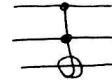


Lecture 12 <sup>2/25/10</sup> Universal set of fault-tolerant gate gadgets

① Fault-tolerant Toffoli [Shor, 9605011]

$$T|a, b, c\rangle = |a, b, c \oplus ab\rangle$$

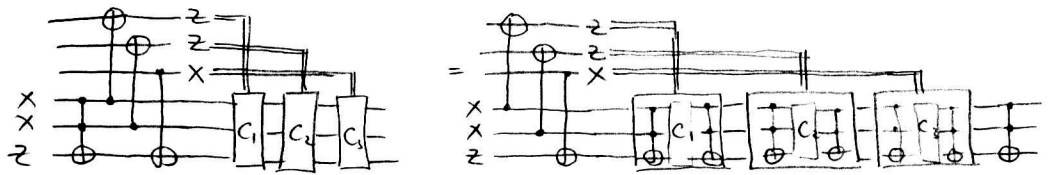


- construction using computation + teleportation (using adaptive classical control)

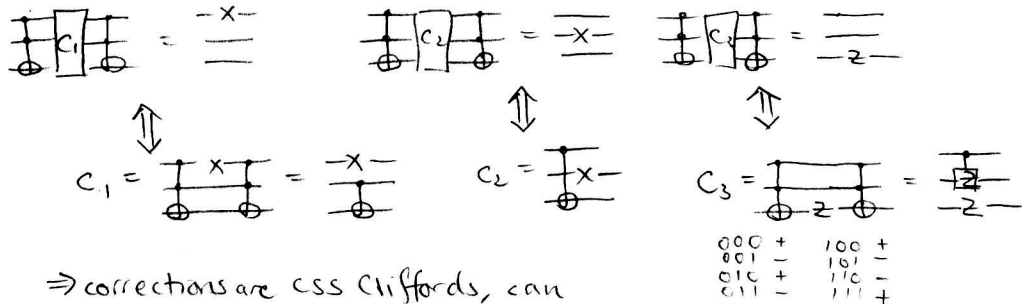
1. prepare the ancilla  $\sum_{a,b} T|a, b, 0\rangle = \sum_{a,b} |a, b, a \cdot b\rangle = \begin{matrix} X \\ X \\ Z \end{matrix} \begin{matrix} \oplus \\ \oplus \\ \oplus \end{matrix}$

2. teleport into it

one-qubit teleportation:  $\begin{matrix} X \\ Z \end{matrix} \begin{matrix} \oplus \\ \oplus \end{matrix} = \begin{matrix} X \\ Z \end{matrix} \begin{matrix} \oplus \\ \oplus \end{matrix} = \dots$



thus corrections should satisfy



⇒ corrections are CSS Cliffords, can be applied fault tolerantly

To prepare the encoded ancilla fault tolerantly:

note  $|+\rangle_1 |+\rangle_2 |+\rangle_3 |+\rangle_4 = |+\rangle \otimes (|A\rangle + |B\rangle)$

$$|A\rangle = \frac{1}{2}(|000\rangle + |010\rangle + |100\rangle + |111\rangle)$$

$$|B\rangle = X_3 |A\rangle$$

$$\begin{matrix} \Lambda_1(\Lambda_2(z_3)) \\ \Lambda_1(z_4) \end{matrix} |0\rangle \otimes (|A\rangle + |B\rangle)_{2,3,4} + |1\rangle \otimes \Lambda_2(z_3) z_4 (|A\rangle + |B\rangle)_{2,3,4} = |+\rangle \otimes |A\rangle + |-\rangle \otimes |B\rangle$$

↳ measure & correct repeatedly

assume self-dual CSS code with  $\bar{X} = X^{\otimes n}$ ,  $\bar{Z} = Z^{\otimes n}$

⇒  $\Lambda(\bar{Z})$ ,  $\bar{Z}$  are both transversal

⇒ encode qubits 2,3,4 into the code

use a verified cat state  $|0^n\rangle + |1^n\rangle$  for the first qubit

note:  $\Lambda_1(\Lambda_2(z_3))$  uses a Toffoli

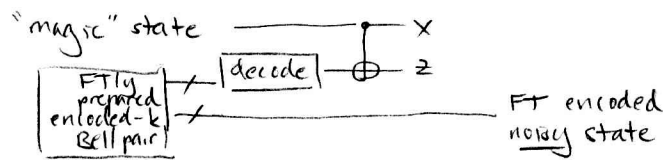
A similar technique will work for any  $C_3$  gate, e.g.,  $\begin{pmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{pmatrix}$ .  
 The corresponding state can be teleported into, requiring only Clifford corrections.  
 The ancilla is an eigenstate of a Clifford gate, can measure the eigenvalue using controlled Cliffords from a verified cat state.

② Magic states distillation [Knill Laflamme Zurek 9610011, Bravyi Kitaev 0403025]  
 so far: Clifford gates + prep. of certain nonstabilizer states  $\Rightarrow$  universality  
 (with classical control)

universality at level 0  $\Rightarrow$  FT universality at level 1

$\Rightarrow \dots \Rightarrow$  FT universality at level  $k$

new idea: Clifford gates + prep. of certain noisy nonstabilizer states  $\Rightarrow$  universality



Advantage: Through the highest level of encoding, all operations are Cliffords<sup>(CSS)</sup>  
 $\rightarrow$  easy to simulate (with Pauli noise channels)

$\rightarrow$  one less exec to count malignant sets within, simplified analysis

Disadvantage: Typically higher overhead (necessarity?)

Note: If Cliffords + noisy magic state prep.  $\Rightarrow$  universality, c.c.  
 then  $\Rightarrow$  prep. of noiseless magic states

$\therefore$  Suffices to study distillation problem:

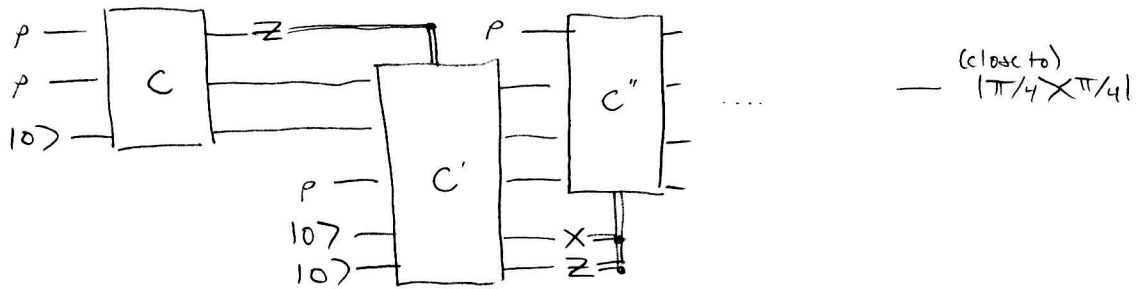
From which noisy states can we asymptotically prepare a magic state,

$$\text{eg. } |A\rangle = \sum_{a,b=0}^1 |a,b,a,b\rangle$$

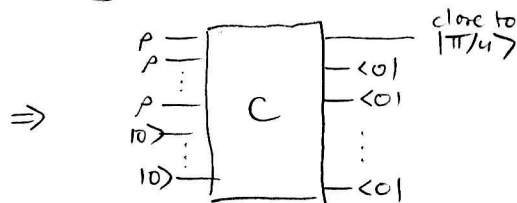
$$\text{or } |\pi/4\rangle = \frac{1}{\sqrt{2}}(|0\rangle + e^{i\pi/4}|1\rangle)$$

using adaptive classical control of stabilizer operations?

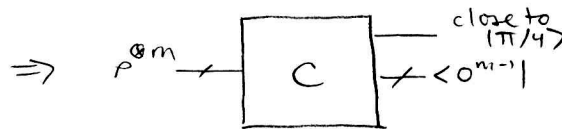
Model. Can repeatedly prepare identical copies of  $\rho$ , a single-qubit state.



Claim 1: May assume all measurements are postselected +1.



Claim 2: May assume no stabilizer preparations used.

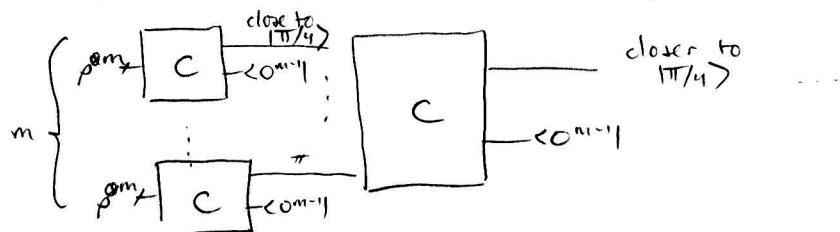


(Corollary of Gottesman-Knill, any  $m$  qubits and  $n$   $|0\rangle$  ancillas can be implemented with  $k$  ancilla qubits,  $l = m + n - k$  extra ancilla qubits,  $\max\{0, k - m\}$ )

$\Rightarrow C'$  encodes into a stabilizer code

stabilizer codes  $\Leftrightarrow$  distillation schemes

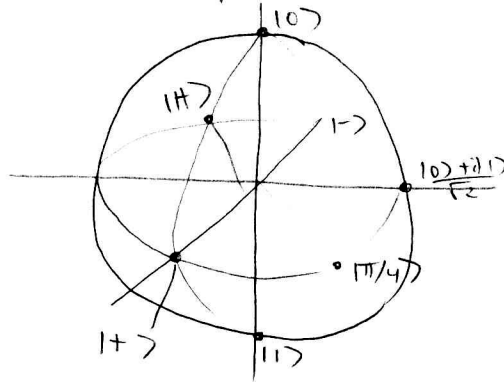
code concatenation  $\Leftrightarrow$  recursive distillation



Good codes can be found by trial and error (no general theory).

Note: Need overhead to be polynomial, preferably polylogarithmic.

Bloch sphere



1. symmetrize  $\rho = \rho(x, y, z)$   
 $= \frac{1}{2}(I + xX + yY + zZ)$   
 about  $x=z$  axis by applying  $I$  or  $H$   
 with 50/50 prob. at random  
 $\Rightarrow \rho' = \frac{1}{2}(\rho + H\rho H)$   
 $= \frac{1}{2}\left(\frac{1}{2}(1+xX+yY+zZ) + \frac{1}{2}(1+xZ-yY+zX)\right)$   
 $= \frac{1}{2}\left(1 + \frac{x+z}{2}X + \frac{x+z}{2}Z\right)$   
 $= \rho\left(\frac{x+z}{2}, 0, \frac{x+z}{2}\right)$   
 $= \frac{1}{2}\begin{pmatrix} 1 + \frac{x+z}{2} & \frac{x+z}{2} \\ \frac{x+z}{2} & 1 - \frac{x+z}{2} \end{pmatrix} = \frac{1}{2}\begin{pmatrix} 1+x' & x' \\ x' & 1-x' \end{pmatrix}$

2. Decode the code

$$\Rightarrow \rho_{out} = \begin{pmatrix} \langle 0 | \rho^{otm} | 0 \rangle & \langle 0 | \rho^{otm} | 1 \rangle \\ \langle 1 | \rho^{otm} | 0 \rangle & \langle 1 | \rho^{otm} | 1 \rangle \end{pmatrix} / \text{trace}$$

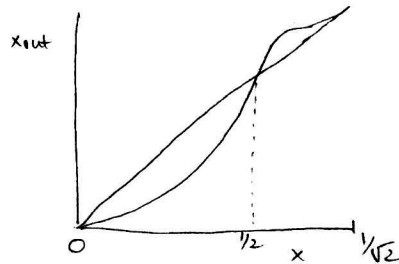
Eg. for the Steane  $[[7, 1, 3]]$  code,

$$|0\rangle = \frac{1}{\sqrt{8}} \left( |0^7\rangle + |0001111\rangle + |0110011\rangle + |1010101\rangle + |1011100\rangle + |11011010\rangle + |1100110\rangle + |1101001\rangle \right)$$

$$\Rightarrow \langle 0 | \rho^{otm} | 0 \rangle = \frac{1}{8} (p_{00}^7 + 7p_{00}^3(p_{01}^4 + p_{10}^4) + 7p_{00}^3 p_{11}^4 + 7(p_{00} p_{01}^4 p_{11}^2))$$

similarly for the other terms, since  $|1\rangle = X^{07}|0\rangle$

$$\Rightarrow \rho_{out} = \rho(x_{out}, 0, x_{out}), \quad x_{out} = \frac{x^2(7 + 8x^4)}{1 + 14x^4}$$



Note: No  $\rho$  inside octahedron can be distilled, convex combination of stabilizer states.

Application to  $v = 5/2$  anym FQHE.

Open problems: Distillable region, asymptotic distillation rate,  $n$ -qubit states?

(Tight) tolerable noise threshold upper bounds

based on Gottesman-Knill theorem, magic states distillation

Vidhani Huelga Plenio 0408076

Buhrman, Cleve, Laurent, Lindur, Schrijver, Unger '06

R?

Howard, van Dam 0907.3189

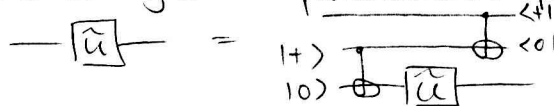
Model: Assume perfect classical control of perfect stabilizer operations (Clifford gates, prep, meas.), and a noisy non-Clifford gate.

Example:  $E_{\text{depolarizing}}(p) = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{pmatrix}$   
randomize qubit w/ prob  $4p/3$

Claim: If  $p \geq \frac{3}{4} \cdot \frac{6-2\sqrt{2}}{7} \approx 34\%$ , then computation is classically simulatable.

Proof:

1. May assume  $\tilde{u}$  only used in postselected teleportation



2. If  $(I \otimes \tilde{u})(|00\rangle + |11\rangle)$  is a (known) mixture of stabilizer states,

$\uparrow$   
 Janczowski isomorphism of  $\tilde{u}$

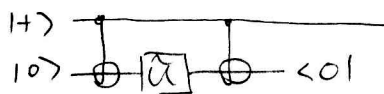
then Gottesman-Knill gives efficient classical simulation.

$$= (1 - \frac{4}{3}p) \frac{1}{2} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ e^{i\pi/4} & 0 & 0 & 1 \end{pmatrix} + \frac{4}{3}p \cdot \frac{1}{4} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

$$= \frac{3}{4} \frac{6-2\sqrt{2}}{7} \left( \frac{4+\sqrt{2}}{14} \right) J(I) + \left( \frac{4+\sqrt{2}}{14} \right) J \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix} + \left( \frac{3-\sqrt{2}}{14} \right) J(X) + \left( \frac{3-\sqrt{2}}{14} \right) J(Y) \quad \square$$

Claim: If  $p < \frac{3}{4} \frac{6-2\sqrt{2}}{7}$ , then the model allows for efficient qv. universality

Proof: Applying



$$\text{gives } \frac{1}{2} \begin{pmatrix} 1 & e^{-i\pi/4} \frac{3-4p}{3-2p} \\ e^{i\pi/4} \frac{3-4p}{3-2p} & 1 \end{pmatrix}$$

$$= p(x, x, 0), \quad x = \frac{1}{\sqrt{2}} \frac{3-4p}{3-2p}$$

$$= \frac{1}{2} \quad \text{at } p = \frac{3}{4} \frac{6-2\sqrt{2}}{7} \quad \checkmark \quad \square$$