

Review: MIPS Register Usage

See also: `kern/arch/mips/include/asmdefs.h`

R0, zero = ## zero (always returns 0)

R1, at = ## (assembler temporary)

R2, v0 = ## return value / system call number

R3, v1 = ## return value

R4, a0 = ## 1st argument (to subroutine)

R5, a1 = ## 2nd argument

R6, a2 = ## 3rd argument

R7, a3 = ## 4th argument

Review: MIPS Register Usage

R08-R15, t0-t7 = ## temp variables
R24-R25, t8-t9 = ## temp variables
 ## (use without saving)
R16-R23, s0-s7 = ## subroutine variables
 ## (saved/restored by subroutines)
R26-27, k0-k1 = ## reserved for interrupt handler
R28, gp = ## global pointer
 ## (for easy access to some variables)
R29, sp = ## stack pointer
R30, s8/fp = ## 9th subroutine reg / frame pointer
R31, ra = ## return addr (used by jal)

How a System Call Works

```
## In OS/161 all user programs are linked with  
## lib/crt0/mips-crt0.S which contains __start, the starting  
## point for OS/161 programs. See lib/crt0/mips-crt0.S  
## for comments on what this code is doing.
```

```
004000b0 <__start>:
```

```
4000b0: 3c1c1000    lui  gp,0x1000  
4000b4: 279c7fff0   addiu gp,gp,32752  
4000b8: 3c08fffff   lui  t0,0xfffff  
4000bc: 3508ffff8   ori  t0,t0,0xffff8  
4000c0: 03a8e824    and  sp,sp,t0  
4000c4: 27bdfbff0   addiu sp,sp,-16  
4000c8: 3c011000    lui  at,0x1000  
4000cc: ac250004    sw   a1,4(at)
```

How a System Call Works

```
4000d0: 0c100040    jal 400100 <main>    # Call main
4000d4: 00000000    nop
4000d8: 00408021    move    s0,v0
4000dc: 0c100050    jal 400140 <exit>
4000e0: 02002021    move    a0,s0
4000e4: 0c100069    jal 4001a4 <_exit>
4000e8: 02002021    move    a0,s0
4000ec: 02002021    move    a0,s0
4000f0: 24020000    li      v0,0
4000f4: 0000000c    syscall
4000f8: 0810003b    j 4000ec <__start+0x3c>
4000fc: 00000000    nop
```

How a System Call Works

```
/* See how a function/system call happens. */
#include <unistd.h>
#include <errno.h>

int
main()
{
    int x;
    int y;
    x = close(999);
    y = errno;

    return x;
}
```

How a System Call Works

```
% cs350-objdump -d syscall > syscall.out
```

```
% cat syscall.out
```

```
00400100 <main>:
```

400100:	27bdf fe0	addiu	sp, sp, -32
400104:	afb f001c	sw	ra, 28(sp)
400108:	afbe 0018	sw	s8, 24(sp)
40010c:	03a0 f021	move	s8, sp
400110:	0c10 007b	jal	4001ec <close>
400114:	2404 03e7	li	a0, 999
400118:	afc2 0010	sw	v0, 16(s8)
40011c:	3c02 1000	lui	v0, 0x1000
400120:	8c42 0000	lw	v0, 0(v0)
400124:	0000 0000	nop	

How a System Call Works

<main> continued

400128:	afc20014	sw	v0, 20(s8)
40012c:	8fc20010	lw	v0, 16(s8)
400130:	03c0e821	move	sp, s8
400134:	8fbf001c	lw	ra, 28(sp)
400138:	8fbe0018	lw	s8, 24(sp)
40013c:	03e00008	jr	ra
400140:	27bd0020	addiu	sp, sp, 32

How a System Call Works

```
## See lib/libc/syscalls.S for details/comments */  
## At bit easier to understand from disassembled code.
```

```
% cs350-objdump -d syscall > syscall.S
```

```
% cat syscall.S
```

```
...
```

```
0040022c <close>:
```

```
    40022c: 08100074    j 4001d0 <__syscall>
```

```
    400230: 24020007    li v0,7
```

```
00400234 <reboot>:
```

```
    400234: 08100074    j 4001d0 <__syscall>
```

```
    400238: 24020008    li v0,8
```

```
...
```

How a System Call Works

From lib/libc/syscalls.S

```
.set noreorder
.text
.type __syscall,@function
.ent __syscall

__syscall:
    syscall          /* make system call */
    beq a3, $0, 1f    /* if a3 is zero, call succeeded */
    nop              /* delay slot */
    sw v0, errno      /* call failed: store errno */
    li v1, -1         /* and force return value to -1 */
    li v0, -1
1:
    j ra             /* return */
    nop             /* delay slot */
.end __syscall
.set reorder
```

How a System Call Works

From `lib/libc/syscalls.S`

`SYSCALL(close, 7)`

`SYSCALL(reboot, 8)`

`SYSCALL(sync, 9)`

`SYSCALL(sbrk, 10)`

`SYSCALL(getpid, 11)`

How a System Call Works

syscall instruction generates an exception.

Processor begins execution at virtual address 0x8000 0080

From: kern/arch/mips/mips/exception.S

Q: where does this address live? A: kseg0

exception:

```
    move k1, sp          /* Save previous stack pointer in k1 */
    mfc0 k0, c0_status    /* Get status register */
    andi k0, k0, CST_KUp  /* Check the we-were-in-user-mode bit */
    beq k0, $0, 1f        /* If clear, from kernel, already have stack
    nop                   /* delay slot */
```

```
/* Coming from user mode - load kernel stack into sp */
    la k0, curkstack      /* get address of "curkstack" */
    lw sp, 0(k0)          /* get its value */
    nop                   /* delay slot for the load */
```

```
1:
    mfc0 k0, c0_cause      /* Now, load the exception cause. */
    j common_exception    /* Skip to common code */
    nop                   /* delay slot */
```

How a System Call Works

From: kern/arch/mips/mips/exception.S

common_exception:

- o saves the contents of the registers
- o calls mips_trap (C code in kern/arch/mips/mips/trap.c)
- o restores the contents of the saved registers
- o rfe (return from exception)

From: kern/arch/mips/mips/trap.c

mips_trap:

- o figures out the exception type/cause
- o calls the appropriate handing function
(for system call this is mips_syscall).

How a System Call Works

From: kern/arch/mips/mips/syscall.c

```
mips_syscall(struct trapframe *tf)
{
    assert(curspl==0);
    callno = tf->tf_v0; retval = 0;

    switch (callno) {
        case SYS_reboot:
            /* is in kern/main/main.c */
            err = sys_reboot(tf->tf_a0);
            break;

            /* Add stuff here */

        default:
            kprintf("Unknown syscall %d\n", callno);
            err = ENOSYS;
            break;
    }
```

How a System Call Works

```
if (err) {
    tf->tf_v0 = err;
    tf->tf_a3 = 1;          /* signal an error */
} else {
    /* Success. */
    tf->tf_v0 = retval;
    tf->tf_a3 = 0;          /* signal no error */
}

/* Advance the PC, to avoid the syscall again. */
tf->tf_epc += 4;

/* Make sure the syscall code didn't forget to lower spl */
assert(curspl==0);
}
```

C Code for Sections Example

```
#include <unistd.h>
```

```
#define N    (200)
```

```
int x = 0xdeadbeef;
```

```
int y1;
```

```
int y2;
```

```
int y3;
```

```
int array[4096];
```

```
char const *str = "Hello World\n";
```

```
const int z = 0xabcddcba;
```

```
struct example {
```

```
    int ypos;
```

```
    int xpos;
```

```
};
```

C Code for Sections Example (cont'd)

```
int
main()
{
    int count = 0;
    const int value = 1;
    y1 = N;
    y2 = 2;
    count = x + y1;
    y2 = z + y2 + value;

    reboot(RB_POWEROFF);
    return 0; /* avoid compiler warnings */
}
```

cs350-readelf Output: Sections and Segments

```
% cs350-readelf -a segments > readelf.out
```

```
% cat readelf.out
```

Section Headers:

[Nr]	Name	Type	Addr	Off	Size	ES	Flg	...	Al
[0]		NULL	00000000	000000	000000	00		...	0
[1]	.reginfo	MIPS_REGINFO	00400094	000094	000018	18	A	...	4
[2]	.text	PROGBITS	004000b0	0000b0	000250	00	AX	...	16
[3]	.rodata	PROGBITS	00400300	000300	000020	00	A	...	16
[4]	.data	PROGBITS	10000000	001000	000010	00	WA	...	16
[5]	.sbss	NOBITS	10000010	001010	000014	00	WAp	...	4
[6]	.bss	NOBITS	10000030	00101c	004000	00	WA	...	16
[7]	.comment	PROGBITS	00000000	00101c	000036	00		...	1

...

Key to Flags:

W (write), A (alloc), X (execute), M (merge), S (strings)

I (info), L (link order), G (group), x (unknown)

0 (extra OS processing required) o (OS specific), p (processor s

Size = 0x250 = 592 bytes

Off = offset into the ELF file

Addr = virtual address

cs350-readelf Output: Sections and Segments

Program Headers:

Type	Offset	VirtAddr	PhysAddr	FileSiz	MemSiz	Flg	Align
REGINFO	0x000094	0x00400094	0x00400094	0x00018	0x00018	R	0x4
LOAD	0x000000	0x00400000	0x00400000	0x00320	0x00320	R E	0x1000
LOAD	0x001000	0x10000000	0x10000000	0x00010	0x04030	RW	0x1000

Section to Segment mapping:

Segment Sections...

00	.reginfo
01	.reginfo .text .rodata
02	.data .sbss .bss

.reginfo = register info (not used)
.text = code/machine instructions
.rodata = read-only data (e.g., string literals, consts)
.data = data (i.e., global variables)
.bss = Block Started by Symbol
has a symbol but no value (i.e, uninitialized data)
.sbss = ?small bss?

cs350-objdump -s output: Sections and Segments

```
% cs350-objdump -s segments > objdump.out
% cat objdump.out
...
```

Contents of section .text:

```
4000b0 3c1c1001 279c8000 3c08ffff 3508ffff <...'...<...5...
...
```

Decoding 3c1c1001 to determine instruction

0x3c1c1001 = binary 11110000001110000001000000000001

0011 1100 0001 1100 0001 0000 0000 0001

001111 | 00000 | 11100 | 0001 0000 0000 0001

instr | rs | rt | immediate

LUI | 0 | reg 28 | 0x1001

LUI | unused | reg 28 | 0x1001

Load unsigned immediate into rt (register target)

lui gp, 0x1001

cs350-objdump -s output: Sections and Segments

Contents of section .rodata:

```
400300 48656c6c 6f20576f 726c640a 00000000 Hello World.....
400310 abcddcba 00000000 00000000 00000000 .....
...
## 0x48 = 'H' 0x65 = 'e' 0x0a = '\n' 0x00 = '\0'
## Align next int to 4 byte boundary
## const int z = 0xabcddcba
## If compiler doesn't prevent z from being written/hardware could
## Size = 0x20 = 32 bytes "Hello World\n\0" = 13 + 3 padding = 16
##           + const int z = 4 = 20
## Then align to the next 16 byte boundry at 32 bytes.
```

cs350-objdump -s output: Sections and Segments

Contents of section .data:

```
10000000 deadbeef 00400300 00000000 00000000 .....@.....  
...  
## Size = 0x10 bytes = 16 bytes  
## int x = deadbeef (4 bytes)  
## char const *str = "Hello World\n"; (4 bytes)  
## value stored in str = 0x00400300.  
## NOTE: that that is the address of the start  
## of the .rodata section (i.e., address of 'H').
```

cs350-objdump -s output: Sections and Segments

```
% cs350-nm -n segments > nm.out
```

```
% cat nm.out
```

```
...
```

```
10000010 A __bss_start
```

```
10000010 A _edata
```

```
10000010 A _fbss
```

```
10000010 S y3
```

```
10000014 S y2
```

```
10000018 S y1
```

```
1000001c S errno
```

```
10000020 S __argv
```

```
...
```

```
## .bss Size = 0x4000 = 16384 = 4096 * sizeof(int)
```

```
10000030 B array
```

```
10004030 A __end
```