

EQ trans (k). $t_1 = t_2 \quad t_2 = t_3 \quad \dots \quad t_k = t_{k+1}$ for any terms t_1, \dots, t_{k+1} .

$$t_1 = t_{k+1}$$

"applying transitivity k times".

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Peano Arithmetic

EQ subs (r): $t_1 = t_2$

for any variable z ,
and terms r, t_1 and t_2

Axioms for Equality

$$\uparrow [E t_1 / z] = r [E t_2 / z]$$

ED * ED

EQ 1: $(\forall x (x = x))$ (reflexivity)

EQ 2: For each formula α and variable z ,

$$\text{--- } ((\forall x (\forall y ((x = y) \rightarrow (\alpha[x/z] \rightarrow \alpha[y/z])))))$$

=T * =T

EQ symm: $(\forall x (\forall y ((x = y) \rightarrow (y = x))))$ (symmetry)

EQ trans: $(\forall w (\forall x (\forall y ((x = y) \rightarrow ((y = w) \rightarrow (x = w))))))$
(transitivity)

Peano Arithmetic:

Constant symbols: 0

Function symbols: $s^{(1)}$, $+^{(2)}$, $\times^{(2)}$

Interpretation I:

Domain: \mathbb{N} the set of natural numbers $\{0, 1, 2, \dots\}$

$0^I = \text{zero}$.

$s^{(1)}$: the successor function. $s(x) = x + 1$, $\forall x \in \mathbb{N}$.

$+^{(2)}$: the addition function.

$\times^{(2)}$: the multiplication function

Every natural number has a corresponding term:

$$0 \equiv 0$$

$$1 \equiv s(0)$$

$$2 \equiv s(s(0))$$

$$3 \equiv s(s(s(0)))$$

Peano Axioms:

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$$\text{PA1: } (\forall x (\neg (s(x) = 0)))$$

- Every natural number's successor is not zero.
- Zero is not a successor.

$$\text{PA2: } (\forall x (\forall y ((s(x) = s(y)) \rightarrow (x = y))))$$

- If two numbers are the same, they must have the same predecessor $P(x)$ $((x = y) \rightarrow (P(x) = P(y)))$.
- No number has two different predecessors.

Two properties of addition:

$$\text{PA3: } (\forall x (x + 0 = x))$$

- The addition of any number and zero is the same number.
- Adding zero to any number yields the same number.

$$\text{PA4: } (\forall x (\forall y (x + s(y) = s(x + y))))$$

- The addition of one number and the successor of another number is the successor of the sum of the two numbers.
- Adding a successor yields the successor of adding the number.

Two properties of multiplication:

$$\text{PA5: } (\forall x (x \times 0 = 0))$$

- Multiplying any number by zero yields zero.

$$\text{PA6: } (\forall x (\forall y (x \times s(y) = (x \times y) + x)))$$

- Multiplying one number and the successor of another number is the product of the two numbers plus the first number.

① Proving $(\forall x \varphi)$ by induction.

PA7: For each formula φ and variable x ,
induction hypothesis

$$(\underbrace{\varphi[0/x]}_{\text{base case}} \rightarrow (\underbrace{(\forall x (\varphi[x/x] \rightarrow \varphi[S(x)/x])}_{\text{induction step}})) \rightarrow (\underbrace{(\forall x \varphi)}_{\text{conclusion}}))$$

This axiom allows us to prove properties of natural numbers by induction.

Base case: We need to prove $\varphi[0/x]$.

Induction step: We need to prove
 $(\forall x (\varphi[x/x] \rightarrow \varphi[S(x)/x]))$

Consider an arbitrary natural number k

Induction hypothesis: assume $\varphi[k/x]$ is true

We need to prove $\varphi[S(k)/x]$ is true

...

By the principle of mathematical induction, $(\forall x \varphi)$ is true
 QED

② Proving $(\forall x \varphi)$ by $\forall i$ (direct proof).

Proof:	1	u fresh assumption
	\vdots	\equiv
	\vdots	\equiv
	n	$\varphi[u/x]$
	$n+1$	$(\forall x \varphi) \quad \forall i: 1-n$

Properties of Natural Numbers

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Theorem: $(\forall x (\neg (s(x) = x)))$

Proof by induction:

Base case: We need to prove that $(\neg (s(0) = 0))$

- 1 $(\forall x (\neg (s(x) = 0)))$ PA1
- 2 $(\neg (s(0) = 0))$ $\forall e: 1 [0/x]$

Induction step: We need to prove that

$\phi \vdash (\forall x ((\neg (s(x) = x)) \rightarrow (\neg (s(s(x)) = s(x))))))$

Hint: use $\rightarrow i$ and PA2.

3	u fresh	assumption.
4	$(\neg (s(u) = u))$	assumption.
5	$(s(s(u)) = s(u))$	assumption.
6	$(\forall x (\forall y ((s(x) = s(y)) \rightarrow (x = y))))$	PA2
7	$(s(s(u)) = s(u)) \rightarrow (s(u) = u)$	
8	$(\forall y ((s(s(u)) = s(y)) \rightarrow (s(u) = y)))$	$\forall e: 6 [s(u)/x]$
9	$((s(s(u)) = s(u)) \rightarrow (s(u) = u))$	$\forall e: 8 [u/y]$
10		
11		
12		
13	$(s(u) = u)$	$\rightarrow e: 5, 9$
14	\perp	$\perp i: 4, 13$
15	$(\neg (s(s(u)) = s(u)))$	$\neg i: 5-14$
16	$((\neg (s(u) = u)) \rightarrow (\neg (s(s(u)) = s(u))))$	$\rightarrow i: 4-15$
17	$(\forall x ((\neg (s(x) = x)) \rightarrow (\neg (s(s(x)) = s(x))))))$	$\forall i: 3-16$

Theorem: $(\forall x (\neg (S(x) = x)))$

Proof by induction continued:

Let's finish the proof.

- 15 $((\neg (S(0) = 0)) \rightarrow ((\forall x (\neg (S(x) = x)) \rightarrow (\neg (S(S(x)) = S(x))))))$
 $\rightarrow (\forall x (\neg (S(x) = x)))$ PA7: $[\psi = (\neg (S(x) = x))]$
- 16 $((\forall x (\neg (S(x) = x)) \rightarrow (\neg (S(S(x)) = S(x)))) \rightarrow (\forall x (\neg (S(x) = x))))$
 $\rightarrow e: 2, 15$
- 17 $(\forall x (\neg (S(x) = x)))$ $\rightarrow e: 14, 16$

QED

Theorem: $\emptyset \vdash (\forall x (\forall y (x+y = y+x)))$

"Addition is commutative."

Proof by induction on x .

(x) Base case: We need to prove that $(\forall y (0+y = y+0))$.

We prove this by induction on y .

(y) Base case: We need to prove that $(0+0 = 0+0)$.

1 $(\forall x (x=x))$ EQ1.

2 $(0+0 = 0+0)$ $\forall e: 1 [(0+0)/x]$

(y) Induction step: We need to prove that

$(\forall y ((0+y = y+0) \rightarrow (0+S(y) = S(y)+0)))$

Hint: use PA4, PA3 and EQsubs(r), EQsymm,

3	u fresh	assumption
4	$(0+u = u+0)$	assumption.
5		
6	★ $(0+S(u) = S(0+u))$	PA4 + $\forall e \times 2$
7	★ $(S(0+u) = S(u+0))$	EQsubs($S(x)$): 4
8	$(u+0 = u)$	PA3 + $\forall e$
9	★ $(S(u+0) = S(u))$	EQsubs($S(x)$): 8.
10	$(S(u)+0 = S(u))$	PA3 + $\forall e$ + $\forall e \times 2$
11	$((S(u)+0 = S(u)) \rightarrow (S(u) = S(u)+0))$	EQsymm
12	★ $(S(u) = S(u)+0)$	$\rightarrow e: 10, 11$
13		
14		
15	$(0+S(u) = S(u)+0)$	EQtrans(3): 6, 7, 9, 12
16	$((0+u = u+0) \rightarrow (0+S(u) = S(u)+0))$	$\rightarrow i: 4-15$
17	$(\forall y ((0+y = y+0) \rightarrow (0+S(y) = S(y)+0)))$	$\forall i: 3-16$

Thought process of proof:

$(0 + u = u + 0)$ assumption.

6 $0 + s(u) = s(0 + u)$ PA4.

7 $s(0 + u) = s(u + 0)$ assumption + EQsubs.

9 $s(u + 0) = s(u)$ PA3 + EQsubs.

12 $s(u) = s(u) + 0$ PA3 + EQsymm

$= s(u) + 0$

15 $0 + s(u) = s(u) + 0$ EQtrans.

Proof by induction on x continued.

Base case: We need to prove that $(\forall y (0 + y = y + 0))$.

$$18 ((0 + 0 = 0 + 0) \rightarrow ((\forall y (0 + y = y + 0) \rightarrow (0 + S(y) = S(y) + 0))) \rightarrow (\forall y (0 + y = y + 0))) \quad \text{PA7.}$$

$$19 ((\forall y (0 + y = y + 0) \rightarrow (0 + S(y) = S(y) + 0)) \rightarrow (\forall y (0 + y = y + 0))) \quad \rightarrow e: 2, 18$$

$$20 (\forall y (0 + y = y + 0)) \quad \rightarrow e: 17, 19$$

This completes the proof of the base case.

(X) Induction step: We need to prove that

$$(\forall x (\forall y (x + y = y + x)) \rightarrow (\forall y (S(x) + y = y + S(x))))$$

Consider an arbitrary $x \in \mathbb{N}$. We will prove that

$$((\forall y (x + y = y + x)) \rightarrow (\forall y (S(x) + y = y + S(x))))$$

Assume that $(\forall y (x + y = y + x))$ (X)

We will prove that $(\forall y (S(x) + y = y + S(x)))$

We prove this by induction on y .

(Y) Base case: We will prove that $S(x) + 0 = 0 + S(x)$.

$$21 (\forall y (0 + y = y + 0)) \quad \text{by the base case of induction on } x$$

$$22 (0 + S(x) = S(x) + 0) \quad \forall e: 21$$

23

$$24 ((0 + S(x) = S(x) + 0) \rightarrow (S(x) + 0 = 0 + S(x))) \quad \text{EqSymm} + \forall e x 2$$

$$25 (S(x) + 0 = 0 + S(x)) \quad \rightarrow e: 22, 24$$

Properties of Natural Numbers.

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Proof by induction on x continued:

Induction step continued:

Induction on y continued:

(y) Induction step: We will prove that

$$(\forall y ((S(x) + y = y + S(x)) \rightarrow (S(x) + S(y) = S(y) + S(x))))$$

Choose a fresh variable u .

$$\text{We will prove } ((S(x) + u = u + S(x)) \rightarrow (S(x) + S(u) = S(u) + S(x)))$$

Assume $(S(x) + u = u + S(x))$. (**)

We will prove $(S(x) + S(u) = S(u) + S(x))$.

30★ $S(x) + S(u) = S(S(x) + u)$

31★ $S(x) + u = u + S(x)$

32★ $S(S(x) + u) = S(u + S(x))$

33★ $S(u + S(x)) = S(S(u + x))$

34 $x + u = u + x$

35 $((x + u = u + x) \rightarrow (u + x = x + u))$

36 $u + x = x + u$

37★ $S(S(u + x)) = S(S(x + u))$

38 $x + S(u) = S(x + u)$

39 $((x + S(u) = S(x + u)) \rightarrow (S(x + u) = x + S(u)))$ EQsymm + $\forall x 2$

40 $S(x + u) = x + S(u)$

41★ $S(S(x + u)) = S(x + S(u))$

42 $x + S(u) = S(u) + x$

43★ $S(x + S(u)) = S(S(u) + x)$

44 $S(u) + S(x) = S(S(u) + x)$

45 $(S(u) + S(x) = S(S(u) + x)) \rightarrow (S(S(u) + x) = S(u) + S(x))$ EQsymm + $\forall x$.

46★ $S(S(u) + x) = S(u) + S(x)$

47 $(S(x) + S(u) = S(u) + S(x))$

PA4 + $\forall x 2$.

reflexivity on (**).

EQsubs($S(x)$): 31.

PA4 + $\forall x 2$.

(*) + $\forall e$.

EQsymm + $\forall x 2$.

$\rightarrow e$: 34, 35.

EQsubs($S(S(x))$): 36.

PA4 + $\forall x 2$

EQsymm + $\forall x 2$

$\rightarrow e$: 38, 39.

EQsubs($S(x)$): 40.

(*) + $\forall e$.

EQsubs($S(x)$): 42.

PA4 + $\forall x 2$

EQsymm + $\forall x$.

$\rightarrow e$: 44, 45.

EQtrans(6): 30, 32, 33, 37, 41, 43, 46

Thought process of the proof:

$$(\forall y (x+y = y+x)) \quad (*)$$

$$s(x) + u = u + s(x) \quad (**)$$

$$30 \quad s(x) + s(u) = s(s(x) + u) \quad \text{PA4}$$

$$32 \quad s(s(x) + u) = s(u + s(x)) \quad (**) + \text{EQsubs}$$

$$33 \quad s(u + s(x)) = s(s + u(x)) \quad \text{PA4} + \text{EQsubs}$$

$$37 \quad s(s(u+x)) = s(s(x+u)) \quad (*) + \text{EQsubs}$$

$$41 \quad s(s(x+u)) = s(x+s(u)) \quad \text{PA4} + \text{EQsubs} + \text{EQsymm}$$

$$43 \quad s(x+s(u)) = s(s(u)+x) \quad (*) + \text{EQsubs}$$

$$46 \quad s(s(u)+x) = s(u) + s(x) \quad \text{PA4} + \text{EQsymm}$$

$$47 \quad s(x) + s(u) = s(u) + s(x) \quad \text{EQtrans()}$$