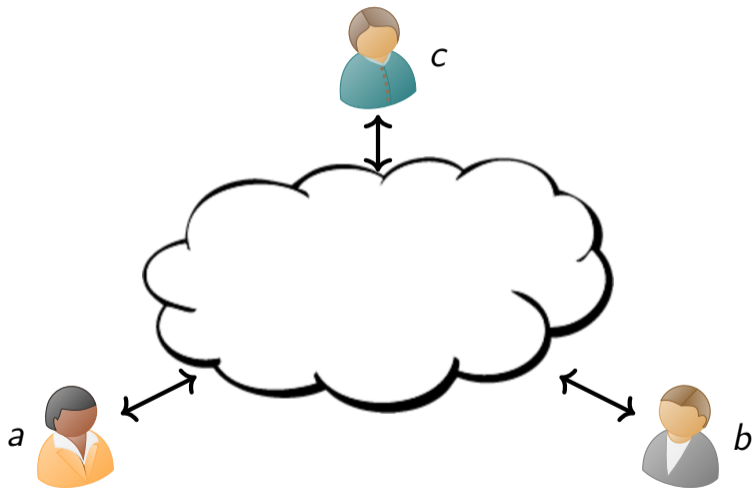


# PRAC: Round-Efficient 3-Party MPC for Dynamic Data Structures

Sajin Sasy, Adithya Vadapalli, **Ian Goldberg**

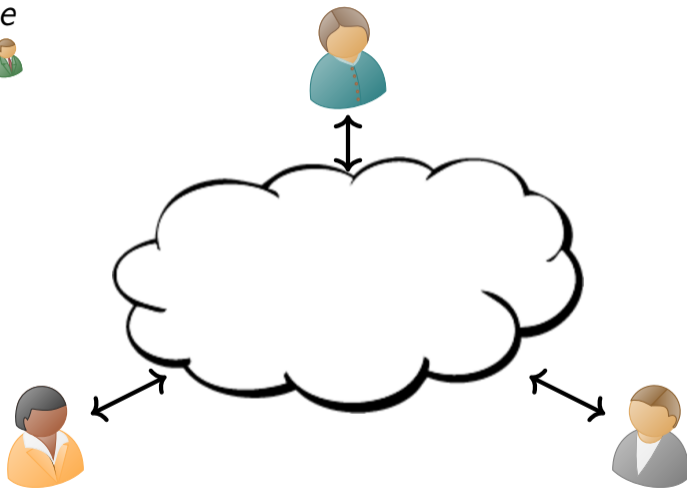
CrySP Group Meeting, 16 Oct 2023

# MPC: multiparty computation



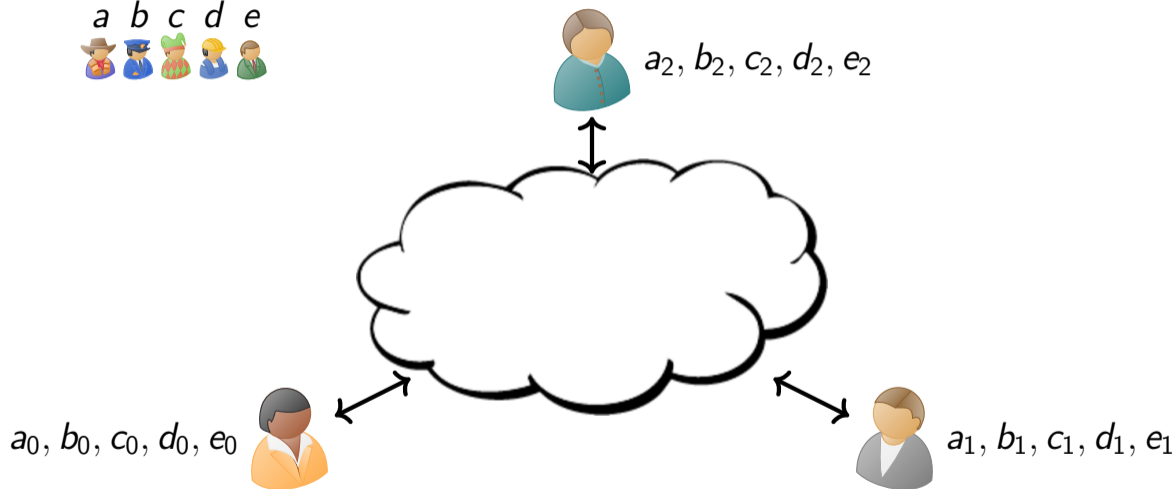
# MPC: multiparty computation

*a b c d e*  

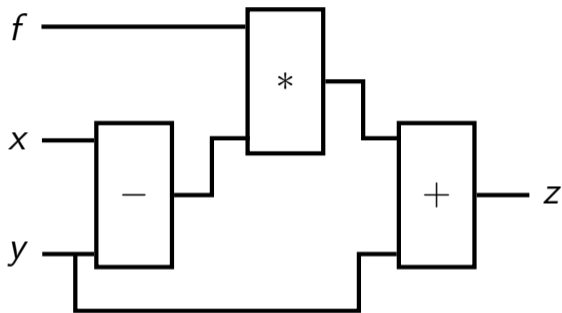



# MPC: multiparty computation

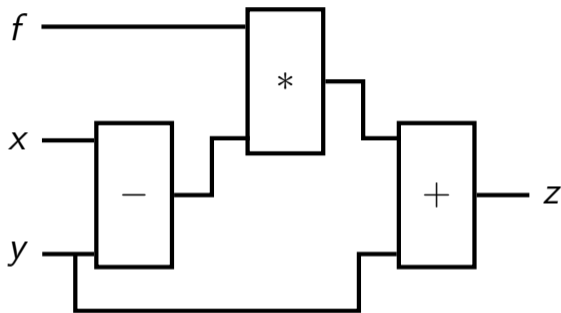
$a$   $b$   $c$   $d$   $e$



# Computation



# Computation



$$z = \begin{cases} y & \text{if } f = 0 \\ x & \text{if } f = 1 \end{cases}$$

# The hard part: memory

$$A = A_0 \oplus A_1, \quad i = i_0 \oplus i_1$$

$$(A_0, i_0), (A_1, i_1) \xrightarrow{\text{cloud}} z_0, z_1$$

$$z_0 \oplus z_1 = (A_0 \oplus A_1)[i_0 \oplus i_1]$$

# The hard part: memory

$$(A_0, i_0, v_0), (A_1, i_1, v_1) \xrightarrow{\text{cloud}} A'_0, A'_1$$

$$(A'_0 \oplus A'_1)[j] = \begin{cases} (A_0 \oplus A_1)[j] & j \neq i_0 \oplus i_1 \\ v_0 \oplus v_1 & j = i_0 \oplus i_1 \end{cases}$$



# The hard part: memory

$$(A_0, i_0, v_0), (A_1, i_1, v_1) \xrightarrow{\text{cloud}} A'_0, A'_1$$

$$(A'_0 \oplus A'_1)[j] = \begin{cases} (A_0 \oplus A_1)[j] & j \neq i_0 \oplus i_1 \\ v_0 \oplus v_1 & j = i_0 \oplus i_1 \end{cases}$$

DORAM: Distributed Oblivious Random Access Memory

# DORAMs: two approaches

## Computation-efficient:

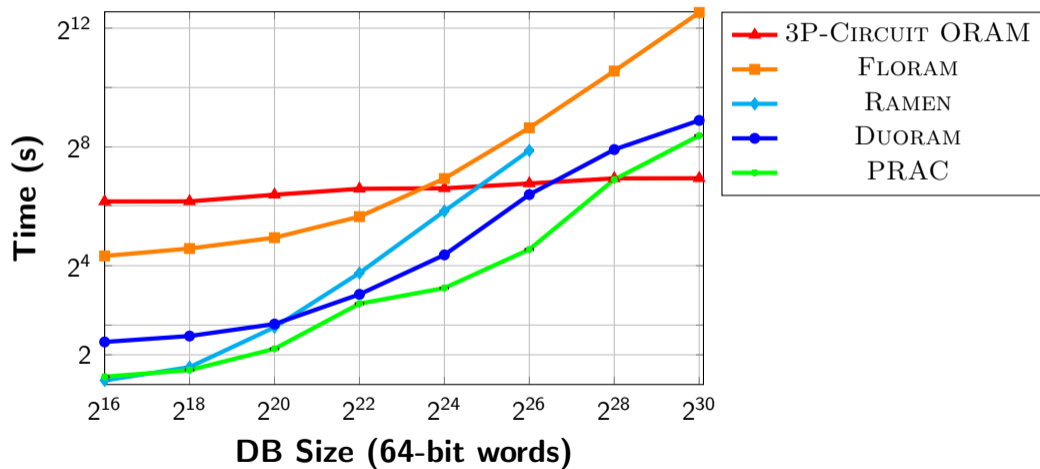
e.g., 3P-Circuit ORAM, Ramen,  
GigaDORAM

## Communication-efficient:

e.g., Floram, Duoram, **PRAC**

# DORAMs: two approaches

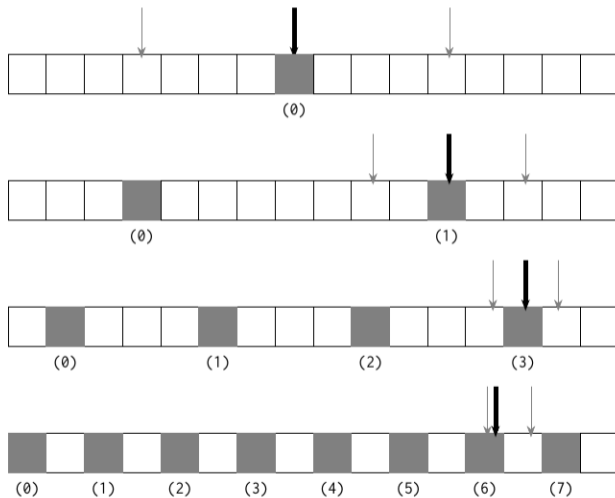
10 READ operations



# PRAC: Private Random Access Computations

- General MPC framework for random-access computations
- Underlying DORAM based on Duoram  
(Vadapalli, Henry, Goldberg, USENIX Security 2023)
- Implementation improvements
- Algorithmic improvements
  - use of incremental DPFs and wide DPFs
  - three oblivious data protocols:
    - binary search
    - heaps
    - AVL trees

# Binary search



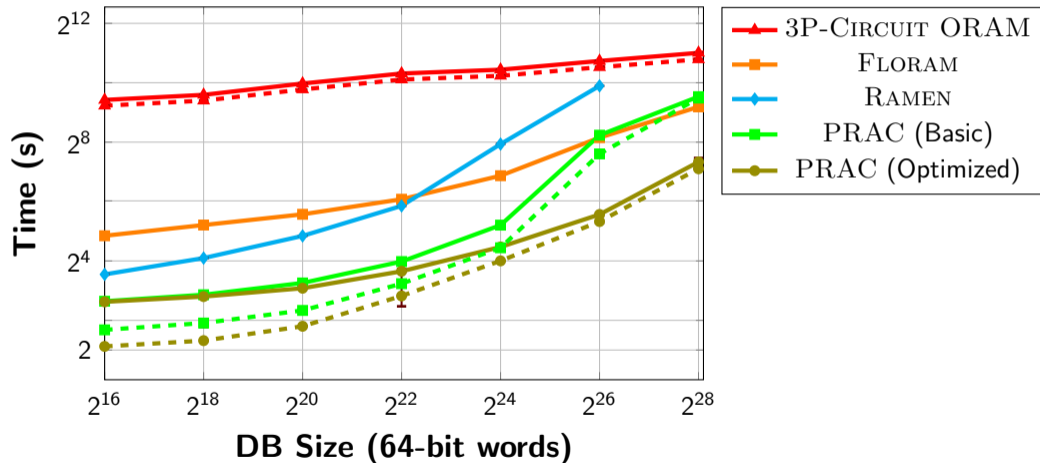
# Binary search

Analytical comparison of MPC binary search

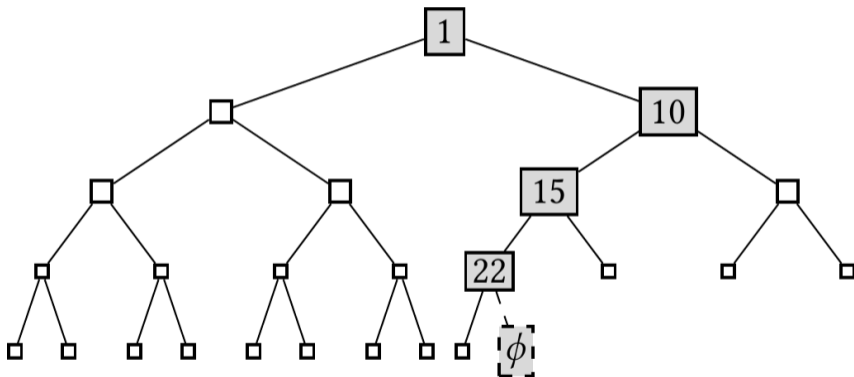
Protocol	Communication		Computation
	Bandwidth	Rounds	
Floram	$\mathcal{O}(\lg^2 n)$	$\mathcal{O}(\lg^2 n)$	$\mathcal{O}(n \lg n)$
Blanton & Yuan	$\mathcal{O}(\sqrt{n})$	$\mathcal{O}(\lg n)$	$\mathcal{O}(n)$
PRAC (Basic)	$\mathcal{O}(\lg^2 n)$	$\mathcal{O}(\lg n)$	$\mathcal{O}(n \lg n)$
PRAC (Optimized)	$\mathcal{O}(\lg n)$	$\mathcal{O}(\lg n)$	$\mathcal{O}(n)$

# Binary search

## 1 Binary Search



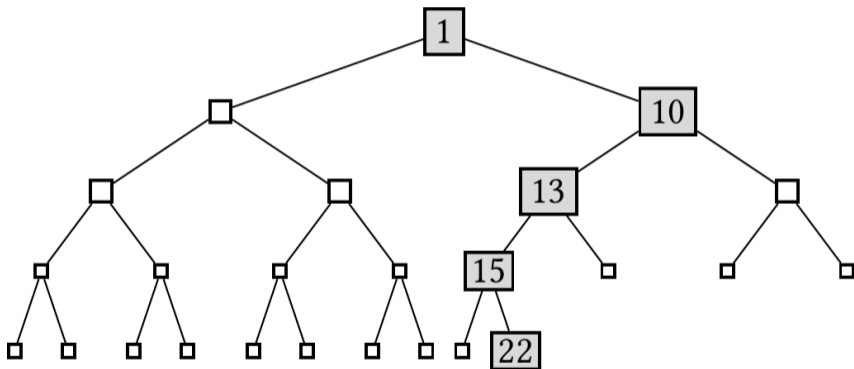
# Heaps: insertion



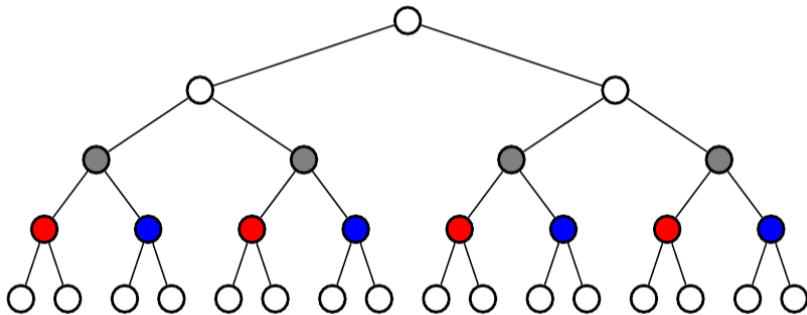




# Heaps: insertion



# Heaps: extraction



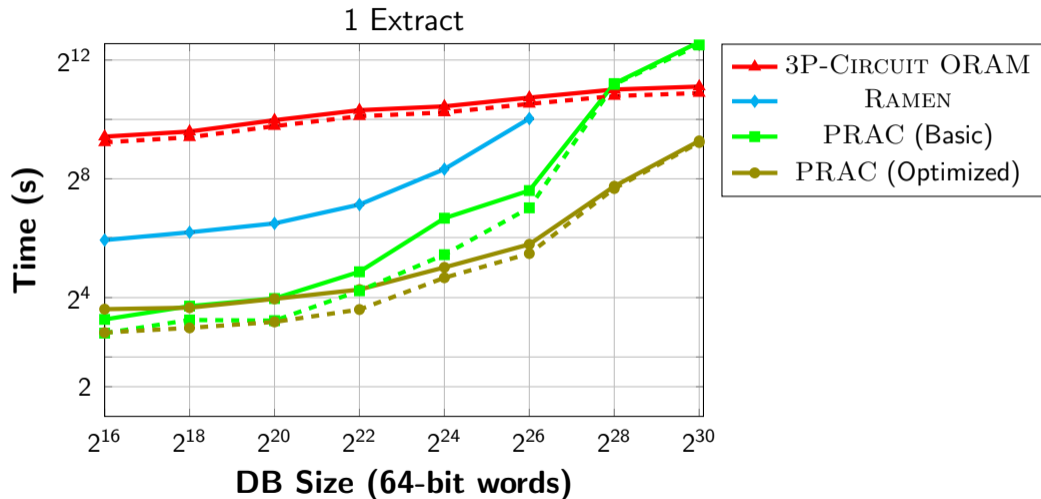
## Analytical comparison of MPC Heap protocols

Protocol	Insert		Computation
	Communication Bandwidth	Rounds	
Mazloom et al.	$\mathcal{O}(\lg^2 n)$	$\mathcal{O}(\lg^2 n)$	$\mathcal{O}(\lg^2 n)$
PRAC (Basic)	$\mathcal{O}(\lg n)$	$\mathcal{O}(\lg n)$	$\mathcal{O}(\lg n)$
PRAC (Optimized)	$\mathcal{O}(\lg n)$	$\mathcal{O}(\lg \lg n)$	$\mathcal{O}(\lg n)$

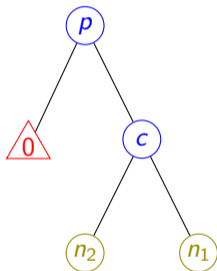
## Analytical comparison of MPC Heap protocols

Protocol	Extract		Computation
	Communication Bandwidth	Rounds	
Mazloom et al.	$\mathcal{O}(1)$	$\mathcal{O}(1)$	$\mathcal{O}(n)$
PRAC (Basic)	$\mathcal{O}(\lg^2 n)$	$\mathcal{O}(\lg n)$	$\mathcal{O}(n \lg n)$
PRAC (Optimized)	$\mathcal{O}(\lg n)$	$\mathcal{O}(\lg n)$	$\mathcal{O}(n)$

# Heaps



# AVL Trees

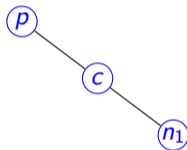


Case 1:  $n_1$  inserted  
Case 2:  $n_2$  inserted

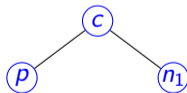
Case 1: L(R) Rotation

$$F_{dr} = d_{pc} \oplus d_{cn} = 0$$

After first ORotate:



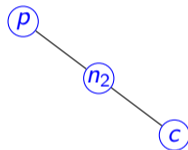
After second ORotate:



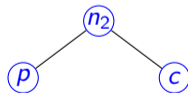
Case 2: RL(LR) Rotation

$$F_{dr} = d_{pc} \oplus d_{cn} = 1$$

After first ORotate:



After second ORotate:



# Takeaways

- Generally for MPC in reasonable deployments, minimizing **communication** (especially **rounds**) is much more important than minimizing **(parallelizable) computation**
- PRAC provides communication-optimized MPC framework for computations requiring random-access memory and dynamic data structures, including:
  - oblivious binary search (improved over prior work)
  - oblivious heaps (improved over prior work)
  - oblivious AVL trees (for the first time!)