# Minimal Elements for the Prime Numbers 

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## Motivation

Fact
The following 26 numbers are prime:
$2,3,5,7,11,19,41,61,89,409,449,499,881,991,6469,6949$, 9001, 9049, 9649, 9949, 60649, 666649, 946669, 60000049, 66000049, 66600049

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The following 26 numbers are prime:
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## Claim

Give me a prime number and I can remove some of its digits to obtain a prime on this list!

## Minimal Primes

- The primes in this list are known as the minimal primes because this the smallest list of numbers for which this claim holds.


## Minimal Sets

- More generally, any language (set of strings over a finite alphabet) has its own minimal set of elements and the minimal primes are the minimal set of the language

$$
\{2,3,5,7,11,13,17,19,23, \ldots\}
$$

## Definitions

- $x$ is a subword of $y$ if one can strike out zero or more symbols of $y$ to get $x$.
- A string of symbols $s$ is minimal for a language $L$ if

1. $s$ is a member of $L$ and
2. $s$ does not contain another member of $L$ as a subword.

- $M(L)$ denotes the set of minimal elements of $L$.


## Higman-Haines Theorem

- $M(L)$ is finite for every language $L$.


## Computation of Minimal Sets

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- Can lead to some strange behaviour...
- The minimal set for primes of the form $4 n+1$ has 146 elements, the largest of which has 79 digits.
- The minimal set for primes of the form $4 n+3$ has 113 elements, the largest of which has 19,153 digits!


## Computation of Minimal Sets

## Proposed Computation Process

- The following process will determine $M(L)$ if it can be implemented:

1. $M:=\emptyset$
2. while $L \neq \emptyset$ do
3. choose $x$, a shortest string in $L$
4. add $x$ to $M$
5. remove from $L$ all words containing the subword $x$ 6. return $M$

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- Caveat: We might not have a nice way of performing operations on $L$.


## Computation of Minimal Sets

Using Overapproximations

- This process also works if $L$ is replaced with an overapproximation $L^{\prime}$, so long as once no more minimal elements remain to be found we can show that $L^{\prime}=\emptyset$.


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## Using Overapproximations

- This process also works if $L$ is replaced with an overapproximation $L^{\prime}$, so long as once no more minimal elements remain to be found we can show that $L^{\prime}=\emptyset$.
- In practice, we choose $L^{\prime}$ to be a regular language, e.g.,

$$
\{2,5\} \cup \Sigma^{*}\{1,3,7,9\}
$$

is a regular overapproximation to the set of primes over the alphabet $\Sigma:=\{0, \ldots, 9\}$.

## Computation of Minimal Sets

Sample Language

- We will work with overapproximations of the form $x L^{*} z$ where $x$ and $z$ are strings of digits and $L$ is a set of digits.
- To be able to apply the process previously described, we need to be able to test if $x L^{*} z$ contains a prime or not.


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- To be able to apply the process previously described, we need to be able to test if $x L^{*} z$ contains a prime or not.
- It is unknown if this problem is decidable.


## Computation of Minimal Sets

Necessary Operations

- In order to perform the process previously described, we need to perform the following operations on the language $x L^{*} z$ :

1. Determine if the language contains a prime.
2. If so, determine the smallest prime(s) in the language.
3. If a prime is found, shrink the language under consideration so that it no longer contains that prime.

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2. If so, determine the smallest prime(s) in the language.
3. If a prime is found, shrink the language under consideration so that it no longer contains that prime.

- And any strings which contain that prime as a subword.


## Proving that $x L^{*} z$ contains no primes

Method 1: Find a common divisor

Theorem. If $N$ divides $x z$ and all numbers of the form $x L z$ then $N$ divides all numbers of the form $x L^{*} z$.

## Proving that $x L^{*} z$ contains no primes

Method 1: Find a common divisor

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Example. 7 divides 49 and 469 so 7 divides 4669, 46669, and all numbers of the form $46^{*} 9$.

## Proof

$N$ divides $x z$ and all $x L z$ implies $N$ divides all $x L^{*} z$

Say $y \in L^{*}$ contains the digits $y_{1}, \ldots, y_{n}$ and $z$ is a digit. By telescoping,

$$
\begin{aligned}
x y z-x z & =\sum_{i=1}^{n}\left(x y_{i} y_{i+1} \cdots y_{n} z-x y_{i+1} \cdots y_{n} z\right) \\
& =\sum_{i=1}^{n} 10^{n-i}\left(x y_{i}-x\right) \\
& =\sum_{i=1}^{n} 10^{n-i-1}\left(x y_{i} z-x z\right)
\end{aligned}
$$

$N$ must divide $x y z$ since it divides every other term in this equation.

## Proving that $x L^{*} z$ contains no primes

## Method 2: Use an algebraic factorization

Let $[x]_{b}$ represent the evaluation of the string $x$ in base $b$; the following are some example algebraic factorizations:

$$
\begin{aligned}
& {[\overbrace{4 \cdots 4}^{n} 1]_{16}=\left(8 \cdot 4^{n}+7\right)\left(8 \cdot 4^{n}-7\right) / 15} \\
& {[1 \overbrace{0 \cdots 0}^{n} 1]_{8}=\left(2^{n+1}+1\right)\left(4^{n+1}-2^{n+1}+1\right)}
\end{aligned}
$$

Once $n$ is large enough the right side obviously factors and cannot be prime.

## Proving that $x L^{*} z$ contains no primes

Combination method

The family 19* in base 17 contains no primes, because

$$
[1 \overbrace{9 \cdots 9}^{2 n}]_{17}=\left(5 \cdot 17^{n}+3\right)\left(5 \cdot 17^{n}-3\right) / 16
$$

and all $[1 \overbrace{9 \cdots 9}^{2 n+1}]_{17}$ are even, since $[19]_{17}$ and $[1999]_{17}$ are even.

## Proving that $x L^{*} z$ contains a prime

- In practice, if $x L^{*} z$ could not be ruled out as only containing composites and $|L|>1$ then a relatively small prime could always be found in the language.
- Intuitively, this is because there are a large number of small strings in such a language, and at least one is likely to be prime.
- For example, there are $2^{n-2}$ strings of length $n$ in the language $1\{2,3\}^{*} 1$.


## Searching for primes in $x y^{*} z$

- In the case $|L|=1$ the family is of the form $x y^{*} z$, and there is only a single string of each length $\geqslant|x z|$.
- Some families $x y^{*} z$ could not be ruled out as only containing composites and no primes could be found in the family, even after searching through numbers with over 100,000 digits.


## Does $x y^{*} z$ contain large primes?

- The prime number theorem tells us that the chance that a random $n$-digit number is prime is approximately $1 / n$. If one conjectures the numbers $x y^{*} z$ behave similarly you would expect $\sum_{n=2}^{\infty} 1 / n=\infty$ primes of the form $x y^{*} z$.


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- The prime number theorem tells us that the chance that a random $n$-digit number is prime is approximately $1 / n$. If one conjectures the numbers $x y^{*} z$ behave similarly you would expect $\sum_{n=2}^{\infty} 1 / n=\infty$ primes of the form $x y^{*} z$.
- Of course, this doesn't always happen, but it's at least a reasonable conjecture in the absence of evidence to the contrary.


## In Practice. . .

- Many $x y^{*} z$ families contain no small primes even though they do contain very large primes.
- For example, the smallest prime in the base 23 family $9 E^{*}$ is $9 \mathrm{E}^{800873}$ which when written in decimal contains 1,090,573 digits.


## In Practice. . .

- Many $x y^{*} z$ families contain no small primes even though they do contain very large primes.
- For example, the smallest prime in the base 23 family $9 E^{*}$ is $9 E^{800873}$ which when written in decimal contains 1,090,573 digits.
- Technically, probable primality tests were used to show this (which have a very small chance of making an error) because all known primality tests run far too slowly to run on a number of this size.


## Shrinking the Language

- Recall that once a minimal prime has been found we want to shrink the language being searched while still keeping it large enough that it contains all remaining minimal primes.


## Shrinking $x L^{*} z$

- Say that $x y z$ is discovered to be prime with $y \in L$. Then $x L^{*} z$ can be replaced with

$$
x(L \backslash\{y\})^{*} z
$$

## Shrinking $x L^{*} z$

- Say that xyyz is discovered to be prime with $y \in L$. Then $x L^{*} z$ can be replaced with

$$
x(L \backslash\{y\})^{*} z \quad \cup \quad x(L \backslash\{y\})^{*} y(L \backslash\{y\})^{*} z
$$

## Shrinking $x L^{*} z$

- Say that $x y \hat{y} z$ and $x \hat{y} y z$ are discovered to be prime with $y, \hat{y} \in L$ and $y \neq \hat{y}$. Then $x L^{*} z$ can be replaced with

$$
x(L \backslash\{y\})^{*} z \quad \cup \quad x(L \backslash\{\hat{y}\})^{*} z .
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## Shrinking $x L^{*} z$

- Say that $x y \hat{y} z$ is discovered to be prime with $y, \hat{y} \in L$ and $y \neq \hat{y}$. Then $x L^{*} z$ can be replaced with

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x(L \backslash\{y\})^{*}(L \backslash\{\hat{y}\})^{*} z .
$$

## Exploring $x L^{*} z$

- If the methods we've discussed cannot be used to rule out or shrink $x L^{*} z$ where $L=\left\{y_{1}, \ldots, y_{n}\right\}$ then we can replace it by

$$
x L^{*} y_{1} z \quad \cup \quad x L^{*} y_{2} z \quad \cup \quad \cdots \quad \cup \quad x L^{*} y_{n} z
$$

and re-run the methods on this new language.

## Experimental Results

- There is no guarantee that the techniques discussed will ever terminate, but in practice they often do.
- They are able to determine the minimal primes of the form $4 n+1$ and $4 n+3$ and the minimal primes expressed in the bases $b$ for $2 \leqslant b \leqslant 16$ and $b=18,20,22,23,24$, and 30 .


## Experimental Results

- There is no guarantee that the techniques discussed will ever terminate, but in practice they often do.
- They are able to determine the minimal primes of the form $4 n+1$ and $4 n+3$ and the minimal primes expressed in the bases $b$ for $2 \leqslant b \leqslant 16$ and $b=18,20,22,23,24$, and 30 .
- The bases $b=17,19,21$, and $25 \leqslant b \leqslant 29$ are solved with the exception of 37 families of the form $x y^{*} z$.


## Summary of Results for Bases up to 30

| Base | \# elements | Max. length | \# unsolved families |
| :---: | :---: | :---: | :---: |
| 2 | 2 | 2 |  |
| 3 | 3 | 3 |  |
| 4 | 3 | 2 |  |
| 5 | 8 | 5 |  |
| 6 | 7 | 5 |  |
| 7 | 9 | 5 |  |
| 8 | 15 | 9 |  |
| 9 | 12 | 4 |  |
| 10 | 26 | 8 |  |
| 11 | 152 | 45 |  |
| 12 | 17 | 8 |  |
| 13* | 228 | 32,021 |  |
| 14 | 240 | 86 |  |
| 15 | 100 | 107 |  |
| 16 | 483 | 3545 |  |
| 17* | $\geqslant 1279$ | $\geqslant 111,334$ | 1 |
| 18 | 50 | 33 |  |
| 19* | $\geqslant 3462$ | $\geqslant 110,986$ | 1 |
| 20 | 651 | 449 |  |
| $21^{*}$ | $\geqslant 2600$ | $\geqslant 479,150$ | 1 |
| 22 | 1242 | 764 |  |
| 23* | 6021 | 800,874 |  |
| 24 | 306 | 100 |  |
| $25^{*}$ | $\geqslant 17,597$ | $\geqslant 136,967$ | 12 |
| 26 | $\geqslant 5662$ | $\geqslant 8773$ | 2 |
| $27^{*}$ | $\geqslant 17,210$ | $\geqslant 109,006$ | 5 |
| 28* | $\geqslant 5783$ | $\geqslant 94,538$ | 1 |
| 29* | $\geqslant 57,283$ | $\geqslant 174,240$ | 14 |
| 30 | 220 | 1024 |  |

## Unsolved Families

| Base | Family | Algebraic form | Base | Family | Algebraic form |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 17 | F19* | $\left(5 \cdot 821 \cdot 17^{n}-3^{2}\right) / 16$ | 29 | $1 \mathrm{~A}^{*}$ | $\left(19 \cdot 29^{n}-5\right) / 14$ |
| 19 | EE16* | $\left(2^{2} \cdot 13 \cdot 307 \cdot 19^{n}-1\right) / 3$ |  | 68LO* 6 | $7 \cdot 757 \cdot 29^{n+1}+2 \cdot 3$ |
| 21 | G0*FK | $2^{4} \cdot 21^{n+2}+5 \cdot 67$ |  | AMP* | $\left(8761 \cdot 29^{n}-5^{2}\right) / 28$ |
| 25 | $6 \mathrm{MF}^{*} 9$ | $\left(1381 \cdot 25^{n+1}-53\right) / 8$ |  | $C^{*} \mathrm{FK}$ | $\left(3 \cdot 29^{n+2}+2 \cdot 331\right) / 7$ |
|  | CM1* | $\left(59 \cdot 131 \cdot 25^{n}-1\right) / 24$ |  | $\mathrm{F}^{*}$ OPF | $\left(3 \cdot 5 \cdot 29^{n+3}+139 \cdot 1583\right) / 28$ |
|  | EE1* | $\left(8737 \cdot 25^{n}-1\right) / 24$ |  | FKI* | $\left(6379 \cdot 29^{n}-3^{2}\right) / 14$ |
|  | E1*E | $\left(337 \cdot 25^{n+1}+311\right) / 24$ |  | $\mathrm{F}^{*}$ OP | $\left(3 \cdot 5 \cdot 29^{n+2}+7573\right) / 28$ |
|  | EFO* | $2 \cdot 3 \cdot 61 \cdot 25^{n}-1$ |  | LP09* | $\left(31 \cdot 16607 \cdot 29^{n}-3^{2}\right) / 28$ |
|  | F1*F1 | $\left(19^{2} \cdot 25^{n+2}+37 \cdot 227\right) / 24$ |  | OOPS* A | $2 \cdot 10453 \cdot 29^{n+1}-19$ |
|  | FO* KO | $3 \cdot 5 \cdot 25^{n+2}+2^{2} \cdot 131$ |  | PC* | $\left(2 \cdot 89 \cdot 29^{n}-3\right) / 7$ |
|  | FOK* 0 | $\left(5 \cdot 11 \cdot 41 \cdot 25^{n+1}+19\right) / 6$ |  | PPPL* 0 | $\left(87103 \cdot 29^{n+1}+3^{2}\right) / 4$ |
|  | LOL* 8 | $\left(53 \cdot 83 \cdot 25^{n+1}-3 \cdot 37\right) / 8$ |  | Q* GL | $\left(13 \cdot 29^{n+2}-3 \cdot 1381\right) / 14$ |
|  | M1*F1 | $\left(23^{2} \cdot 25^{n+2}+37 \cdot 227\right) / 24$ |  | Q* LO | $\left(13 \cdot 29^{n+2}-19 \cdot 109\right) / 14$ |
|  | M10* 8 | $19 \cdot 29 \cdot 25^{n+1}+2^{3}$ |  | RM* G | $\left(389 \cdot 29^{n+1}-5 \cdot 19\right) / 14$ |
|  | OL* 8 | $\left(199 \cdot 25^{n+1}-3 \cdot 37\right) / 8$ |  |  |  |
| 26 | A* 6 F | $\left(2 \cdot 26^{n+2}-7 \cdot 71\right) / 5$ |  |  |  |
|  | I* GL | $\left(2 \cdot 3^{2} \cdot 26^{n+2}-11 \cdot 113\right) / 25$ |  |  |  |
| 27 | 80*9A | $2^{3} \cdot 27^{n+2}+11 \cdot 23$ |  |  |  |
|  | 999G* | $\left(101 \cdot 877 \cdot 27^{n}-2^{3}\right) / 13$ |  |  |  |
|  | CL*E | $\left(3^{2} \cdot 37 \cdot 27^{n+1}-7 \cdot 29\right) / 26$ |  |  |  |
|  | EI*F8 | $\left(191 \cdot 27^{n+2}-2^{3} \cdot 149\right) / 13$ |  |  |  |
|  | $\mathrm{F}^{*} 9 \mathrm{FM}$ | (3.5.27 ${ }^{n+3}-113557$ )/26 |  |  |  |
| 28 | OA* F | $\left(2 \cdot 7 \cdot 47 \cdot 28^{n+1}+5^{3}\right) / 27$ |  |  |  |

