

Real Number Proving in PVS

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Why Real Number Proving in PVS ?

- ▶ There are more numbers than integers in *real* life (despite what model-checkers are telling you).
- ▶ Conceptually, it is easier to reason on a continuous framework than on a discrete one.
- ▶ A lot of classical results in calculus, trigonometry, and continuous mathematics.
- ▶ Sometimes you cannot avoid them: hybrid systems, engineering applications, etc.

I Use a CAS, Why Should I Bother with PVS?

Computer Algebra Systems (CAS):

- ▶ Mathematica, Maple, Matlab, Scilab, . . . offer very powerful symbolic and numerical engines.
- ▶ CAS do not aim *soundness*. Singularities and exceptions are well-known problems of CAS.
- ▶ CAS do not support specification languages but programming languages.

CAS vs. Theorem Provers

- ▶ Real analysis is not a traditional strength of theorem provers.
- ▶ Theorem provers and CAS can be integrated in useful ways:
 - ▶ Computer algebra systems can be used to perform mechanical simplifications and find potential solutions.
 - ▶ Theorem prover are then used to verify the correctness of a particular solution.

I. Real Numbers in PVS

- ▶ Reals are defined as an uninterpreted subtype of number in the prelude library:

```
real: TYPE+ FROM number
```

- ▶ All numeric constants are real:
 - ▶ naturals: $0, 1, \dots$
 - ▶ integers: $\dots, -1, 0, 1, \dots$
 - ▶ rationals: $\dots, -1/10, \dots, 3/2, \dots$
- ▶ Decimal notation is supported in PVS 4: The decimal number **3.141516** is syntactic sugar for the rational number $31416/10000$.

PVS's real numbers are \mathbb{R}

(Rather than floating point numbers)

- ▶ All the **standard properties**: infinite, non-enumerable, $\mathbb{N} \subseteq \mathbb{Z} \subseteq \mathbb{Q} \subseteq \mathbb{R}, \dots$
- ▶ **Exact** arithmetic: $1/3 + 1/3 + 1/3 = 1$.
- ▶ The type real is **unbounded**:

```
googol      : real = 10^100  
googolplex : real = 10^googol
```

```
googol_prop : LEMMA  
  googolplex > googol * googol
```

PVS's real is Built-in

- ▶ Numerical expressions can be **automatically** reduced by the theorem prover (no need to prove $1+1=2$), ...
- ▶ ...except for *machine physical limitations*: Try to prove `googol_prop` with **grind**.

You can still prove `googol_prop` using analytical methods.

Subtypes of real

```
nzreal  : TYPE+ = {r:real | r /= 0} % Nonzero reals
nnreal  : TYPE+ = {r:real | r >= 0} % Nonnegative reals
npreal  : TYPE+ = {r:real | r <= 0} % Nonpositive reals
negreal : TYPE+ = {r:real | r < 0} % Negative reals
posreal : TYPE+ = {r:real | r > 0} % Positive reals

rat      : TYPE+ FROM real
int      : TYPE+ FROM rat
nat      : TYPE+ FROM int
```

The uninterpreted type `number` is the only `real`'s supertype predefined in PVS: no complex numbers, no hyper-reals, no \mathbb{R}^∞ , ...

Predefined Operations

`+, -, *: [real, real -> real]`

`/: [real, nzreal -> real]`

`-: [real -> real]`

`sgn(x:real) : int = IF x >= 0 THEN 1 ELSE -1 ENDIF`

`abs(x:real) : {nny: nreal | nny >= x} = ...`

`max(x,y:real): {z: real | z >= x AND z >= y} = ...`

`min(x,y:real): {z: real | z <= x AND z <= y} = ...`

`^(x: real,i:{i:int | x /= 0 OR i >= 0}): real = ...`

...and what about $\sqrt{\quad}$, \int , \log , \exp , \sin , \cos , \tan , π , \lim , ... ?

NASA PVS Libraries

Among many others, the following libraries are available at <http://shemesh.larc.nasa.gov/ftp/larc/PVS-library/pvslib.html>

- ▶ `reals`: Square, square root, quadratic formula, polynomials.
- ▶ `analysis`: Real analysis, limits, continuity, derivatives, integrals.
- ▶ `series`: Power series, Taylor's theorem.
- ▶ `lnexp` and `lnexp_fnd`: Logarithm, exponential, and hyperbolic functions.
- ▶ `trig` and `trig_fnd`: Trigonometry.
- ▶ `complex`: Complex numbers.
- ▶ `float`: Floating point numbers.

To Be Or Not To Be (Foundational) ?

- ▶ Axiomatic theories `trig` and `lnexp` typecheck faster.
- ▶ Foundational theories `trig_fnd` and `lnexp_fnd` have **no** axioms, and are updated regularly.
- ▶ Careful what you wish for:

```
|-----  
{1}  sin(pi / 2) > 1 / 2
```

Rule? (`grind`)

```
Integral rewrites Integral[real](0, 1, atan_deriv_fn)  
  to integral(0, 1, atan_deriv_fn)  
atan_value rewrites atan_value(1)  
  to integral(0, 1, atan_deriv_fn)  
atan rewrites atan(1)  
  to integral(0, 1, atan_deriv_fn)  
pi rewrites pi  
  to 4 * integral(0, 1, atan_deriv_fn)  
sin_value rewrites sin_value  
  to ...
```

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  to integral(0, 1, atan_deriv_fn)  
pi rewrites pi  
  to 4 * integral(0, 1, atan_deriv_fn)  
sin_value rewrites sin_value  
  to ...
```

II. Low Level Real Number Proving

Real numbers in PVS are axiomatically defined in the PVS prelude:

- ▶ Theory `real_axioms`:
Commutativity, associativity, identity, etc. These properties are known to the decision procedures, so they rarely need to be cited.

- ▶ Theory `real_props`:
Order and cancellation laws. These lemmas are **not** used automatically by the standard decision procedures.

If You Really Want to Know ...

```
real_props: THEORY
BEGIN
  both_sides_plus_le1: LEMMA  $x + z \leq y + z$  IFF  $x \leq y$ 
  both_sides_plus_le2: LEMMA  $z + x \leq z + y$  IFF  $x \leq y$ 
  both_sides_minus_le1: LEMMA  $x - z \leq y - z$  IFF  $x \leq y$ 
  both_sides_minus_le2: LEMMA  $z - x \leq z - y$  IFF  $y \leq x$ 
  both_sides_div_pos_le1: LEMMA  $x/pz \leq y/pz$  IFF  $x \leq y$ 
  both_sides_div_neg_le1: LEMMA  $x/nz \leq y/nz$  IFF  $y \leq x$ 
  ...
  abs_mult: LEMMA  $\text{abs}(x * y) = \text{abs}(x) * \text{abs}(y)$ 
  abs_div: LEMMA  $\text{abs}(x / n0y) = \text{abs}(x) / \text{abs}(n0y)$ 
  abs_abs: LEMMA  $\text{abs}(\text{abs}(x)) = \text{abs}(x)$ 
  abs_square: LEMMA  $\text{abs}(x * x) = x * x$ 
  abs_limits: LEMMA  $-(\text{abs}(x) + \text{abs}(y)) \leq x + y$  AND
                $x + y \leq \text{abs}(x) + \text{abs}(y)$ 
END real_props
```

Tip 1: Avoid real_props

|-----
{1} nnx / (nnx + 1) <= 1

Rule? (grind)

|-----
{1} nnx / (1 + nnx) <= 1

Rule? (grind :theories "real_props")

div_mult_pos_le1 rewrites nnx / (1 + nnx) <= 1
to TRUE

Q.E.D.

A Toy Example

toy :

|-----
{1} $x * (1 - x) \leq 1$

Rule? (grind :theories "real_props")

toy :

|-----
{1} $x - x * x \leq 1$

A Toy Example

toy :

|-----
{1} $x * (1 - x) \leq 1$

Rule? (grind :theories "real_props")

toy :

|-----
{1} $x - x * x \leq 1$

Tip 2: Use both-sides To Operate Both Sides of a Formula

(But **only** to **add/subtract** both sides of a formula)

toy :

$$\{1\} \quad x - x * x \leq 1$$

Rule? (both-sides "-" "1/4")

Applying $- 1 / 4$ to both sides of an inequality/equality conjunction, this simplifies to:

toy :

$$\{1\} \quad x - x * x - 1 / 4 \leq 1 - 1 / 4$$

Rule?

Tip 2: Use both-sides To Operate Both Sides of a Formula

(But **only** to **add/subtract** both sides of a formula)

toy :

$$\{1\} \quad x - x * x \leq 1$$

Rule? (both-sides "-" "1/4")

Applying $- 1 / 4$ to both sides of an inequality/equality conjunction, this simplifies to:

toy :

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Rule?

Tip 3: Use case to Prove What You Want ...

(No what PVS offers you)

Rule? (case "x - x * x - 1 / 4 <= 0")

this yields 2 subgoals:

toy.1 :

{-1} x - x * x - 1 / 4 <= 0

|-----

[1] x - x * x - 1 / 4 <= 1 - 1 / 4

Rule? (assert)

This completes the proof of toy.1.

toy.2 :

|-----

{1} x - x * x - 1 / 4 <= 0

[2] x - x * x - 1 / 4 <= 1 - 1 / 4

Rule? (hide 2)

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(No what PVS offers you)

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toy.1 :

{-1} x - x * x - 1 / 4 <= 0

|-----

[1] x - x * x - 1 / 4 <= 1 - 1 / 4

Rule? (assert)

This completes the proof of toy.1.

toy.2 :

|-----

{1} x - x * x - 1 / 4 <= 0

[2] x - x * x - 1 / 4 <= 1 - 1 / 4

Rule? (hide 2)

Tip 3: Use case to Prove What You Want ...

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Rule? (case "x - x * x - 1 / 4 <= 0")

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toy.1 :

{-1} x - x * x - 1 / 4 <= 0

|-----

[1] x - x * x - 1 / 4 <= 1 - 1 / 4

Rule? (assert)

This completes the proof of toy.1.

toy.2 :

|-----

{1} x - x * x - 1 / 4 <= 0

[2] x - x * x - 1 / 4 <= 1 - 1 / 4

Rule? (hide 2)

... And Arrange Expressions With case-replace

toy.2 :

|-----

[1] $x - x * x - 1 / 4 \leq 0$

Rule? (case-replace

" $x - x * x - 1 / 4 = -(x-1/2)*(x-1/2)$ "

:hide? t)

this yields 2 subgoals:

toy.2.1 :

|-----

{1} $-(x - 1 / 2) * (x - 1 / 2) \leq 0$

Rule? (grind :theories "real_props")

this simplifies to:

toy.2.1 :

|-----

{1} $-(x - 1 / 2) * x - (1 * -(x - 1 / 2)) / 2 \leq 0$

... And Arrange Expressions With case-replace

toy.2 :

|-----

[1] $x - x * x - 1 / 4 \leq 0$

Rule? (case-replace

" $x - x * x - 1 / 4 = -(x-1/2)*(x-1/2)$ "

:hide? t)

this yields 2 subgoals:

toy.2.1 :

|-----

{1} $-(x - 1 / 2) * (x - 1 / 2) \leq 0$

Rule? (grind :theories "real_props")

this simplifies to:

toy.2.1 :

|-----

{1} $-(x - 1 / 2) * x - (1 * -(x - 1 / 2)) / 2 \leq 0$

... And Arrange Expressions With case-replace

toy.2 :

|-----

[1] $x - x * x - 1 / 4 \leq 0$

Rule? (case-replace

" $x - x * x - 1 / 4 = -(x-1/2)*(x-1/2)$ "

:hide? t)

this yields 2 subgoals:

toy.2.1 :

|-----

{1} $-(x - 1 / 2) * (x - 1 / 2) \leq 0$

Rule? (grind :theories "real_props")

this simplifies to:

toy.2.1 :

|-----

{1} $-(x - 1 / 2) * x - (1 * -(x - 1 / 2)) / 2 \leq 0$

Tip 4: Introduce New Names to Avoid Distribution Laws

toy.2.1 :

|-----

{1} $-(x - 1 / 2) * (x - 1 / 2) \leq 0$

Rule? (name-replace "X" "(x-1/2)" :hide? nil)

this simplifies to:

toy.2.1 :

{-1} $(x - 1 / 2) = X$

|-----

{1} $-X * X \leq 0$

Tip 4: Introduce New Names to Avoid Distribution Laws

toy.2.1 :

|-----

{1} $-(x - 1 / 2) * (x - 1 / 2) \leq 0$

Rule? (name-replace "X" "(x-1/2)" :hide? nil)

this simplifies to:

toy.2.1 :

{-1} $(x - 1 / 2) = X$

|-----

{1} $-X * X \leq 0$

Finally ...

```
toy.2.1 :  
{-1} (x - 1 / 2) = X  
  |-----  
{1}  -X * X <= 0
```

Rule? (`grind :theories "real_props"`)

This completes the proof of toy.2.1.

```
toy.2.2 :
```

```
  |-----  
{1}  x - x * x - 1 / 4 = -(x - 1 / 2) * (x - 1 / 2)  
{2}  x - x * x - 1 / 4 <= 0
```

Rule? (`assert`)

Q.E.D.

Tip 5: Don't Reinvent the Wheel

(Look into the NASA libraries first!)

Theory reals@quadratic:

quadratic_le_0 : LEMMA

$a \cdot x^2 + b \cdot x + c \leq 0$ IFF

$((\text{discr}(a,b,c) \geq 0$ AND

$((a > 0$ AND $x_2(a,b,c) \leq x$ AND $x \leq x_1(a,b,c)$) OR

$(a < 0$ AND $(x \leq x_1(a,b,c)$ OR $x_2(a,b,c) \leq x))$) OR

$(\text{discr}(a,b,c) < 0$ AND $c \leq 0)$)

A Simpler Proof

|-----
{1} x * (1 - x) <= 1

Rule? (lemma "quadratic_le_0"
("a" "-1" "b" "1" "c" "-1" "x" "x"))
(grind)

Trying repeated skolemization, instantiation, and
if-lifting,

Q.E.D.

III. Strategies for Algebraic Manipulations

- ▶ **Manip**: Package for algebraic manipulations of real-valued expressions.
- ▶ Developed by B. Di Vito (NASA LaRC).
- ▶ Included as part of the PVS NASA Libraries.
- ▶ The package consists of:
 - ▶ Strategies.
 - ▶ Extended notations for formulas and expressions.
 - ▶ Emacs extensions.
 - ▶ Support functions for strategy developers.

Manip Strategies: Basic Manipulations

Strategy	Description
(swap-rel fnums)	Swap sides and reverse relations
(swap! expr-loc)	$x \circ y \Rightarrow y \circ x$
(group! expr-loc LR)	$(x \circ y) \circ z \Rightarrow x \circ (y \circ z)$
(flip-ineq fnums)	Negate and move inequalities
(split-ineq fnum)	Split \leq (\geq) into $<$ ($>$) and $=$

Extended Formula Notation

- ▶ Standard
 - ▶ *: All formulas.
 - ▶ -: All formulas in the antecedent.
 - ▶ +: All formulas in the consequent.

- ▶ Extended (Manip strategies only)
 - ▶ $(\hat{\ } n_1 \dots n_k)$: All formulas but n_1, \dots, n_k
 - ▶ $(-\hat{\ } n_1 \dots n_k)$: All antecedent formulas but n_1, \dots, n_k
 - ▶ $(+\hat{\ } n_1 \dots n_k)$: All consequent formulas but n_1, \dots, n_k

(Basic) Extended Expression Notation

- ▶ Term indexes:
 - ▶ L,R: Left- or right-hand side of a formula.
 - ▶ n : n -th term from left to right in a formula.
 - ▶ $-n$: n -th term from right to left in a formula.
 - ▶ $*$: All terms in a formula.
 - ▶ $(\wedge n_1 \dots n_k)$: All terms in a formula but n_1, \dots, n_k .
- ▶ Location references:
 - ▶ $(! \text{ fnum LR } i_1 \dots i_n)$: Term in formula fnum , Left- or Right-hand side, at recursive path location $i_1 \dots i_n$.

Examples

```
{-1}  x * r + y * r + 1 >= r - 1
      |-----
{1}   r = y * 2 * x + 1
```

Rule? (swap-rel -1)

```
{-1}  r - 1 <= x * r + y * r + 1
      |-----
{1}   r = y * 2 * x + 1
```

Rule? (swap! (! -1 R 1))

```
{-1}  r - 1 <= r * x + y * r + 1
      |-----
{1}   r = y * 2 * x + 1
```

Examples

```
{-1}  x * r + y * r + 1 >= r - 1
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Examples

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{-1}  r - 1 <= x * r + y * r + 1
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{1}   r = y * 2 * x + 1
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Rule? (swap! (! -1 R 1))

```
{-1}  r - 1 <= r * x + y * r + 1
      |-----
{1}   r = y * 2 * x + 1
```

Examples

$\{-1\}$ $x * r + y * r + 1 \geq r - 1$
|-----
 $\{1\}$ $r = y * 2 * x + 1$

Rule? (swap-rel -1)

$\{-1\}$ $r - 1 \leq x * r + y * r + 1$
|-----
 $\{1\}$ $r = y * 2 * x + 1$

Rule? (swap! (! -1 R 1))

$\{-1\}$ $r - 1 \leq r * x + y * r + 1$
|-----
 $\{1\}$ $r = y * 2 * x + 1$

```
{-1}  r - 1 <= r * x + y * r + 1
      |-----
[1]   r = y * 2 * x + 1
```

Rule? (group! (! 1 R 1) R)

```
{-1}  r - 1 <= r * x + y * r + 1
      |-----
{1}   r = y * (2 * x) + 1
```

Rule? (flip-ineq -1)

```
      |-----
{1}   r - 1 > r * x + y * r + 1
[2]   r = y * (2 * x) + 1
```

```
{-1}  r - 1 <= r * x + y * r + 1
      |-----
[1]   r = y * 2 * x + 1
```

Rule? (group! (! 1 R 1) R)

```
[-1]  r - 1 <= r * x + y * r + 1
      |-----
{1}   r = y * (2 * x) + 1
```

Rule? (flip-ineq -1)

```
|-----
{1}   r - 1 > r * x + y * r + 1
[2]   r = y * (2 * x) + 1
```

```
{-1}  r - 1 <= r * x + y * r + 1
      |-----
[1]   r = y * 2 * x + 1
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Rule? (group! (! 1 R 1) R)

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[-1]  r - 1 <= r * x + y * r + 1
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```
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{1}   r - 1 > r * x + y * r + 1
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```

```
{-1}  r - 1 <= r * x + y * r + 1
      |-----
[1]   r = y * 2 * x + 1
```

Rule? (group! (! 1 R 1) R)

```
[-1]  r - 1 <= r * x + y * r + 1
      |-----
{1}   r = y * (2 * x) + 1
```

Rule? (flip-ineq -1)

```
|-----
{1}   r - 1 > r * x + y * r + 1
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```

```
|-----
{1}  r - 1 > r * x + y * r + 1
[2]  r = y * (2 * x) + 1
```

Rule? (split-ineq 1)

Splitting off the equality case from formula 1,
this yields 2 subgoals:

```
{-1}  r - 1 = r * x + y * r + 1
      |-----
[1]   r - 1 > r * x + y * r + 1
[2]   r = y * (2 * x) + 1
```

Rule? (postpone)

```
|-----
{1}  r - 1 = r * x + y * r + 1
[2]  r - 1 > r * x + y * r + 1
[3]  r = y * (2 * x) + 1
```

|-----
{1} $r - 1 > r * x + y * r + 1$
{2} $r = y * (2 * x) + 1$

Rule? (split-ineq 1)

Splitting off the equality case from formula 1,
this yields 2 subgoals:

{-1} $r - 1 = r * x + y * r + 1$
|-----
{1} $r - 1 > r * x + y * r + 1$
{2} $r = y * (2 * x) + 1$

Rule? (postpone)

|-----
{1} $r - 1 = r * x + y * r + 1$
{2} $r - 1 > r * x + y * r + 1$
{3} $r = y * (2 * x) + 1$

More Strategies

Strategy	Description
(mult-by fnums term)	Multiply formula by term
(div-by fnums term)	Divide formula by term
(move-terms fnum L R tnums)	Move additive terms left and right
(isolate fnum L R tnum)	Isolate additive terms
(cross-mult fnums)	Perform cross-multiplications
(factor fnums)	Factorize formulas
(factor! expr-loc)	Factorize terms
(mult-eq fnum fnum)	Multiply equalities
(mult-ineq fnum fnum)	Multiply inequalities

More Examples

$$\{-1\} \quad (x * r + y) / pa > (r - 1) / pb$$

|-----

$$\{1\} \quad r - y * 2 * x = 1$$

Rule? (cross-mult -1)

$$\{-1\} \quad pb * r * x + pb * y > pa * r - pa$$

|-----

$$\{1\} \quad r - y * 2 * x = 1$$

Rule? (isolate 1 L 1)

$$\{-1\} \quad pb * r * x + pb * y > pa * r - pa$$

|-----

$$\{1\} \quad (r = 1 + y * 2 * x)$$

More Examples

$$\{-1\} \quad (x * r + y) / pa > (r - 1) / pb$$

$$\{1\} \quad r - y * 2 * x = 1$$

Rule? (cross-mult -1)

$$\{-1\} \quad pb * r * x + pb * y > pa * r - pa$$

$$\{1\} \quad r - y * 2 * x = 1$$

Rule? (isolate 1 L 1)

$$\{-1\} \quad pb * r * x + pb * y > pa * r - pa$$

$$\{1\} \quad (r = 1 + y * 2 * x)$$

More Examples

$$\{-1\} \quad (x * r + y) / pa > (r - 1) / pb$$

$$\{1\} \quad r - y * 2 * x = 1$$

Rule? (cross-mult -1)

$$\{-1\} \quad pb * r * x + pb * y > pa * r - pa$$

$$\{1\} \quad r - y * 2 * x = 1$$

Rule? (isolate 1 L 1)

$$\{-1\} \quad pb * r * x + pb * y > pa * r - pa$$

$$\{1\} \quad (r = 1 + y * 2 * x)$$

More Examples

$$\{-1\} \quad (x * r + y) / pa > (r - 1) / pb$$

$$\{1\} \quad r - y * 2 * x = 1$$

Rule? (cross-mult -1)

$$\{-1\} \quad pb * r * x + pb * y > pa * r - pa$$

$$[1] \quad r - y * 2 * x = 1$$

Rule? (isolate 1 L 1)

$$[-1] \quad pb * r * x + pb * y > pa * r - pa$$

$$\{1\} \quad (r = 1 + y * 2 * x)$$

$$\{-1\} \quad x * y - pa + na < x * na * pa$$

$$\{-2\} \quad r - y * 2 * x = 1$$

|-----

$$\{1\} \quad 2 * pa = 2 * x + 2 * y$$

Rule? (move-terms -1 L (2 3))

$$\{-1\} \quad (x * y < x * na * pa + pa - na)$$

$$\{-2\} \quad r - y * 2 * x = 1$$

|-----

$$\{1\} \quad 2 * pa = 2 * x + 2 * y$$

Rule? (factor 1)

$$\{-1\} \quad (x * y < x * na * pa + pa - na)$$

$$\{-2\} \quad r - y * 2 * x = 1$$

|-----

$$\{1\} \quad 2 * pa = 2 * (x + y)$$

$$\{-1\} \quad x * y - pa + na < x * na * pa$$

$$\{-2\} \quad r - y * 2 * x = 1$$

|-----

$$\{1\} \quad 2 * pa = 2 * x + 2 * y$$

Rule? (move-terms -1 L (2 3))

$$\{-1\} \quad (x * y < x * na * pa + pa - na)$$

$$[-2] \quad r - y * 2 * x = 1$$

|-----

$$[1] \quad 2 * pa = 2 * x + 2 * y$$

Rule? (factor 1)

$$[-1] \quad (x * y < x * na * pa + pa - na)$$

$$[-2] \quad r - y * 2 * x = 1$$

|-----

$$\{1\} \quad 2 * pa = 2 * (x + y)$$

$$\{-1\} \quad x * y - pa + na < x * na * pa$$

$$\{-2\} \quad r - y * 2 * x = 1$$

|-----

$$\{1\} \quad 2 * pa = 2 * x + 2 * y$$

Rule? (move-terms -1 L (2 3))

$$\{-1\} \quad (x * y < x * na * pa + pa - na)$$

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|-----

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Rule? (factor 1)

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|-----

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|-----

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|-----

$$\{1\} \quad 2 * pa = 2 * (x + y)$$

$$[-1] \quad (x * y < x * na * pa + pa - na)$$

$$[-2] \quad r - y * 2 * x = 1$$

|-----

$$\{1\} \quad 2 * pa = 2 * (x + y)$$

Rule? (mult-eq -1 -2)

$$\{-1\} \quad (x * y) * (r - y * 2 * x) < (x * na * pa + pa - na) * 1$$

$$[-2] \quad (x * y < x * na * pa + pa - na)$$

$$[-3] \quad r - y * 2 * x = 1$$

|-----

$$[1] \quad 2 * pa = 2 * (x + y)$$

Rule? (div-by 1 "2")

...

|-----

$$\{1\} \quad (pa = (x + y))$$

$$[-1] \quad (x * y < x * na * pa + pa - na)$$

$$[-2] \quad r - y * 2 * x = 1$$

|-----

$$\{1\} \quad 2 * pa = 2 * (x + y)$$

Rule? (mult-eq -1 -2)

$$\{-1\} \quad (x * y) * (r - y * 2 * x) < (x * na * pa + pa - na) * 1$$

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Rule? (mult-eq -1 -2)

$$\{-1\} \quad (x * y) * (r - y * 2 * x) < (x * na * pa + pa - na) * 1$$

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$$[-3] \quad r - y * 2 * x = 1$$

|-----

$$\{1\} \quad 2 * pa = 2 * (x + y)$$

Rule? (div-by 1 "2")

...

|-----

$$\{1\} \quad (pa = (x + y))$$

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|-----

$$\{1\} \quad 2 * pa = 2 * (x + y)$$

Rule? (mult-eq -1 -2)

$$\{-1\} \quad (x * y) * (r - y * 2 * x) < (x * na * pa + pa - na) * 1$$

$$[-2] \quad (x * y < x * na * pa + pa - na)$$

$$[-3] \quad r - y * 2 * x = 1$$

|-----

$$[1] \quad 2 * pa = 2 * (x + y)$$

Rule? (div-by 1 "2")

...

|-----

$$\{1\} \quad (pa = (x + y))$$

The Field Package

- ▶ **Field**: A simplification procedure for the field of real numbers.
- ▶ Included as part of the PVS NASA Libraries.
- ▶ The package consists of:
 - ▶ The strategy `field`.
 - ▶ Several *extra-`tegies`*.

field

$$\{-1\} \quad \text{vox} > 0$$

$$\{-2\} \quad s * s - D * D > D$$

$$\{-3\} \quad s * \text{vix} * \text{voy} - s * \text{viy} * \text{vox} \neq 0$$

$$\{-4\} \quad ((s * s - D * D) * \text{voy} - D * \text{vox} * \text{sqrt}(s * s - D * D)) / \\ (s * (\text{vix} * \text{voy} - \text{vox} * \text{viy})) * s * \text{vox} \neq 0$$

$$\{-5\} \quad \text{voy} * \text{sqrt}(s * s - D * D) - D * \text{vox} \neq 0$$

|-----

$$\{1\} \quad (\text{viy} * \text{sqrt}(s * s - D * D) - \text{vix} * D) / \\ (\text{voy} * \text{sqrt}(s * s - D * D) - \text{vox} * D) = \\ (D * D - s * s) / (((s * s - D * D) * \text{voy} - D * \text{vox} * \\ \text{sqrt}(s * s - D * D)) / \\ (s * (\text{vix} * \text{voy} - \text{vox} * \text{viy})) * s * \text{vox}) + \\ \text{vix} / \text{vox}$$

Rule? (field 1)

Q.E.D.

field

$$\{-1\} \quad \text{vox} > 0$$

$$\{-2\} \quad s * s - D * D > D$$

$$\{-3\} \quad s * \text{vix} * \text{voy} - s * \text{viy} * \text{vox} \neq 0$$

$$\{-4\} \quad ((s * s - D * D) * \text{voy} - D * \text{vox} * \text{sqrt}(s * s - D * D)) / \\ (s * (\text{vix} * \text{voy} - \text{vox} * \text{viy})) * s * \text{vox} \neq 0$$

$$\{-5\} \quad \text{voy} * \text{sqrt}(s * s - D * D) - D * \text{vox} \neq 0$$

|-----

$$\{1\} \quad (\text{viy} * \text{sqrt}(s * s - D * D) - \text{vix} * D) / \\ (\text{voy} * \text{sqrt}(s * s - D * D) - \text{vox} * D) = \\ (D * D - s * s) / (((s * s - D * D) * \text{voy} - D * \text{vox} * \\ \text{sqrt}(s * s - D * D)) / \\ (s * (\text{vix} * \text{voy} - \text{vox} * \text{viy})) * s * \text{vox}) + \\ \text{vix} / \text{vox}$$

Rule? (field 1)

Q.E.D.

Some Extra-ategies

Strategy	Description
<code>(grind-reals)</code>	grind + real_props
<code>(cancel-by fnum term)</code>	Cancel a common term in a formula
<code>(skoletin fnum)</code>	Skolemize let-in expressions
<code>(skeep fnum)</code>	Skolemize with same variable names
<code>(neg-formula fnum)</code>	Negate a formula
<code>(add-formula fnum fnum)</code>	Add two formulas

Forget Tip 1 and Tip 4, Use grind-reals

|-----
{1} (x - 1 / 2) * (x - 1 / 2) >= 0

Rule? (`grind-reals :nodistrib 1`)

Q.E.D.

Forget Tip 1 and Tip 4, Use grind-reals

|-----
{1} (x - 1 / 2) * (x - 1 / 2) >= 0

Rule? (`grind-reals :nodistrib 1`)

Q.E.D.

cancel-by

$$\begin{array}{l} \{-1\} \quad 4 * (pa * pb) + (pa * 6) * pa = pa * ((c + 1) * 2) \\ \quad |----- \\ \{1\} \quad 2 * pb + 3 * pa = c \end{array}$$

Rule? (cancel-by -1 "2*pa")

$$\begin{array}{l} \{-1\} \quad (3 * pa) + (2 * pb) = 1 + c \\ \quad |----- \\ \{1\} \quad 2 * pa = 0 \\ \{2\} \quad 3 * pa + 2 * pb = c \end{array}$$

cancel-by

$$\begin{array}{l} \{-1\} \quad 4 * (pa * pb) + (pa * 6) * pa = pa * ((c + 1) * 2) \\ \quad |----- \\ \{1\} \quad 2 * pb + 3 * pa = c \end{array}$$

Rule? (cancel-by -1 "2*pa")

$$\begin{array}{l} \{-1\} \quad (3 * pa) + (2 * pb) = 1 + c \\ \quad |----- \\ \{1\} \quad 2 * pa = 0 \\ \{2\} \quad 3 * pa + 2 * pb = c \end{array}$$

PVS's Let-in Expressions

- ▶ Let-in expressions are used in PVS to introduce local definitions.
- ▶ They are automatically unfolded by the theorem prover.

```
|-----  
{1}  LET a = a * y + 2 IN  
      LET b = a + x IN  
      LET c = a + b IN -b + 4 * a * c / 2 = 0
```

Rule? (**assert**)

```
|-----  
{1}  (32 + 8 * (x*x * y*y) + 4 * (x*x*y) + 16 * (x*y) +  
      16 * (x*y) + 8*x) / 2 + -(2 + x*y + x) = 0
```

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- ▶ Let-in expressions are used in PVS to introduce local definitions.
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{1}  (32 + 8 * (x*x * y*y) + 4 * (x*x*y) + 16 * (x*y) +  
      16 * (x*y) + 8*x) / 2 + -(2 + x*y + x) = 0
```

Let-in Expressions Go Wild

```
|-----  
{1} LET a = (x + 1) IN LET b = a * a IN  
    LET c = b * b IN c * c >= a
```

Rule? (assert)

```
|-----  
{1} 1 + x + (x*x*x*x*x*x*x*x*x + x*x*x*x*x*x*x*x*x)  
    + (x*x*x*x*x*x*x*x*x + x*x*x*x*x*x*x*x*x)  
    + (x*x*x*x*x*x*x*x*x + x*x*x*x*x*x*x*x*x)  
    ...  
    + (x*x + x)  
    + (x*x + x)  
    + (x*x + x)  
    >= 1 + x
```

Let-in Expressions Go Wild

```
|-----  
{1} LET a = (x + 1) IN LET b = a * a IN  
    LET c = b * b IN c * c >= a
```

Rule? (assert)

```
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{1} 1 + x + (x*x*x*x*x*x*x*x*x*x + x*x*x*x*x*x*x*x*x*x)  
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      + (x*x*x*x*x*x*x*x*x*x + x*x*x*x*x*x*x*x*x*x)  
      ...  
      + (x*x + x)  
      + (x*x + x)  
      + (x*x + x)  
      >= 1 + x
```

Tip 8. To unfold a let-in, use skoletin

```
|-----  
{1}  LET a = (x + 1) IN LET b = a * a IN  
      LET c = b * b IN c * c >= a
```

Rule? (skoletin 1)

```
{-1}  a = (x + 1)  
      |-----  
{1}  LET b = a * a IN LET c = b * b IN c * c >= a
```

Rule? (skoletin* 1)

```
{-1}  c = b * b  
{-2}  b = a * a  
[-3]  a = (x + 1)  
      |-----  
{1}  c * c >= a
```

Tip 8. To unfold a let-in, use `skoletin`

```
|-----  
{1}  LET a = (x + 1) IN LET b = a * a IN  
      LET c = b * b IN c * c >= a
```

Rule? (`skoletin 1`)

```
{-1}  a = (x + 1)  
      |-----  
{1}  LET b = a * a IN LET c = b * b IN c * c >= a
```

Rule? (`skoletin* 1`)

```
{-1}  c = b * b  
{-2}  b = a * a  
{-3}  a = (x + 1)  
      |-----  
{1}  c * c >= a
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{1}  LET a = (x + 1) IN LET b = a * a IN  
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```
{-1}  a = (x + 1)
```

```
|-----  
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```

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```
{-1}  c = b * b
```

```
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```

```
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Tip 8. To unfold a let-in, use skoletin

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```
{-1}  c = b * b
```

```
{-2}  b = a * a
```

```
[-3]  a = (x + 1)
```

```
|-----  
{1}  c * c >= a
```

More examples

```
|-----  
{1}  FORALL (nnx: nreal, x: real):  
      nnx > x - nnx*nnx AND x + 2 * nnx*nnx >= 4 * nnx  
      IMPLIES nnx > 1
```

Rule? (skeep)

```
{-1}  nnx > x - nnx*nnx  
{-2}  x + 2 *nnx*nnx >= 4 * nnx  
      |-----  
{1}  nnx > 1
```

Rule? (neg-formula -1)

```
{-1}  nnx*nnx - x > -nnx  
[-2]  x + 2 * nnx*nnx >= 4 * nnx  
      |-----  
[1]  nnx > 1
```

More examples

```
|-----  
{1}   FORALL (nnx: nnreal, x: real):  
      nnx > x - nnx*nnx AND x + 2 * nnx*nnx >= 4 * nnx  
      IMPLIES nnx > 1
```

Rule? (*skeep*)

```
{-1}  nnx > x - nnx*nnx  
{-2}  x + 2 *nnx*nnx >= 4 * nnx  
|-----  
{1}   nnx > 1
```

Rule? (*neg-formula -1*)

```
{-1}  nnx*nnx - x > -nnx  
[-2]  x + 2 * nnx*nnx >= 4 * nnx  
|-----  
[1]   nnx > 1
```

More examples

```
|-----  
{1}  FORALL (nnx: nreal, x: real):  
      nnx > x - nnx*nnx AND x + 2 * nnx*nnx >= 4 * nnx  
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More examples

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{1}  FORALL (nnx: nreal, x: real):  
      nnx > x - nnx*nnx AND x + 2 * nnx*nnx >= 4 * nnx  
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{-2}  x + 2 * nnx*nnx >= 4 * nnx  
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```

```
{-1} nnx*nnx - x > -nnx
{-2} x + 2 * nnx*nnx >= 4 * nnx
      |-----
[1]   nnx > 1
```

Rule? (add-formulas -1 -2)

```
{-1} 3 * (nnx*nnx) > -nnx + 4 * nnx
      |-----
[1]   nnx > 1
```

Rule? (cancel-by -1 "nnx")

Q.E.D.

```
{-1} nnx*nnx - x > -nnx
[-2] x + 2 * nnx*nnx >= 4 * nnx
      |-----
[1]   nnx > 1
```

Rule? (add-formulas -1 -2)

```
{-1} 3 * (nnx*nnx) > -nnx + 4 * nnx
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{-1} nnx*nnx - x > -nnx
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Rule? (add-formulas -1 -2)

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{-1} 3 * (nnx*nnx) > -nnx + 4 * nnx
      |-----
[1]   nnx > 1
```

Rule? (cancel-by -1 "nnx")

Q.E.D.

IV. Strategies for Specialized Domains

- ▶ Linear arithmetic via **Yices**.
- ▶ Numerical calculations via **Interval**.

Yices

- ▶ Yices is a SMT (Satisfiability Modulo Theories) solver developed at SRI.
- ▶ Background theories supported by Yices:
 - ▶ **Linear arithmetic (real and integer)**: addition and multiplication by scalar.
 - ▶ Arrays.
 - ▶ Uninterpreted functions.
 - ▶ Datatypes.
 - ▶ Bit vectors.
 - ▶ **Quantifiers**.

yices

```
|-----  
{1} EXISTS (x, y: real): x /= y
```

Rule? (yices)

```
(assert  
  (not (exists ( x_1::real y_2::real) (/= x_1 y_2))))
```

Result = unsat

Logical context is inconsistent. Use (reset) to reset.

unsat

Yices translation of negation is unsatisfiable

Simplifying with Yices,

Q.E.D.

yices

```
|-----  
{1} EXISTS (x, y: real): x /= y
```

Rule? (yices)

```
(assert  
  (not (exists ( x_1::real y_2::real) (/= x_1 y_2))))
```

Result = unsat

Logical context is inconsistent. Use (reset) to reset.

unsat

Yices translation of negation is unsatisfiable

Simplifying with Yices,

Q.E.D.

yices

|-----
{1} EXISTS (x, y: real):
1 <= 3 * x - 3 * y AND 3 * x - 3 * y <= 2

Rule? (yices)

Yices translation of negation is unsatisfiable
Simplifying with Yices,

Q.E.D.

Behind this QED there is no a "real" PVS proof!

yices

|-----
{1} EXISTS (x, y: real):
1 <= 3 * x - 3 * y AND 3 * x - 3 * y <= 2

Rule? (yices)

Yices translation of negation is unsatisfiable
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yices

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{1} EXISTS (x, y: real):
1 <= 3 * x - 3 * y AND 3 * x - 3 * y <= 2

Rule? (yices)

Yices translation of negation is unsatisfiable
Simplifying with Yices,

Q.E.D.

*Behind this QED there is **no** a "real" PVS proof!*

The Interval Package

- ▶ A package for interval analysis.
- ▶ Exact real calculations including trigonometric and transcendental functions.
- ▶ <http://research.nianet.org/~munoz/Interval>

numerical

|-----
{1} sin(6 * pi / 180) + sqrt(2) <= 2.11

Rule? (numerical)

Evaluating formula using numerical approximations,
Q.E.D.

Behind this QED there is a "real" PVS proof!

numerical

|-----
{1} sin(6 * pi / 180) + sqrt(2) <= 2.11

Rule? (numerical)

Evaluating formula using numerical approximations,
Q.E.D.

Behind this QED there is a "real" PVS proof!

numerical

|-----
{1} sin(6 * pi / 180) + sqrt(2) <= 2.11

Rule? (numerical)

Evaluating formula using numerical approximations,
Q.E.D.

Behind this QED there is a "real" PVS proof!

instint

{-1} x ## [| 0, 2 |]

|-----

{1} sqrt(x) + sqrt(3) < 315 / 100

Rule? (instint)

Proving that an expression is in a given interval,

Q.E.D.

instint

{-1} x ## [| 0, 2 |]

|-----

{1} sqrt(x) + sqrt(3) < 315 / 100

Rule? (instint)

Proving that an expression is in a given interval,

Q.E.D.

instint with Splitting

```
{-1} x ## [| 0, 1 |]
```

```
|-----
```

```
{1} x * (1 - x) ## [| 0, 1 / 3 |]
```

Rule? (instint)

```
{-1} x * (1 - x) ## Mult([|0, 1|], Sub([|1|], [|0, 1|]))
```

```
{-2} x ## [|0, 1|]
```

```
|-----
```

```
{1} Mult([|0, 1|], Sub([|1|], [|0, 1|])) < [|0, 1 / 3|]
```

Rule? (undo)

Rule? (instint :splitting 6)

Q.E.D.

instint with Splitting

{-1} x ## [| 0, 1 |]

|-----

{1} x * (1 - x) ## [| 0, 1 / 3 |]

Rule? (instint)

{-1} x * (1 - x) ## Mult([|0, 1|], Sub([|1|], [|0, 1|]))

{-2} x ## [|0, 1|]

|-----

{1} Mult([|0, 1|], Sub([|1|], [|0, 1|])) < [|0, 1 / 3|]

Rule? (undo)

Rule? (instint :splitting 6)

Q.E.D.

instint with Splitting

```
{-1} x ## [| 0, 1 |]
```

```
|-----
```

```
{1} x * (1 - x) ## [| 0, 1 / 3 |]
```

```
Rule? (instint)
```

```
{-1} x * (1 - x) ## Mult([|0, 1|], Sub([|1|], [|0, 1|]))
```

```
{-2} x ## [|0, 1|]
```

```
|-----
```

```
{1} Mult([|0, 1|], Sub([|1|], [|0, 1|])) < [|0, 1 / 3|]
```

```
Rule? (undo)
```

```
Rule? (instint :splitting 6)
```

```
Q.E.D.
```

Essential Tools for the Real Practitioner

- ▶ PVS NASA libraries.
- ▶ Manip/Field packages.
- ▶ Depending on your application: Yices and Interval.