Module 2

Finite Automata

The simplest computers for the simplest languages

CS 360: Introduction to the Theory of Computing

Collin Roberts, University of Waterloo

2.1

Topics of Module 2

- Deterministic finite automata
- Nondeterministic finite automata and the equivalence to DFAs
- ε-NFAs, and their equivalence to DFAs.

2.2

1 Deterministic finite automata

Deterministic finite automata

Finite automaton: a simple computing machine

- Given a finite input word, it moves from one program state to another.
- Each move is based on one input letter
- At the end of the input, the machine either *accepts* or *rejects* the input, depending on the machine state.

Vital limitation of a finite automaton:

- It cannot look back in its input.
- The only memory is in the state; aside from that, it has forgotten everything.

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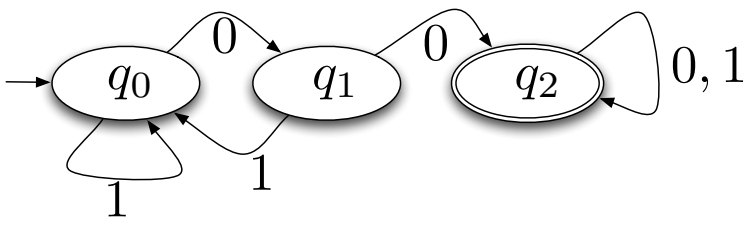
Definitions

Finite automaton described by 5 parameters:

- Q = set of computation states
- Σ = finite input alphabet
- δ = transition function
  - **Important:** In a DFA, there must be a transition defined for *every state* and for *every possible alphabet character*.
- q_0 = start state
- F = accept states



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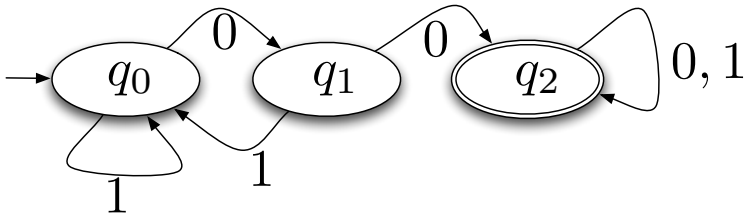
This DFA



In this finite automaton:

- $Q = \{q_0, q_1, q_2\}$,
- $\Sigma = \{0, 1\}$,
- δ is a function from $Q \times \Sigma \to Q$. It includes $(q_0, 0) \to q_1$
- $q_0 = q_0$,
- $F = \{q_2\}$.

This DFA *accepts* the *language* of words with 00 as a substring. **Question:** How would you prove that? 2.5

Acceptance, extension

Given a DFA M, what does "M accepts w" mean?

- Starting at q_0, follow transition function δ for each letter in w, in order.
- String w accepted by M if at the end of w's transitions, we wind up in a state in F.

A more formal definition of acceptance comes by looking at the *extended transition function*, $\hat{\delta}$.

- $\hat{\delta}(q, w)$: state we finish in if we start at state q and follow δ for each letter in the word w in turn.

Function $\hat{\delta}$ is a function from $Q \times \Sigma^* \to Q$. It usually cannot be written down in a closed form, which leads us to the following technique. 2.6

Formal definition of extended transition function

Formally, $\hat{\delta}(q, w)$ is defined recursively:

- $\hat{\delta}(q, \varepsilon) = q$, for all states q.
- If $|w| > 0$, then we can write $w = xa$, where $|a| = 1$.
- Then define $\hat{\delta}(q, w) = \delta(\hat{\delta}(q, x), a)$.

Unpack that:

- Let $q_1 = \hat{\delta}(q, x)$.
- In words, q_1 is the state that we reach starting from q after we have read the prefix x.
- Then, process the single letter a: $\delta(q_1, a) = \delta(\hat{\delta}(q, x), a)$.

This can be defined from the other end:

- If $w = ax$, then $\hat{\delta}(q, w) = \hat{\delta}(\delta(q, a), x)$.
- The textbook gives the first definition, so we will stick with that. 2.7

Lemma about extended transition function

Lemma: Let $D = (Q, \delta, \Sigma, q_0, F)$ be a DFA, with extended transition function $\hat{\delta}$. Let $q \in Q$ and $a \in \Sigma$ be arbitrary. Then $\hat{\delta}(q, a) = \delta(q, a)$, i.e. $\hat{\delta}$ agrees with δ for any state q and on any single alphabet symbol a.

Proof:

$$\begin{aligned} \hat{\delta}(q, a) &= \hat{\delta}(q, \varepsilon a) \\ &= \delta(\hat{\delta}(q, \varepsilon), a) \\ &= \delta(q, a). \end{aligned}$$

□ 2.8

Language of an DFA

The *language* of the DFA M: all words M accepts.

Formally, acceptance of a word:

- $M = (Q, \Sigma, \delta, q_0, F)$ accepts $w \in \Sigma^*$ exactly when $\hat{\delta}(q_0, w) \in F$.

Language of the DFA: all accepted words.

- $L(M) = \{w \in \Sigma^* \text{ where } \hat{\delta}(q_0, w) \in F\}$.
- (or just $\{w \in \Sigma^* \text{ where } M \text{ accepts } w\}$.)

Terminology: $L(M)$ can be called:

- The language of the DFA M
- The language *accepted* by the DFA M
- The language *recognized* by the DFA M

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Do DFAs compute?

In a certain sense, yes.

Example:

- "Is x a multiple of 3?" can be answered by a DFA.
  - The input word w is the binary representation of x.
  - Then the machine accepts w if x is a multiple of 3.
- This language, $L = \{w | w \text{ is the binary representation of a number } x \text{ that is divisible by } 3\}$, is accepted by an DFA.
  - L includes $11, 110, 1111, 0$, and does not include $10, \varepsilon$.
- So in that sense, yes, they compute.

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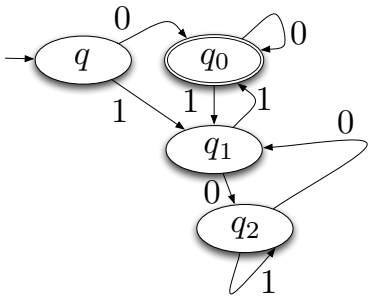
Divisibility DFA

Reminder of conventions:

- start state has an unlabelled arrow
- accept states are double circles
- label arrows with values for δ

We will have a state for each remainder (0, 1, and 2), plus a special start state, so we do not accept ε.

This gives this DFA:



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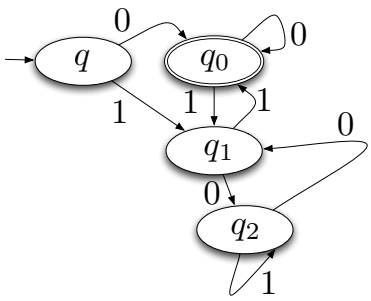
Why does that DFA work?



- Adding a new symbol (0 or 1): double and add the new symbol.
- Double a multiple of 3 and add 0: a new multiple of 3,
- Double a multiple of 3 and add 1: a number with remainder 1,
- Double a number with remainder 1 and add 1: a multiple of 3,
- *etc.*
- We are in state q_i when i is the remainder considering what we have already read.

This machine only accepts multiples of 3. 2.12

Enhancements to DFAs

We are going to prove a theorem soon that FAs accept a specific class of languages, called *regular languages*.

- That should be robust: changes to an FA should not invalidate the property.
- So we will change FAs in a variety of small and large ways.
- The first major change is *nondeterminism*.

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2 Nondeterministic Finite Automata

Nondeterminism

How does a DFA work?

- In a given state, for each input letter, there is exactly one choice for what to do, and
- The machine accepts if after all letters have been read, it is in an accept state.

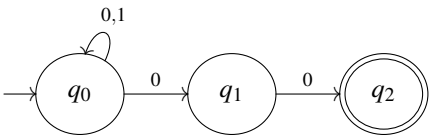
What if there were choices, instead?

- In a given state, for each letter, there may be a *choice* of what to do.
- The machine accepts if *some* sequence of choices results in an accept state.

Nondeterministic FAs allow a huge blow-up in computation: there will be lots going on in parallel. As such they may not be very realistic models of any kind of computers. 2.14

An example

$L = \{$all words with 00 as the last two symbols$\}$



- Nondeterministic FA: Transition from a state, given an alphabet symbol, is to a *set* of possible states.
- *Note*: sometimes a state has no outgoing transition for a given symbol; the second state has no output labelled 1.
- If a thread reaches a state which has no outgoing transition for the given input symbol, then that thread *crashes*; it proceeds no further.

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Formal definition

NFA defined by 5 parameters

- Q = set of computation states
- Σ = finite input alphabet
- δ = transition function
  - **Important:** In an NFA, there need *not* be a transition defined for every state and for every possible alphabet character.
- q_0 = start state

- F = accept states

Differences from DFAs:

- δ: function from $Q \times \Sigma \to \{\text{subsets of } Q\}$.
  - Recall notation: $2^Q = \{\text{subsets of } Q\}$.
  - Based on a state in Q and an input letter from Σ, which states are now active in Q?
  - (Different from DFA, where it is from $Q \times \Sigma \to Q$.)
- Accepts whenever *any* state path from q_0 is to an accept state.

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New form for extended transition function

We must enhance the definition for our extended transition function. We now need to allow the output of the transition function to be a set of states instead of a single state.

- $\hat{\delta}(q,w)$ = all states that we can reach from start state q processing the input word w.
- $\hat{\delta}$: function $Q \times \Sigma^* \to 2^Q$.
- Base case: $\hat{\delta}(q,\varepsilon) = \{q\}$.
- Recursive case: If $|w| > 0$, then we may write $w = xa$, where $|a| = 1$.
- Then can we define $\hat{\delta}(q,xa) = \delta(\hat{\delta}(q,x),a)$?
  - Well, no.
  - The function δ is $Q \times \Sigma \to 2^Q$.
  - But $\hat{\delta}(q,x)$ is in 2^Q instead of in Q, so $\delta(\hat{\delta}(\ldots),a)$ is not allowed.
  - We really need to define $\hat{\delta}(q,xa) = \bigcup_{p \in \hat{\delta}(q,x)} \delta(p,a)$ instead.
  - *Note:* This definition handles threads that crash correctly.
    - A thread that crashes has no outgoing transition from state p for input symbol a.
    - In other words, $\delta(p,a) = \emptyset$.
    - But then, $\delta(p,a)$ contributes nothing to the union of sets of states, reflecting the fact that the thread has crashed and proceeds no further.

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Acceptance by an NFA and the language of an NFA

Acceptance of an NFA:

- NFA M accepts a word w if $\hat{\delta}(q_0,w) \cap F$ is nonempty.
- There is a path from q_0, labelled by letters of w, winding up in an accept state from F.

Language of an NFA:

- The language of the NFA $M = (Q,\Sigma,\delta,q_0,F)$ is:

$$L(M) = \{w \in \Sigma^* | \hat{\delta}(q_0,w) \cap F \neq \emptyset\}.$$

- That is, all words accepted by the NFA.

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Where are we?

DFAs:

- Model of computation: finite states, follow single transitions
- Acceptance of a word: transitions lead to an accept state from M
- Language: All accepted words

NFAs:

- Model of computation: finite states, possibly many transitions per letter, or possibly none.
- Acceptance: any path of transitions leads to an accept state.
- The extended transition function is more complicated.
- Language: All accepted words.

Are these different from each other in terms of power?
Both are limited in that they only read each input letter once, left-to-right.

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NFAs are no more powerful than DFAs

Theorem: Let L be a language that is accepted by an NFA. Then L is accepted by a DFA.

- *Note*: The opposite direction is easy: DFAs are NFAs!
  (Well, must change δ trivially, so that the NFA transitions to the 1-element set corresponding to the transition in the DFA.)
- Need to show: Given an NFA, we can construct a DFA that accepts its language.

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Outline of proof

Here is how we will do this:

- Given an NFA N, we will construct a DFA, D.
- Then, we will have to show that $L(D) = L(N)$.
- As this is an equality of sets, we need to show that every word in $L(N)$ is in $L(D)$, and every word in $L(D)$ is in $L(N)$.

You have seen a little of this in CS 241.

First, let's build the DFA D, using the same alphabet Σ that N uses.

Remember: The machine D does *not* have to have the same states as N, just the same alphabet and language!

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Computation in an NFA

How does the NFA N work?

- If N is in state q, after processing one letter a, N could be in any state from the set $\delta(q,a)$.
- Then, N processes the next letter and winds up in any of another set of states.
  (That is what is built into the extended transition function, $\hat{\delta}_N$.)
- In the DFA D, a single state represents a *set of states* in the NFA N.
- When D reads a new letter in, we jump from one state in the DFA to another (corresponding to the appropriate sets of states in the NFA N).

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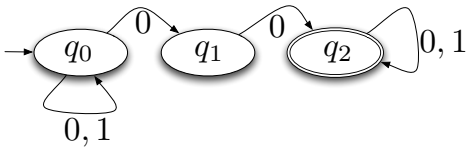
Sketch of the process of the proof

This NFA accepts the language $L = \{\text{words with 00 as a substring}\}$. How would you prove that?



- Its only accept state is q_2.
- Suppose we have processed some letters, and the NFA could be in either state q_0 or q_1, and the next input letter is 1.
  - From q_0 we go to q_0.
  - From q_1, no transition labelled 1.
- The new DFA, if it is in the *set* state $\{q_0, q_1\}$ and reads in a 1, must go to the state corresponding to the set $\{q_0\}$.

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Sketch of the process of the proof

Here is idea of the algorithm for the *subset construction* for the transition function of D:

- For each subset $S \subseteq Q$ of states from N:
  - Recall that S corresponds to a single state in D.
  - For each alphabet symbol, a:
    - For each state $p \in S$, consider $\delta(p,a)$ (recall, this is a set of states).
    - Gather all of these together. Let $T = \bigcup_{p \in S} \delta(p,a)$.
    - Then T also corresponds to a single state in D.
    - Add the transition $S \xrightarrow{a} T$ to the transition function δ_D for D.

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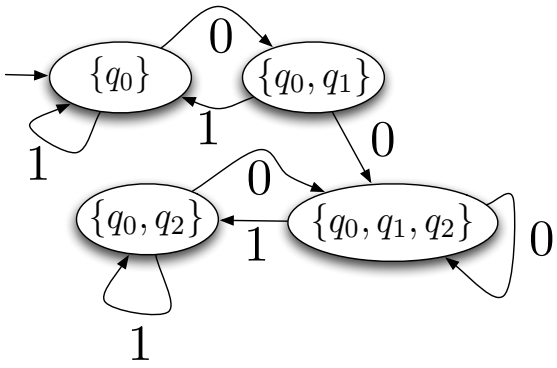
This procedure may make many new states

This 3-state NFA turns into a 4-state DFA.

- In general, if the NFA N has k states, the DFA D could have 2^k states.

Here is the DFA:



And what are accept states?

- $F_D = \{$States in D that represent sets of states in N that include at least one accept state from $F_N\}$.

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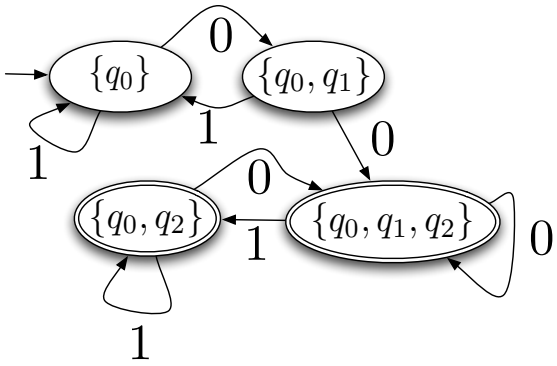
The full DFA



- Convince yourself that accepts the same language as the original construction.
- Interestingly, the DFA state $\{q_0, q_2\}$ is not needed, since we only get to it from another accept state.
- We really only needed 3 states.

Now, let's generalize this idea.

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Subset construction (formal)

Given NFA $N = (Q_N, \Sigma, \delta_N, q_0, F_N)$, construct a new DFA $D = (Q_D, \Sigma, \delta_D, \{q_0\}, F_D)$, with these parameters:

- $Q_D = 2^{Q_N}$.
- Let $S \in Q_D$, i.e. think of S as a set of states from the definition of N. Then
- $F_D = \{S \in Q_D \mid S \cap F_N \neq \emptyset\}$
- That is, each acceptance state in D corresponds to a set of states in N with at least one accept state.
- And a more complicated transition function:

$$\begin{aligned} & \delta_D(S,a) \\ = \; & \{\text{all states reachable in } N \text{ from } S \text{ when we read } a\} \\ = \; & \bigcup_{p \in S} \delta_N(p,a). \end{aligned}$$

- Let D's initial state be $\{q_0\}$, where q_0 is the initial state of N.

This is a DFA, not an NFA. We know which state D is in after reading any alphabet symbol. 2.27

The languages are equal

Now we must show that the languages of the NFA N and the DFA D equal. Think about which words are in the languages of N and of D.

- $w \in L(N) \Leftrightarrow \hat{\delta}_N(q_0, w) \cap F_N \neq \emptyset$.
- $w \in L(D) \Leftrightarrow \hat{\delta}_D(\{q_0\}, w) \in F_D$.

By the definition of F_D, the second statement is equivalent to:

- $w \in L(D) \Leftrightarrow \hat{\delta}_D(\{q_0\}, w) \cap F_N \neq \emptyset$.

We must show these are the same: that if $w \in L(D)$, then $w \in L(N)$, and vice versa. 2.28

Proof

Must show: $w \in L(N) \Leftrightarrow w \in L(D)$.

- That is: $\hat{\delta}_N(q_0, w) \cap F_N \neq \emptyset \Leftrightarrow \hat{\delta}_D(\{q_0\}, w) \cap F_N \neq \emptyset$.
- It suffices to show: $\hat{\delta}_N(q, w) = \hat{\delta}_D(\{q\}, w)$, for any state q from the definition of N.
- (If one of these sets has a non-empty intersection with F_N, then so does the other)

Proof: By induction on $|w|$.

- Base case ($|w| = 0$): Thus $w = \varepsilon$. Then $\hat{\delta}_N(q, \varepsilon) = \{q\} = \hat{\delta}_D(\{q\}, \varepsilon)$.
- Inductive case ($|w| > 0$): The inductive hypothesis is that, for every x with $|x| < |w|$, we have $\hat{\delta}_N(q, x) = \hat{\delta}_D(\{q\}, x)$, for any state q from N.
- Since $|w| > 0$, we may write $w = xa$, where $|a| = 1$.
- Then the induction hypothesis applies to x, so $\hat{\delta}_N(q, x) = \hat{\delta}_D(\{q\}, x)$.

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Inductive case of the proof

- The induction hypothesis is that $\hat{\delta_N}(q, x) = \hat{\delta_D}(\{q\}, x)$.
- We need to show that $\hat{\delta_N}(q, xa) = \hat{\delta_D}(\{q\}, xa)$.
- So now we process the last character, a, on both sides.

$$
\begin{aligned}
\hat{\delta}_D(\{q\}, xa) & \underbrace{=}_{\text{Definition of } \hat{\delta}_D \text{ for the DFA } D} && \delta_D(\hat{\delta}_D(\{q\}, x), a) \\
& \underbrace{=}_{\text{induction hypothesis}} && \delta_D(\hat{\delta_N}(q, x), a) \\
& \underbrace{=}_{\text{Definition of } \delta_D} && \bigcup_{p \in \hat{\delta}_N(q, x)} \delta_N(p, a) \\
& \underbrace{=}_{\text{Definition of } \hat{\delta}_N} && \hat{\delta_N}(q, xa).
\end{aligned}
$$

Reading one more symbol in N, the set of states of N we can be in corresponds with the state we will be in in D (recall that a single state of D represents a set of states of N).

- The NFA N accepts the same language as the DFA D.

The class of languages accepted by NFAs is the same as for DFAs. 2.30

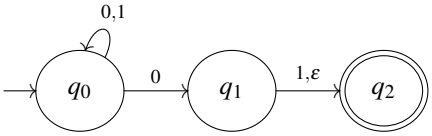
3 epsilon-NFAs

Another expansion to NFA's: ε-transitions

Sometimes in designing an NFA, it is handy to have transitions that happen automatically, without reading any letters of the input.

- Machine models that allow this are ε-*NFAs*.
- An ε-NFA is a 5-tuple, like an NFA.
- The only difference is transition function:
- δ is now a function: $Q \times (\Sigma \cup \{\varepsilon\}) \to 2^Q$.
- (Transitions may exist that do not consume input letters.)

Example: The ε-NFA



Accepts binary words ending with 0 or with 01. (We could do this without ε-transitions, by making the second state an accept state.)

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Formality about ε-NFAs

How do their extended transition functions look?

- Which states could I be in after processing the word x?
- Any states I could be in after just x, with no ε-transitions, plus
- Any states after just x, and *one* ε transition, plus
- Any states after just x, and *two* ε transitions, and so on ...
- ...until we exhaust all possible ε-transitions.

Unroll this to get a recursive definition for $\hat{\delta}$.

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Defining the ε-closure of a state

The definitions should be fairly similar:

- $\hat{\delta}(q, y)$: states we can get to after processing y and having ε-transitions.

Let's start with ε-transitions:

- Let $\text{ECLOSE}(p)$ = all states reachable starting from p only using ε-transitions.
- We can define $\text{ECLOSE}(p)$ recursively:
  - $p \in \text{ECLOSE}(p)$
  - If $q \in \text{ECLOSE}(p)$, then so are all of the states in $\delta(q, \varepsilon)$.

The set $\text{ECLOSE}(p)$ is called the ε-*closure* of the state p.
We will also allow $\text{ECLOSE}(S)$ to be defined for a set S of states via:

$$\text{ECLOSE}(S) = \bigcup_{s \in S} \text{ECLOSE}(s).$$

2.33

Defining the ε-closure of a state

Lemma: For an ε-NFA E, subset $S \subseteq Q$ and decomposition $S = \bigcup_i S_i$,

$$\bigcup_i \text{ECLOSE}(S_i) = \text{ECLOSE}\left(\bigcup_i S_i\right),$$

(i.e. taking ε-closure commutes with taking set unions). **Proof:**

$$\begin{aligned}
\bigcup_i \textsc{Eclose}(S_i) &\underbrace{=}_{\text{Definition of } \textsc{Eclose}(S_i)} \bigcup_i \left(\bigcup_{s \in S_i} \textsc{Eclose}(s) \right) \\
&\underbrace{=}_{S = \bigcup_i S_i} \bigcup_{s \in S} \textsc{Eclose}(s) \\
&\underbrace{=}_{\text{Definition of } \textsc{Eclose}(S)} \textsc{Eclose}(S) \\
&\underbrace{=}_{S = \bigcup_i S_i} \textsc{Eclose}\left(\bigcup_i S_i \right). \square
\end{aligned}$$

2.34

The definition of the extended transition function

We can define $\hat{\delta}$ for ε-NFA's, also recursively.

- Base case: $\hat{\delta}(q, \varepsilon) = \textsc{Eclose}(q)$: states reachable from q with ε-transitions
- Inductive case: Suppose that $w = xa$, where $a \in \Sigma$. (*Note*: *a cannot be* ε, which is not a member of Σ.)
  - We know that $P = \hat{\delta}(q, x)$ is the set of all states in Q that we can get to by following either edges for the letters of x or ε-transitions (including ε-transitions at the end of x).
  - Then, we must follow the transitions for the alphabet symbol a: Let $R = \bigcup_{p \in P} \delta(p, a)$: then R has all of the states we can get to from P after following a transition for a.
  - Last, we might have some more ε-transitions.
- So, $\hat{\delta}(q, w) = \textsc{Eclose}(R) = \textsc{Eclose}\left(\bigcup_{p \in \hat{\delta}(q,x)} \delta(p, a) \right)$

2.35

Languages and power of ε-NFAs

- Language of an ε-NFA: $L = \left\{ w \in \Sigma^* \mid \hat{\delta}(q_0, w) \cap F \neq \emptyset \right\}$
- That is, $\hat{\delta}(q_0, x)$ includes an accept state.

Are ε-NFAs more powerful?

- No.
- Theorem: Given an ε-NFA E, there exists an ordinary DFA D such that $L(D) = L(E)$.
- This is not very surprising: we must show that we can include the ε-transitions of E in the transition function δ_D for D.
- (The other direction is just by definition: a DFA is an ε-NFA, once we make some trivial changes to the structure of δ so that it produces a 1-element set when it reads a symbol and the empty set when it reads in ε.)

To prove the theorem, we must construct a DFA D accepting language $L(E)$.

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The equivalent DFA

- Both machines use the same alphabet, of course.
- We use the subset construction, as when we built the DFA for an NFA.
- Starting state is $q_D = \{\textsc{Eclose}(q_0)\}$. Thus, we start having implicitly taken ε-transitions from the starting state q_0 of E.

The complexity comes in the transition function and the accept states.

- Transition function:
  - From one DFA state, S (corresponding to a set of states in E), if we process one letter a in the new DFA, we should mimic this behaviour from E:

    - follow any edges labelled a
    - take any ε-transitions
  - From any one state from E, say q, this then takes us to:
    - $\delta_E(q,a)$
    - $\text{ECLOSE}(\delta_E(q,a))$
- And we therefore want the union over all states $q \in S$: $\delta_D(S,a) = \bigcup_{q \in S} \text{ECLOSE}(\delta_E(q,a))$.

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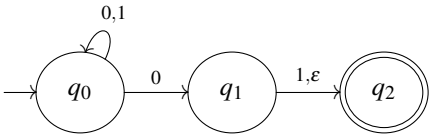
Example of Subset Construction

- Here we demonstrate one step in the subset construction for the earlier small example of an ε-NFA:



- The subset construction gives us

$$\begin{aligned} Q_D &= \{\emptyset, \{q_0\}, \{q_1\}, \{q_2\}, \{q_0,q_1\}, \{q_0,q_2\}, \{q_1,q_2\}, \{q_0,q_1,q_2\}\} \\ q_D &= \text{ECLOSE}(\{q_0\}) = \{q_0\} \end{aligned}$$

- Now we determine $\delta_D(S,a)$, where $S = \{q_0,q_1,q_2\}$ (the state we reach from $\{q_0\}$ upon reading 0) and $a = 1$.
- Computing $\text{ECLOSE}(\delta_E(p,a))$ for each $p \in S$ gives

$$\begin{aligned} \text{ECLOSE}(\delta_E(q_0,1)) &= \text{ECLOSE}(\{q_0\}) = \{q_0\} \\ \text{ECLOSE}(\delta_E(q_1,1)) &= \text{ECLOSE}(\{q_2\}) = \{q_2\} \\ \text{ECLOSE}(\delta_E(q_2,1)) &= \text{ECLOSE}(\emptyset) = \emptyset \end{aligned}$$

2.38

Example of Subset Construction

- Hence the target state coming from the subset construction is $\{q_0,q_2\}$.
- The construction says that we need to add to the transition function for D: $\delta_D(\{q_0,q_1,q_2\},1) = \{q_0,q_2\}$.
- Now to complete the transition function for D, we do this same construction for each of the 8 choices for S (on the previous slide) and each alphabet symbol form $\Sigma = \{0,1\}$.

2.39

Picking the accept states, equality of languages

- We need to figure out which states are accepting states:
  - $F_D = \{S \mid S \in Q_D \text{ and } S \cap F_E \neq \emptyset\}$.
- Declare a word accepted by D if D is in an accept state (according to this recipe) when D finishes processing the word.
- Now, to show equality of the languages of the two automata, we must show that if x is accepted by E, then x is accepted by D, and vice versa.
  - (One concern: do we do the right thing for the word ε?)
  - To show $L(E) = L(D)$, can show that $\hat{\delta}_E(q_0,x) = \hat{\delta}_D(q_D,x)$?
  - If so, then we will be in the same set of ε-NFA states after reading x, and our definitions of F_D and F_E will guarantee that both machines will accept exactly the same words.

2.40

Transition functions

We want to show: $\hat{\delta}_E(q_0,w) = \hat{\delta}_D(q_D,w)$, for all strings w. The proof is by induction on $|w|$.

Base case ($|w| = 0$): In this case, $w = \varepsilon$. We have $\hat{\delta}_E(q_0,\varepsilon) = \text{ECLOSE}(\{q_0\})$, by definition of $\hat{\delta}_E$ in the ε-NFA.

- We therefore have $\hat{\delta}_D(q_D,\varepsilon) \underbrace{=}_{D \text{ is a } DFA} q_D \underbrace{=}_{\text{construction}} \text{ECLOSE}(\{q_0\}) \underbrace{=}_{\text{Definition}} \hat{\delta}_E(q_0,\varepsilon)$.
- So the base case holds.

Inductive case ($|w| > 0$):

- We need to argue that $\hat{\delta}_E(q_0,w) = \hat{\delta}_D(q_D,w)$.
- The induction hypothesis is that $\hat{\delta}_E(q_0,x) = \hat{\delta}_D(q_D,x)$, for all strings x where $|x| < |w|$.
- Write $w = xa$, where a is a single character.
- The induction hypothesis applies to x.
- Thus we may let $\hat{\delta}_E(q_0,x) = \hat{\delta}_D(q_D,x) = S$.

2.41

Transition functions

- Now we compute

$$
\begin{aligned}
\hat{\delta}_E(q_0,xa) &\underbrace{=}_{\text{Definition of } \hat{\delta}_E} \text{ECLOSE}\left(\bigcup_{p \in \hat{\delta}_E(q_0,x)} \delta_E(p,a)\right) \\
&\underbrace{=}_{\text{Definition of } S} \text{ECLOSE}\left(\bigcup_{p \in S} \delta_E(p,a)\right) \\
&\underbrace{=}_{\text{Lemma}} \bigcup_{p \in S} \text{ECLOSE}(\delta_E(p,a)) \\
&\underbrace{=}_{\text{Definition of } \delta_D} \delta_D(S,a) \\
&\underbrace{=}_{\text{Definition of } S} \delta_D(\hat{\delta}_D(\text{ECLOSE}(q_0),x),a) \\
&\underbrace{=}_{\text{Definition of } \hat{\delta}_D} \hat{\delta}_D(\text{ECLOSE}(q_0),xa), \text{ as required.}
\end{aligned}
$$

2.42

Transition functions

- **Remark:** You should convince yourself that the base case handles the input word ε correctly.
- We have proved by induction that the two extended transition functions agree on all input words.
- Therefore, as we argued earlier, this shows that the two automata accept precisely the same languages.
- So we are done.

2.43

End of module 2

Hence, the class of languages accepted by ε-NFAs is the same as the class accepted by ordinary NFAs, which is the same as the class of languages accepted by DFAs.

We have now seen a collection of types of automata:

- Deterministic finite automata
- Nondeterministic finite automata
- ε-NFAs

All three accept the same class of languages. But what is that class?

They are the *regular languages*.

2.44