Logical Approach to Physical Data Independence and Query Compilation

Query Compilation

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The Story So Far...

1. Physical Data Independence (OBDA, Data Exchange, ...)
2. Logic-based formalization (Relational model, constraints)
3. Queries and Answers

$$\text{cert}_{\Sigma,D}(\varphi) = \{\vec{a} \mid \Sigma \cup D \models \varphi(\vec{a})\} = \bigcap_{I \models \Sigma \cup D} \{\vec{a} \mid I \models \varphi(\vec{a})\}$$

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Difficulties

1. $\vec{a} \in \text{cert}_{\Sigma,D}(\varphi)$ undecidable for FOL
   - $\Rightarrow$ *Data Complexity* (in $|D|$) is high for most decidable fragments
   - $\Rightarrow$ *Lite Description Logics* and *CQ*—sacrificing expressiveness
2. Unintuitive Answers
   - $\Rightarrow$ "does John have a phone #?" $\rightarrow$ yes; "what is john's phone #?" $\rightarrow$ { }
   - $\Rightarrow$ non-compositional (no algebra for certain answers)
3. Efficient algorithm only for range-restricted queries over *closed KB*.

Query Compilation

Equivalent Range-restricted Queries

IDEA

Restrict allowed *user queries* to those that are *logically equivalent* to a *range-restricted query over* S_A under Σ.

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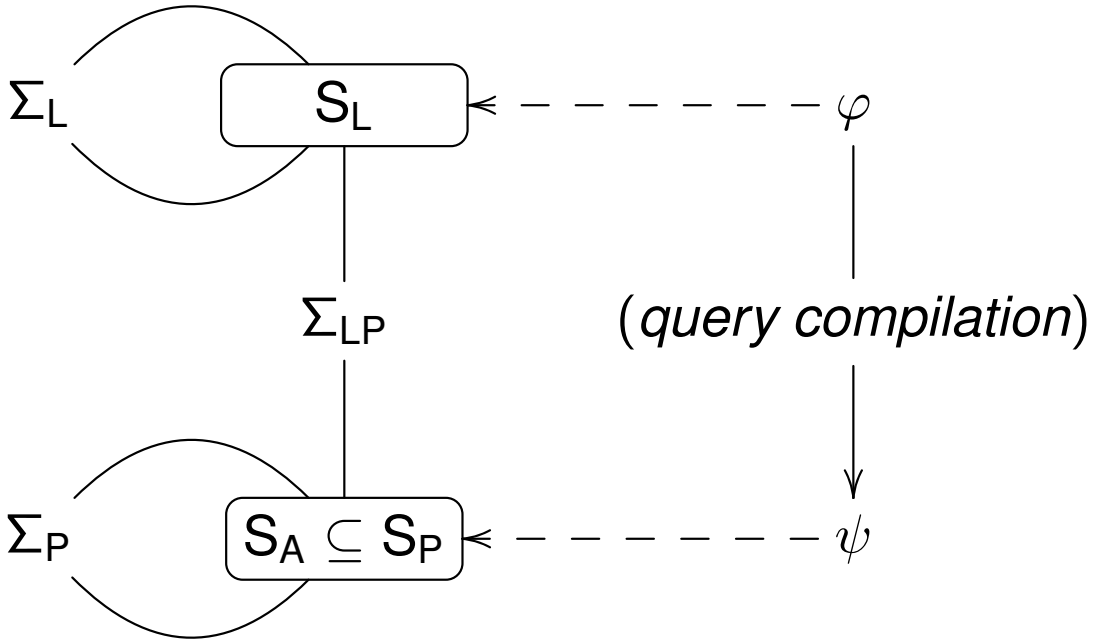
Intuition(s)

- The user *has the illusion* there is *exactly one* explicit *closed-world* database instance over S_L;
- The system has a choice of *different* D_A instances to support this: record id-s/pointers, record distribution/fragmentation, ...

Physical Design and Query Compilation: Overview

$$\Sigma = (\Sigma_L \cup \Sigma_{LP} \cup \Sigma_P)$$



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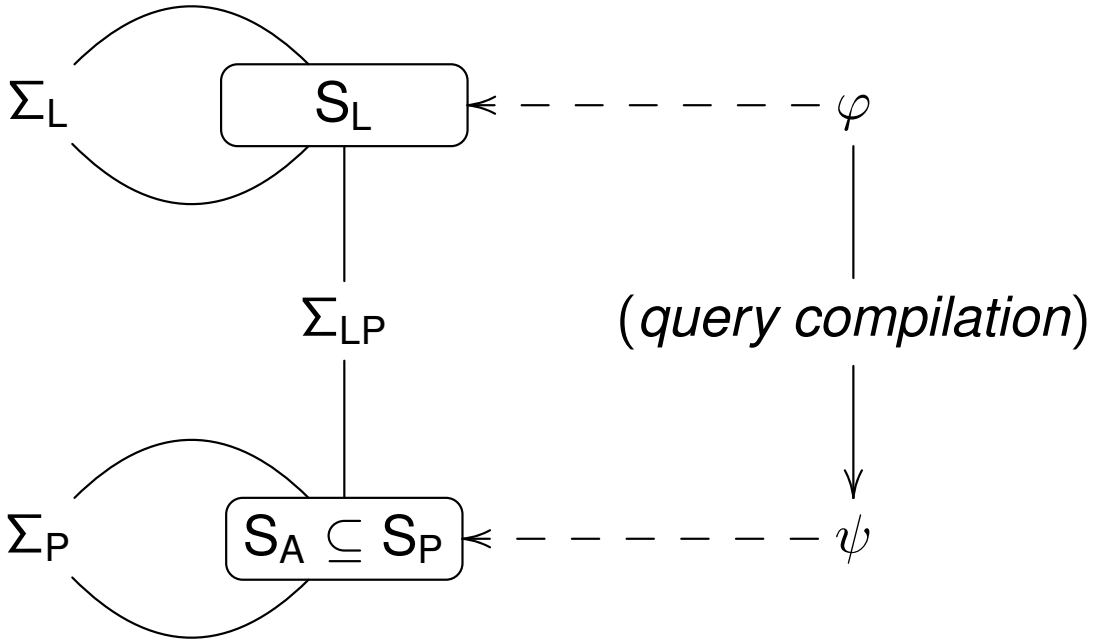

Issues

1. how to find ψ, given φ, Σ, and S_A? (and how hard is this?)
2. how closely does ψ describe actual execution? (and what is left out?)

Case Study: Relational Systems Implementation

Case Study: RDBMs Internals

Physical Design Desiderata (v0)

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- indexing (primary and secondary)
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IDEA: Record IDs

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- vertical fragmentation: fragments contain RID (losses join);
- primary index: search key "clusterd" with RID (typically stores tuples);
- secondary index: search keys and RIDs

$\Rightarrow$ all translates to a *selection of access paths* (and constraints)

Case Study: RDBMs Internals

Example

Logical Schema:

- ternary user relation $A/3$.

Physical Schema:

- a 4-ary *base file* $EAs/4/0$ (A-scan)
  $\Rightarrow A(x, y, z) \leftrightarrow (\exists a.EAs(a, x, y, z))$
  $\Rightarrow a$ is a key for EAs.
- a 4-ary *base file* $EAr/4/1$ (A-ref)
  $\Rightarrow EAs(a, x, y, z) \leftrightarrow EAr(a, x, y, z)$
- an $EAIx/2/1$ index on A's attribute x
  $\Rightarrow EAIx(x, a) \leftrightarrow (\exists y, z.EAs(a, x, y, z))$
- an $EAIy/2/1$ index on A's attribute y
  $\Rightarrow EAIx(y, a) \leftrightarrow (\exists x, z.EAs(a, x, y, z))$

Example: Access Paths

```
(defap A-scan :logical AB  :name EAs
    :in   []          :out  [a :int x :int y :int z :int]
    :cost ["10n" "10n"])

(defap A-ref  :logical AB  :name EAr
    :in   [a :int]  :out  [x :int y :int z :int]
    :cost ["10000log(n)" "1"])

(defap A-idx1 :logical AI1 :name EAIx
    :in   [x :int]  :out  [a :int]
    :cost ["log(n)" "log(n)"])

(defap A-idx2 :logical AI2 :name EAIy
    :in   [y :int]  :out  [a :int]
    :cost ["log(n)" "log(n)"])
```

Example: Constraints

```
(defschema schema
  :constraints [[
; base table
   (fol (-> (A ?x ?y ?z) (exists [?a] (AB ?a ?x ?y ?z))))
   (fol (-> (AB ?a ?x ?y ?z) (A ?x ?y ?z )))
; index on "x"
   (fol (-> (AB ?a ?x ?y ?z) (AI1 ?x ?a)))
   (fol (-> (AI1 ?x ?a) (exists [?y ?z] (AB ?a ?x ?y ?z))))
; index on "y"
   (fol (-> (AB ?a ?x ?y ?z) (AI2 ?y ?a)))
   (fol (-> (AI2 ?y ?a) (exists [?x ?z] (AB ?a ?x ?y ?z))))
 ]
 (primary-key [a] (AB a x y z))
 ]
 :access-paths [A-scan A-ref A-idx1 A-idx2]
 :cost-model ->Complexity
)
```

Case Study: RDBMs Internals

Example

User queries (and expected execution patterns):

1. $A(x, y, z)$
2. $A(x, y, z)$ given value for x
3. $A(x, y, z)$ given value for x and y
4. $\exists y, z.A(x, y, z)$ given x
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Case Study: RDBMs Internals

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   $\Rightarrow$ lookup x in *EAIx*, use RID for lookup in *EAr*, and compare y values.
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5. $\exists z.A(x, y, z)$ given value for x and y
   $\Rightarrow$ lookup x in *EAIx*, y in *EAIy* and compare retriever RIDs.

Example: Running the System

```
scan (A x y z) []
Plan: 0 (10n,10n)
E?x1.EAs(?x1,x,y,z)

lookup x: (A x y z) [x]
Plan: 1 (10000log(n)^2 + 2log(n),1log(n))
E?x1.(EAIx(x,?x1)^E?s0.(EAr(?x1,?s0,y,z)))

lookup x,y: (A x y z) [x,y]
Plan: 22 (10000log(n)^2 + 3log(n),1log(n))
E?x1.(EAIx(x,?x1)^E?s1.E?s0.(EAr(?x1,?s0,?s1,z)^Cmp(y,?s1)))

lookup x index only: (exists [?y ?z] (A x ?y ?z)) [x]
Plan: 0 (1log(n),1log(n))
E?x1.EAIx(x,?x1)

lookup x,y index only: (exists [?z] (A x y ?z)) [x,y]
Plan: 90 (2log(n)^2 + 1log(n),1log(n)^2)
E?x1.(EAIy(y,?x1)^E?s0.(EAIx(x,?s0)^Cmp(?x1,?s0)))
```

Execution Model, part 1

How can we say that, e.g.,

$$\exists x_1.(EAIy(y, x_1) \wedge \exists s_0.(EAIx(x, s_0) \wedge Cmp(x_1, s_0)))$$

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- implements:
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- efficiently—depends on mapping to actual operations:

access paths	$\mapsto$	access path access
conjunction	$\mapsto$	nested loops join
existential quantification	$\mapsto$	projection

and a *cost model* (estimates cost of executing the query from summaries)

Example: Nested Loops Join

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All operators implement an iterator (cursor) protocol:

- `get-first`: gets/searches for the first applicable tuple
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Nested Loops Join (NLJ):

Open and get first tuple:

```
function (Q1 ∧ Q2)-first
  if not Q1-first return false
  while not Q2-first do
    if not Q1-next return false
  return true
```

Get next tuple:

```
function (Q1 ∧ Q2)-next
  if Q2-next return true
  while Q1-next do
    if Q2-first return true
  return false
```

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$$\mathrm{In}(Q) = \begin{cases} \mathrm{In}(Q_1) \cup (\mathrm{In}(Q_2) - \mathrm{Out}(Q_1)) & \text{if } Q = \text{“}(Q_1 \wedge Q_2)\text{”}, \\ \mathrm{In}(Q_1) & \text{if } Q = \text{“}\exists x.Q_1\text{”}, \\ \mathrm{In}(Q_1) \cup \mathrm{In}(Q_2) & \text{if } Q = \text{“}(Q_1 \vee Q_2)\text{”}, \text{ and} \\ \mathrm{In}(Q_1) & \text{if } Q = \text{“}\neg\, Q_1\text{”}. \end{cases}$$

$$\mathrm{Out}(Q) = \begin{cases} \mathrm{Out}(Q_1) \cup \mathrm{Out}(Q_2) & \text{if } Q = \text{“}(Q_1 \wedge Q_2)\text{”}, \\ \mathrm{Out}(Q_1) \setminus \{x\} & \text{if } Q = \text{“}\exists x.Q_1\text{”}, \\ \mathrm{Out}(Q_1) \cap \mathrm{Out}(Q_2) & \text{if } Q = \text{“}(Q_1 \vee Q_2)\text{”}, \text{ and} \\ \emptyset & \text{if } Q = \text{“}\neg\, Q_1\text{”}. \end{cases}$$

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No *selection* operator: realized by

- supplying a value to a index lookup, or
- by an additional nested-loops join with a built-in table `Cmp/2/2`.

Standard RDBMs Summary

Standard (relational) Physical Design

`CREATE TABLE foo` DDL command causes

1. a *logical symbol* `foo` to be created;
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4. . . . numerous pages of additional options these days

- Record IDs and Indexing
- Multi-level store
- Horizontal partitioning
- Data replication
- Delegation to other database engines
- Materialized views and cached query results