# The Longest Path in a Random Graph: Directed Case

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- ► random digraph with (1 − ε)n edges has longest directed path of length O(log n) a.a.s
- random digraph with n edges has longest directed path of length O(\sqrt{n \log n}) a.a.s
- ▶ Note: We can just work with  $D_{n,p}$  instead of  $D'_{n,N}$  by coupling

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# Main Theorem $(D_{n,p})$

For  $\alpha > 1$  there are positive numbers c, K and  $\theta < 1$  such that

 $\Pr[D_{n,\frac{\alpha}{n}} \text{ contains a directed path of length } cn] > 1 - K\theta^n.$ 

Here and onwards we let  $p = \frac{\alpha}{n}$ 

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- 2. The process constructs the graph in a Depth First Search like manner
- 3. Use a branching process to analyze it (can have infinitely many points)
- 4. Show that *cn* generations occur in the branching process (i.e. a cn dipath in the first process) well before we have used up all the points in the first process

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#### 4. ...

Issue: The distributions keep changing. Ex. the probability that  $v_1$  has k children is  $\binom{n-1}{k}p^k(1-p)^{n-1-k}$  while the probability that its first child has l is  $\binom{n-1-k}{l}p^l(1-p)^{n-1-k-l}$  and so on...

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# Directed Graph Process

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- Recursively repeat on the first child until branch dies, then repeat on the sibling of the node that went extinct (Essentially a depth first search!)
- 5. If tree dies out, pick unused vertex with smallest index and restart the process until no vertices remain unused

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Note: BFS can only give us at most a log *n* path but *DFS* can potentially give us much longer paths

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### Galton-Watson Branching Process

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- At each level assign a random order to the branches so the "leftmost infinite path" (LIP) is well defined
- Recall: Since λ > 1 there exists Q = Q(λ, m) such that the probability of extinction is Q < 1</p>

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Claim 1

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#### Claim 2

There exists  $c = c(\alpha) > 0$  such that  $\Pr[L_{cn} > \frac{\delta}{2}n]$  is exponentially small.

#### Claim 3

For all *m*, extinction probability  $Q = Q(\lambda, m) < Q_0 = Q_0(\lambda) < 1$ .

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# Proof of Main Result

- ▶ Claim 3 and independence  $\implies$  first  $\epsilon n$  processes go extinct with probability at most  $Q_0^{\epsilon n}$
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- We had a surplus of  $\delta n$  vertices in the original graph process so we are done!

Let  $f(x) = \sum_{k\geq 0} \lambda_k x^k$ . Since  $\mathbb{E}[\mathbf{Ch}] = \lambda > 1$  it is well known that Q is the unique root of the equation Q = f(Q) for  $0 \leq Q < 1$ .

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Note that  $f(x) = (1 - p + px)^m = (1 - \frac{\lambda(1-x)}{m})^m$  and is monotone increasing in *m* since for  $C = \lambda(1-x)$  by AM-GM

$$1 \cdot \left(1 - \frac{C}{m}\right) \dots \left(1 - \frac{C}{m}\right) \le \left(\frac{1 + m(1 - \frac{C}{m})}{m+1}\right)^{m+1} = \left(1 - \frac{C}{m+1}\right)^{m+1}$$

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Also,  $f(x) \leq e^{-\lambda(1-x)}$  and thus  $Q < Q_0 < 1$  for all m where  $Q_0$  is the unique solution on (0,1) to  $Q_0 = e^{-\lambda(1-Q_0)}$ . Equivalently,  $Q_0 = \frac{x}{\lambda}$  where x is unique value in (0,1) for which  $xe^{-x} = \lambda e^{-\lambda}$ .

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By independence of  $T_i$ ,  $\mathbb{E}[\exp(tS_{\epsilon n})] = \prod_{i=1}^{\epsilon n} \mathbb{E} \exp(tT) \leq K^{\epsilon n}$ Pick  $\epsilon$  small (in terms of  $\alpha, \delta$ ) and pick  $t = t_0$  to get exponentially small probability We want  $\Pr[T_1 + ... + T_{\epsilon n} > \frac{\delta}{2}n]$  exponentially small  $\Pr[S_{\epsilon n} > \frac{\delta}{2}n] = \Pr[\exp(tS_{\epsilon n}) > \exp(t\delta n/2)] \le \frac{\mathbb{E}[\exp(tS_{\epsilon n})]}{\exp(t\delta n/2)}$  (Markov)

By independence of  $T_i$ ,  $\mathbb{E}[\exp(tS_{\epsilon n})] = \prod_{i=1}^{\epsilon n} \mathbb{E}\exp(tT) \leq K^{\epsilon n}$ Pick  $\epsilon$  small (in terms of  $\alpha, \delta$ ) and pick  $t = t_0$  to get exponentially small probability

Lastly show  $\mathbb{E}[e^{t_0 T}] \leq K$  for  $t_0 = t_0(\lambda) > 0$  and  $K = K(\lambda) > 0$ 

#### 

Let  $\overline{T}$  be the total population and  $\overline{t}_l = \Pr[\overline{T} = l]$ , then we have  $\overline{t}_0 = 0$  and  $\overline{t}_{\infty} = 1 - Q$ , and for  $l \ge 1$ 

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Then  $\overline{g}(x) = \sum_{l=1} \overline{t}_l x^l$  satisfies  $\overline{g}(x) = x f(\overline{g}(x))$  so it is defined uniquely for x < 1 (By LIFT) and so  $\overline{g}(x) < Q$ 

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By a continuity argument  $\overline{g}(x)$  can be defined up to  $y_0 = \frac{x_0}{f(x_0)}$ where  $x_0 = \frac{1-p}{(m-1)p}$  is the maximum of  $\frac{x}{f(x)}$  under the assumption  $\overline{g} < x_0$   ${\mathcal T}$  is total population conditioned of the process going extinct. Thus

$$t_{l} = \Pr[T = l] = \Pr[\overline{T} = l | \overline{T} < \infty] = \frac{\overline{t}_{l}}{Q}$$
$$\mathbb{E}T = g'(1) = \frac{1}{1 - f'(Q)} \sim \frac{1}{1 - \lambda Q} \text{ where } g(x) = \sum_{l=0} t_{l} x^{l} = \overline{g}(x)/Q$$

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$$\text{Now } \mathbb{E}e^{t_{0}T} = g(e^{t_{0}}) \leq g(y_{0}) = \frac{x_{0}}{Q} < \frac{1}{Q}.$$

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Now  $\mathbb{E}e^{t_0 T} = g(e^{t_0}) \leq g(y_0) = \frac{x_0}{Q} < \frac{1}{Q}.$ 

We are done as Q is an increasing function in m so  $\frac{1}{Q}$  is uniformly bounded in m as needed

The proof of claim 2 can be done by a similar approach to the proof of claim 1. That is, we define an appropriate RV, bound the MGF and use this to derive the concentration bound.

# Thank you!