# Carmen's Core Concepts (Math 135)

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Week 8

- The following are equivalent (TFAE)
- 2 Inverses
- 3 More on Multiplicative Inverses
- 4 Linear Congruence Theorem 2 [LCT2]
- 5 Fermat's Little Theorem  $[F\ell T]$
- 6 Example of Fermat's Little Theorem
- 🕜 Important Corollaries to F $\ell$ T
- 8 Chinese Remainder Theorem [CRT]
- Ohinese Remainder Theorem Example
- 10 Splitting the Modulus [SM]
- Introduction to Cryptography
- Public Key Cryptography
- 13 Square and Multiply Algorithm

- $a \equiv b \pmod{m}$
- *m* | (*a* − *b*)
- $\exists k \in \mathbb{Z}, a-b=km$
- $\exists k \in \mathbb{Z}, a = km + b$
- a and b have the same remainder when divided by m
- [a] = [b] in  $\mathbb{Z}_m$ .

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For example, solving [10][x] = [1] is the exact same as solving  $10x \equiv 1 \pmod{m}$ .

#### Inverses

- [-a] is the additive inverse of [a], that is, [a] + [-a] = [0].
- If there exists an element [b] ∈ Z<sub>m</sub> such that
   [a][b] = [1] = [b][a], we call [b] the multiplicative inverse of
   [a] and write [b] = [a]<sup>-1</sup> or b ≡ a<sup>-1</sup> mod m.

## More on Multiplicative Inverses

Proposition: Let a ∈ Z and m ∈ N.
(a) [a]<sup>-1</sup> exists in Z<sub>m</sub> if and only if gcd(a, m) = 1.
(a) [a]<sup>-1</sup> is unique if it exists.
Proof:

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$$\begin{array}{ll} [a]^{-1} \text{ exists} & \Leftrightarrow & [a][x] = [1] \text{ is solvable in } \mathbb{Z}_m \\ \Leftrightarrow & ax + my = 1 \text{ is a solvable [LDE]} \\ \Leftrightarrow & \gcd(a, m) = 1 \text{ [GCDOO]} \end{array}$$

completing the proof.

② Assume  $[a]^{-1}$  exists. Suppose there exists a  $[b] \in \mathbb{Z}_m$  such that [a][b] = [1] = [b][a]. Then

$$[a]^{-1}[a][b] = [a]^{-1}[1]$$
  
 $[1][b] = [a]^{-1}$   
 $[b] = [a]^{-1}$ 

**Theorem:** Let  $a, c \in \mathbb{Z}$  and let  $m \in \mathbb{N}$ . Let gcd(a, m) = d. The equation [a][x] = [c] in  $\mathbb{Z}_m$  has a solution if and only if  $d \mid c$ . Moreover, if  $[x] = [x_0]$  is one particular solution, then the complete solution is

$$\left\{ [x_0], [x_0 + \frac{m}{d}], [x_0 + 2\frac{m}{d}], \dots, [x_0 + (d-1)\frac{m}{d}] \right\}$$

# Fermat's Little Theorem $[F\ell T]$

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**Theorem:** If *p* is prime and  $p \nmid a$  then  $a^{p-1} \equiv 1 \mod p$ . Equivalently,  $[a^{p-1}] = [1]$  in  $\mathbb{Z}_p$ . **Proof:** Major Ideas:

• Lemma: Let gcd(a, p) = 1. Let

$$S = \{a, 2a, ..., (p-1)a\} \qquad T = \{1, 2, ..., p-1\}.$$

Then the elements of S are unique modulo p and for all  $s \in S$ , there exists a unique element  $t \in T$  such that  $s \equiv t \mod p$ .

$$\prod_{x \in S} x \equiv \prod_{y \in T} y \mod p \Longleftrightarrow \prod_{k=1}^{p-1} ka \equiv \prod_{j=1}^{p-1} j \mod p$$
$$\iff a^{p-1} \prod_{k=1}^{p-1} k \equiv \prod_{j=1}^{p-1} j \mod p \iff a^{p-1} \equiv 1 \mod p$$

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$$7^{92} \equiv 7^{9(10)+2} \mod 11$$
  
 $\equiv (7^{10})^9 7^2 \mod 11$   
 $\equiv 1^9 \cdot 7^2 \mod 11$  By F $\ell$ T since  $11 \nmid 7$   
 $\equiv 49 \mod 11$   
 $\equiv 5 \mod 11$ 

- **Corollary:** If p is a prime and  $a \in \mathbb{Z}$ , then  $a^p \equiv a \mod p$ .
- Corollary: If p is a prime number and [a] ≠ [0] in Z<sub>p</sub>, then there exists a [b] ∈ Z<sub>p</sub> such that [a][b] = [1], namely [b] = [a<sup>p-2</sup>] = [a]<sup>p-2</sup>.
- Corollary: If r = s + kp, then  $a^r \equiv a^{s+k} \mod p$  where p is a prime and  $a \in \mathbb{Z}$  and  $r, s, k \in \mathbb{N}$ .
- Corollary: Prove that if  $p \nmid a$  and  $r \equiv s \mod (p-1)$ , then  $a^r \equiv a^s \mod p$ .

**Theorem:** If  $gcd(m_1, m_2) = 1$ , then for any choice of integers  $a_1$  and  $a_2$ , there exists a solution to the simultaneous congruences

 $n \equiv a_1 \pmod{m_1}$  $n \equiv a_2 \pmod{m_2}$ 

Moreover, if  $n = n_0$  is one integer solution, then the complete solution is  $n \equiv n_0 \pmod{m_1 m_2}$ .

#### Chinese Remainder Theorem Example

Solve the simultaneous congruence

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**Solution:** Write x = 2 + 7k for some  $k \in \mathbb{Z}$ . Into the second eqn:

$$2 + 7k \equiv 7 \mod 11$$
$$7k \equiv 5 \mod 11$$

Multiplying both sides by 3 gives

$$3 \cdot 7k \equiv 15 \mod 11 \iff 21k \equiv 4 \mod 11$$
  
 $\iff -k \equiv 4 \mod 11 \iff k \equiv 7 \mod 11$ 

Therefore,  $k = 7 + 11\ell$  for some  $\ell \in \mathbb{Z}$ . Thus, since x = 2 + 7kand  $k = 7 + 11\ell$ , we have

$$x = 2 + 7k = 2 + 7(7 + 11\ell) = 51 + 77\ell$$

Therefore,  $x \equiv 51 \mod 77$  is the solution.

**Theorem:** Let m and n be coprime positive integers. Then, for any integers x and a, we have

 $x \equiv a \mod m$  $x \equiv a \mod n$ 

simultaneously if and only if  $x \equiv a \mod mn$ .

- What is Cryptography?
- Private vs Public Key Cryptography (Pad Lock analogy)

- Alice produces a private key d and a public key e.
- Bob uses the public key e to take a message M and encrypt it to some ciphertext C
- **③** Bob then sends *C* over an insecure channel to Alice.
- Alice decrypts C to M using d.
  - Encryption and decryption are inverses to each other.
  - d and e are different,
  - Only *d* is secret.

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Solution: First, we compute successive square powers of 5:

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- $5^{16} \equiv (58)^2 \equiv 31 \mod 101$

$$5^{32} \equiv (31)^2 \equiv 52 \mod 101$$

$$5^{64} \equiv (52)^2 \equiv 78 \mod 101$$

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Now, in binary,  $99 = 64 + 32 + 2 + 1 = 2^6 + 2^5 + 2^1 + 2^0$ . Hence,

$$5^{99} \equiv 5^{64} \cdot 5^{32} \cdot 5^2 \cdot 5^1 \mod 11$$
  
 $\equiv 78 \cdot 52 \cdot 25 \cdot 5 \mod 11$   
 $\equiv 81 \mod 11$