

Last time

- Finished introduction and overview:
 - ◆ Network access and physical media
 - ◆ Internet structure and ISPs
 - ◆ Delay & loss in packet-switched networks
 - ◆ Protocol layers, service models

This time

- Link layer
 - ◆ Overview
 - ◆ Error detection and correction
 - ◆ PPP

Chapter 5: The Data Link Layer

Our goals:

- Understand principles behind data link layer services:
 - ◆ error detection, correction
 - ◆ sharing a broadcast channel: multiple access
 - ◆ link layer addressing
 - ◆ reliable data transfer, flow control

- Instantiation and implementation of various link layer technologies

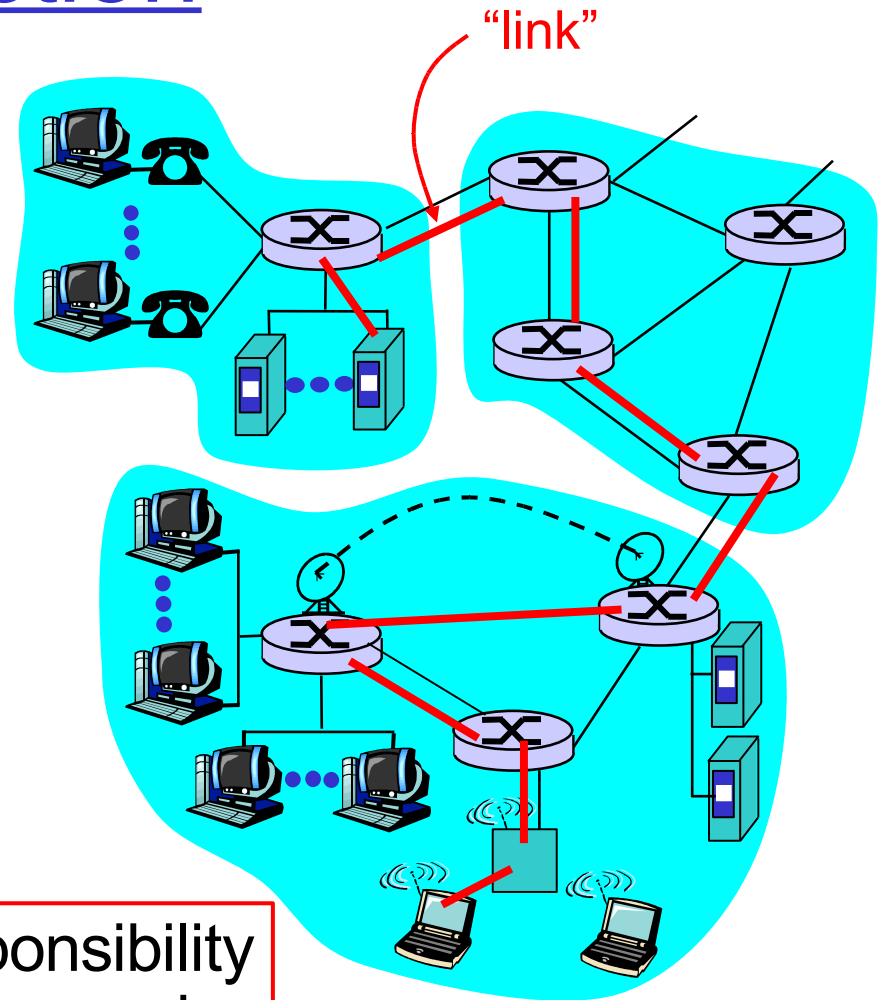
Link Layer

- 5.1 Introduction and services
- 5.2 Error detection and correction
- 5.3 Multiple access protocols
- 5.4 Link-Layer Addressing
- 5.5 Ethernet
- 5.6 Hubs and switches
- 5.7 PPP
- 5.8 Link Virtualization: ATM and MPLS

Link Layer: Introduction

Some terminology:

- Hosts and routers are **nodes**
- Communication channels that connect adjacent nodes along communication path are **links**
 - ◆ wired links
 - ◆ wireless links
 - ◆ LANs
- A layer-2 packet is a **frame**, and encapsulates a layer-3 datagram



The data-link layer has the responsibility of transferring datagrams from one node to an adjacent node over a link

Link layer: context

- Datagram transferred by different link protocols over different links:
 - ◆ e.g., Ethernet on first link, frame relay on intermediate links, 802.11 on last link
- Each link protocol provides different services
 - ◆ e.g., may or may not provide reliable delivery over link

Transportation analogy

- Trip from Princeton to Lausanne
 - ◆ limo: Princeton to JFK
 - ◆ plane: JFK to Geneva
 - ◆ train: Geneva to Lausanne
- Tourist = **datagram**
- Transport segment = **communication link**
- Transportation mode = **link layer protocol**
- Travel agent = **routing algorithm**

Link Layer Services

Different link layers offer different combinations of these services:

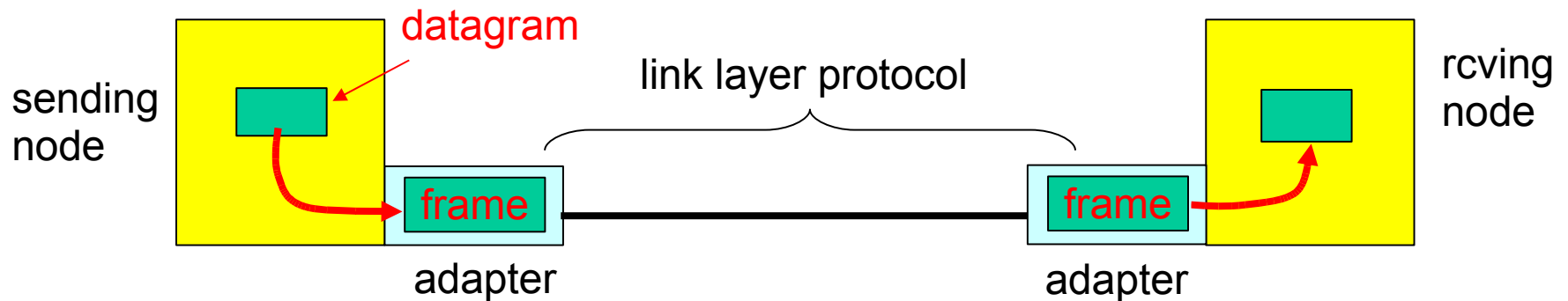
- **Framing, link access:**
 - ◆ encapsulate datagram into frame, adding header, trailer
 - ◆ channel access if shared medium
 - ◆ “MAC” addresses used in frame headers to identify source, destination

- **Reliable delivery between adjacent nodes**
 - ◆ based on acknowledgements and retransmissions
 - details later in the course (transport layer)
 - ◆ seldom used on low bit error link (fiber, some twisted pair)
 - ◆ wireless links: high error rates
 - Q: why both link-level and end-to-end reliability?

Link Layer Services (more)

- **Flow Control:**
 - ◆ pacing between adjacent sending and receiving nodes
- **Error Detection:**
 - ◆ errors caused by signal attenuation, noise.
 - ◆ receiver detects presence of errors:
 - signals sender for retransmission or drops frame
- **Error Correction:**
 - ◆ receiver identifies *and corrects* bit error(s) without resorting to retransmission
- **Half-duplex and full-duplex:**
 - ◆ with half duplex, nodes at both ends of link can transmit, but not at same time

Adapters Communicating



- Link layer implemented in “adapter” (aka NIC)
 - ◆ Ethernet card, PCMCIA card, 802.11 card
- Sending side:
 - ◆ encapsulates datagram in a frame
 - ◆ adds error checking bits, flow control, etc.
 - ◆ medium access control
- Receiving side:
 - ◆ looks for errors, flow control, etc
 - ◆ extracts datagram, passes to receiving node
- Adapter is semi-autonomous
- Link & physical layers

Link Layer

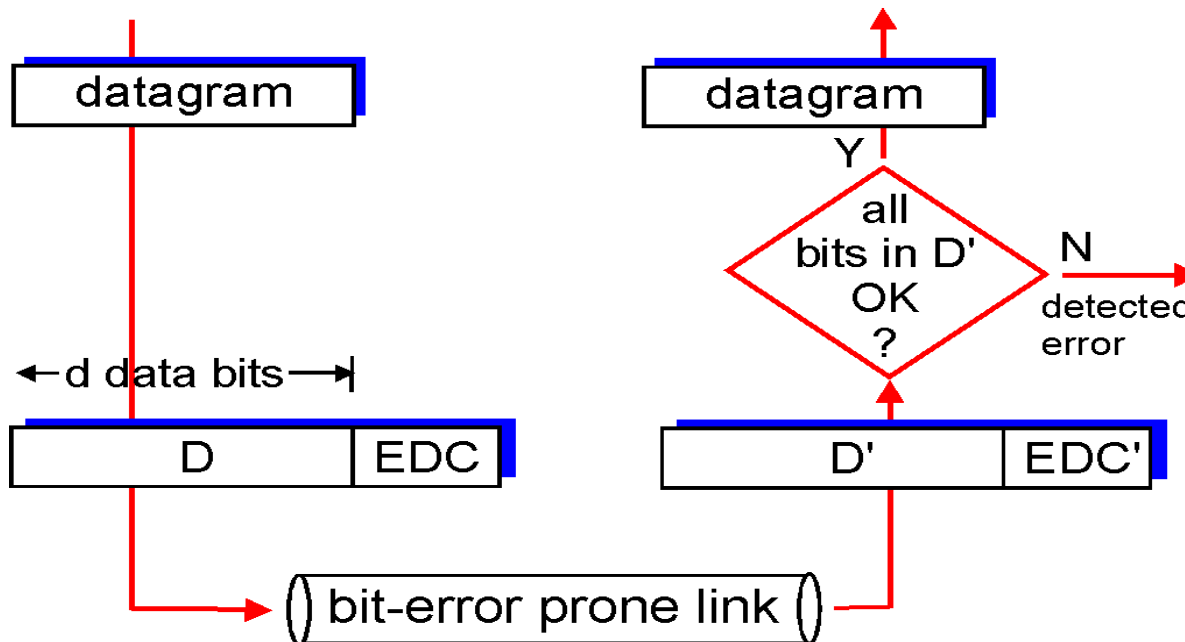
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Error Detection

EDC= Error Detection and Correction bits (redundancy)

D = Data protected by error checking, may include header fields

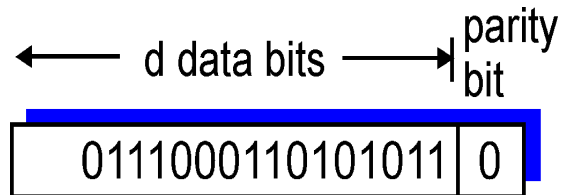
- Error detection not 100% reliable!
 - protocol may miss some errors, but rarely
 - larger EDC field yields better detection and correction



Parity Checking

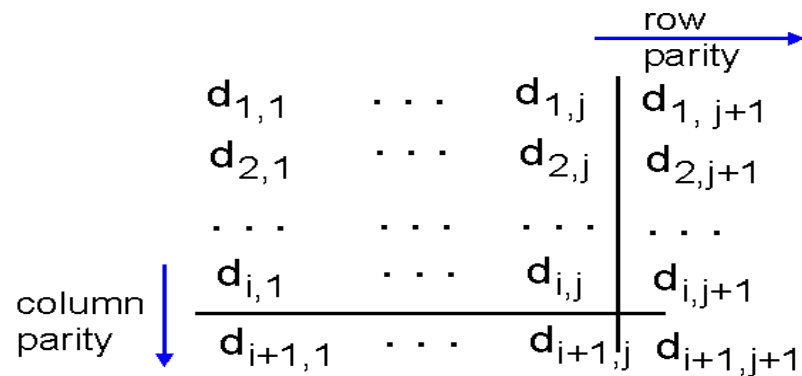
Single Bit Parity:

Detect single bit errors



Two Dimensional Bit Parity:

Detect *and correct* single bit errors



1	0	1	0	1	1
1	1	1	1	0	0
0	1	1	1	0	1
⏟	⏟	⏟	⏟	⏟	⏟
1	0	1	0	1	0

no errors

1	0	1	0	1	1
1	1	1	1	0	0
0	1	1	1	0	1
⏟	⏟	⏟	⏟	⏟	⏟
1	0	1	0	1	0

parity error

parity error

correctable
single bit error

Internet checksum

Goal: detect “errors” (e.g., flipped bits) in transmitted segment; easy to implement in software

Sender:

- Treat segment contents as sequence of 16-bit integers in 1's complement notation
- Checksum:
 - ◆ add integers
 - ◆ invert sum
- Put checksum value into checksum field

Receiver:

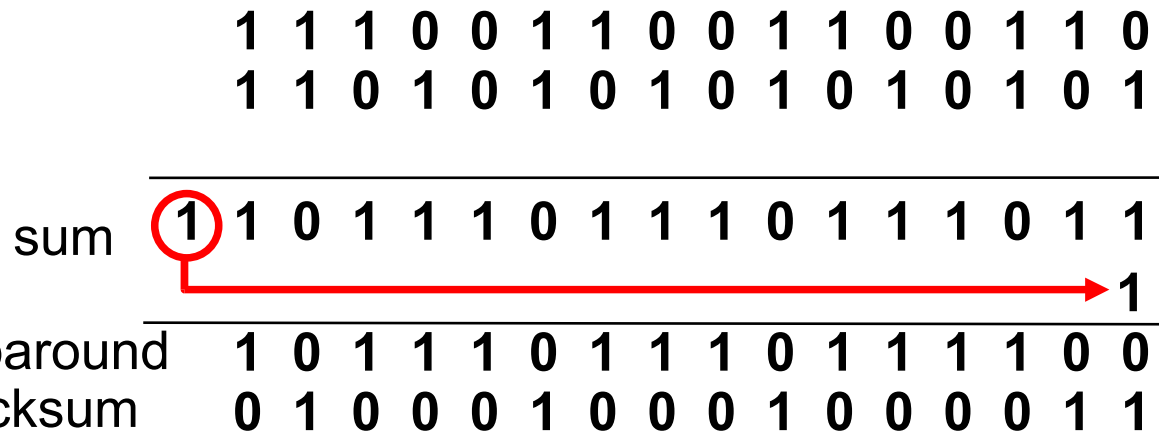
- Compute checksum of received segment
- Check if computed checksum equals checksum field value:
 - ◆ NO - error detected
 - ◆ YES - no error detected. *But maybe errors nonetheless?*

Internet Checksum Example

- Note

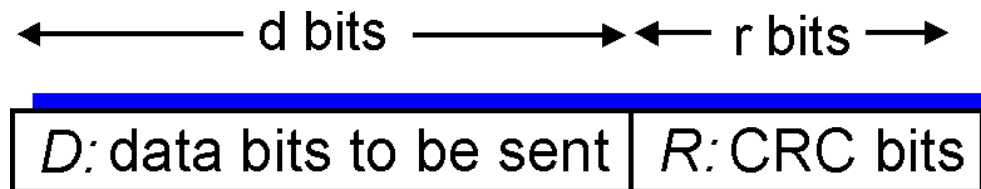
- ◆ When adding numbers in 1's complement notation, a carry-over from the most significant bit needs to be added to the result

- Example: add two 16-bit integers



Checksumming: Cyclic Redundancy Check

- Sender has data **D**, which is **d** bits long.
- Sender computes **R**, an **r**-bit CRC of those **d** bits
 - ◆ **r** is usually 8, 12, 16, or 32
- Sender sends **D** followed by **R** to the receiver
- Receiver computes the CRC of the received data
 - ◆ If no match, there were certainly errors
 - ◆ If match, no error detected



Checksumming: Cyclic Redundancy Check

- Better at detecting errors than other methods we've seen
 - ◆ Can detect all burst errors up to r bits
 - ◆ But much more complicated to compute

- Since the link layer is handled semi-autonomously in adapter hardware, CRCs are more appropriate for the link layer than for higher layers in the protocol stack.

- The mathematical details are in the text.

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Point to Point Data Link Control

- This is the simplest kind of link layer
 - ◆ One sender
 - ◆ One receiver
 - ◆ One link

- The most popular such protocol is **PPP** (point-to-point protocol)
 - ◆ You probably use this at home to connect to the Internet
 - Dialup: plain PPP
 - ADSL: PPP over Ethernet (PPPoE)

PPP Design Requirements [RFC 1557]

- **packet framing:** encapsulation of network-layer datagram in data link frame
 - ◆ carry network layer data of any network layer protocol (not just IP) *at the same time*
 - ◆ ability to demultiplex upwards
- **bit transparency:** must carry any bit pattern in the data field
- **error detection** (no correction)
- **connection liveness:** detect, signal link failure to network layer
- **network layer address negotiation:** endpoint can learn/configure each other's network address

PPP non-requirements

- no error correction/recovery
- no flow control
- out of order delivery OK
- no need to support multipoint links (e.g., polling)

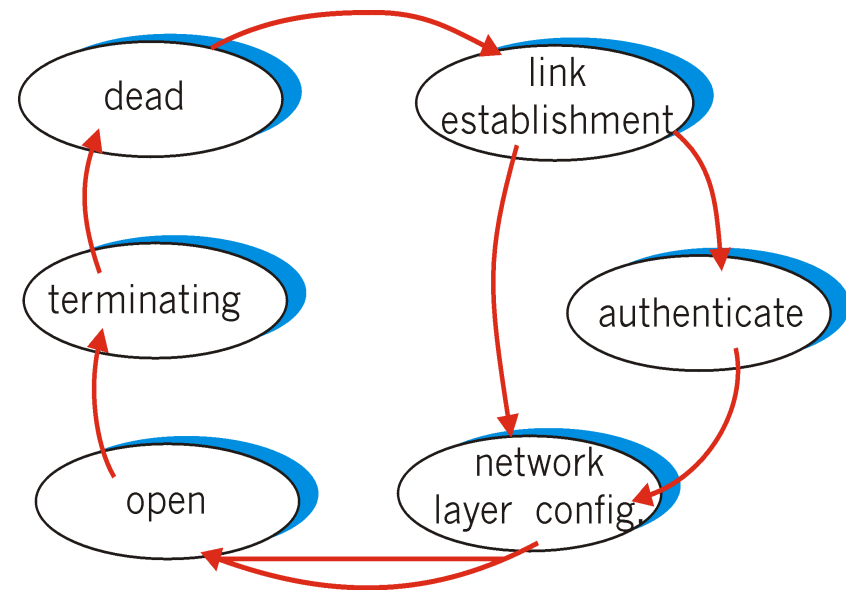
Error recovery, flow control, data re-ordering
all relegated to higher layers!

PPP Link Control Protocol

Before exchanging network-layer data, data link peers must

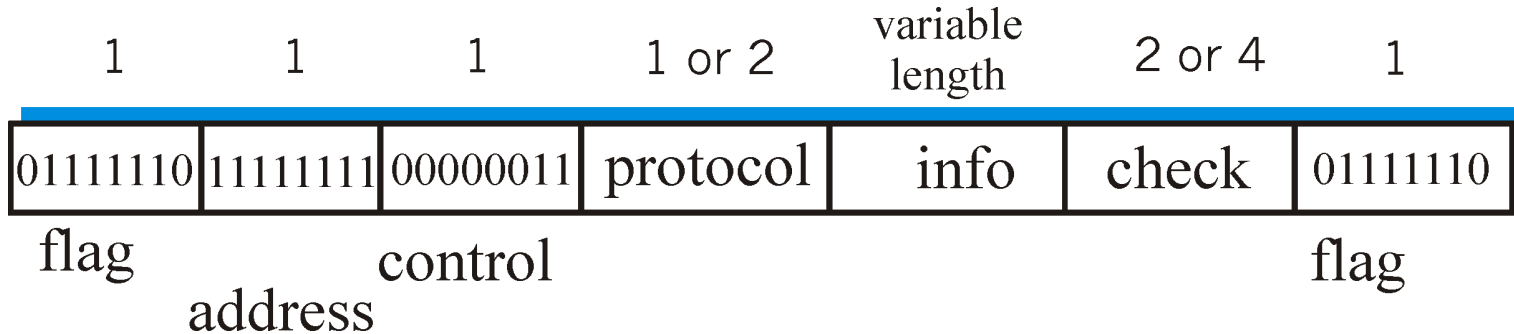
- **configure PPP link** (max. frame length, authentication)
- **learn/configure network layer information**

These are straightforward two-party protocols, because PPP is point-to-point



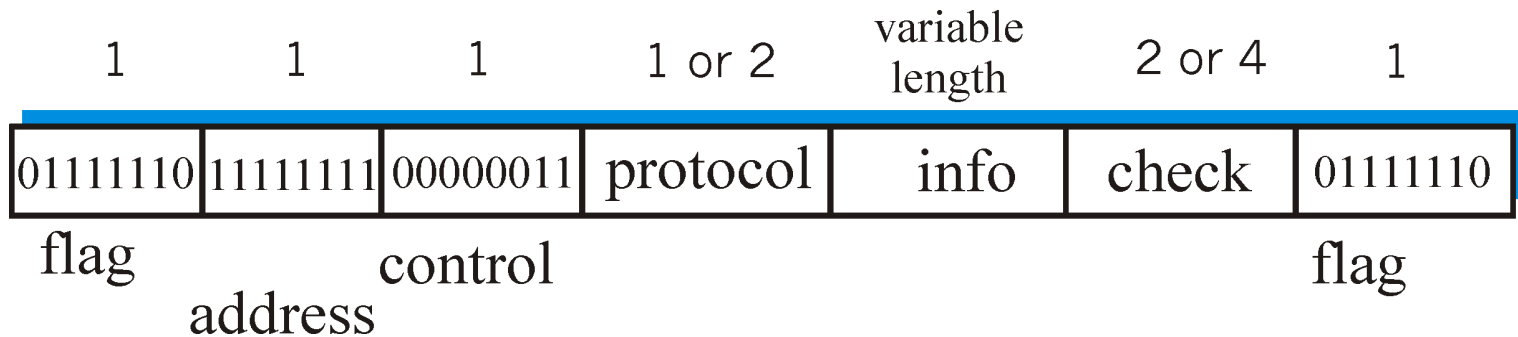
PPP Data Frame

- **Flag:** delimiter (framing)
- **Address:** does nothing (only one option)
- **Control:** does nothing; in the future possible multiple control fields
- **Protocol:** upper layer protocol to which frame delivered (e.g. PPP-LCP, IP, Appletalk, etc.)



PPP Data Frame

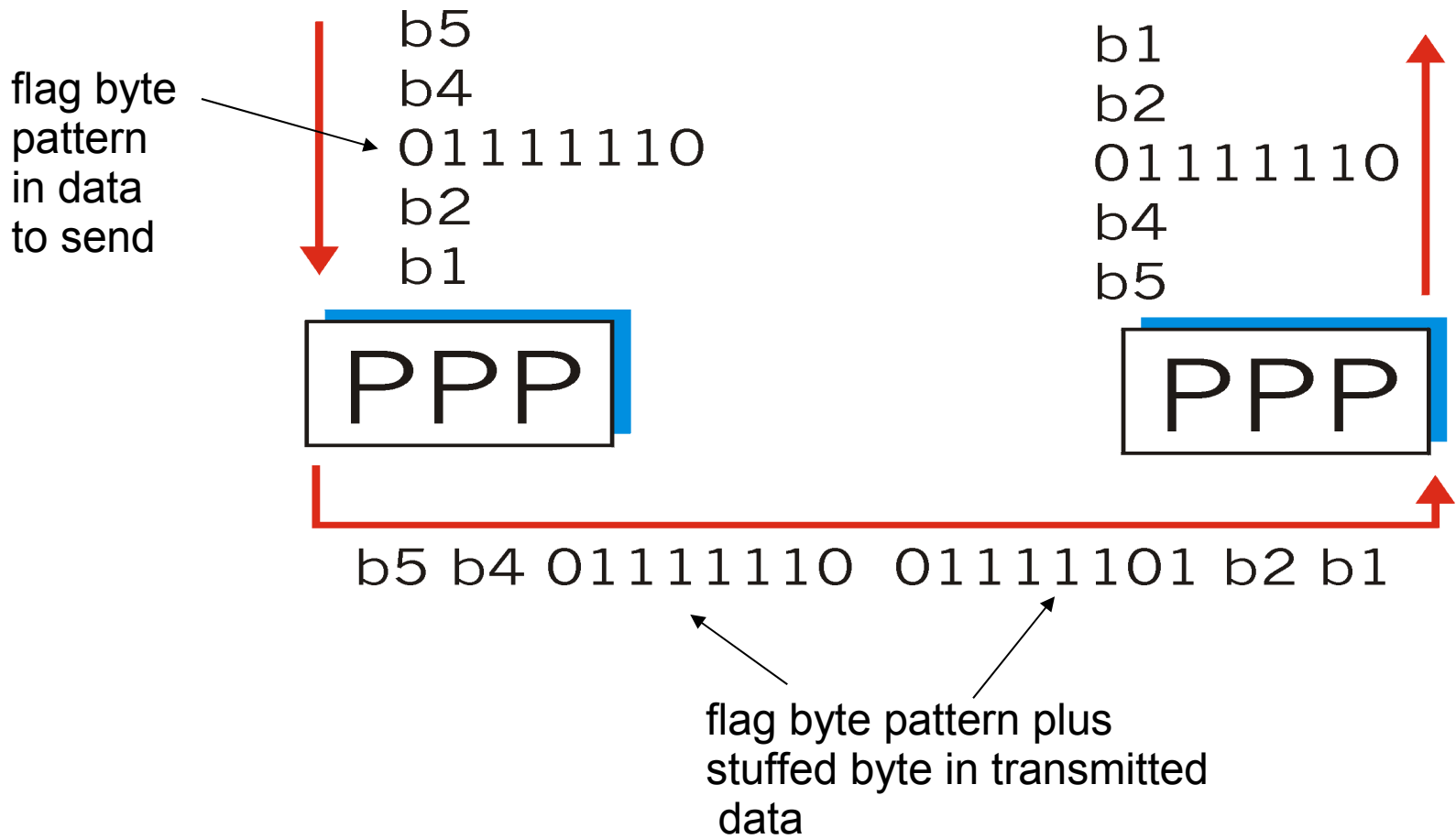
- **Info:** upper layer data being carried
- **Check:** cyclic redundancy check for error detection



Byte Stuffing

- “Data transparency” requirement: protocol, info, checksum fields must be allowed to include flag pattern <01111110>
 - **Q:** is received <01111110> the closing flag, or part of some other field?
- **Sender:** adds (“stuffs”) escape sequence <01111101> byte before:
 - each **non-flag** <01111110> byte
 - each <01111101> byte
- **Receiver:**
 - discard <01111101> byte and treat next as data
 - single <01111110>: flag byte

Byte Stuffing



Recap

- Link layer overview
 - ◆ Services
 - ◆ Adapters
- Error detection and correction
 - ◆ Parity check
 - ◆ Internet checksum
 - ◆ CRC
- PPP
 - ◆ Byte stuffing

Next time

- Multiple access protocols
- Link-layer addressing
- Ethernet