

Modeling with ClaferIG

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Part I

Quick glance at Clafer

Clafer 1: Clafer model

- 1 Processor
 - 2 Core 1..2
-

Model vs instances

Clafer 4: Clafer model

- 1 Processor
 - 2 Core 1..2
-

Clafer 5: First instance

- 1 Processor
 - 2 Core
-

Clafer 6: Second instance

- 1 Processor
 - 2 Core1
 - 3 Core2
-

Clafers 7: Model with constraints

- 1 `Person *`
 - 2 `Age : integer`
 - 3 `[Age \geq 0]`
-

Part II

What is the problem?

Satisfiability

$$\text{satisfiable}(\text{model}) = \begin{cases} \text{true} & \exists x \in \text{instances}(\text{model}) \\ \text{false} & \text{otherwise} \end{cases}$$

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Theorem

Satisfiability for first-order logic is undecidable.



Figure: Alonzo Church

$$\text{satisfiable}(\text{model}, \text{scope}) = \begin{cases} \text{true} & \exists x \in \text{instances}(\text{model}, \text{scope}) \\ \text{false} & \text{otherwise} \end{cases}$$

$$\text{satisfiable}(\text{model}, \text{scope}) = \begin{cases} \text{true} & \exists x \in \text{instances}(\text{model}, \text{scope}) \\ \text{false} & \text{otherwise} \end{cases}$$

Clafer 9: Fermat's last theorem

```
1 A 1..*
2 B 1..*
3 C 1..*
4 [ (#A × #A × #A) + (#B × #B × #B) = #C × #C × #C ]
```

$$\text{satisfiable}(\text{model}, \text{scope}) = \begin{cases} \text{true} & \exists x \in \text{instances}(\text{model}, \text{scope}) \\ \text{false} & \text{otherwise} \end{cases}$$

Clafer 10: Fermat's last theorem

```
1 A 1..*
2 B 1..*
3 C 1..*
4 [ (#A × #A × #A) + (#B × #B × #B) = #C × #C × #C ]
```

$$\text{scope} = \left\{ \begin{array}{l} A \rightarrow 10 \\ B \rightarrow 10 \\ C \rightarrow 10 \end{array} \right\}$$

Meaning of true

How to interpret $\text{satisfiable}(\text{model}, \text{scope}) = \mathbf{true}$?

Perfect: accepts all intended instances.

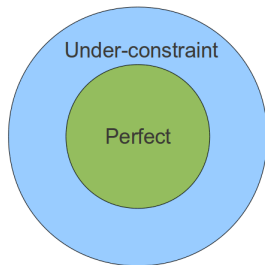


Figure: Satisfiable models

Meaning of true

How to interpret $\text{satisfiable}(\text{model}, \text{scope}) = \mathbf{true}$?

Perfect: accepts all intended instances.

Under-constraint: accepts all intended instances plus more.

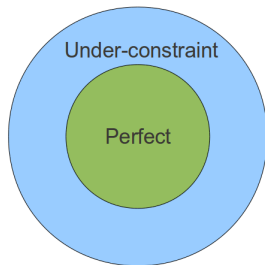


Figure: Satisfiable models

Under-constraint example

Clafer 11: Spouses model

```
1 abstract Person
2     marriedTo → Person ?
3     [this = marriedTo.marriedTo]
4
5 Alice : Person
6 Bob : Person
```

satisfiable(spouses model, scope) = true

Under-constraint example

Clafer 12: Spouses model

```
1 abstract Person
2     marriedTo → Person ?
3     [this = marriedTo.marriedTo]
4
5 Alice : Person
6 Bob : Person
```

$satisfiable(spouses\ model, scope) = \mathbf{true}$

But it's wrong.

Meaning of false

How to interpret $\text{satisfiable}(\text{model}, \text{scope}) = \mathbf{false}$?

Over-constraint

Clafer 13: $1 = 2$?

```
1 A : integer
2 [ A = 1 ]
3 [ A = 2 ]
4 [ A ≠ 0 ]
```

Meaning of false

How to interpret $\text{satisfiable}(\text{model}, \text{scope}) = \text{false}$?

Over-constraint

Clafer 15: $1 = 2$?

```
1 A : integer
2 [ A = 1 ]
3 [ A = 2 ]
4 [ A ≠ 0 ]
```

Insufficient scope

Clafer 16: Not enough A

```
1 A 5..*
```

$\text{scope} = \{A \rightarrow 4\}$

Part III

ClaferIG

$$\text{satisfiable}(model, scope) = \begin{cases} true & \exists x \in \text{instances}(model, scope) \\ false & \text{otherwise} \end{cases}$$

$$\text{find}(model, scope) = \text{instances}(model, scope)$$

Detecting under-constraint

Clafer 17: Spouses model

```
1 abstract Person
2     marriedTo → Person ?
3     [this = marriedTo.marriedTo]
4
5 Alice : Person
6 Bob : Person
```

$satisfiable(spouses\ model, scope) = \mathbf{true}$

$find(spouses\ model, scope) = \dots$

Detecting under-constraint

Clafer 18: ClaferIG instance of spouses model

```
1 Alice
2   marriedTo1 = Alice
3 Bob
4   marriedTo2 = Bob
```

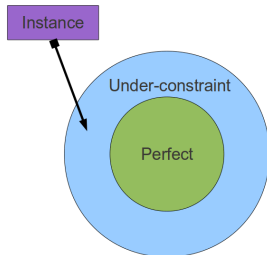


Figure: Proof of under-constraint

Clafer 19: $1 = 2?$

```
1 A : integer
2 [ A = 1 ]
3 [ A = 2 ]
4 [ A ≠ 0 ]
```

$satisfiable(model, scope) = \mathbf{false}$

Cannot simultaneously satisfy

- $[A = 1]$
- $[A = 2]$

$$\text{unsatCore}(\text{model}, \text{scope}) = \text{minimal}(c)$$

where

$$c \subseteq \text{constraints}(\text{model}) \wedge \text{unsatisfiable}(c, \text{scope})$$

Clafer 20: Over model

```
1 A : integer
2 [ A = 1 ]
3 [ A = 2 ]
4 [ A ≠ 0 ]
```

unsatCore(over model, scope) = ...

Clafer 21: ClaferIG UNSAT core of over model

The following set of constraints cannot be satisfied in the current scope.

- 1) $A = 1$
- 2) $A = 2$

Clafer 22: Over model

```
1 A : integer
2 [ A = 1 ]
3 [ A = 2 ]
4 [ A ≠ 0 ]
```

Clafer 23: ClaferIG fix for over model

Altering the following constraints produced a counterexample.

1) removed $A = 1$
 $A = 2$

$$\text{minimalScope}(\text{model}) = \text{minimal}(\text{scope}) \wedge \text{satisfiable}(\text{model}, \text{scope})$$

Clafer 24: Fermat's last theorem

```
1 A 1..*
2 B 1..*
3 C 1..*
4 [ (#A × #A × #A) + (#B × #B × #B) = #C × #C × #C ]
```

minimalScope(removeConstraints(model))

Clafer 25: Mesopotamian automotive systems

```
1 Chariot           --scope(Chariot) = 1
2   Wheel 2..4      --scope(Wheel) = scope(Chariot) × 2 = 2
3     Spoke 10..16  --scope(Spoke) = scope(Wheel) × 10 = 20
```

Scope dependency

Definition

B depends on A if the scope analysis of B requires the scope analysis of A.

Parent-child dependency

Clafer 26: Parents

1	A		
2	B		
3		C	
4	D		
5			
6	E		
7	F		

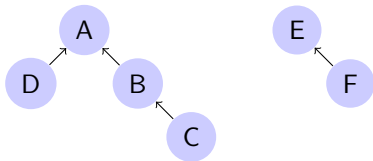


Figure: Dependency graph of parents model

Scope dependency

Subtype dependency

Clafer 27: Subtypes

```
1 abstract A
2   B
3
4 C
5   D : A 2..*
```

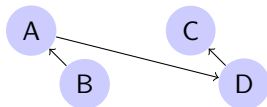


Figure: Dependency graph of subtypes model

Reference dependency

Clafer 28: References

```
1 A
2   B
3
4 C
5   D → A 2..*
```

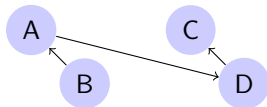
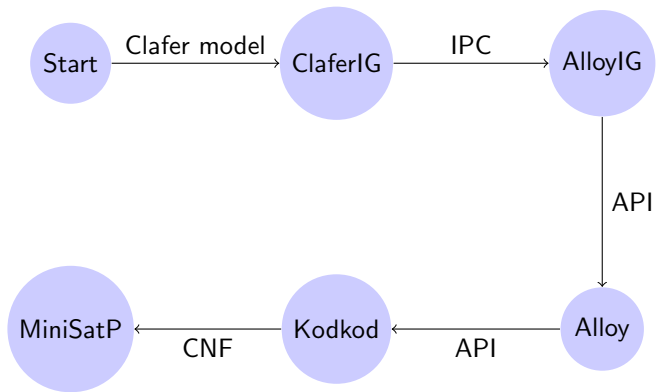


Figure: Dependency graph of references model

Implementation

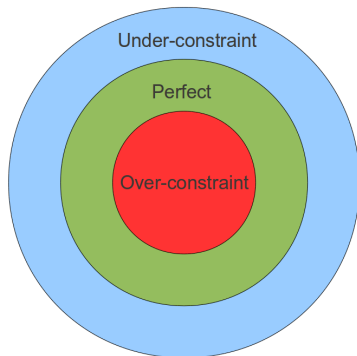


Part IV

Future work

Over-constraint satisfiable models

- Detecting under-constraint is easy
- Detecting over-constraint is hard



Clafer 29: Three's company

```
1 Crowd
2     Member → Person 3..*
3
4 abstract Person
5
6 Jack : Person
7 Janet : Person
```

Analyze the constraints for a tighter bound.

Clafer 30: Processor feature

```
1 Processor
2     VideoCard ?
3
4 [ some VideoCard ]
```

Isomorphism

Clafer 31: Leadership v1

```
1 Party
2   Person1
3     Age = 3
4   Person2
5     Age = 4
```

Clafer 32: Leadership v2

```
1 Party
2   Person1
3     Age = 4
4   Person2
5     Age = 3
```

Can ClaferIG detect isomorphic instances?

$$\mathit{canonical}(\mathit{model1}) = \mathit{canonical}(\mathit{model2}) \iff \mathit{model1} \cong \mathit{model2}$$

$$\mathit{canonical}(/home/zan/../jayna) = /home/jayna$$

$$\mathit{canonical}(/home/jayna) = /home/jayna$$

canonical function for Clafer instances?

<https://github.com/gsdlab/claferIG>