

# Searching for Kochen–Specker Systems With Orderly Generation and Satisfiability Solving

Curtis Bright

with Zhengyu Li and Vijay Ganesh

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# The Free Will Theorem

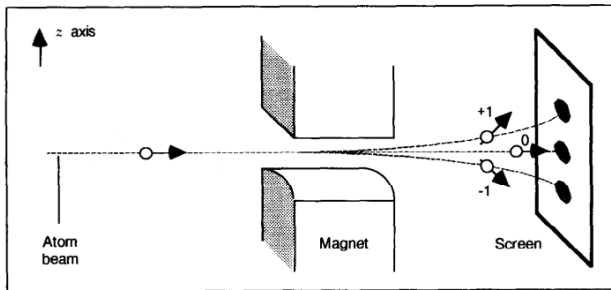
In 2006, John Conway and Simon Kochen proved the *Free Will Theorem*—if humans have free will then so do atoms.



Their proof crucially relies on a configuration of three dimensional vectors called a Kochen–Specker (KS) system.

# The Stern–Gerlach Experiment (1922)

Shoot an atom of orthohelium through a magnetic field:

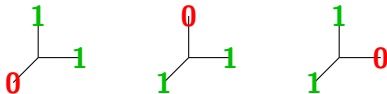


The *spin* of the atom (in the direction of the field) is  $+1$ ,  $-1$ , or  $0$ .

## The SPIN Axiom

Suppose the  $\pm 1$  beams are combined producing the *squared spin* which is either 1 or 0.

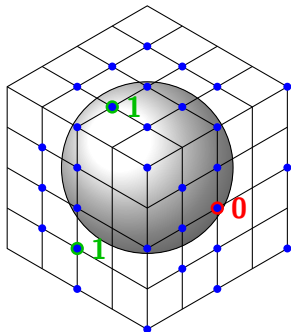
If you measure this in the  $x$ ,  $y$  and  $z$  axes it will be zero **in exactly one of these directions**.



The 101 conspiracy

# The KS Theorem

It is impossible to assign  $\{0, 1\}$  values to the following 31 vectors in a way that maintains the 101 conspiracy.



31 vector KS system of Conway and Kochen

The atom cannot have a predetermined spin in every direction!

# Can We Do Better Than 31?

The best known result is that at least 22 vectors are required.<sup>1</sup>

This was shown by translating a hypothetical 21-vector KS system into a 21-vertex graph and performing an exhaustive search.

There are a huge number of such graphs and the computation took 75 CPU years using the best graph enumeration algorithms.

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<sup>1</sup>S. Uijlen, B. Westerbaan. A Kochen-Specker System Has at Least 22 Vectors. *New Generation Computing*, 2016.

## Reduction to Satisfiability (SAT)

With some cleverness, many restrictive properties a “KS graph” must satisfy can be encoded in Boolean logic.

A SAT approach outperformed the previously used graph enumeration approach. However, a SAT solver generates many isomorphic copies of the same graph.

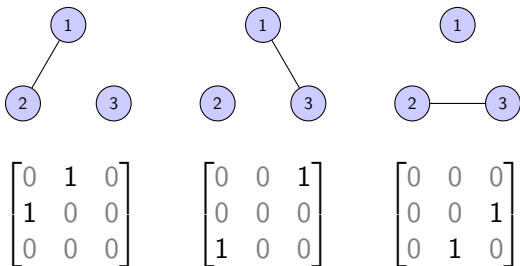
Thus, we combine SAT with isomorph-free exhaustive generation (also previously used to solve Lam’s problem).<sup>2</sup>

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<sup>2</sup>C. Bright, K. Cheung, B. Stevens, I. Kotsireas, V. Ganesh. A SAT-based Resolution of Lam’s Problem. *AAAI 2021*.

# Isomorphisms

When generating combinatorial objects we only care about generating them *up to isomorphism*. Unfortunately, objects usually have many isomorphic representations.



A graph with  $n$  vertices has up to  $n!$  distinct isomorphic adjacency matrices. For efficiency, these should be detected and removed.



# SAT Symmetry Breaking

A typical SAT approach is to add “symmetry breaking” constraints that remove as many isomorphic solutions as possible.

For example, you can order the rows of an adjacency matrix of a graph lexicographically.<sup>3</sup> However, many distinct isomorphic representations still exist, like

$$\begin{bmatrix} 0 & 0 & 1 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \end{bmatrix} \text{ and } \begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 0 \\ 1 & 0 & 0 \end{bmatrix}.$$

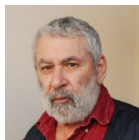
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<sup>3</sup>M. Codish, A. Miller, P. Prosser, P. Stuckey. Constraints for symmetry breaking in graph representation. *Constraints*, 2019.

# Isomorph-free Orderly Generation

Only “canonical” intermediate objects are recorded. The notion of canonicity is defined so that:

1. Every isomorphism class has exactly one canonical representative.
2. If an object is canonical then its parent in the search tree is also canonical.



Developed independently by Faradžev and Read in 1978.

## Canonicity Example

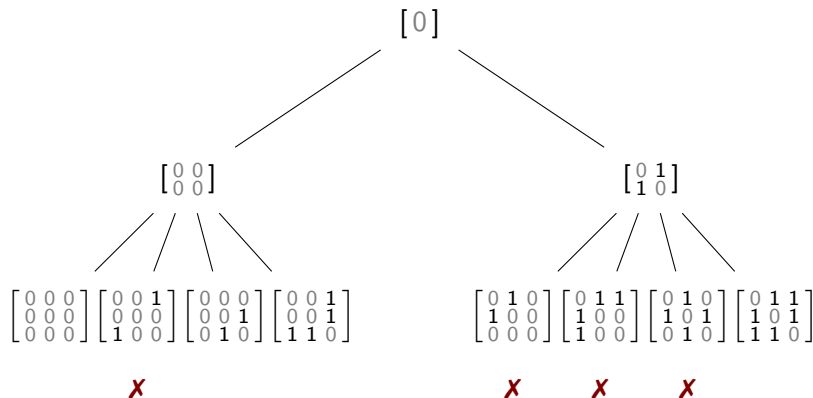
An adjacency matrix of a graph is *canonical* if the vector of its entries below the diagonal is lexicographically smallest (among all matrices in the same isomorphism class).

For example,

$$\begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix}, \begin{bmatrix} 0 & 0 & 1 \\ 0 & 0 & 0 \\ 1 & 0 & 0 \end{bmatrix}, \text{ and } \begin{bmatrix} 0 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$$

are isomorphic adjacency matrices but only the last is canonical.

## Orderly Generation of Graphs



Canonical testing introduces overhead, but every negative test prunes a large part of the search space.

## Orderly Generation in Practice

Each canonical test is independent, making the method easy to parallelize.

Verifying a matrix is *noncanonical* is often fast—it requires finding a single permutation of the vertices giving a lex-smaller matrix.

# SAT and Isomorph-free Generation

There have been surprisingly few attempts at combining isomorph-free generation and SAT solving.<sup>4,5</sup>

The “SAT modulo symmetry” paradigm also uses a canonicity test.<sup>6</sup>

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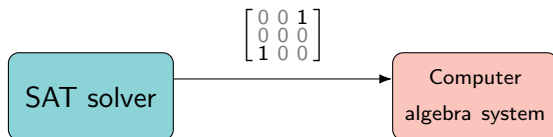
<sup>4</sup>T. Junttila, M. Karppa, P. Kaski, J. Kohonen. An adaptive prefix-assignment technique for symmetry reduction. *Journal of Symbolic Computation*, 2020.

<sup>5</sup>J. Savela, E. Oikarinen, M. Järvisalo. Finding periodic apartments via Boolean satisfiability and orderly generation. *LPAR 2020*.

<sup>6</sup>M. Kirchweger, S. Szeider. SAT Modulo Symmetries for Graph Generation. *CP 2021*.

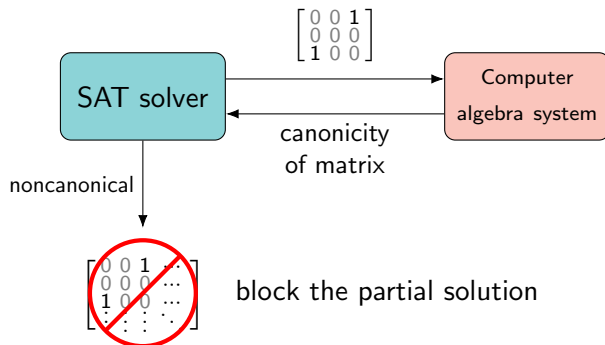
## Orderly Generation in SAT

During the search the SAT solver will find partial solutions (complete definitions for the edges in some subgraphs)...



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## KS Search Results

The speedup factor that we found when using SAT-based orderly generation in the search for KS systems of a given order:

order	speedup factor
16	6.5
17	13.6
18	37.8
19	104.5

The order 21 case was resolved in 25.7 CPU days (over 1000 times faster than the previous search).

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The order 22 case was resolved in 5.3 CPU years. No KS system was found, so a KS system *must have at least 23 directions*.

## A Promising Future

Isomorph-free generation and SAT can be combined to produce a hybrid solver capable of exponential speedups over a pure SAT or computer algebra approach.

The approach is very general and can be applied to many combinatorial generation problems. I believe it has yet to be used to its full potential.

Thank You!

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