An Introduction to Structured Query Language

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Outline

1 The SQL Standard **2** SQL DML **Basic** Queries Data Modification **Complex** Queries Set and Multiset Operations Unknown values Subqueries Table Expressions Outer joins Ordering results Grouping and Aggregation Having clauses

3 SQL DDL

Tables Integrity Constraints Views Triggers

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Structured Query Language (SQL) is made up of two sub-languages:

- SQL Data Manipulation Language (DML)
 - SELECT statements perform queries
 - INSERT, UPDATE, DELETE statements modify the instance of a table
- SQL Data Definition Language (DDL)
 - CREATE, DROP statements modify the database schema
 - GRANT, REVOKE statements enforce the security model

The SQL Standard



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select LastName,HireDate
from Employee
where Salary > 100000

Find the last names and hire dates of employees who make more than \$100000.

Note

SQL is declarative (non-navigational)

- Relational model: relations are sets
- SQL standard: tables are multisets (a.k.a. bags)
 - Duplicate tuples may be stored.
 - SQL queries may result in duplicates even if none of the input tables themselves contain duplicates.
 - Select distinct is used to eliminate duplicates from the result of a query.

select distinct LastName,HireDate from Employee

```
where Salary > 100000
```

```
select P.ProjNo, E.LastName
from Employee E, Project P
where P.RespEmp = E.EmpNo
and P.DeptNo = 'E21'
```

For each project for which department E21 is responsible, find the name of the employee in charge of that project.

```
select attribute-expression-list
from relation-list
[where condition]
```

Note

The result of such a query is a relation which has one attribute for each element of the query's attribute-expression-list.

The SQL "Where" Clause

Conditions may include

- arithmetic operators +, -, *, /
- comparisons =, <>, <, <=, >, >=
- logical connectives and, or and not

```
select E.LastName
from Employee E,
        Department D,
        Employee Emgr
where E.WorkDept = D.DeptNo
    and D.MgrNo = Emgr.EmpNo
    and E.Salary > Emgr.Salary
```

List the last names of employees who make more than their manager.

- Return the difference between each employee's actual salary and a base salary of \$40000
 select E.EmpNo, E.Salary - 40000 as SalaryDiff
 from Employee E
- As above, but report zero if the actual salary is less than the base salary

SQL DML: Insertion & Deletion

Insert a single tuple into the Employee relation.

delete from Employee;

delete from Employee
where WorkDept = 'A00';

Delete all employees from the Employee table.

Delete all employees in department A00 from the Employee table. update Employee Increase the salary of every
set Salary = Salary * 1.05; employee by five percent.

update Employee
set WorkDept = 'E01'
where WorkDept = 'E21';

Move all employees in department E21 into department E01.

Set Operations

• SQL defines UNION, INTERSECT and EXCEPT operations (EXCEPT is set difference)

select empno

from employee

except

select mgrno

from department

- These operations result in sets
 - Q1 UNION Q2 includes any tuple that is found (at least once) in Q1 or in Q2
 - Q_1 INTERSECT Q_2 includes any tuple that is found (at least once) in both Q_1 and Q_2
 - Q_1 EXCEPT Q_2 includes any tuple that is found (at least once) in Q_1 and is not found Q_2

- SQL provides a multiset version of each of the set operations: UNION ALL, INTERSECT ALL, EXCEPT ALL
- suppose Q_1 includes n_1 copies of some tuple t, and Q_2 includes n_2 copies of the same tuple t.
 - Q_1 UNION ALL Q_2 will include $n_1 + n_2$ copies of t
 - Q_1 INTERSECT ALL Q_2 will include $\min(n_1, n_2)$ copies of t
 - Q_1 EXCEPT ALL Q_2 will include $\max(n_1 n_2, 0)$ copies of t

NULL values

- the value NULL can be assigned to an attribute to indicate unknown or missing data
- NULLs are a necessary evil lots of NULLs in a database instance suggests poor schema design
- NULLs can be prohibited for certain attributes by schema constraints, e.g., NOT NULL, PRIMARY KEY
- predicates and expressions that involve attributes that may be NULL may evaluate to NULL
 - x + y evaluates to NULL if either x or y is NULL
 - x > y evaluates to NULL if either x or y is NULL
 - how to test for NULL? Use is NULL or is not NULL

Note

SQL uses a three-valued logic: TRUE, FALSE, NULL

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Logical Expressions in SQL

AND	TRUE	FALSE	NULL
TRUE	TRUE	FALSE	NULL
FALSE	FALSE	FALSE	FALSE
NULL	NULL	FALSE	NULL

OR	TRUE	FALSE	NULL
TRUE	TRUE	TRUE	TRUE
FALSE	TRUE	FALSE	NULL
NULL	TRUE	NULL	NULL
		I	I
NOT	TRUE	FALSE	NULL
	FALSE	TRUE	NULL

NULL and the SQL Where Clause

• The query:

select *
from employee
where hiredate <> '05/05/1947'

will *not* return information about employees whose hiredate is NULL.

Note

The condition in a where clause filters out any tuples for which the condition evaluates to FALSE or to NULL.

Subqueries

• These two queries are equivalent.

```
select deptno, deptname
from department d, employee e
where d.mgrno = e.empno and e.salary > 50000
```

select deptno, deptname
from department
where mgrno in
 (select empno
 from employee
 where salary > 50000)

Subquery Constructs in SQL

- SQL supports the use of the following predicates in the where clause. A is an attribute, Q is a query, op is one of >, <, <>, =, <=, >=.
 - A IN (Q)
 - A NOT IN (Q)
 - *A op* SOME (*Q*)
 - A op ALL (Q)
 - EXISTS (Q)
 - NOT EXISTS (Q)
- For the first four forms, the result of Q must have a single attribute.

Another Subquery Example

• Find the name(s) and number(s) of the employee(s) with the highest salary.

```
select empno, lastname
from employee
where salary >= all
        ( select salary
            from employee )
```

Note

Is this query correct if the schema allows the salary attribute to contain NULLs?

- This query also returns the employee(s) with the largest salary: select empno, lastname from employee E1 where salary is not null and not exists
 (select * from employee E2 where E2.salary > E1.salary)
- This query contains a *correlated* subquery the subquery refers to an attribute (E1.salary) from the outer query.

Scalar Subqueries

- Subquery that returns an atomic value (one row / one column)
- in the where clause:

```
select empno, lastname
 from employee
 where salary >
          (select salary
           from employee e2
           where e2.empno = '000190')
• in the select clause:
 select projno,
         (select deptname
          from department d
          where e.workdept = d.deptno)
 from project p, employee e
```

```
where p.respemp = e.empno
```

Table Expressions

```
• in the from clause:
 select projno, projname
 from project p,
    (select mgrno
    from department, employee
    where mgrno = empno and salary > 100000) as m
 where respemp = mgrno
• in a with clause:
 with Mgrs(empno) as
    (select mgrno
    from department, employee
    where mgrno = empno and salary > 100000)
 select projno, projname
 from project, Mgrs
 where respemp = empno
```

• List the manager of each department. Include in the result departments that have no manager.

Note

SQL supports left, right, and full outer joins.

Ordering Results

- No particular ordering on the rows of a table can be assumed when queries are written. (This is important!)
- No particular ordering of rows of an intermediate result in the query can be assumed either.
- However, it is possible to order the final result of a query, using the order by clause.

select distinct e.empno, emstdate, firstnme, lastname
from employee e, emp_act a
where e.empno = a.empno and a.projno = 'PL2100'
order by emstdate

Grouping and Aggregation: An Example

• For each department, list the number of employees it has and their combined salary.

Grouping and Aggregation: Operational Semantics

- The result of a query involving grouping and aggregation can be determined as follows:
 - 1 form the cross product of the relations in the from clause
 - eliminate tuples that do not satisy the condition in the where clause
 - 3 form the remaining tuples into groups, where all of the tuples in a group match on all of the grouping attributes
 - generate one tuple per group. Each tuple has one attribute per expression in the select clause.
- aggregation functions are evaluated separately for each group

Grouping and Aggregation Example

• Apply where

DEPTNO	DEPTNAME	SALARY
 A00	SPIFFY COMPUTER SERVICE DIV.	52750.00
A00	SPIFFY COMPUTER SERVICE DIV.	46500.00
B01	PLANNING	41250.00
C01	INFORMATION CENTER	38250.00
D21	ADMINISTRATION SYSTEMS	36170.00
D21	ADMINISTRATION SYSTEMS	22180.00
D21	ADMINISTRATION SYSTEMS	19180.00
D21	ADMINISTRATION SYSTEMS	17250.00
D21	ADMINISTRATION SYSTEMS	27380.00
E01	SUPPORT SERVICES	40175.00
E11	OPERATIONS	29750.00
E11	OPERATIONS	26250.00
E11	OPERATIONS	17750.00
E11	OPERATIONS	15900.00
E21	SOFTWARE SUPPORT	26150.00

Grouping and Aggregation Example (cont'd)

• Apply where, then group by

DEPTNO	DEPTNAME	SALARY
A00	SPIFFY COMPUTER SERVICE DIV.	52750.00
A00	SPIFFY COMPUTER SERVICE DIV.	46500.00
B01	PLANNING	41250.00
C01	INFORMATION CENTER	38250.00
D21	ADMINISTRATION SYSTEMS	36170.00
D21	ADMINISTRATION SYSTEMS	22180.00
D21	ADMINISTRATION SYSTEMS	19180.00
D21	ADMINISTRATION SYSTEMS	17250.00
D21	ADMINISTRATION SYSTEMS	27380.00
E01	SUPPORT SERVICES	40175.00
E11	OPERATIONS	29750.00
E11	OPERATIONS	26250.00
E11	OPERATIONS	17750.00
E11	OPERATIONS	15900.00
E21	SOFTWARE SUPPORT	26150.00

Grouping and Aggregation Example (cont'd)

• Finally project and aggregate

DEPTNO	DEPTNAME	TOTALSALARY	EMPLOYEES
A00	SPIFFY COMPUTER SERVICE DIV.	99250.00	2
B01	PLANNING	41250.00	1
C01	INFORMATION CENTER	38250.00	1
D21	ADMINISTRATION SYSTEMS	122160.00	5
E01	SUPPORT SERVICES	40175.00	1
E11	OPERATIONS	89650.00	4
E21	SOFTWARE SUPPORT	26150.00	1

Aggregation Functions in SQL

- count(*): number of tuples in the group
- $\operatorname{count}(E)$: number of tuples for which E (an expression that may involve non-grouping attributes) is non-NULL
- count(distinct E): number of distinct non-NULL E values
- sum(E): sum of non-NULL E values
- sum(distinct E): sum of distinct non-NULL E values
- $\operatorname{avg}(E)$: average of non-NULL E values
- avg(distinct E): average of distinct non-NULL E values
- $\min(E)$: minimum of non-NULL E values
- $\max(E)$: maximum of non-NULL E values

The Having Clause

• List the average salary for each large department.

```
select deptno, deptname, avg(salary) as MeanSalary
from department d, employee e
where e.workdept = d.deptno
group by deptno, deptname
having count(*) >= 4
```

Note

The where clause filters tuples before they are grouped, the having clause filters groups.

Grouping and Aggregation: Operational Semantics

- The result of a query involving grouping and aggregation can be determined as follows:
 - 1 form the cross product of the relations in the from clause
 - eliminate tuples that do not satisy the condition in the where clause
 - 3 form the remaining tuples into groups, where all of the tuples in a group match on all of the grouping attributes
 - eliminate any groups of tuples for which the having clause is not satisfied
 - **5** generate one tuple per group. Each tuple has one attribute per expression in the select clause.
- aggregation functions are evaluated separately for each group

Grouping and Aggregation with Having

Apply where, then group by		
DEPTNO	DEPTNAME	SALARY
A00	SPIFFY COMPUTER SERVICE DIV.	52750.00
A00	SPIFFY COMPUTER SERVICE DIV.	46500.00
в01	PLANNING	41250.00
C01	INFORMATION CENTER	38250.00
D21	ADMINISTRATION SYSTEMS	36170.00
D21	ADMINISTRATION SYSTEMS	22180.00
D21	ADMINISTRATION SYSTEMS	19180.00
D21	ADMINISTRATION SYSTEMS	17250.00
D21	ADMINISTRATION SYSTEMS	27380.00
E01	SUPPORT SERVICES	40175.00
E21	SOFTWARE SUPPORT	26150.00
E11	OPERATIONS	29750.00
E11	OPERATIONS	26250.00
E11	OPERATIONS	17750.00
E11	OPERATIONS	15900.00

Grouping and Aggregation with Having (cont'd)

• After grouping, apply having

DEPTNO	DEPTNAME		SALARY
D21	ADMINISTRATION	SYSTEMS	36170.00
D21	ADMINISTRATION	SYSTEMS	22180.00
D21	ADMINISTRATION	SYSTEMS	19180.00
D21	ADMINISTRATION	SYSTEMS	17250.00
D21	ADMINISTRATION	SYSTEMS	27380.00
E11	OPERATIONS		29750.00
E11	OPERATIONS		26250.00
E11	OPERATIONS		17750.00
E11	OPERATIONS		15900.00
Finally	r project and a	raregate	

rinally project and aggregate

DEPTNO	DEPTNAME		MEANSALARY
D21	ADMINISTRATION	SYSTEMS	24432.00
E11	OPERATIONS		22412.50

Selecting Non-Grouping Attributes

db2 => select deptno, deptname, sum(salary) \
db2 (cont.) => from department d, employee e \
db2 (cont.) => where e.workdept = d.deptno \
db2 (cont.) => group by deptno
SQL0119N An expression starting with "DEPTNAME"
specified in a SELECT clause, HAVING clause, or
ORDER BY clause is not specified in the GROUP BY
clause or it is in a SELECT clause, HAVING clause,
or ORDER BY clause with a column function and no
GROUP BY clause is specified. SQLSTATE=42803

Note

Non-grouping attributes may appear in the select clause only in aggregate expressions. (Why?)

SQL DDL: Tables

```
create table Employee (
   EmpNo char(6),
   FirstName varchar(12),
   MidInit char(1),
   LastName varchar(15),
   WorkDept char(3),
   HireDate date
```

```
alter table Employee
  add column Salary decimal(9,2)
```

drop table Employee

SQL DDL: Data Types

Some of the attribute domains defined in SQL:

- INTEGER
- DECIMAL(p,q): p-digit numbers, with q digits right of decimal
- FLOAT(p): p-bit floating point numbers
- CHAR(n): fixed length character string, length n
- VARCHAR(n): variable length character string, max. length n
- DATE: describes a year, month, day
- TIME: describes an hour, minute, second
- TIMESTAMP: describes and date and a time on that date
- YEAR/MONTH INTERVAL: time interval
- DAY/TIME INTERVAL: time interval

• . . .

Integrity Constraints in SQL

Most commonly-used SQL schema constraints:

- Not Null
- Primary Key
- Unique
- Foreign Key
- Column or Tuple CHECK

Note

Recent SQL standards also allows more powerful integrity constraints. However, they are not supported by all commercial DBMSs.

SQL DDL: Integrity Constraints

```
create table Employee (
   EmpNo char(6) not null primary key,
   FirstName varchar(12) not null,
   MidInit char(1),
   LastName varchar(15) not null,
   WorkDept char(3) not null references Department
                                on delete cascade,
   HireDate date,
   Salary decimal(9,2) check (Salary >= 10000),
   constraint unique name dept
       unique (FirstName, LastName, WorkDept)
alter table Employee
  add column StartDate date
  add constraint hire before start
      check (HireDate <= StartDate);
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```

Another SQL Constraint Example

```
create table registeredin (
  coursenum char(5) not null,
  term char(3) not null,
  id char(8) not null references student
                       on delete no action,
  sectionnum char(2) not null,
  mark integer,
  constraint mark check check (
      mark \geq 0 and mark \leq 100 ).
  primary key (coursenum, term, id),
  foreign key (coursenum, sectionnum, term)
    references section
```

More Powerful SQL Integrity Constraints

```
create assertion balanced_budget check (
not exists (
    select deptno
    from department d
    where budget <
        (select sum(salary)
        from employee
        where workdept = d.deptno)))</pre>
```

Note

General assertions are not supported by current versions of DB2.

Views

Recall the three-level schema architecture:

- External schema
- 2 Conceptual schema
- 3 Physical schema

Definition (View)

A *view* is a relation in the external schema whose instance is determined by the instances of the relations in the conceptual schema.

A view has many of the same properties as a base relation in the conceptual schema:

- its schema information appears in the database schema
- access controls can be applied to it
- other views can be defined in terms of it

Updating Views

Demacra

- Modifications to a view's instance must be propagated back to instances of relations in conceptual schema.
- Some views cannot be updated unambiguously.

Conceptual Schema

Persons		
NAME	CITIZENSHIP	
Ed	Canadian	
Dave	Canadian	
Wes	American	

NationalPastimes

CITIZENSHIP	PASTIME
Canadian	Hockey
Canadian	Curling
American	Hockey
American	Baseball

External Schema

PersonalPastimes

NAME	PASTIME
Ed	Hockey
Ed	Curling
Dave	Hockey
Dave	Curling
Wes	Hockey
Wes	Baseball

- 1 What does it mean to insert (Darryl, Hockey)?
- 2 What does it mean to delete (Dave, Curling)?

Customizing the schema for a particular user/application:

```
create view ManufacturingProjects as
  ( select projno, projname, firstnme, lastname
    from project, employee
    where respemp = empno and deptno = 'D21' )
```

Once defined, SQL DML can be used to query a view like any other table:

select * from ManufacturingProjects

According to SQL-92, a view is updatable only if its definition satisfies a variety of conditions:

- The query references exactly one table
- The query only outputs simple attributes (no expressions)
- There is no grouping/aggregation/distinct
- There are no nested queries
- There are no set operations

These rules are more restrictive than necessary.

Definition

A *trigger* is a procedure executed by the database in response to a change to the database instance.

Basic components of a trigger description:

Event: Type of change that should cause trigger to fire

Condition: Test performed by trigger to determine whether further action is needed

Action: Procedure executed if condition is met

```
create trigger log addr
   after update of addr, phone on person
   referencing OLD as o NEW as n
   for each row
   mode DB2SOL /* DB2-specific syntax */
   when (o.status = 'VIP' or n.status = 'VIP')
      insert into VIPaddrhist(pid, oldaddr, oldphone,
         newaddr, newphone, user, modtime)
      values (o.pid, o.addr, o.phone,
         n.addr, n.phone, user, current timestamp)
```