

ASSIGNMENT 4

ACKNOWLEDGE YOUR SOURCES.

1. [5 marks] **Please send your answer to this question as text in an email message**, so I can easily and quickly approve or make suggestions.

Decide on your course project. Details about the project are now on the course web page. Please specify:

- (a) which paper(s) you will present
- (b) why you chose this paper (just a line or two)

Note: In case you wish to change your project topic later on, just send me a revised proposal and the reason or the change.

I will ask you to choose a presentation date on Piazza.

2. [10 marks] The *farthest point Voronoi diagram* of point set $P = \{p_1, \dots, p_n\}$, in the plane, is the locus of points equally far from more than one site of P . More precisely, define the farthest region for p_i to be $\{x \in \mathbb{R}^2 : d(x, p_i) \geq d(x, p_j), \forall p_j \neq p_i\}$.
 - (a) Prove that the farthest region for p_i is convex.
 - (b) Characterize the sites with unbounded farthest regions.
 - (c) Characterize the sites with empty regions.
3. [5 marks] Show how to compute the Gabriel graph from the Delaunay triangulation in linear time.
4. [10 marks] Give an $O(n \log n)$ time algorithm for the following problem. The input is a set S of n points in the plane and a natural number k . Partition S into a set of k disjoint subsets S_1, S_2, \dots, S_k that are “as far from each other as possible”—to be precise, you must maximize the minimum distance between two point from different subsets. Keep your algorithm high-level. You may use algorithms and results from class. Make sure you prove that your approach is correct.

HINT: this whole assignment is about Voronoi diagrams and Delaunay triangulations.