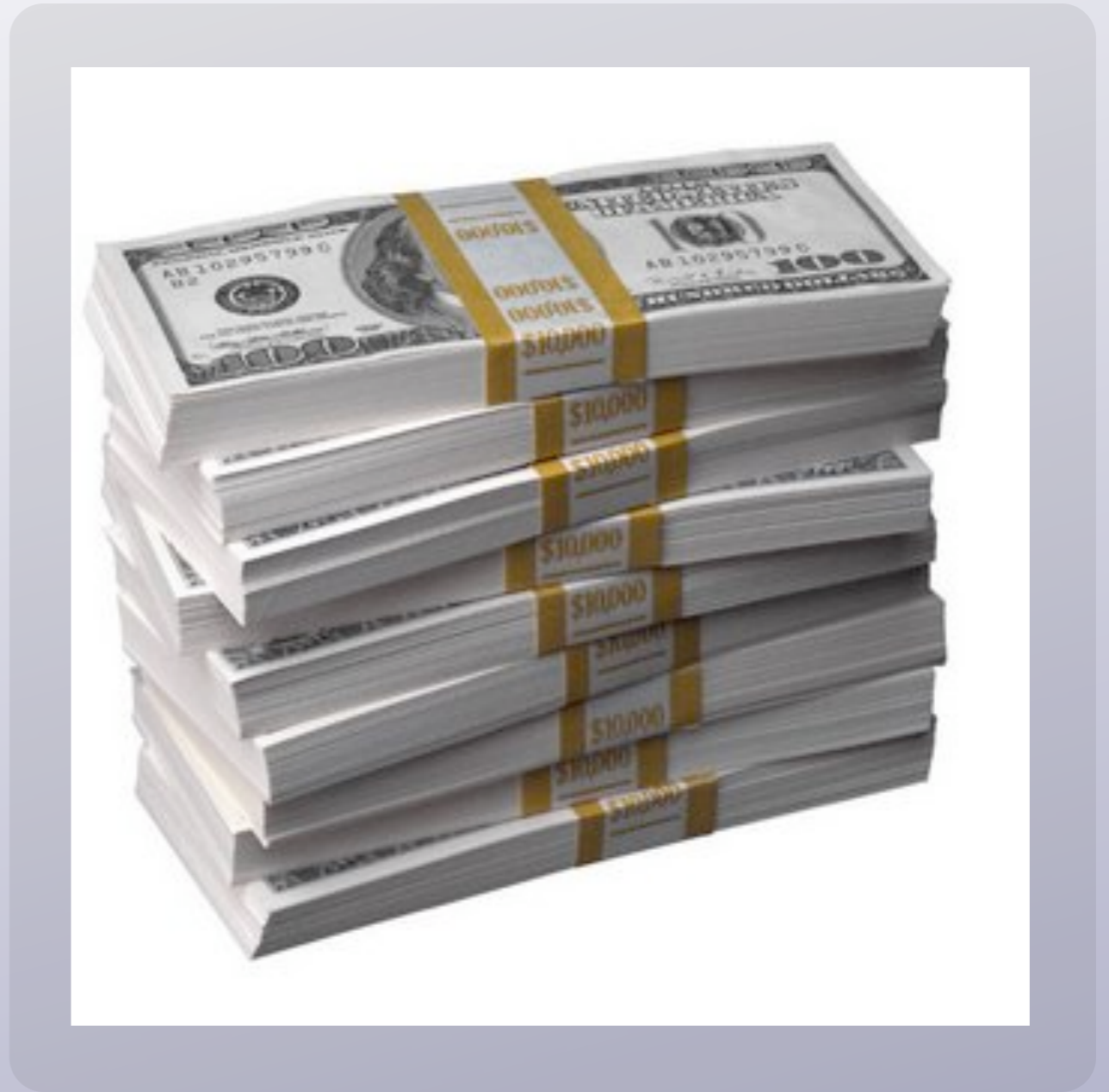
Lossless Abstraction of Imperfect Informations Games

Andrew Gilpin
Tuomas Sandholm

OR



How to Play Poker Perfectly



That means money.



Potentially, even a lot of money.



Due to time constraints,
Details about making profits
will be left as an exercise.

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FLASHBACK

Imperfect Information Games

- Sometimes agents have not observed everything, or else can not remember what they have observed

Imperfect information games: Choice nodes H are partitioned into *information sets*.

- If two choice nodes are in the same information set, then the agent can not distinguish between them.
- Actions available to an agent must be the same for all nodes in the same information set

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This is not my work

The Problem : Finding Nash

Sequential imperfect information game can be expressed in normal matrix form.

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→ **Exponential cost**

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Better : use the sequence form.

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Sequential imperfect information game can be expressed in normal matrix form.

→ **Exponential cost**

Better : use the sequence form.

→ **Linear cost**

Goal of the Article

Create an smaller game equivalent to the initial one.

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→ **Automatically**

Using abstractions.

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Create an smaller game **equivalent** to the initial one.

→ **Automatically**

Using **abstractions**

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Plan

Introduction (just done)

1. Rhode Island Hold'em

2. Games with ordered signals

3. Filtered Signal Tree

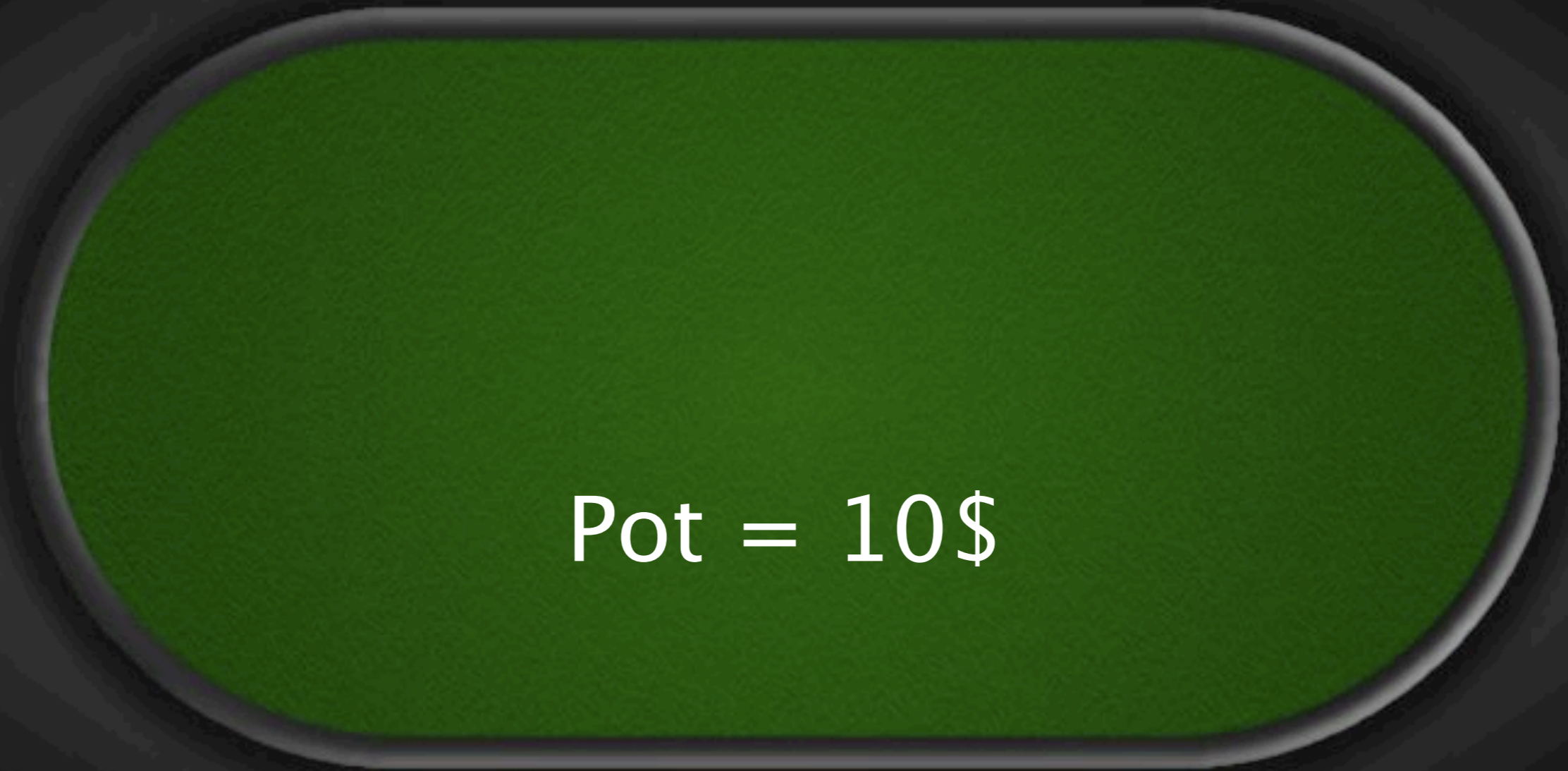
4. Main Theorem

Discussion & Conclusion

1. Rhode Island Hold'em



Opponent : Ante : +5\$



Pot = 10\$

Before 1st round

Me : Ante : +5\$



Pot = 10\$



1st round



Pot = 20\$



1st round

1. Me : Bet : +10\$

2. Opponent : Call : +10\$



Pot = 30\$



1st round

1. Me : Bet : +10\$



Pot = 30\$



End of 1st round



Pot = 30\$



2nd round



Pot = 50\$



2nd round

1. Me : Bet : +20\$

2. Opponent : Raise : +40\$



Pot = 90\$



2nd round

1. Me : Bet : +20\$



Pot = 110\$



2nd round

3. Me : Call : +20\$



Pot = 110\$



End of 2nd round



Pot = 110\$



3rd round



Pot = 130\$



3rd round

1. Me : Bet : +20\$

2. Opponent : Call : +20\$



3rd round

1. Me : Bet : +20\$

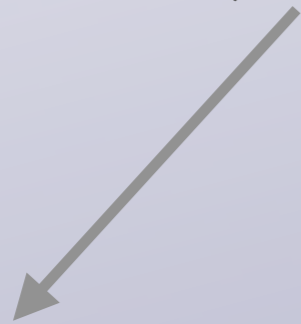


Showdown : I won 150\$

2. Games with ordered signals

$$\Gamma = (I, G, L, \Theta, \kappa, \gamma, p, \lambda, \omega, u)$$

$$\Gamma = (I, G, L, \Theta, \kappa, \gamma, p, \succeq, \omega, u)$$



Players

Tree describing
how the game
proceeds

$$\Gamma = (I, G, L, \Theta, \kappa, \gamma, p, \succeq, \omega, u)$$

Players

Player's
turns

Tree describing
how the game
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$$\Gamma = (I, G, L, \Theta, \kappa, \gamma, p, \succeq, \omega, u)$$

Players

Player's
turns Set of
cards

Tree describing
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Players

Player's
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Players

Number of
private cards
for each turn

Player's
turns Set of
cards

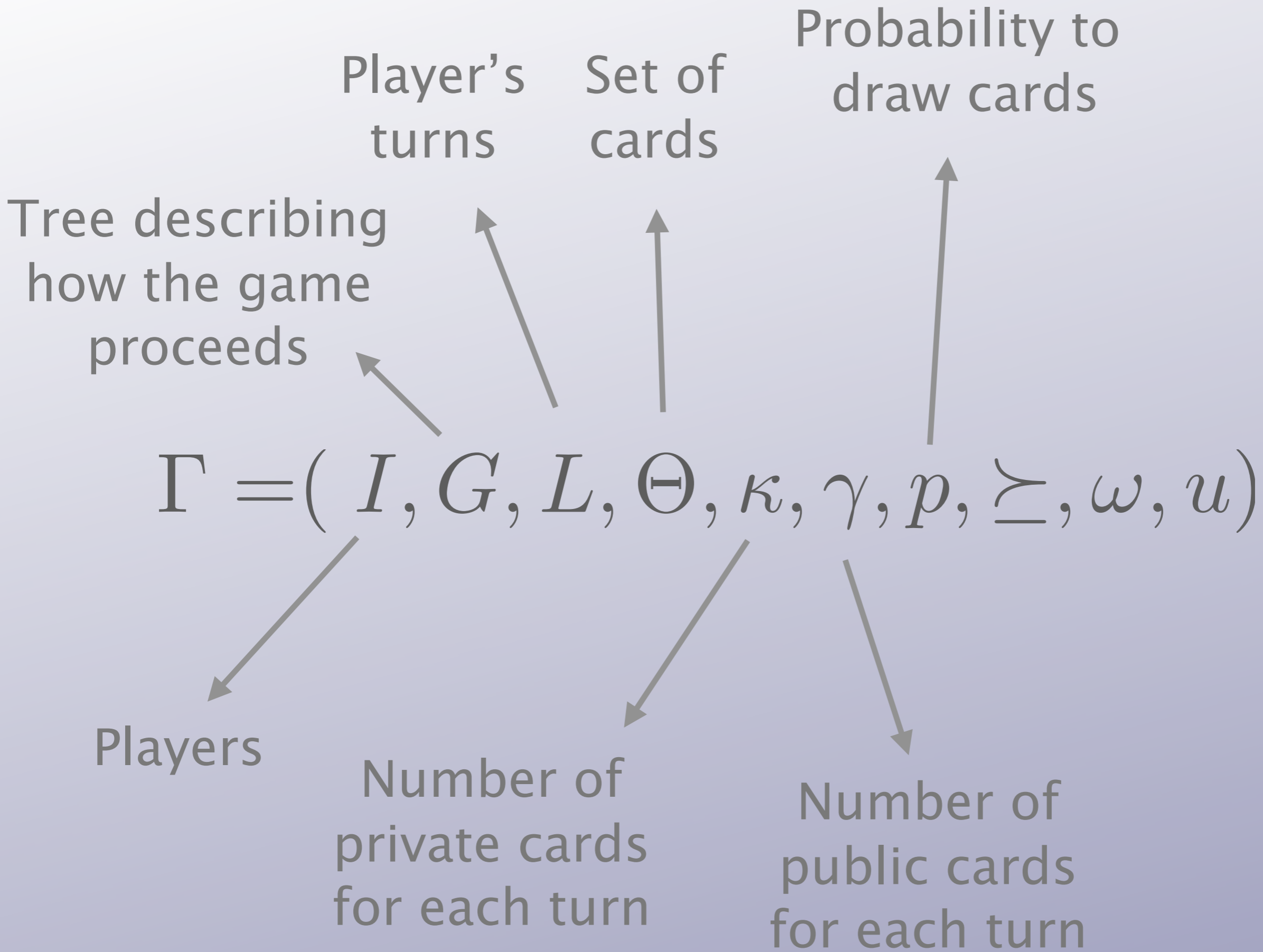
Tree describing
how the game
proceeds

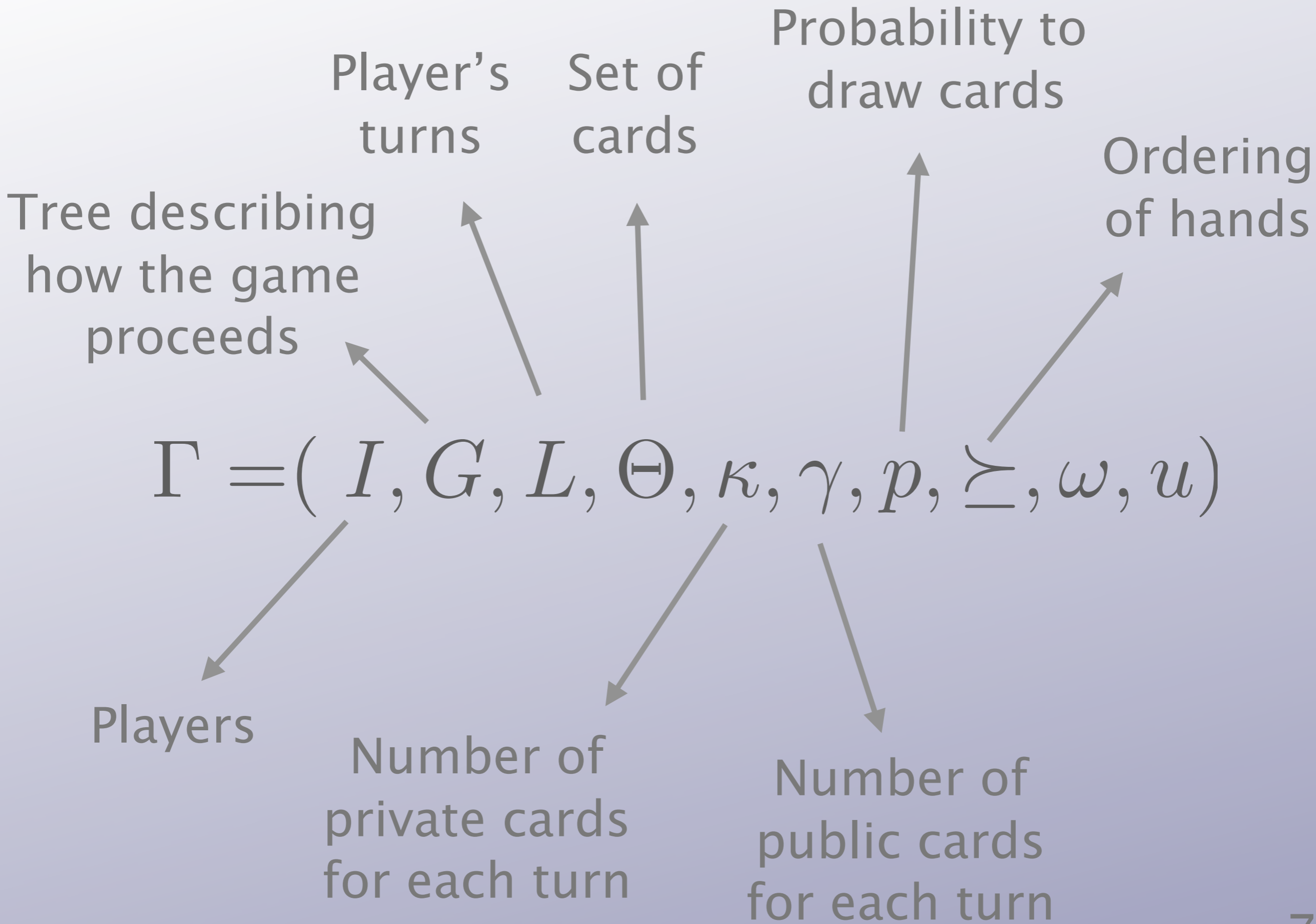
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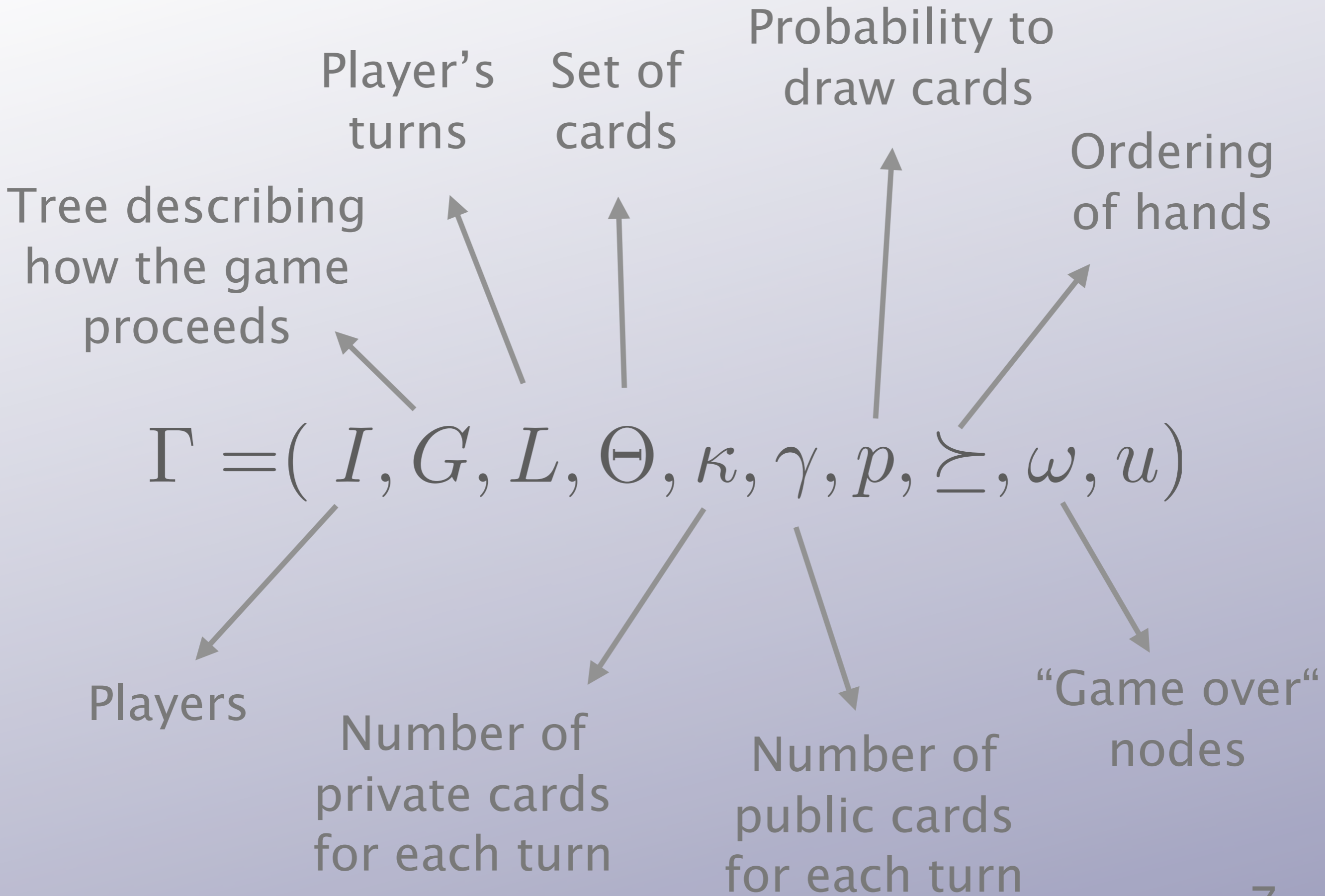
Players

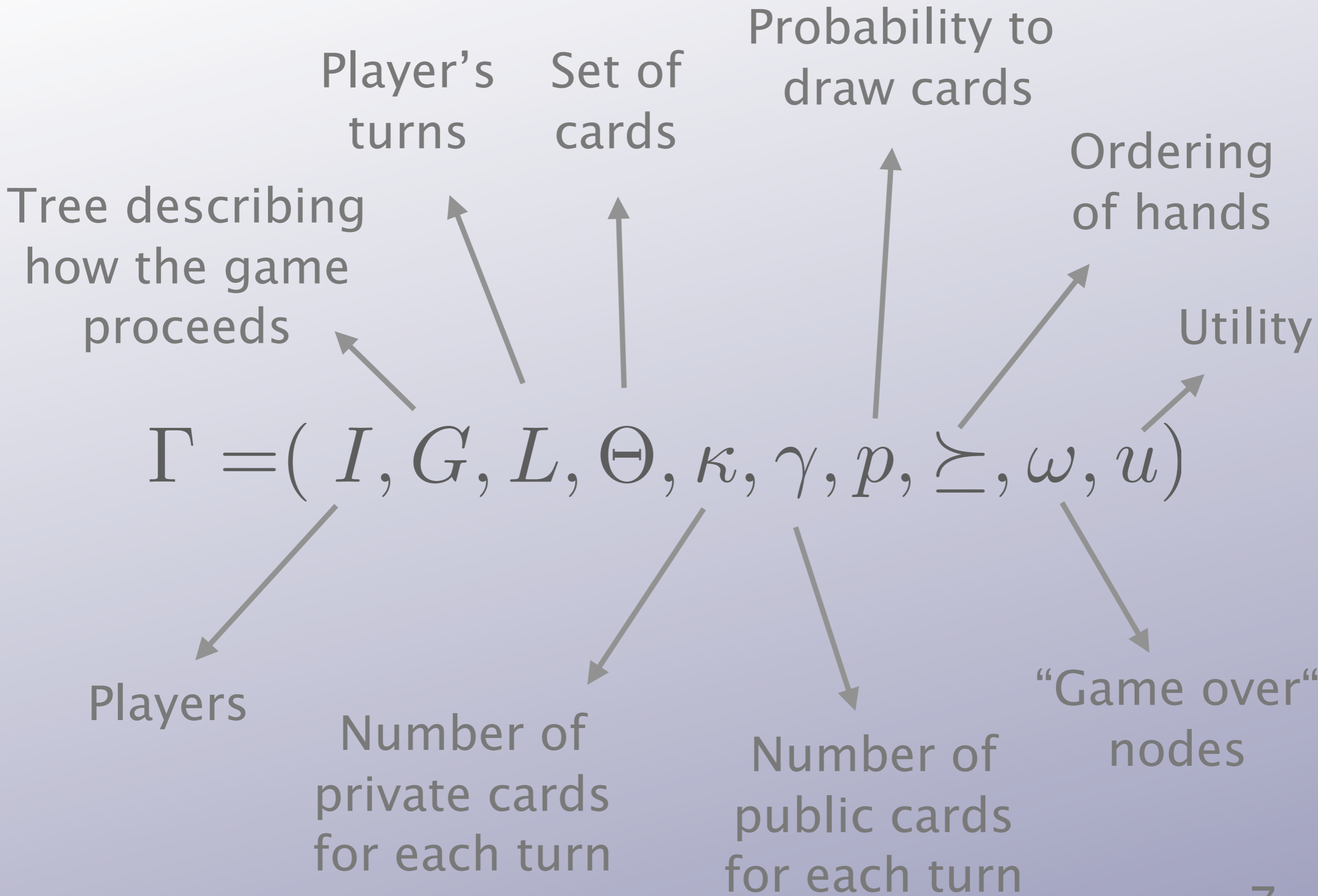
Number of
private cards
for each turn

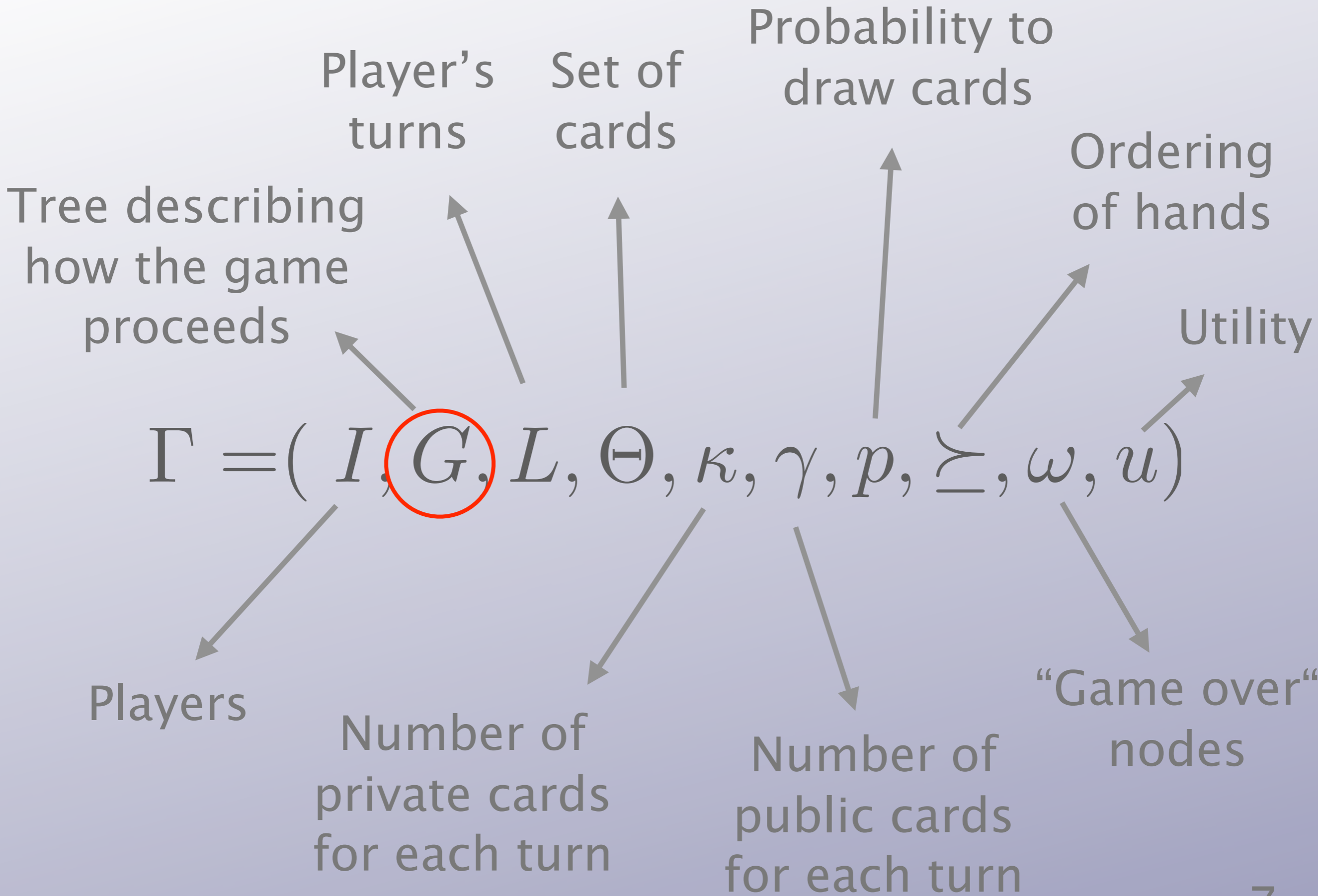
Number of
public cards
for each turn

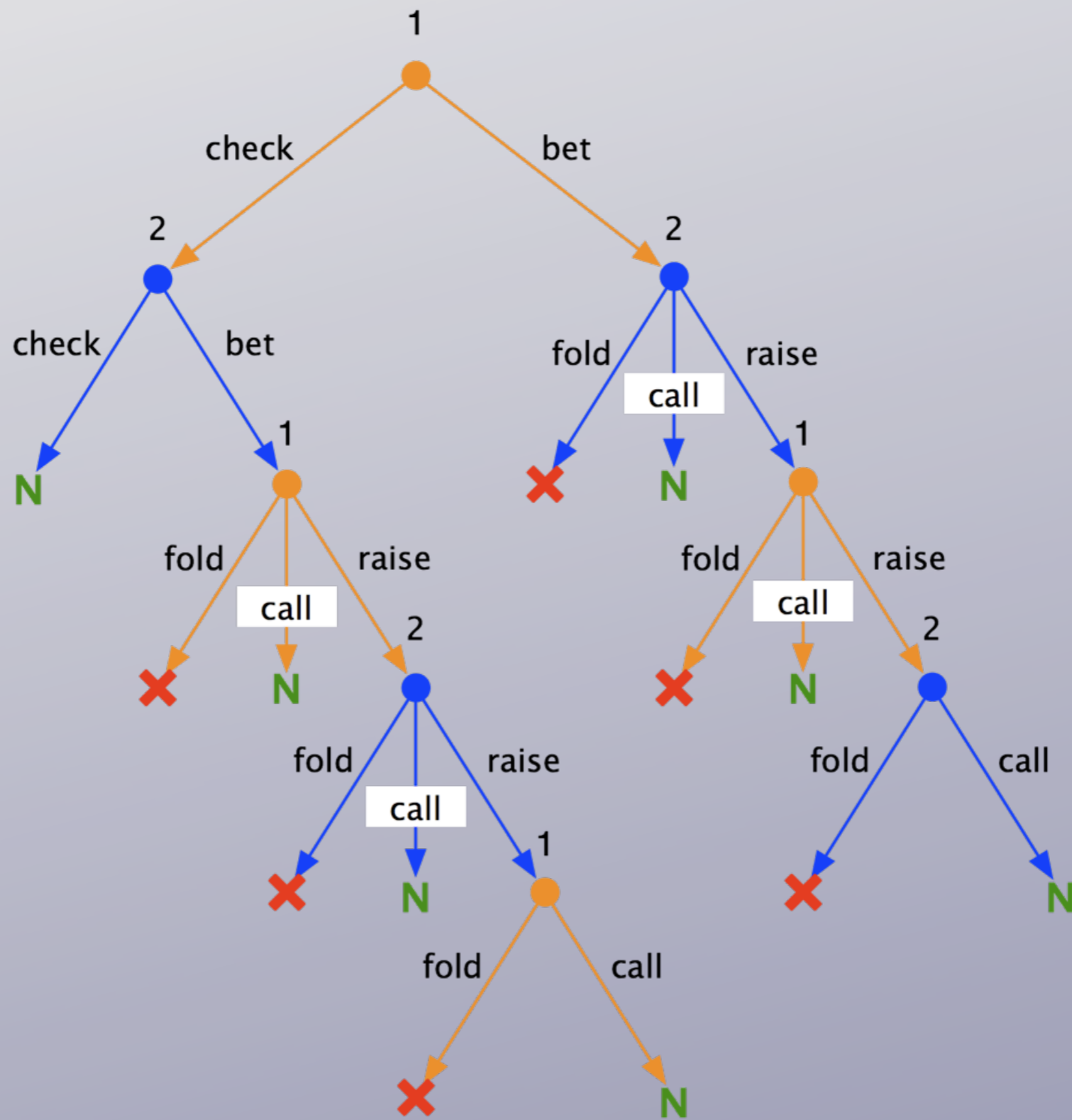


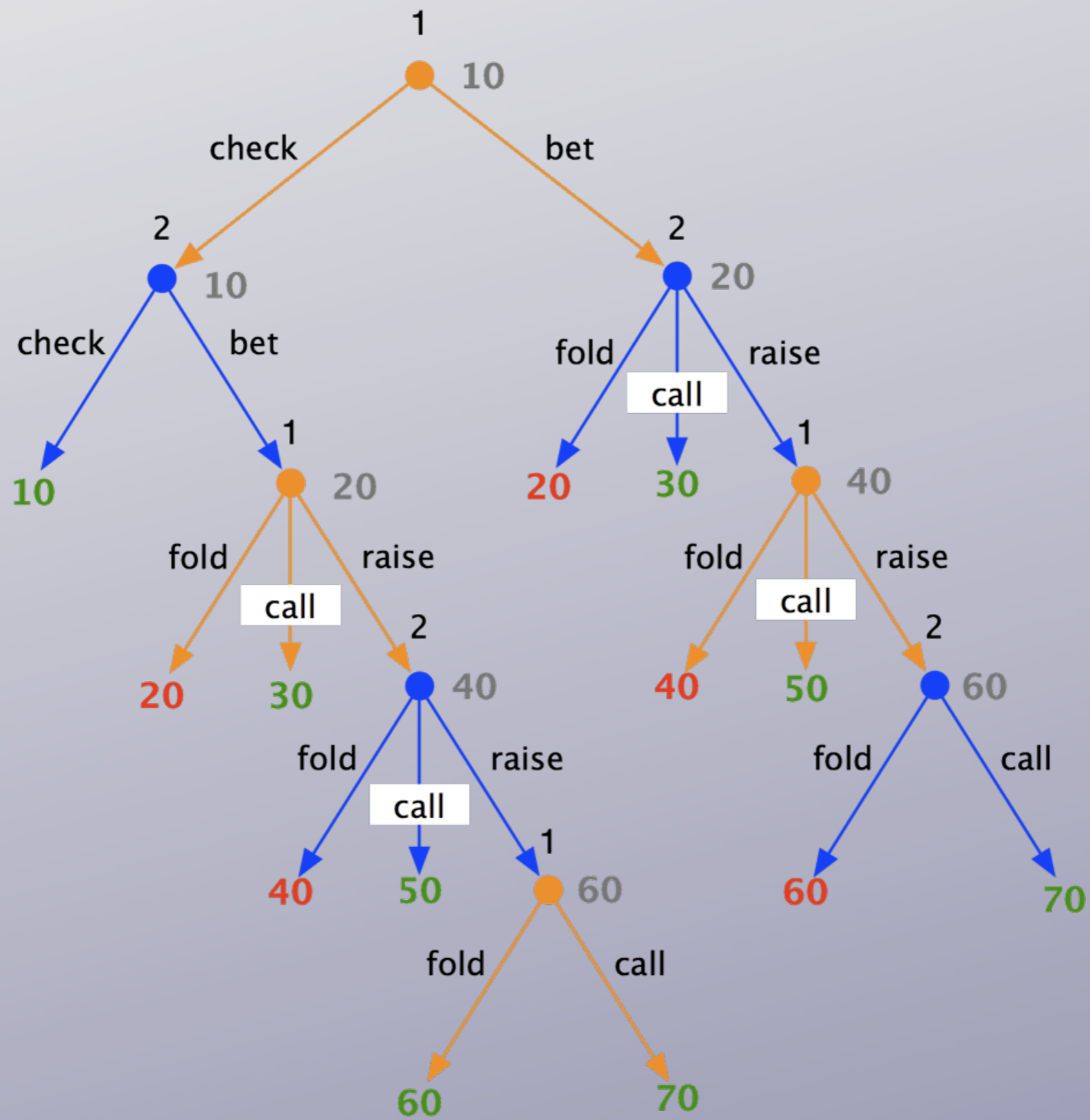


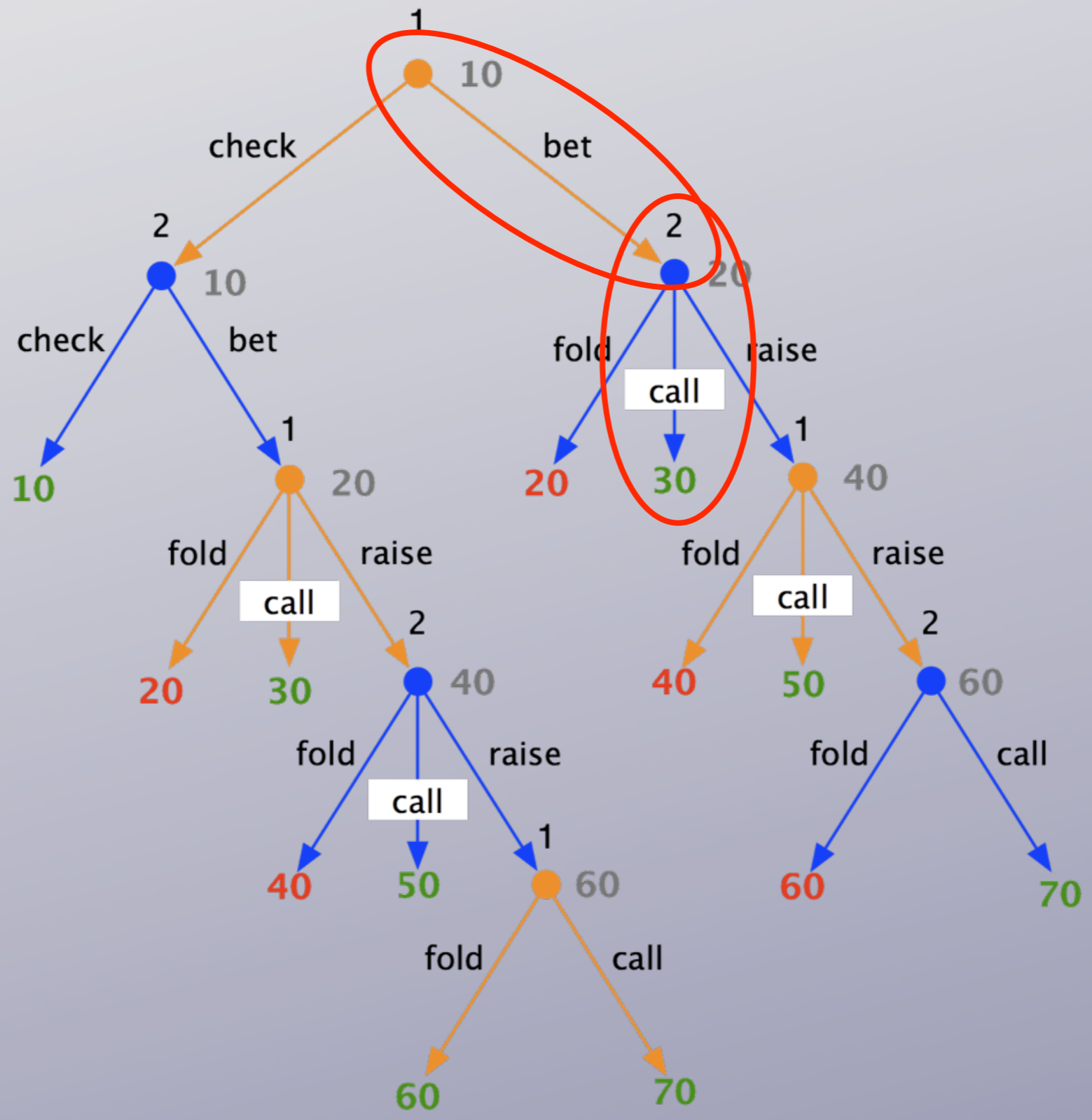












“A filtered ordered game is an extensive form game satisfying perfect recall.”

from the article

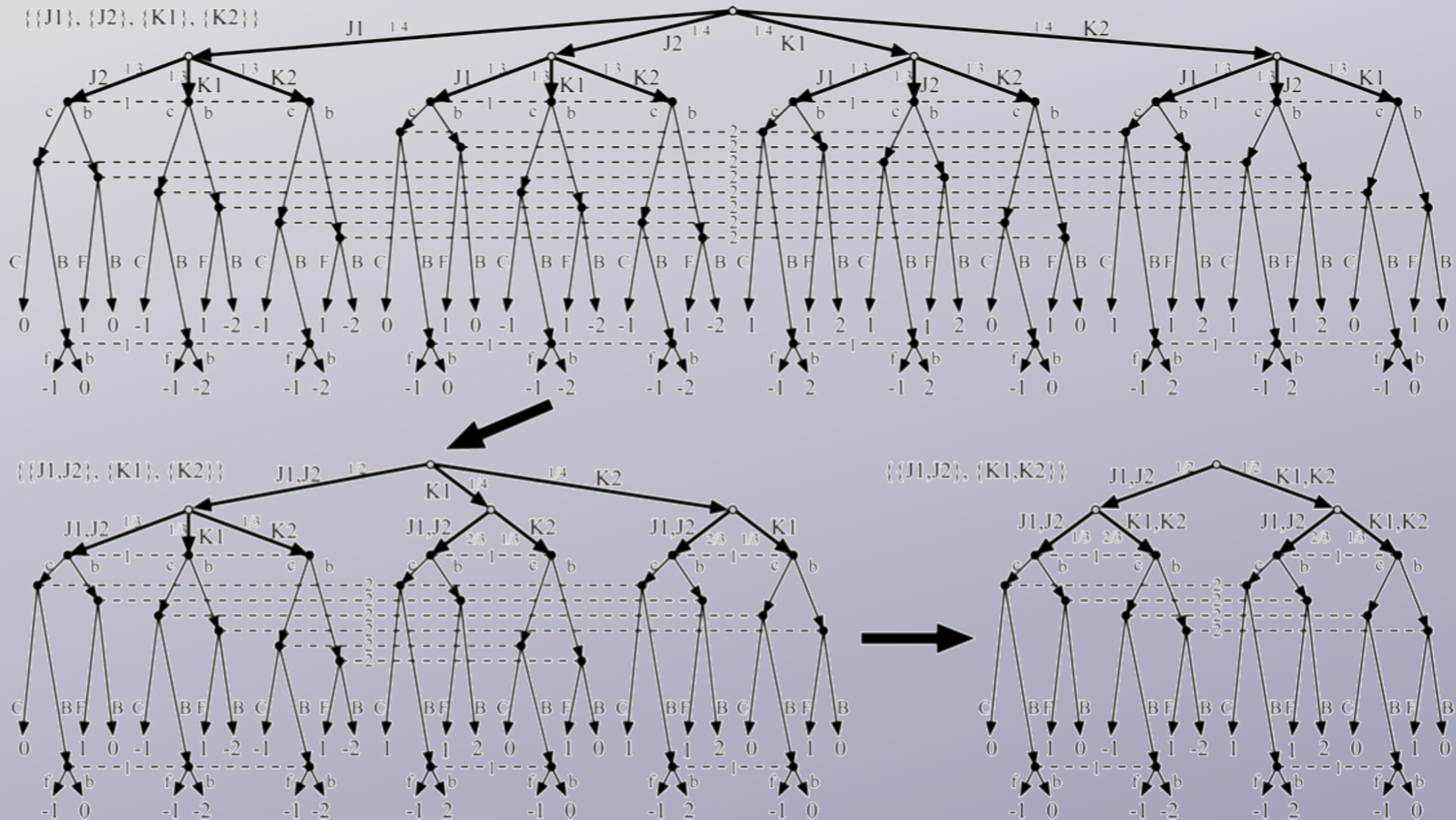
It means that we can use
behavior strategies

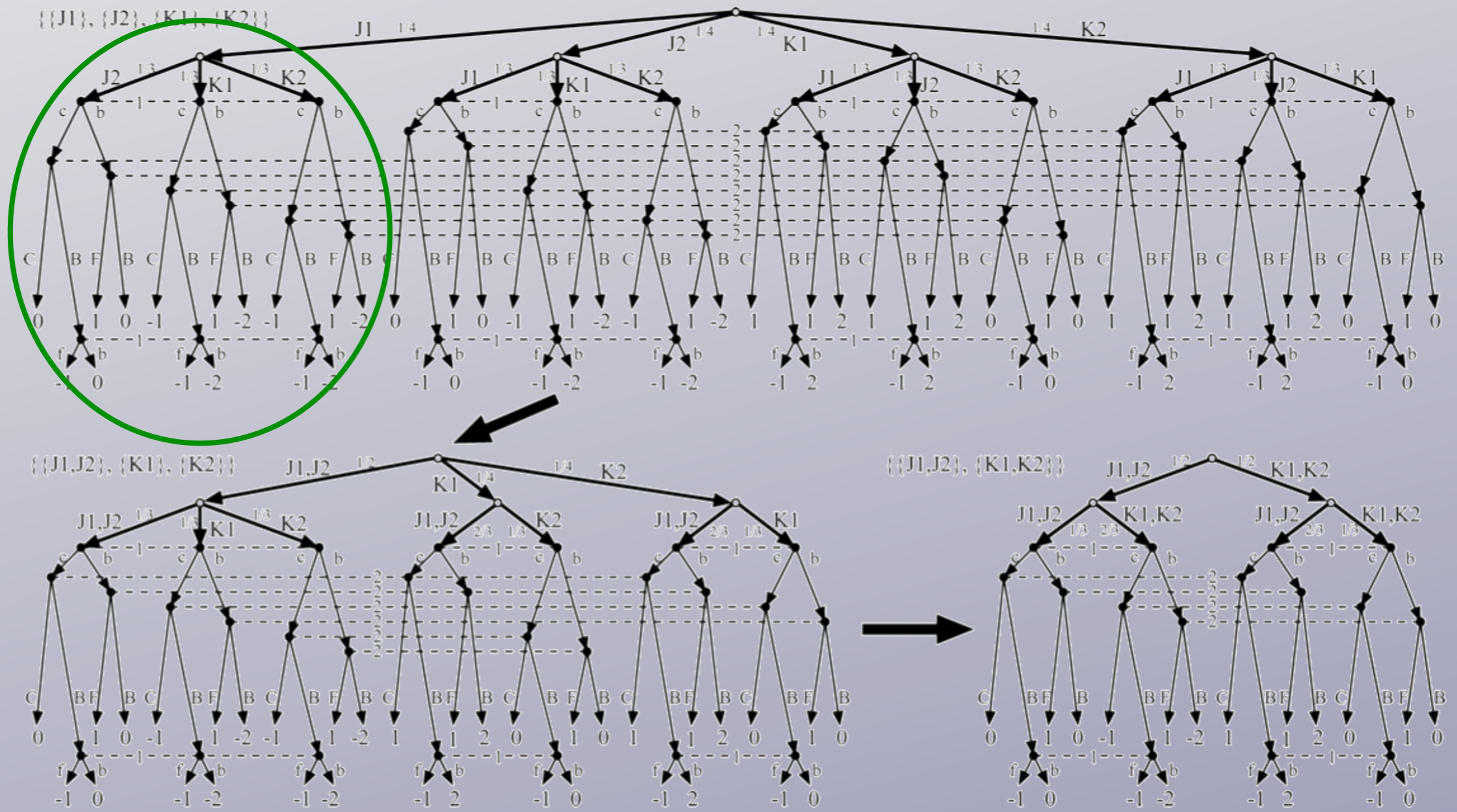
Two limitations
in generality, though.

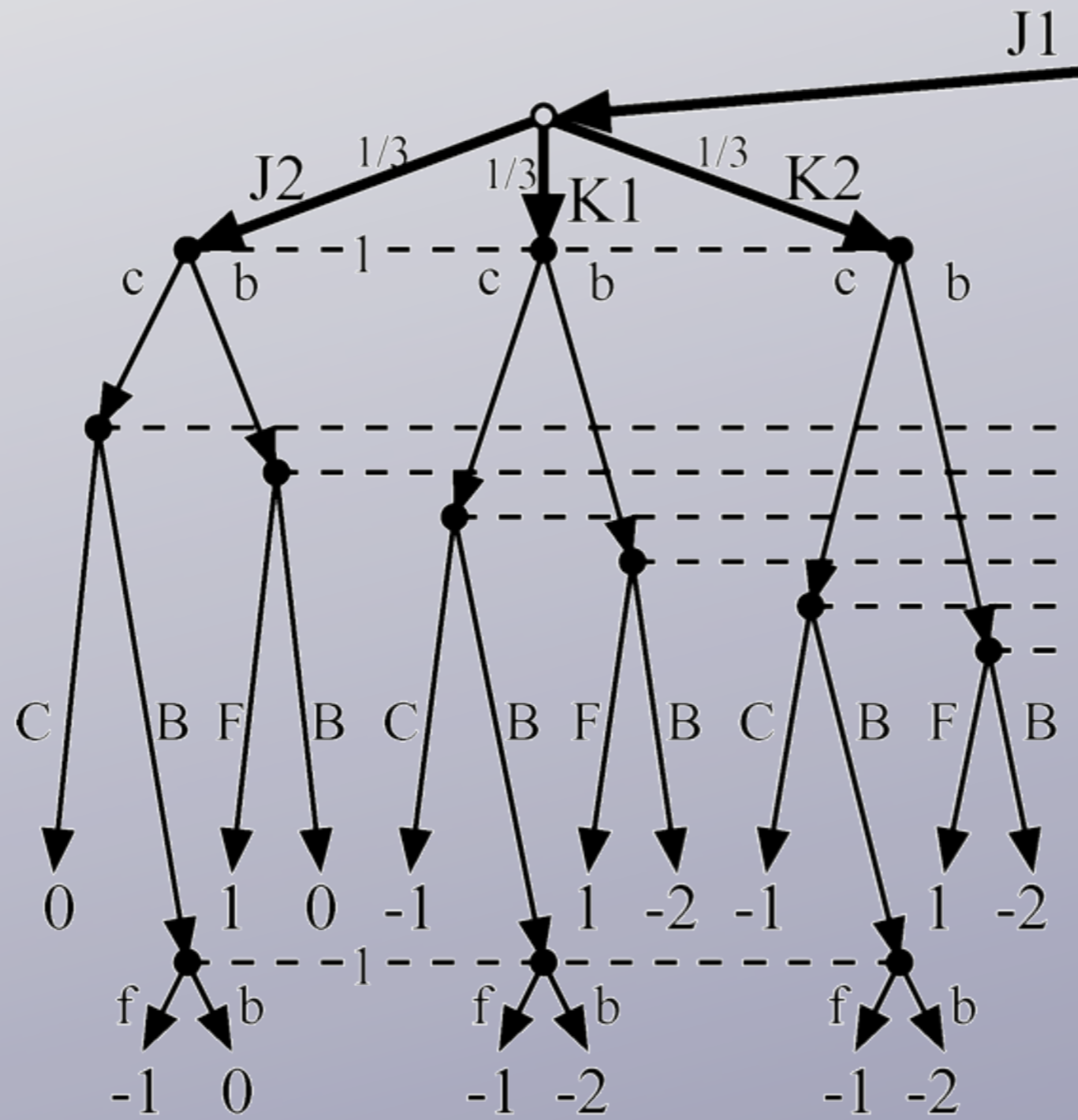
First, structure of player
actions and chance action

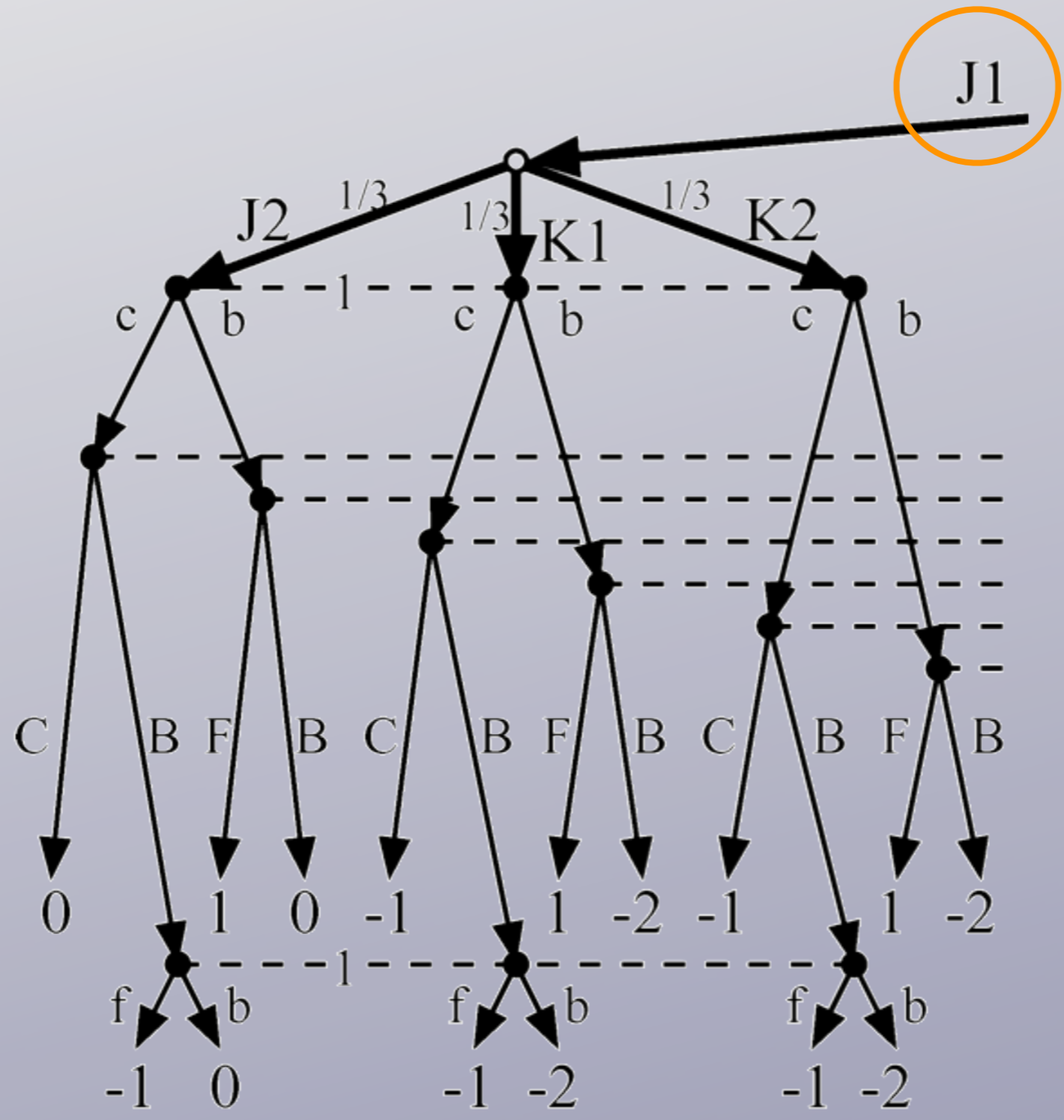
Second, the rank of hands is
the same for everyone.

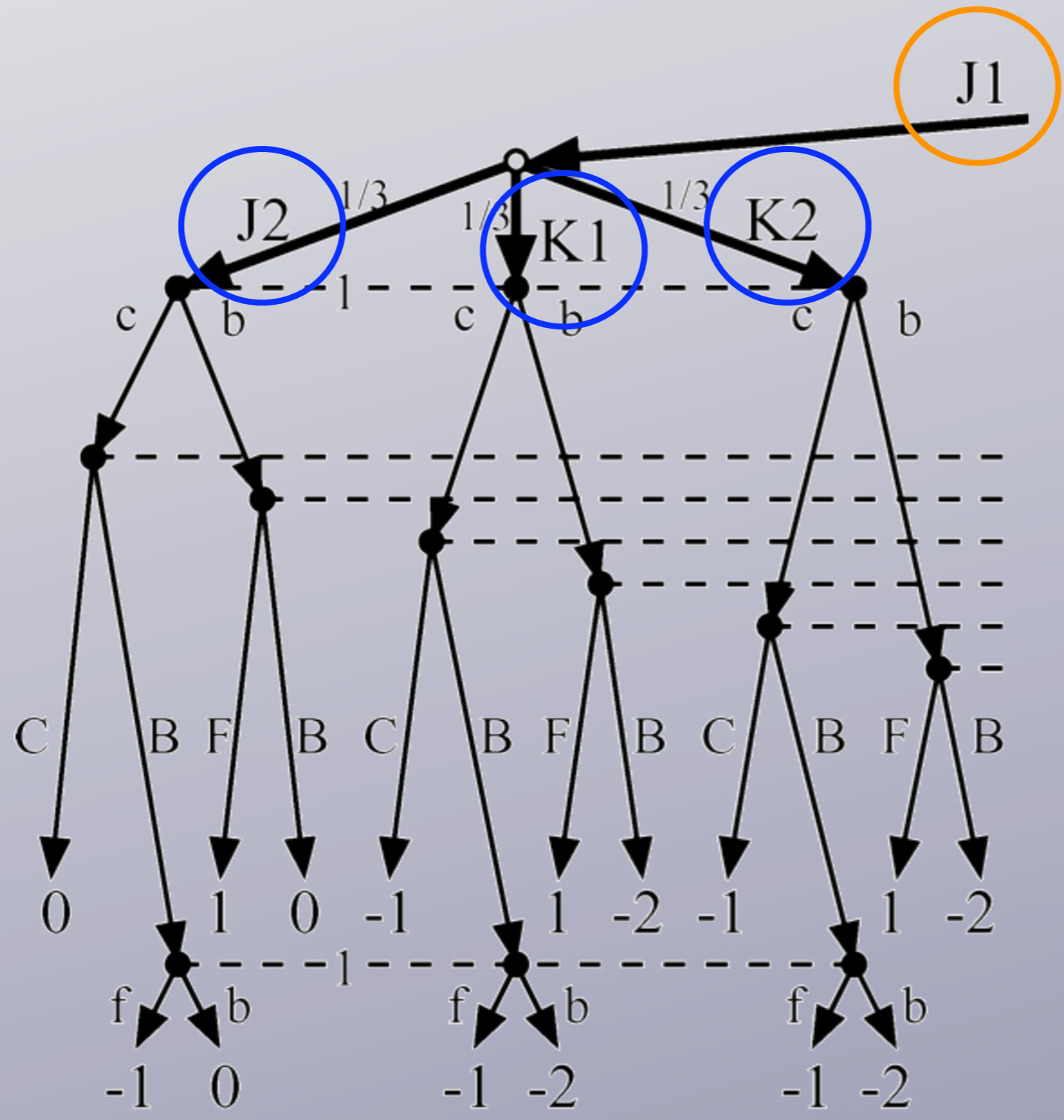
3. Filtered Signal Tree

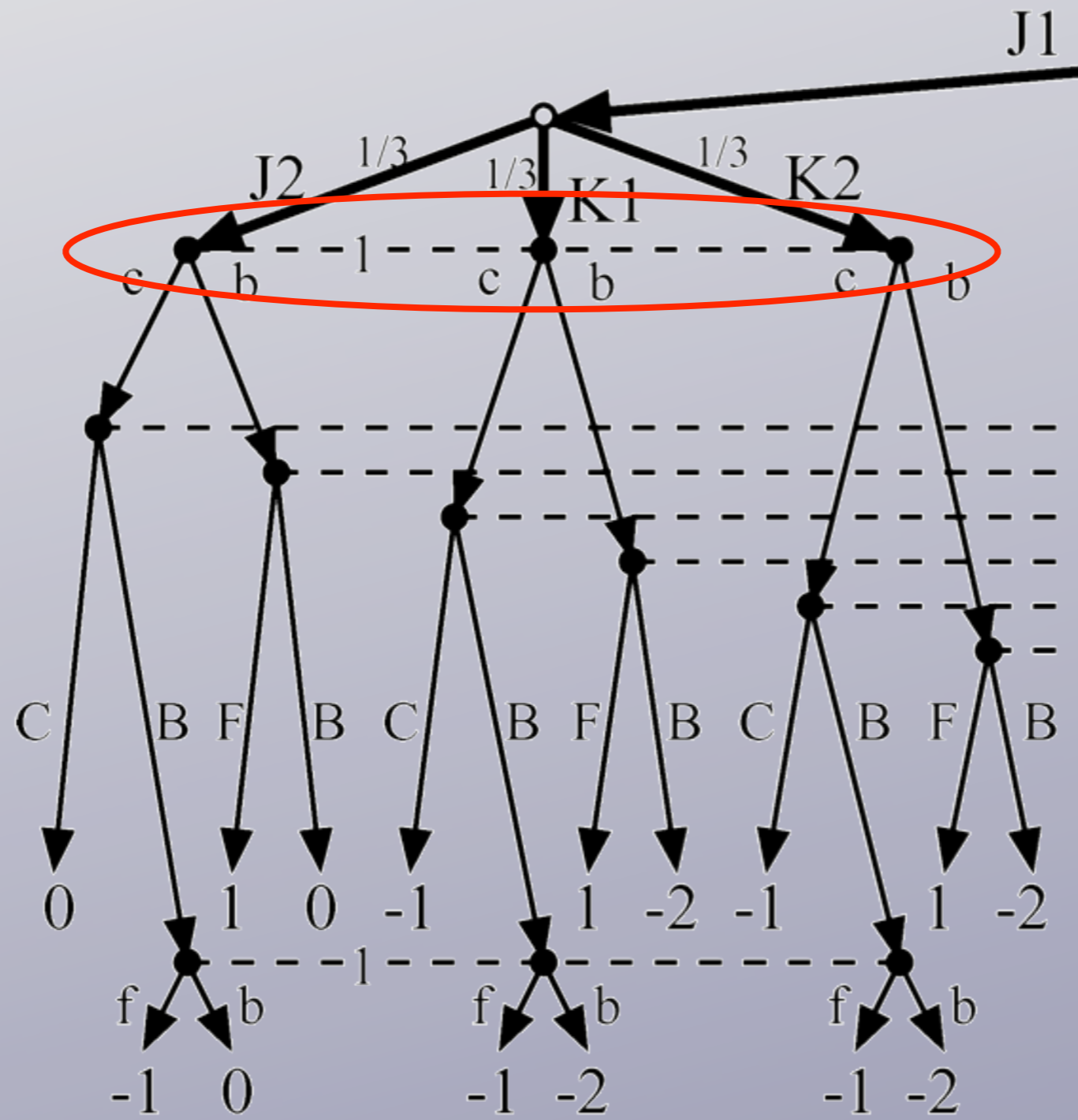


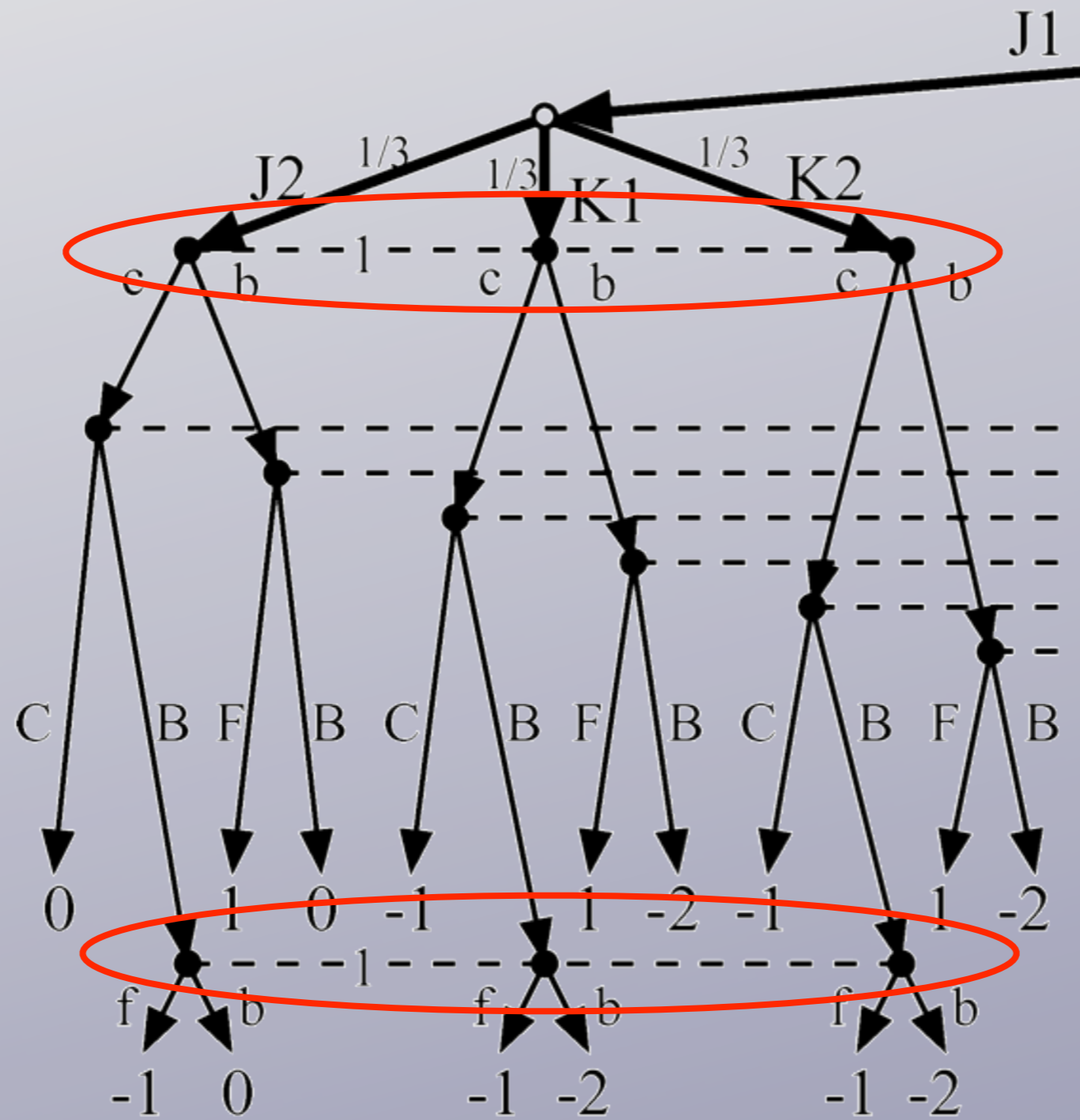




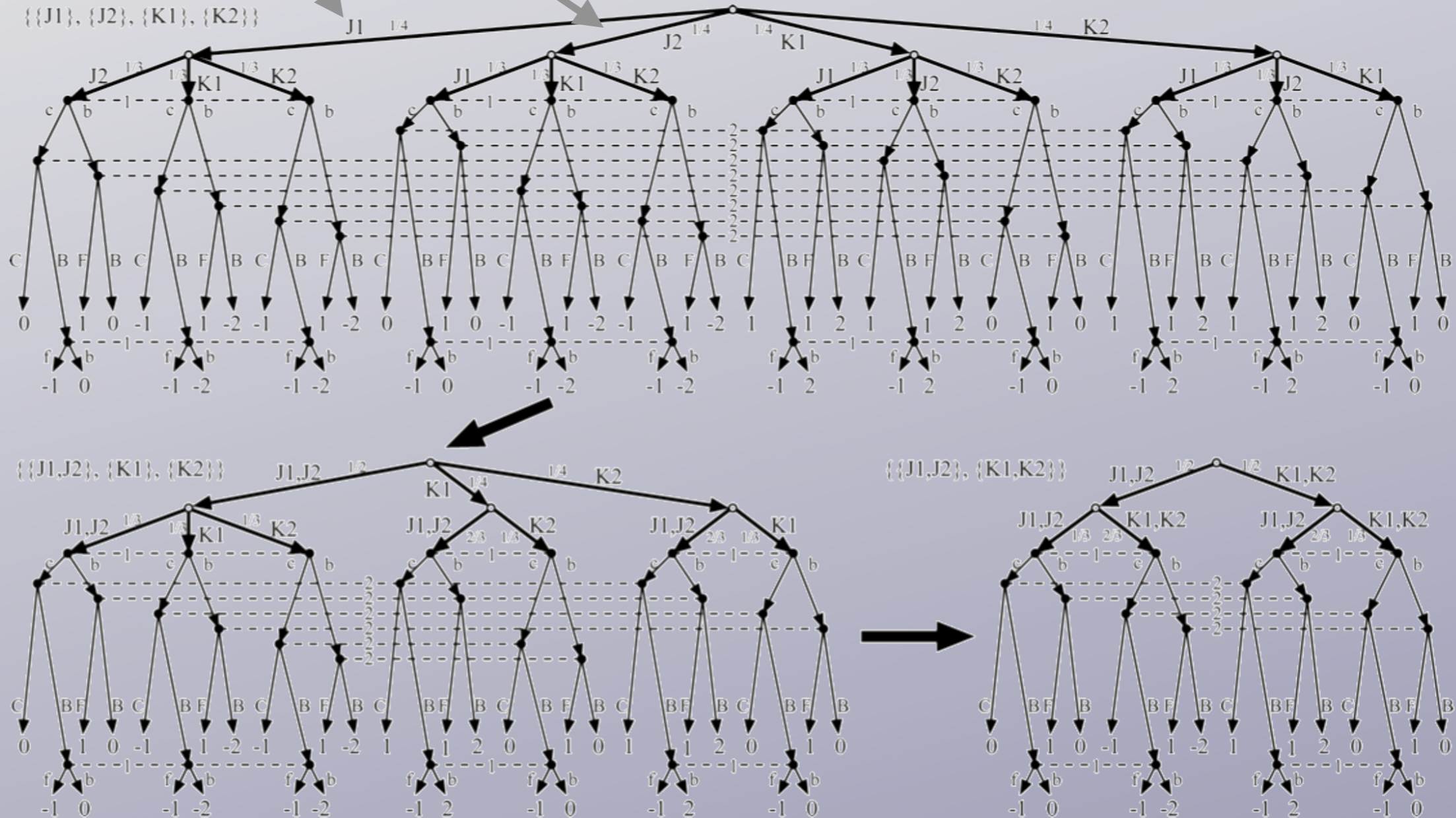




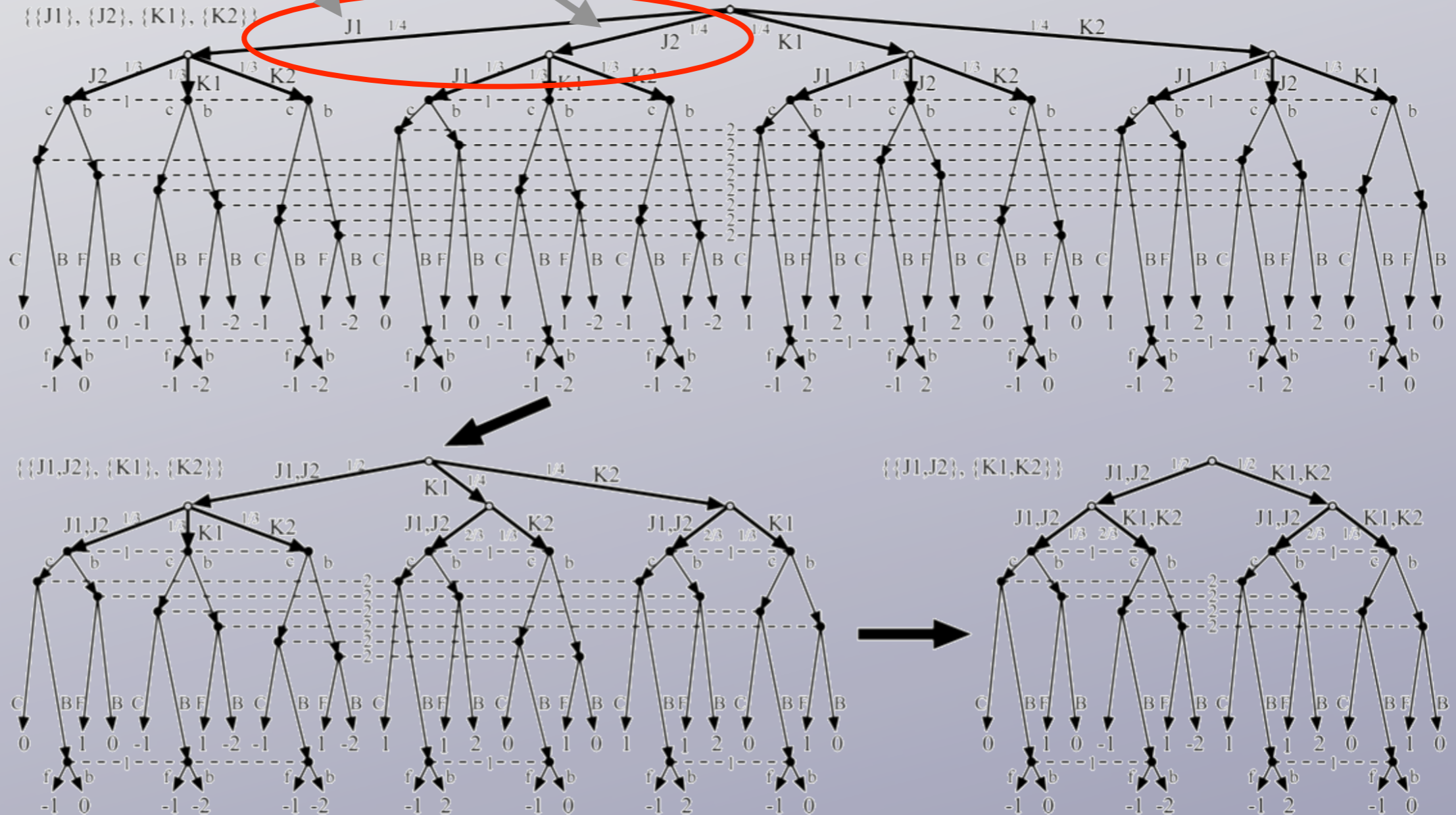




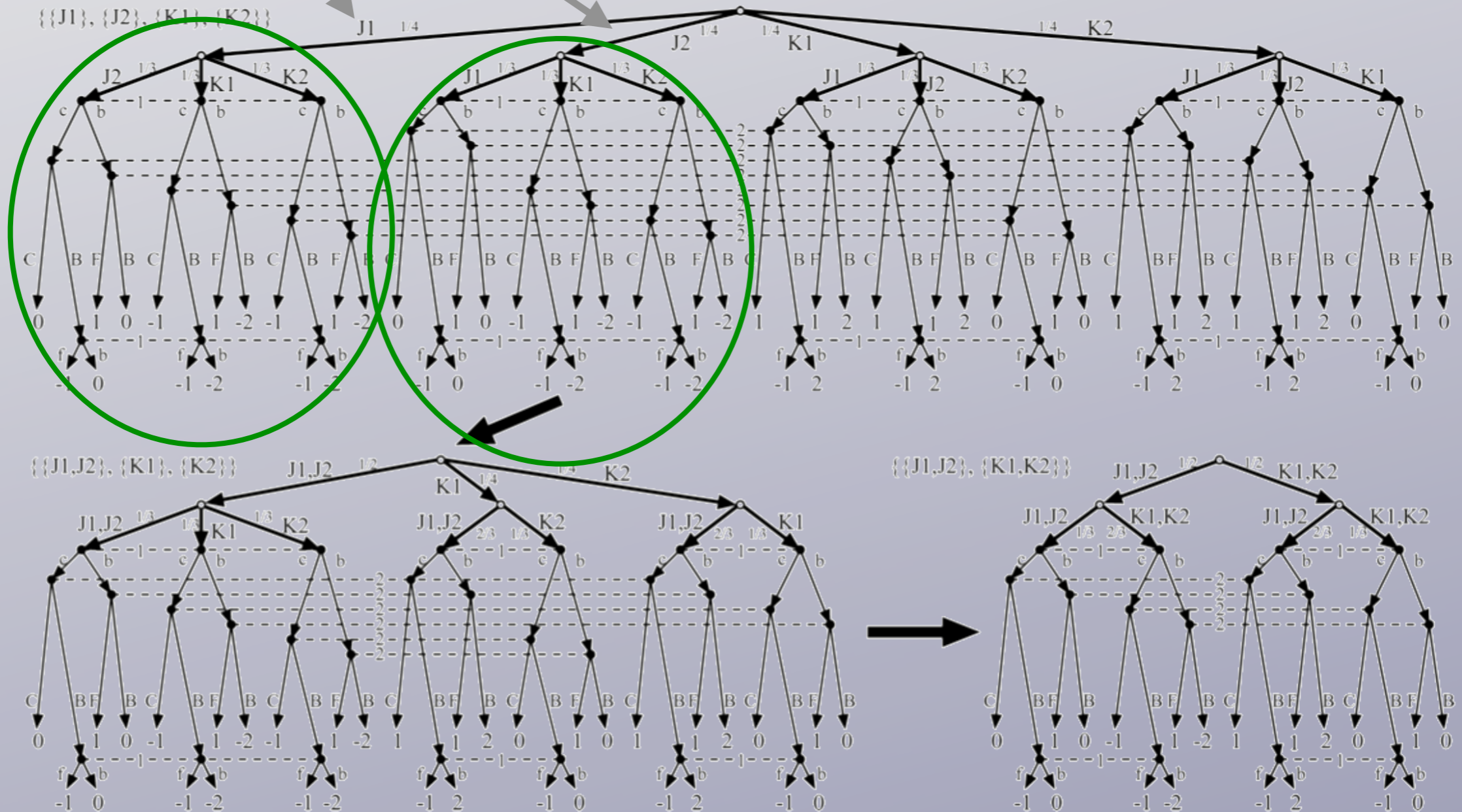
J1 J2



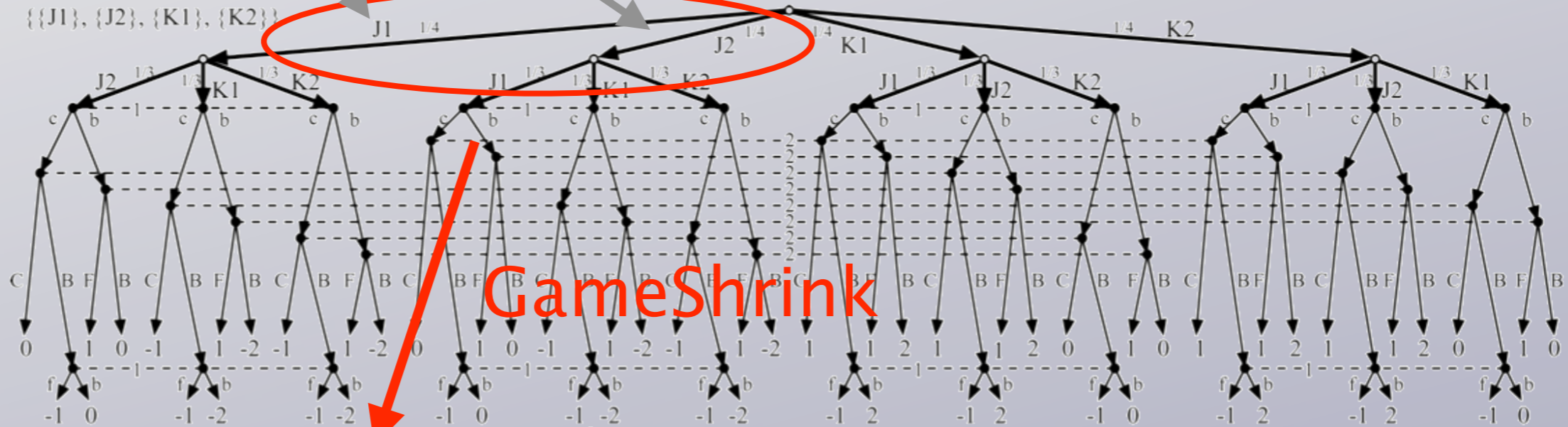
J1 J2



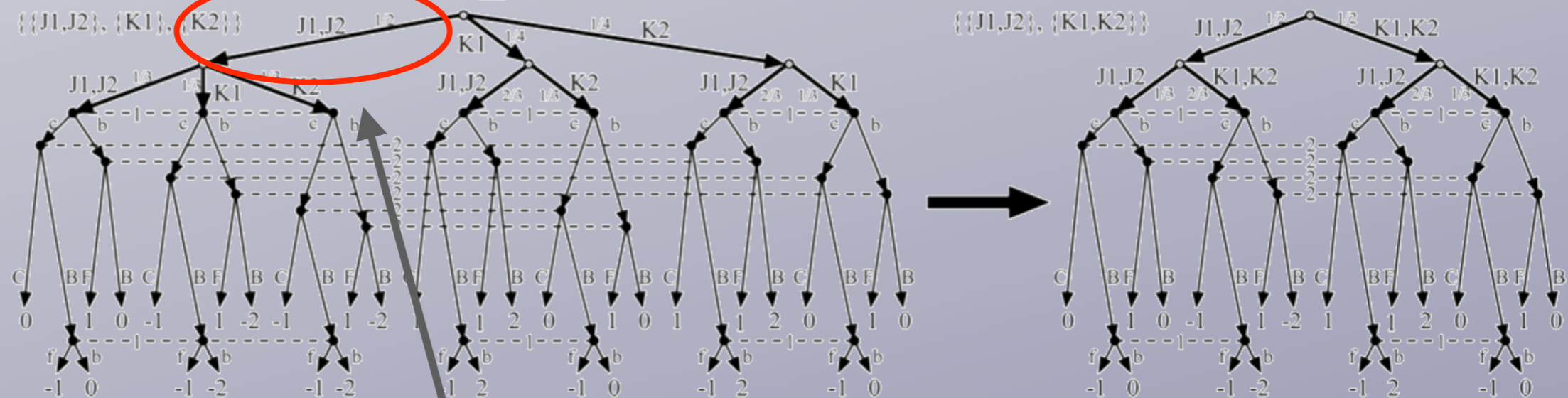
J1 J2



J1 J2



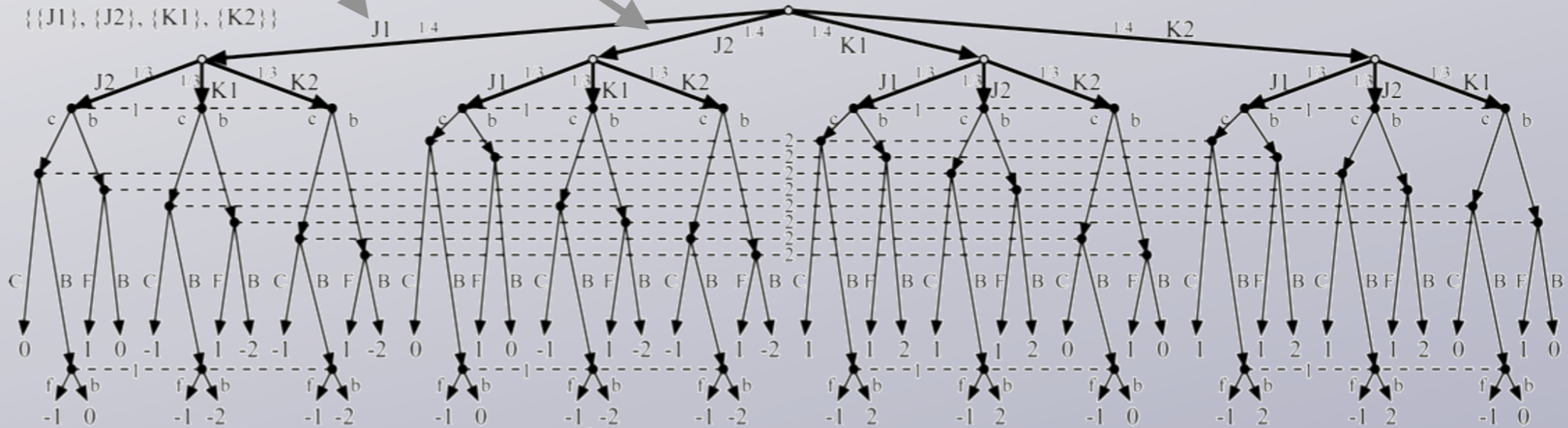
GameShrink



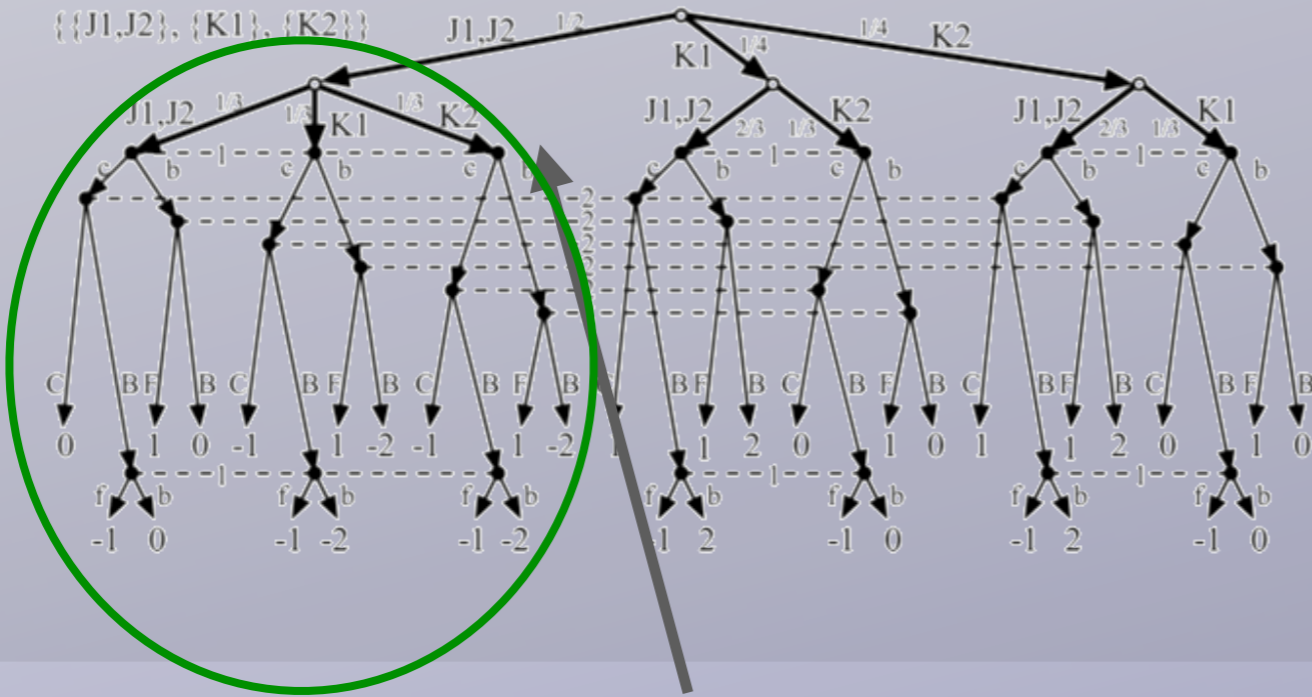
J1, J2

J1 J2

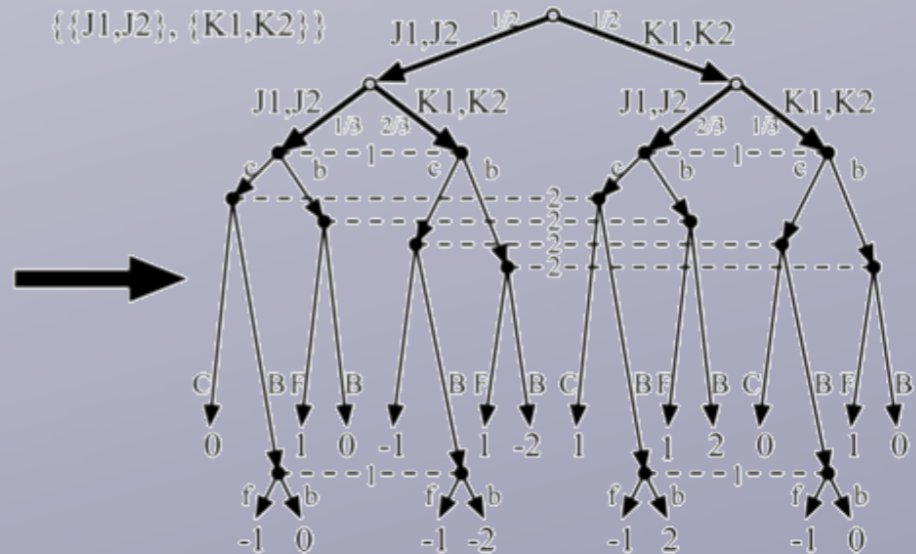
$\{\{J1\}, \{J2\}, \{K1\}, \{K2\}\}$



$\{\{J1, J2\}, \{K1\}, \{K2\}\}$

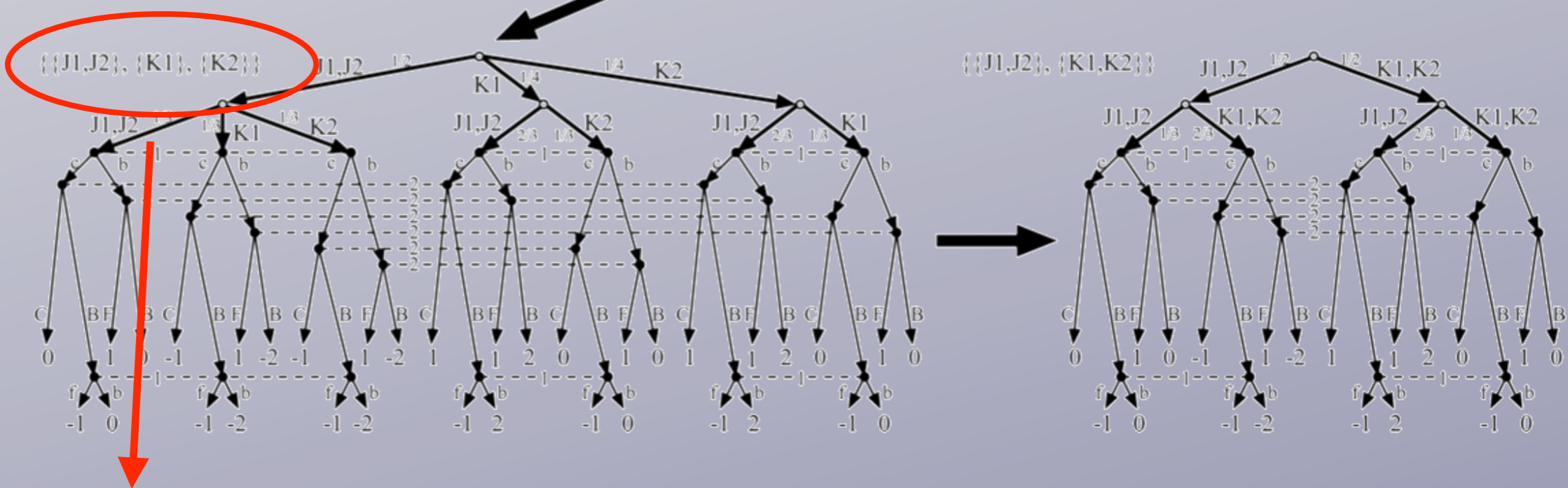
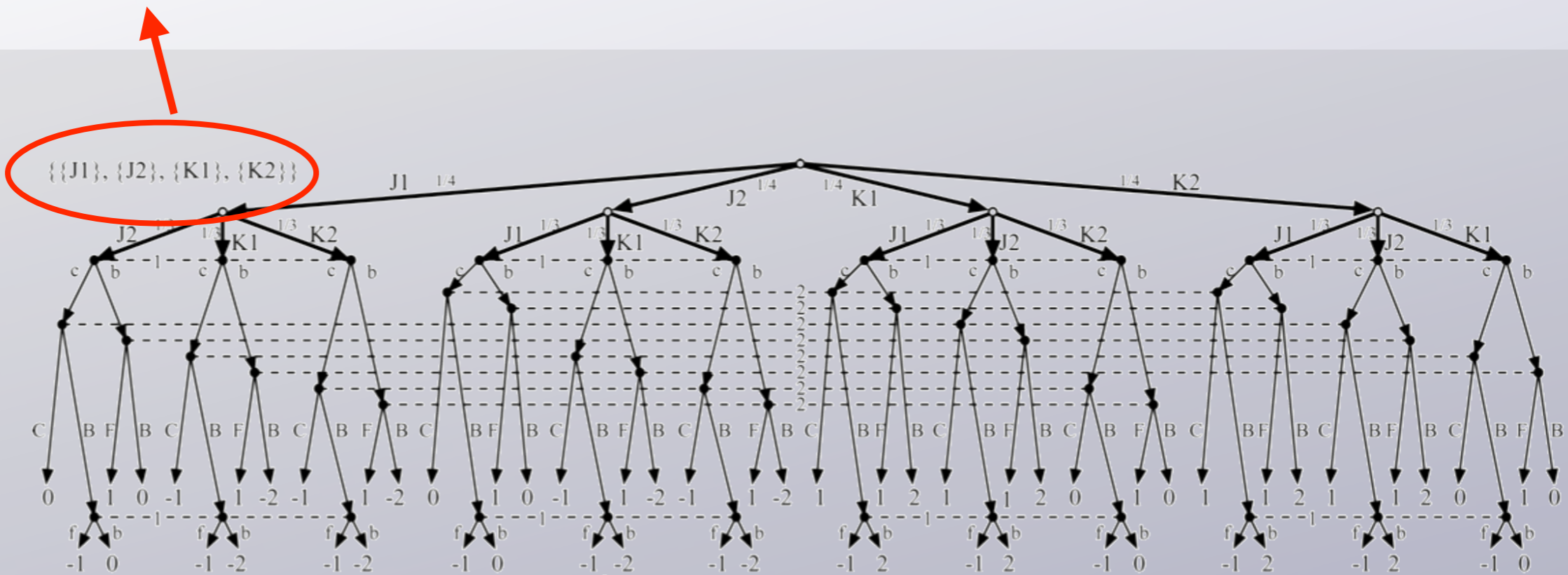


$\{\{J1, J2\}, \{K1, K2\}\}$



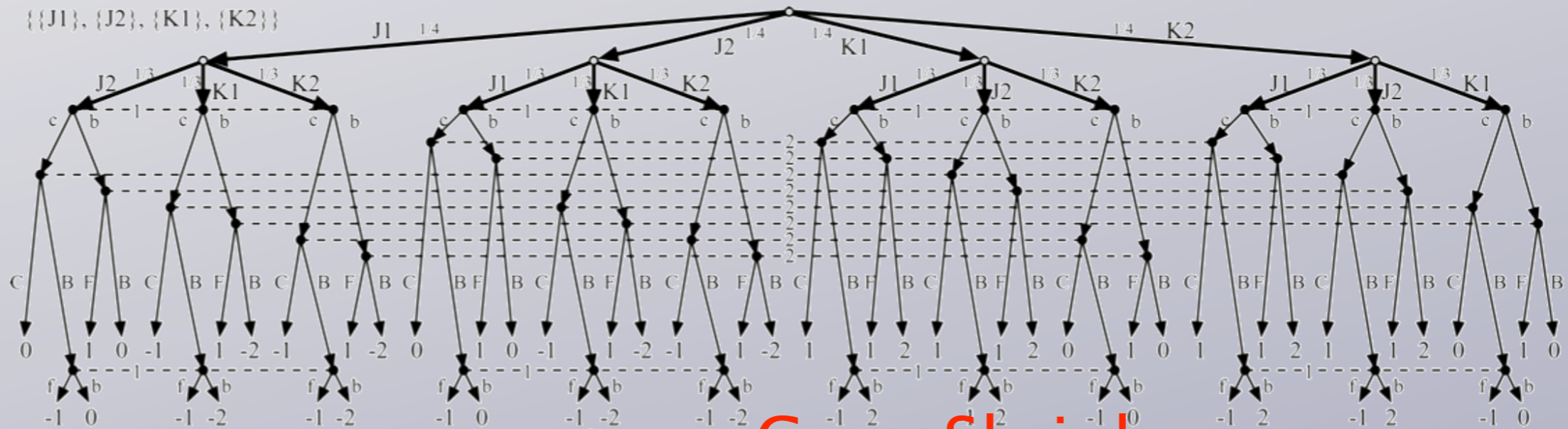
J1, J2

$\{J1, J2, K1, K2\}$



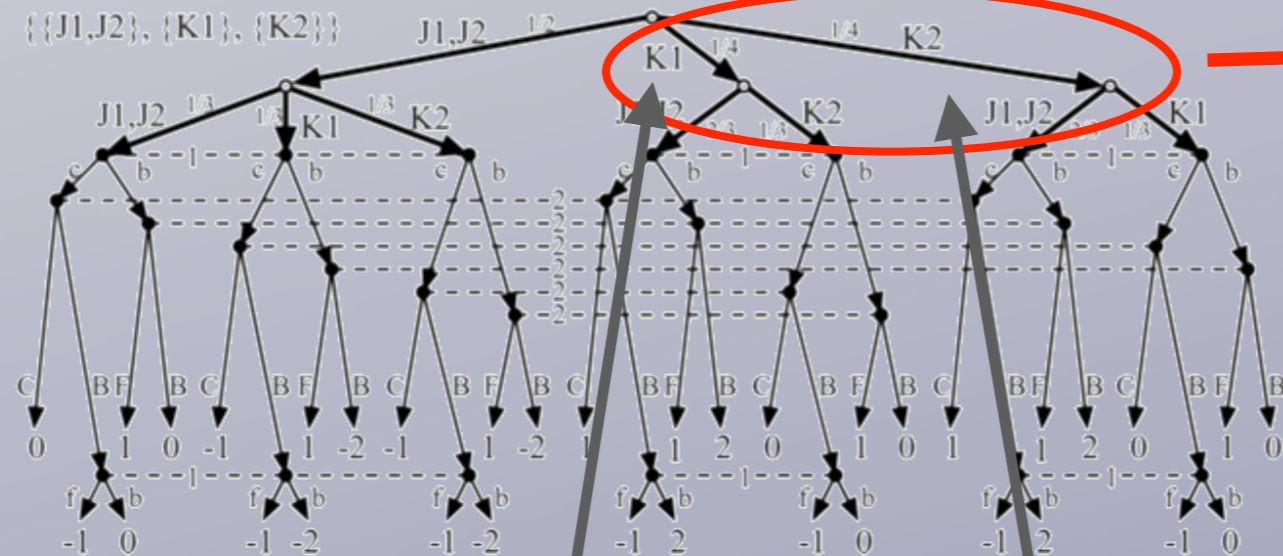
$\{\{J1, J2\}, K1, K2\}$

$\{\{J1\}, \{J2\}, \{K1\}, \{K2\}\}$

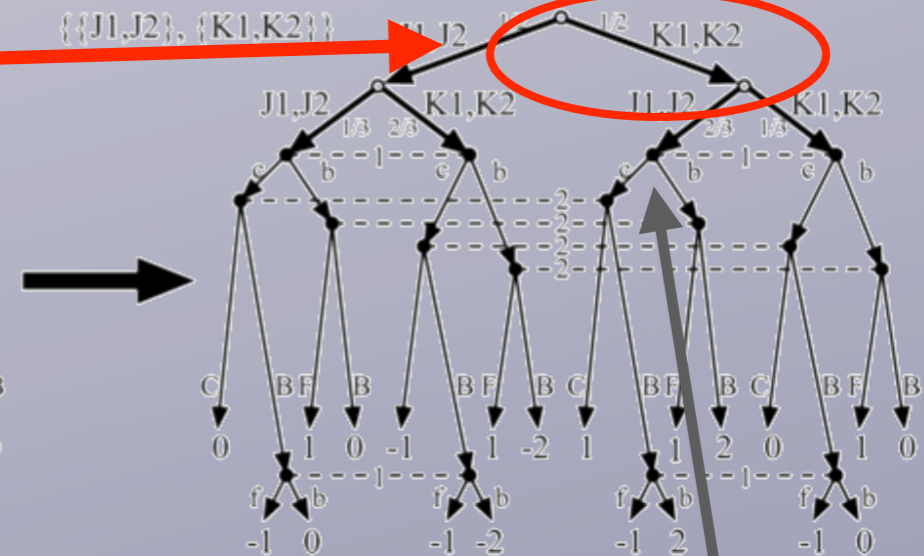


GameShrink

$\{\{J1, J2\}, \{K1\}, \{K2\}\}$



$\{\{J1, J2\}, \{K1, K2\}\}$

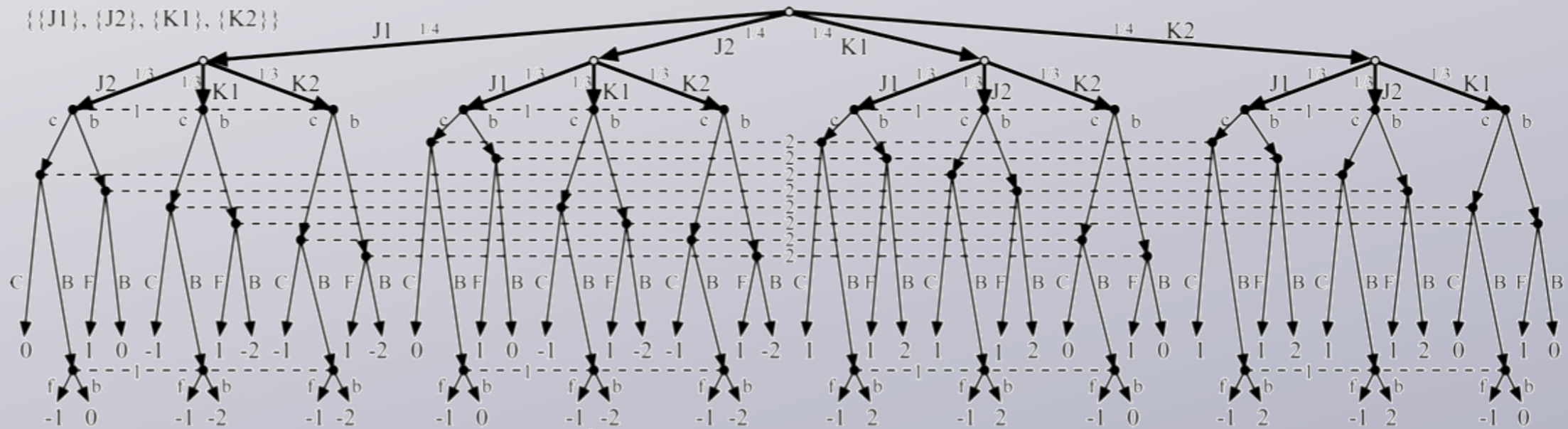


K1

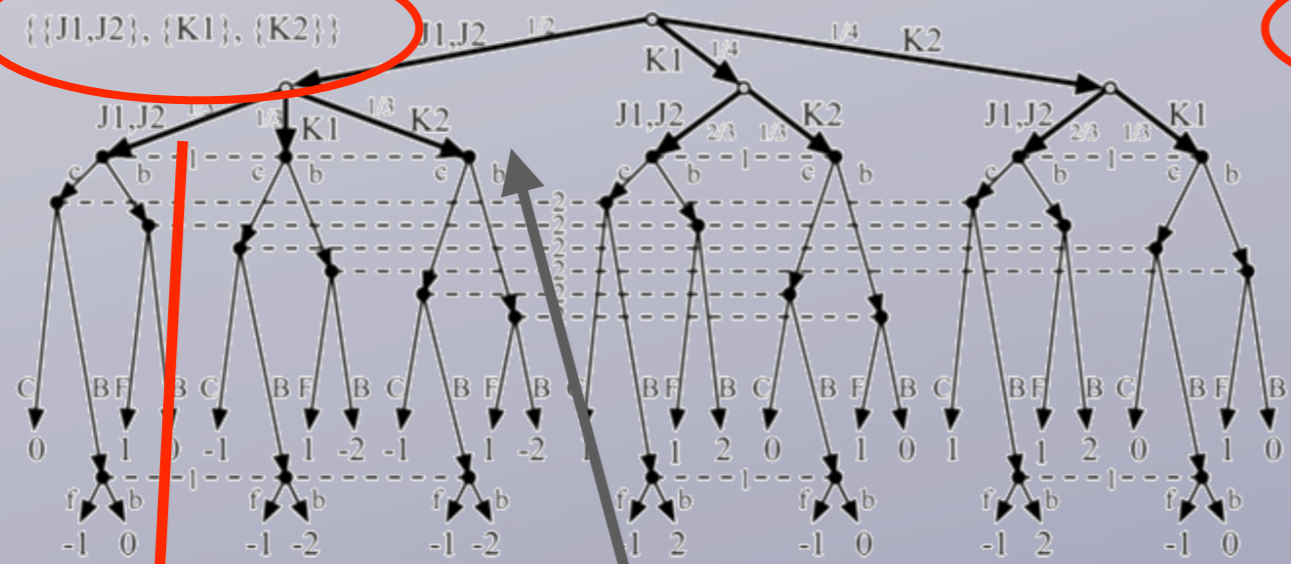
K2

K1, K2

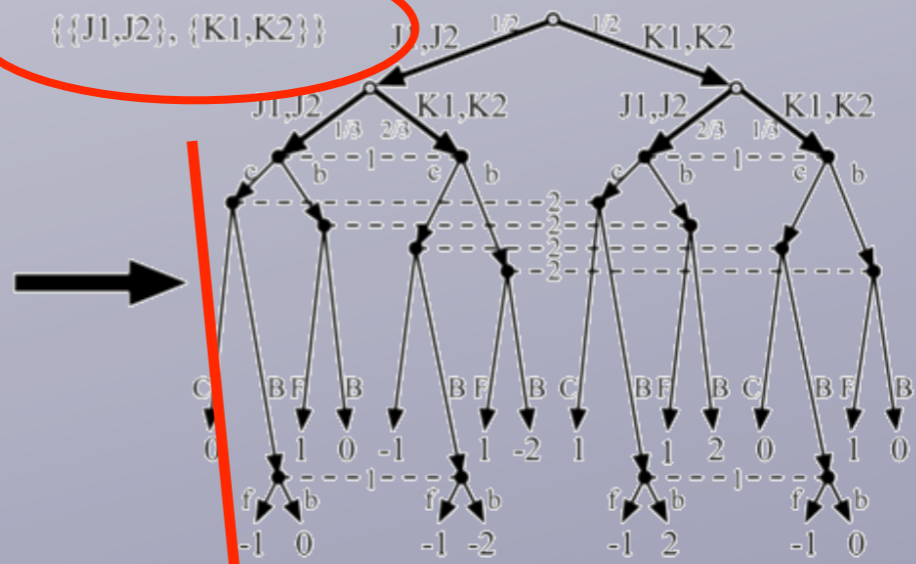
$\{\{J1\}, \{J2\}, \{K1\}, \{K2\}\}$



$\{\{J1, J2\}, \{K1\}, \{K2\}\}$



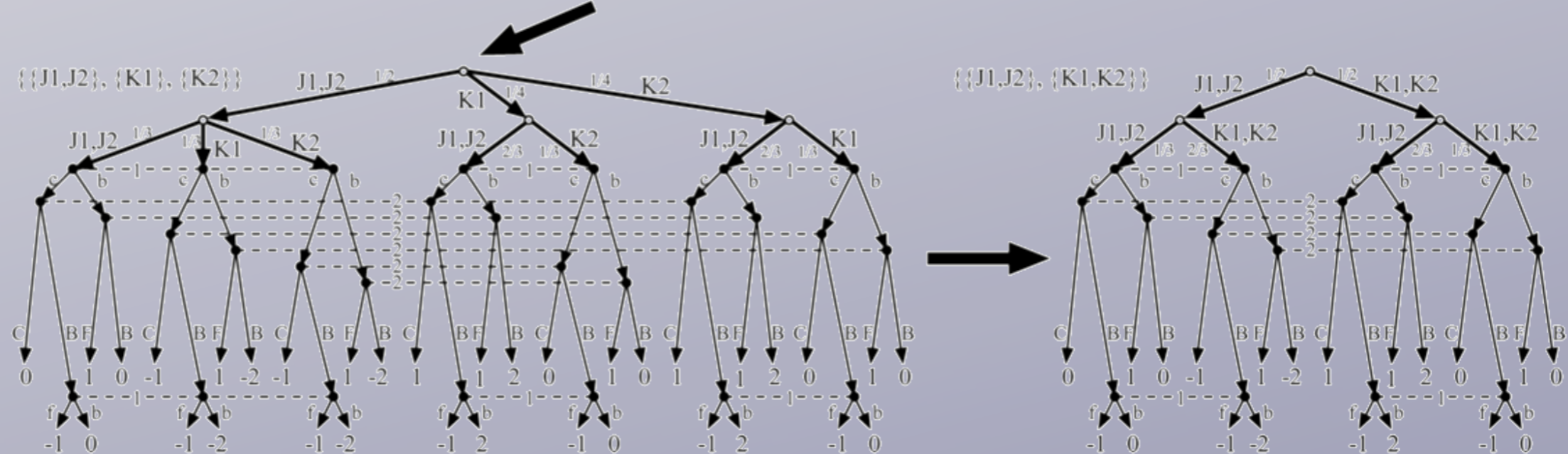
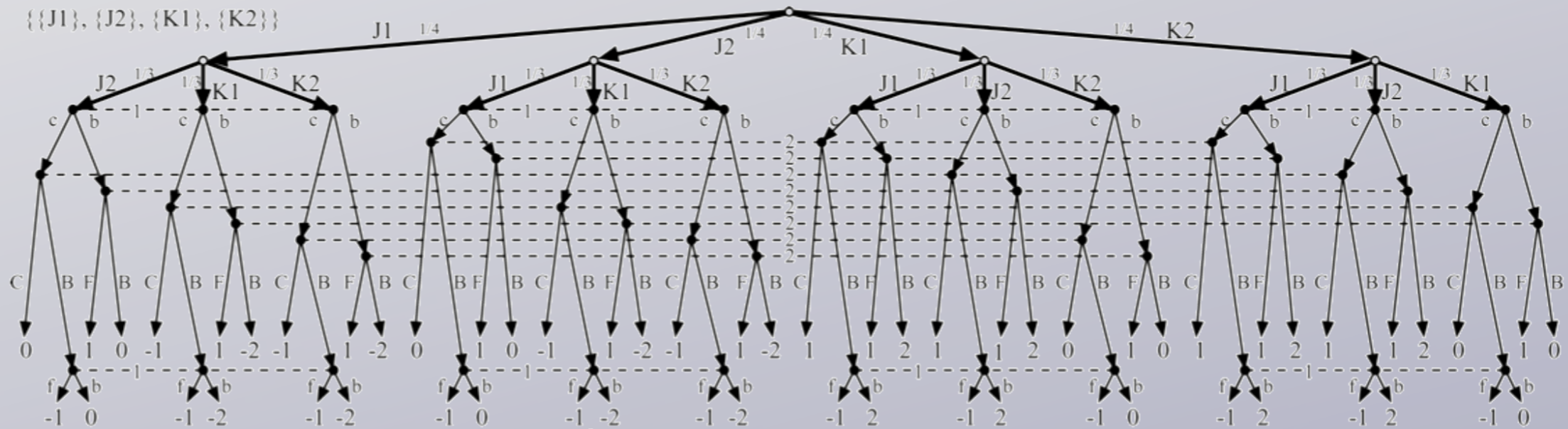
$\{\{J1, J2\}, \{K1, K2\}\}$



$\{\{J1, J2\}, K1, K2\}$

$\{\{J1, J2\}, \{K1, K2\}\}$

113 nodes



39 nodes

r rounds, b nonterminal leaves

size of signal tree is at most

$\frac{1}{b^r}$ size of game tree

in our case, $\frac{1}{b^r} = 0,003$

Algorithm in $O(n^2)$

4. Main Theorem

GameShrink does not
modify Nash equilibria.

GameShrink : algorithm for
ordered game isomorphic
abstraction transformation

Conclusion & Discussion

Main Points (x3)

1. Create a smaller,
equivalent game.

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equivalent game.

3.1 billion to 6 millions

2. Apply on games with
ordered signals

3. Calculated Nash equilibrium for Rhode Island Hold'em

Weaknesses

1. Approximations to
crack larger games.

2. Not all abstractions are used

3. Limits of generality

One last thing

3. Calculated Nash equilibrium for Rhode Island Hold'em

TRY IT!

www.cs.cmu.edu/~gilpin/gsi.html