

$$r(P) = \sum_{Q_i \rightarrow P} \frac{r(Q_i)}{|Q_i|}$$

of outlinks.

$$r(P) = \frac{1 - \delta}{N} + \delta \sum_{Q_i \rightarrow P} \frac{r(Q_i)}{|Q_i|}$$

$$r(A) + r(B) + r(C) + r(D) = 1. \quad \textcircled{1}$$

$$r(A) = r(C) \quad \textcircled{2}$$

$$r(B) = r(A) / 3 = r(C) / 3 \quad \textcircled{3}$$

$$r(C) = \frac{r(A)}{3} + r(B) + r(D) \quad \textcircled{4}$$

$$r(D) = \frac{r(A)}{3} = \frac{r(C)}{3} \quad \textcircled{5}$$

Use eqn ①.

$$r(C) + \frac{r(C)}{3} + r(C) + \frac{r(C)}{3} = 1.$$

$$\frac{8}{3} r(C) = 1$$

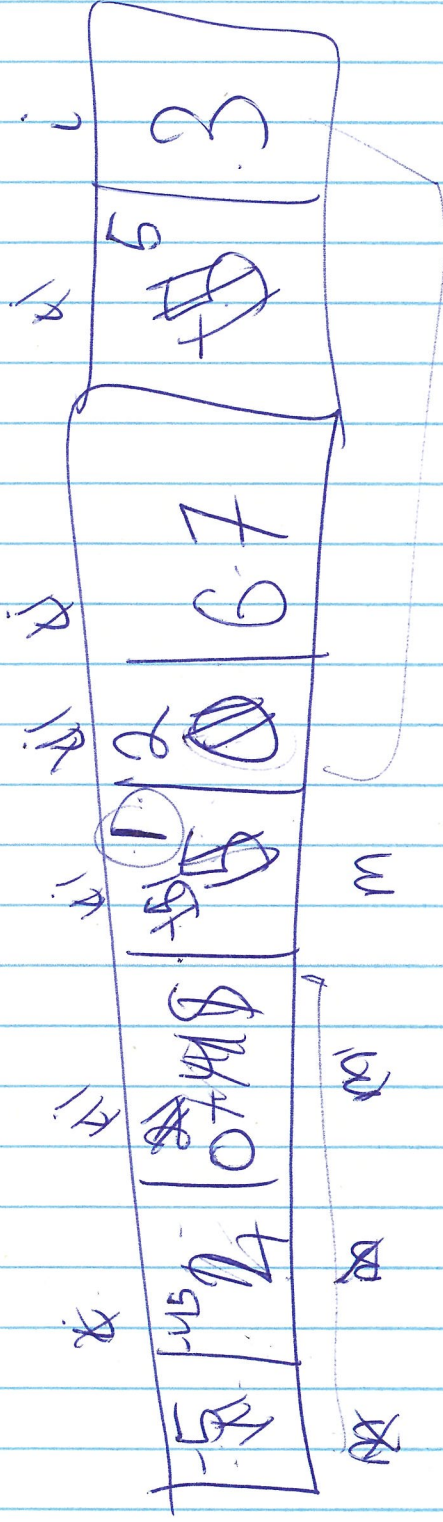
$$r(C) = \frac{3}{8} \quad \textcircled{6}$$

⑤ into ② gives $r(A) = \frac{3}{8}$.

⑤ into ③ and ④ gives $r(B) = r(D) = r(C) / 3 = \frac{1}{8}$.

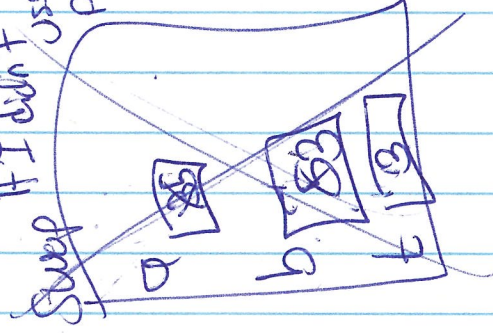
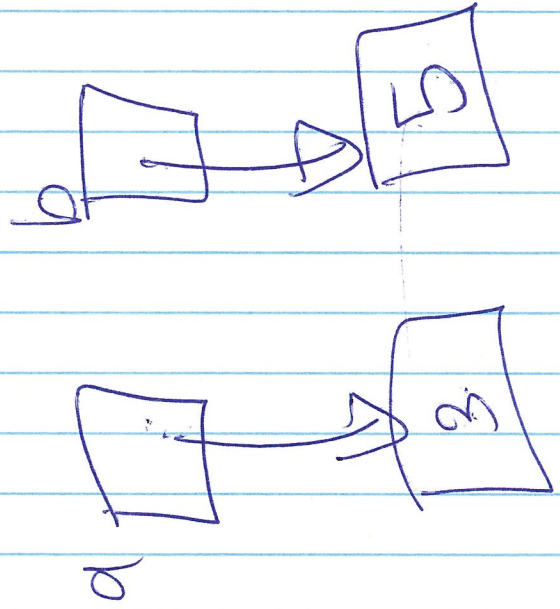
Quick Sort.

②



Last thing, swap $a[l]$ and $a[m]$.

Swap if I didn't use pointers.



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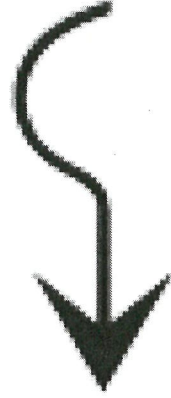
Clicker

Recalling our table below, how many ones in UTF-8 would have been used for our encoding of the arrow U+2B3F?

Code Point Range in Hex	UTF-8 Byte Sequence in Binary
→ 000000-00007F	0xxxxxx <i>7 bits</i>
→ 000080-0007FF	110xxxxx 10xxxxxx <i>6 bits</i> <i>11 bits</i>
→ 000800-00FFFF	1110xxxx 10xxxxxx 10xxxxxx <i>11 bits</i> <i>16 bits</i>
010000-10FFFF	11110xxx 10xxxxxx 10xxxxxx 10xxxxxx <i>21 bits</i>

2B3F = 0010 1011 0011 1111
A ↑ ↑ ↑ B 3 F

7FF
0111 1111 1111
11 ones



- a) 11
- b) 12
- c) 13
- d) 14
- e) None of the above

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```

void fun2(int n) {
    int s=0;
    for (int i=0; i <= sqrt(n); i++) {
        for (int j=1; j < i; j*=4) {
            for (int k=0; k <= pow(3, j); k++) {
                s++;
            }
        }
    }
}

```

for some constant C
 $3^1 + 3^2 + 3^3 + \dots + 3^{\log_4 n}$
 $\frac{3^{(\log_4 n + 1)} - 3}{3 - 1}$

$$\begin{aligned}
 \text{Runtime} &= \sum_{i=0}^{\sqrt{n}} \sum_{j=1}^{\log_4 i} \sum_{k=0}^{3^j-1} C \\
 &= C \sum_{i=0}^{\sqrt{n}} \sum_{j=1}^{\log_4 i} 3^j \\
 &= C \sum_{i=0}^{\sqrt{n}} \frac{3^{\log_4 i + 1} - 3}{3 - 1}
 \end{aligned}$$

```

void fun1(int n) {
    int s=0;
    for (int i=0; i < n; i++) {
        for (int j=0; j <= i; j++) {
            s++;
        }
    }
}

```

$$\begin{aligned}
 \text{Runtime} &= \sum_{i=0}^{n-1} \sum_{j=0}^i C \quad (\text{for some constant } C) \\
 &= C \sum_{i=0}^{n-1} \sum_{j=0}^i 1 \\
 &= C \sum_{i=0}^{n-1} i = C \left(\frac{n(n-1)}{2} \right) \\
 &= O(n^2) \quad \sum_{i=0}^{n-1} i = O(n^2)
 \end{aligned}$$

$$d \log n = \log n^a$$

$$\begin{aligned}
 3 \log_4 3 &= 3 \frac{\log_3 i}{\log_3 4} \\
 &= \cancel{3} \log_3 i \frac{1}{\log_3 4} \\
 &= i \frac{1}{\log_3 4}
 \end{aligned}$$

⑥

How many times does the 'j' loop run? Ans: 2 times where

$4^x < i \leq 4^{x+1}$

$$4^x < i \leq 4^{x+1}$$

\Rightarrow roughly $2 = \log_4 i$

$$O(\log n) = O(\log n)$$

$$O(\log \log n) = O(\log \log n)$$

$$\leq \frac{3}{2} \sum_{i=0}^{\sqrt{n}} 3^{\log_4 i}$$

$$= \frac{3}{2} \sum_{i=0}^{\sqrt{n}} i^{\frac{1}{\log_3 4}}$$

$$= O(\sqrt{n} \left(\frac{1}{\log_3 4} + 1 \right))$$

$$= O(\sqrt{n} \left(\frac{1 + \log_3 4}{2 \log_3 4} \right))$$

$$\frac{1 + \log_3 4}{2 \log_3 4}$$

exponent.

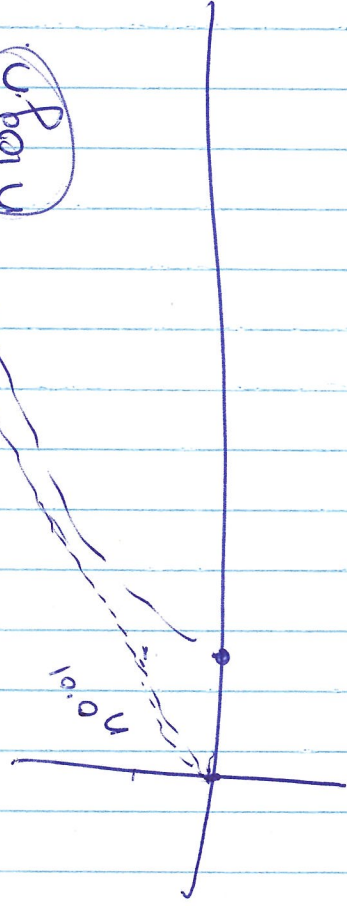
L'Hopital's rule.

Show $n \log n \in O(n^{1.01})$

pf. Recall if $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)}$ is finite then $f(n) = O(g(n))$.

$$\begin{aligned} \text{Evaluate } \lim_{n \rightarrow \infty} \frac{n \log n}{n^{1.01}} &= \lim_{n \rightarrow \infty} \frac{1}{0.01 n^{0.01}} = 0. \end{aligned}$$

\therefore From class, we have that $n \log n \in O(n^{1.01})$.



(7)

$$\lim_{n \rightarrow \infty} \frac{n \log n}{n^{1.01}} \stackrel{\text{L'H}}{=} \lim_{n \rightarrow \infty} \frac{(\frac{1}{n})}{0.01 n^{-0.99}}$$

$$\frac{\log n}{10}$$

$$\frac{n^{0.01}}{0.01 \cdot 10 \cdot 60}$$

$$\frac{60 \cdot 10^{0.6}}{n^{1.01}}$$

$$\frac{10}{60 \cdot 60}$$

$$\frac{10}{60 \cdot 10^{0.6}}$$