# The Longest Path in a Random Graph: Directed Case 

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- $D_{n, N}^{\prime}$ uniform random digraph with $n$ vertices and $N$ edges


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- random digraph with $n$ edges has longest directed path of length $O(\sqrt{n \log n})$ a.a.s
- Note: We can just work with $D_{n, p}$ instead of $D_{n, N}^{\prime}$ by coupling


## Main Result

Main Theorem ( $D_{n, N}^{\prime}$ )
$D_{n, \alpha n}^{\prime}$ contains a path of length cn a.a.s, where $\alpha>1$ and $c=c(\alpha)$ is constant.

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Main Theorem ( $D_{n, p}$ )
For $\alpha>1$ there are positive numbers $c, K$ and $\theta<1$ such that

$$
\operatorname{Pr}\left[D_{n, \frac{\alpha}{n}} \text { contains a directed path of length } c n\right]>1-K \theta^{n} \text {. }
$$

Here and onwards we let $p=\frac{\alpha}{n}$

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1. We define a graph process that generates $D_{n, p}$ and the desired long directed path
2. The process constructs the graph in a Depth First Search like manner
3. Use a branching process to analyze it (can have infinitely many points)
4. Show that $c n$ generations occur in the branching process (i.e. a cn dipath in the first process) well before we have used up all the points in the first process

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Issue: The distributions keep changing. Ex. the probability that $v_{1}$ has $k$ children is $\binom{n-1}{k} p^{k}(1-p)^{n-1-k}$ while the probability that its first child has $I$ is $\binom{n-1-k}{l} p^{\prime}(1-p)^{n-1-k-l}$ and so on...

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5. If tree dies out, pick unused vertex with smallest index and restart the process until no vertices remain unused

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- To guarantee a $D_{n, p}$ we appropriately randomize the remaining directed edges (i.e. edges between children, from one tree to another, from child to ancestor, etc.)
Note: BFS can only give us at most a $\log n$ path but DFS can potentially give us much longer paths


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- At each level assign a random order to the branches so the "leftmost infinite path" (LIP) is well defined
- Recall: Since $\lambda>1$ there exists $Q=Q(\lambda, m)$ such that the probability of extinction is $Q<1$


## Results on Branching Processes

Let $T$ denote the total population until extinction conditioned on process dying. Let $T_{1}, T_{2}, \ldots$ be a sequence of independent RV s distributed according to $T$ and set $S_{k}=T_{1}+. .+T_{k}$

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$L_{k}=\#$ of points to the left of LIP (inclusive) at level $k$.
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Claim 3
For all $m$, extinction probability $Q=Q(\lambda, m)<Q_{0}=Q_{0}(\lambda)<1$.

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- We had a surplus of $\delta n$ vertices in the original graph process so we are done!


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Let $f(x)=\sum_{k \geq 0} \lambda_{k} x^{k}$. Since $\mathbb{E}[\mathbf{C h}]=\lambda>1$ it is well known that $Q$ is the unique root of the equation $Q=f(Q)$ for $0 \leq Q<1$.

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Note that $f(x)=(1-p+p x)^{m}=\left(1-\frac{\lambda(1-x)}{m}\right)^{m}$ and is monotone increasing in $m$ since for $C=\lambda(1-x)$ by AM-GM

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1 \cdot\left(1-\frac{C}{m}\right) \ldots\left(1-\frac{C}{m}\right) \leq\left(\frac{1+m\left(1-\frac{C}{m}\right)}{m+1}\right)^{m+1}=\left(1-\frac{C}{m+1}\right)^{m+1}
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Also, $f(x) \leq e^{-\lambda(1-x)}$ and thus $Q<Q_{0}<1$ for all $m$ where $Q_{0}$ is the unique solution on $(0,1)$ to $Q_{0}=e^{-\lambda\left(1-Q_{0}\right)}$. Equivalently, $Q_{0}=\frac{x}{\lambda}$ where $x$ is unique value in $(0,1)$ for which $x e^{-x}=\lambda e^{-\lambda}$.

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Pick $\epsilon$ small (in terms of $\alpha, \delta$ ) and pick $t=t_{0}$ to get exponentially small probability
Lastly show $\mathbb{E}\left[e^{t_{0} T}\right] \leq K$ for $t_{0}=t_{0}(\lambda)>0$ and $K=K(\lambda)>0$

## (Sketchy) Proof of Claim 1 Cont'd

Let $\bar{T}$ be the total population and $\bar{t}_{l}=\operatorname{Pr}[\bar{T}=l]$, then we have $\bar{t}_{0}=0$ and $\bar{t}_{\infty}=1-Q$, and for $I \geq 1$

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\bar{t}_{l}=\sum_{k=0} \lambda_{k} \sum_{l_{1}+\ldots+l_{k}=l-1: i_{i} \geq 1} \bar{t}_{l_{1}} \cdots \bar{t}_{l_{k}}
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Then $\bar{g}(x)=\sum_{l=1} \bar{t}_{1} x^{\prime}$ satisfies $\bar{g}(x)=x f(\bar{g}(x))$ so it is defined uniquely for $x<1$ (By LIFT) and so $\bar{g}(x)<Q$

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By a continuity argument $\bar{g}(x)$ can be defined up to $y_{0}=\frac{x_{0}}{f\left(x_{0}\right)}$ where $x_{0}=\frac{1-p}{(m-1) p}$ is the maximum of $\frac{x}{f(x)}$ under the assumption $\bar{g}<x_{0}$

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$T$ is total population conditioned of the process going extinct. Thus

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t_{l}=\operatorname{Pr}[T=l]=\operatorname{Pr}[\bar{T}=\| \mid \bar{T}<\infty]=\frac{\bar{t}_{l}}{Q}
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$\mathbb{E} T=g^{\prime}(1)=\frac{1}{1-f^{\prime}(Q)} \sim \frac{1}{1-\lambda Q}$ where $g(x)=\sum_{l=0} t_{l} x^{\prime}=\bar{g}(x) / Q$

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Now $\mathbb{E} e^{t_{0} T}=g\left(e^{t_{0}}\right) \leq g\left(y_{0}\right)=\frac{x_{0}}{Q}<\frac{1}{Q}$.

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Thus

$$
t_{l}=\operatorname{Pr}[T=l]=\operatorname{Pr}[\bar{T}=\| \mid \bar{T}<\infty]=\frac{\bar{t}_{l}}{Q}
$$

$\mathbb{E} T=g^{\prime}(1)=\frac{1}{1-f^{\prime}(Q)} \sim \frac{1}{1-\lambda Q}$ where $g(x)=\sum_{l=0} t_{l} x^{\prime}=\bar{g}(x) / Q$ Now $\mathbb{E} e^{t_{0} T}=g\left(e^{t_{0}}\right) \leq g\left(y_{0}\right)=\frac{x_{0}}{Q}<\frac{1}{Q}$.

We are done as $Q$ is an increasing function in $m$ so $\frac{1}{Q}$ is uniformly bounded in $m$ as needed

## Conclusion

The proof of claim 2 can be done by a similar approach to the proof of claim 1. That is, we define an appropriate RV, bound the MGF and use this to derive the concentration bound.

## Conclusion

Thank you!

