

# Technology Transfer from Security Research Projects

**A Personal Perspective**

**N. Asokan**

**<https://asokan.org/asokan/research>**


# Five examples

- Optimistic Fair Exchange
- Generic Authentication Architecture
- Channel Binding in Protocol Composition
- Secure Device Pairing
- On-board Credentials

**Fair Exchange**

How can two mutually distrusting parties exchange digital "items" on the internet?

Existing solutions:




Gradual Exchange protocols      Trusted Third Party protocols

**Generic Authentication Architecture**

Can we bootstrap a general-purpose global-scale authentication and authorization infrastructure from the existing cellular security infrastructure?

- Need was evident:
  - "Global PKIs will not happen"
- Ad-hoc bootstrapping already in use
  - e.g., Coke vending machine accepting payments via SMS, 1997
- Idea: Bootstrap short-lived certificates from "local PKIs"



**Channel Binding in protocol composition**

Composing two secure authentication protocols carelessly can lead to a man-in-the-middle vulnerability

- Protocol composition can ease deployment
- Examples:
  - Server auth. using TLS + user auth. with password
  - Authentication for VPN access using legacy credentials
  - Bootstrapping a "local PKI"

**Secure Device Pairing**

How can the process of pairing two devices be made easy to use without compromising security or adding to cost?

**On-board Credentials**

Can we safely open up widely deployed secure hardware on mobile devices for use by app developers?

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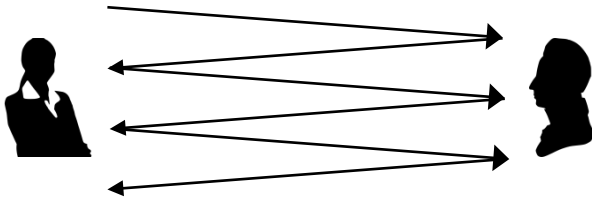
# Fair Exchange

## Optimistic fair exchange - recap

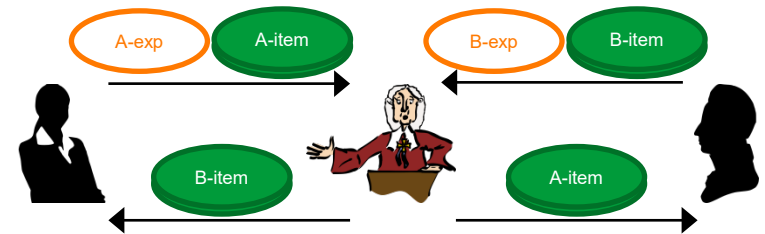
- Optimize for the **common case**
  - Exchange “permits” (e.g., in the form of verifiable escrows) first
  - If the actual exchange fails, take the permits to a third party to recover fairness
- But service providers want to **monetize every exchange**
  - An off-line/on-line computation variant (also more efficient)
- Still no deployment, but lots of citations
  - Renewed interest in the blockchain payment setting

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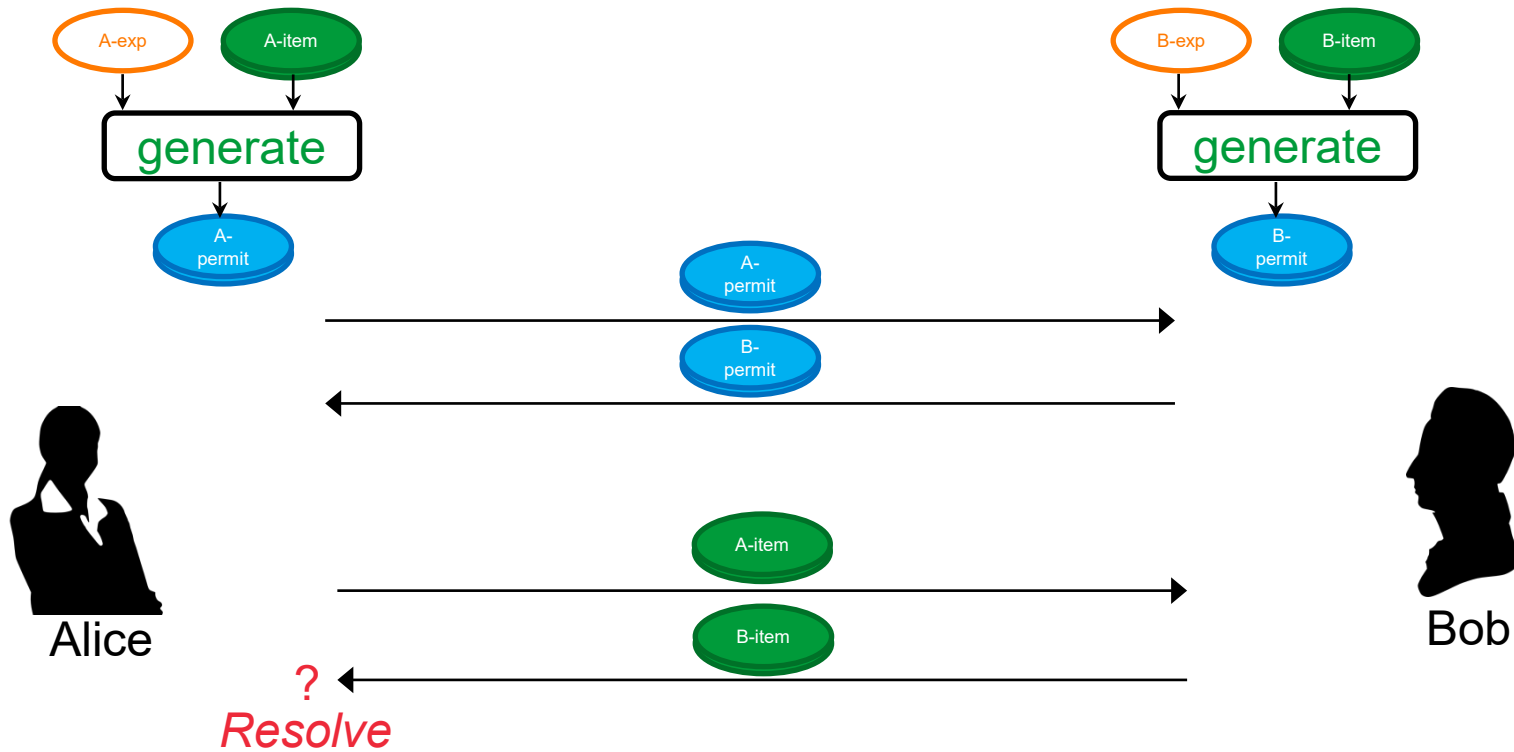


Trusted Third Party protocols

# Fair Exchange: design choices

- Common case: both *want* to complete the exchange
  - design protocol that is efficient for the common case
  - but allows recovery in case of exceptions
- Requirements
  - Effectiveness
  - Fairness
  - Timeliness
  - (Non-invasive)

# Optimistic Fair Exchange



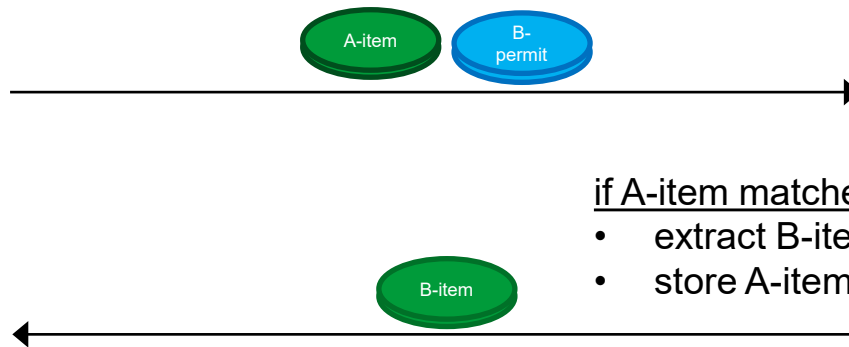
<http://www.semper.org/>

<https://web.archive.org/web/20200627000350/http://www.semper.org/>

<https://semper.schunter.org/>

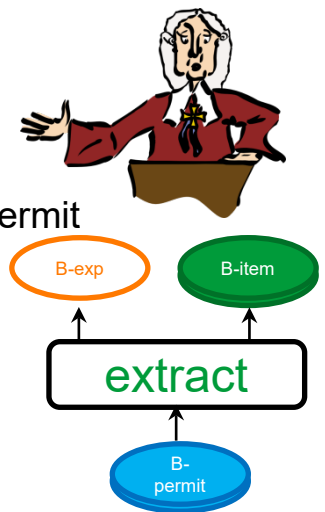
# Optimistic Fair Exchange: Recovery

*Resolve*

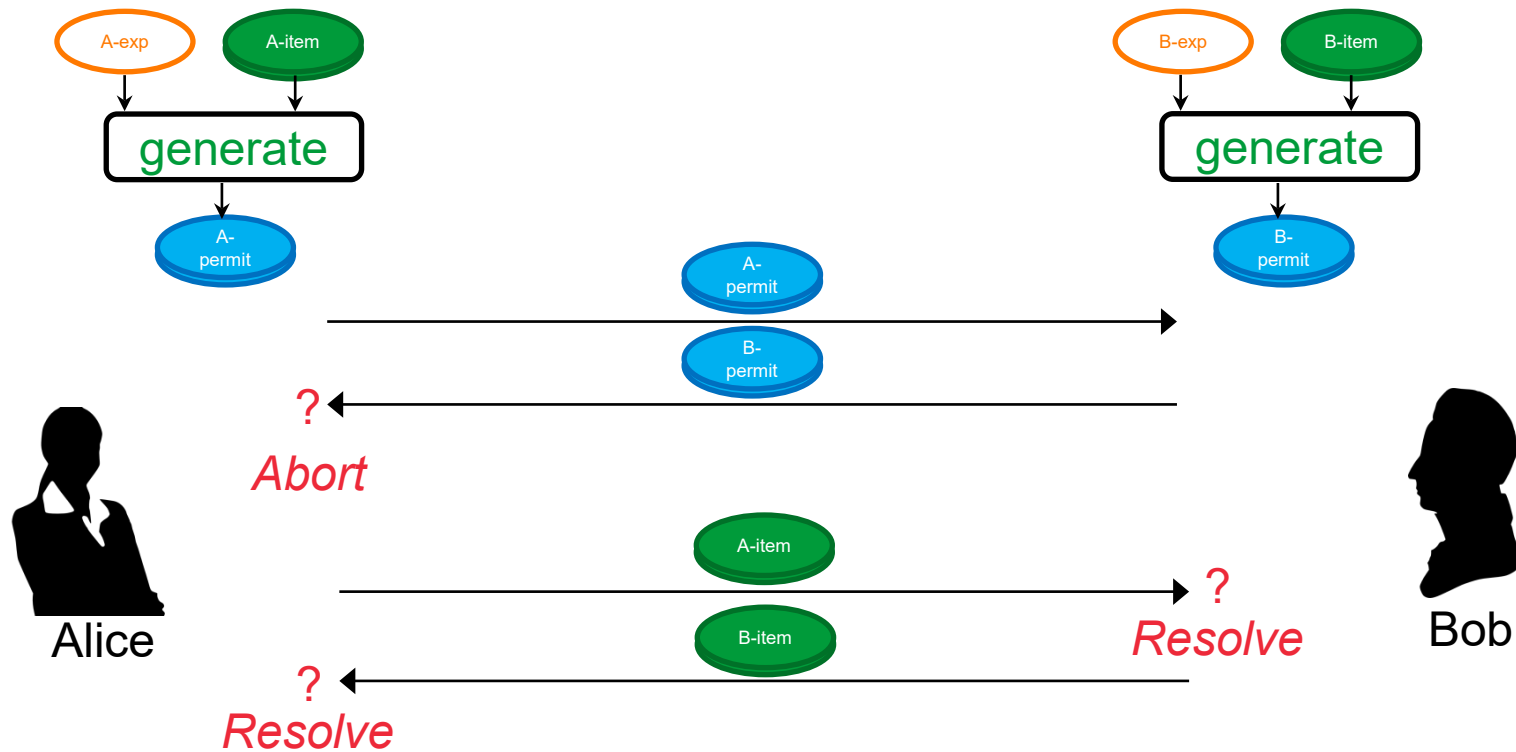


if A-item matches B-exp

- extract B-item from B-permit
- store A-item



# Optimistic Fair Exchange



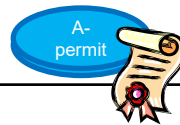


# Optimistic Fair Exchange: Recovery

*Abort*



Alice



If not resolved,  
issue abort token



*Resolve*

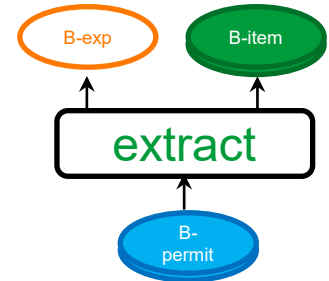


Alice



If not aborted, and  
if A-item matches B-exp

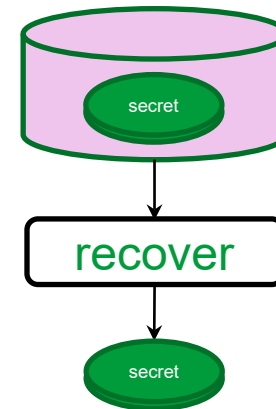
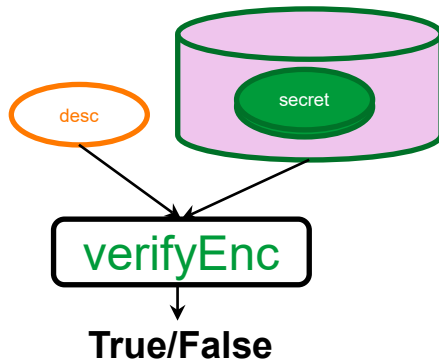
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*Resolve* for Bob is similar

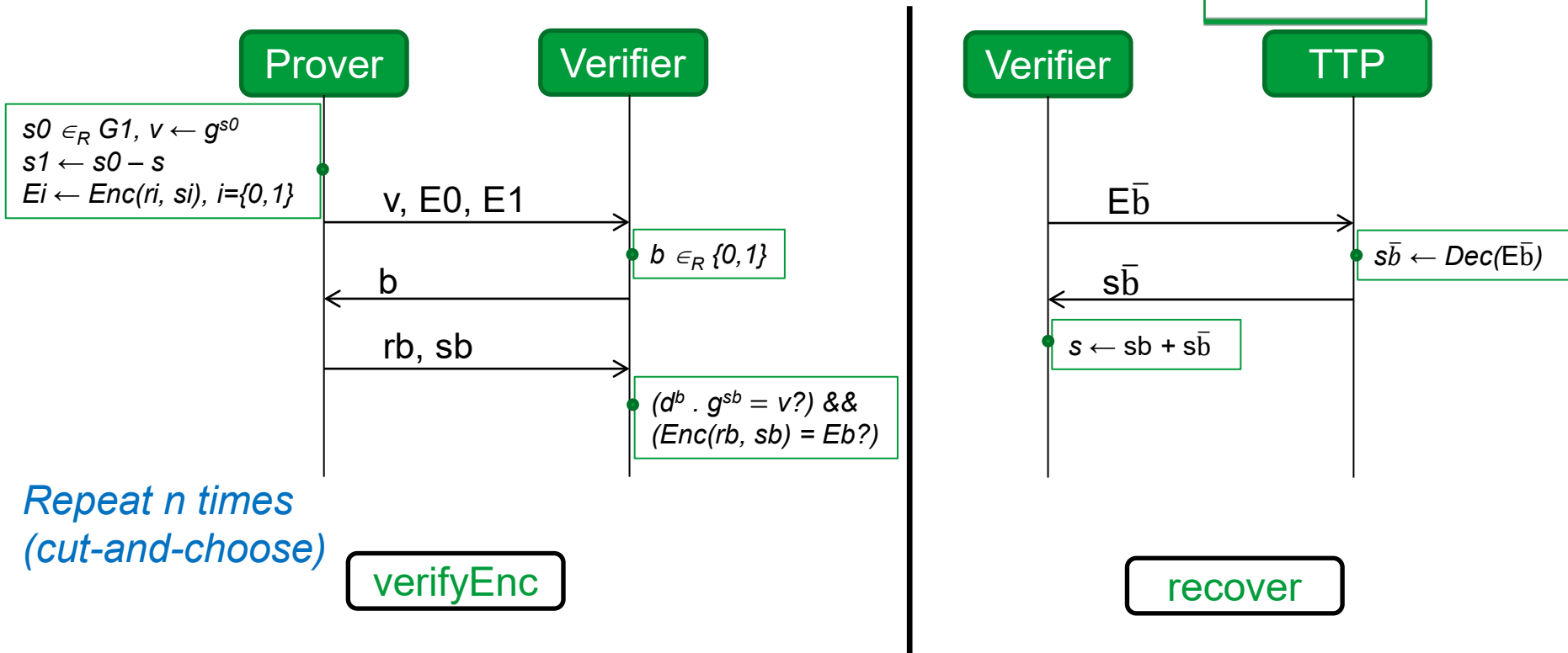
# Verifiable Encryption

Analogy - jewelry in a glass box: can see but can't touch



# Verifiable Encryption of discrete logs

Setting: secret =  $s \in G1$ , desc  $d = g^s$  (in  $G2$ )



# From Verifiable Encryptions to Permits

 A-exp = desc. of  B-item

 A-permit = Verifiable Encryption of  A-exp +  A-item

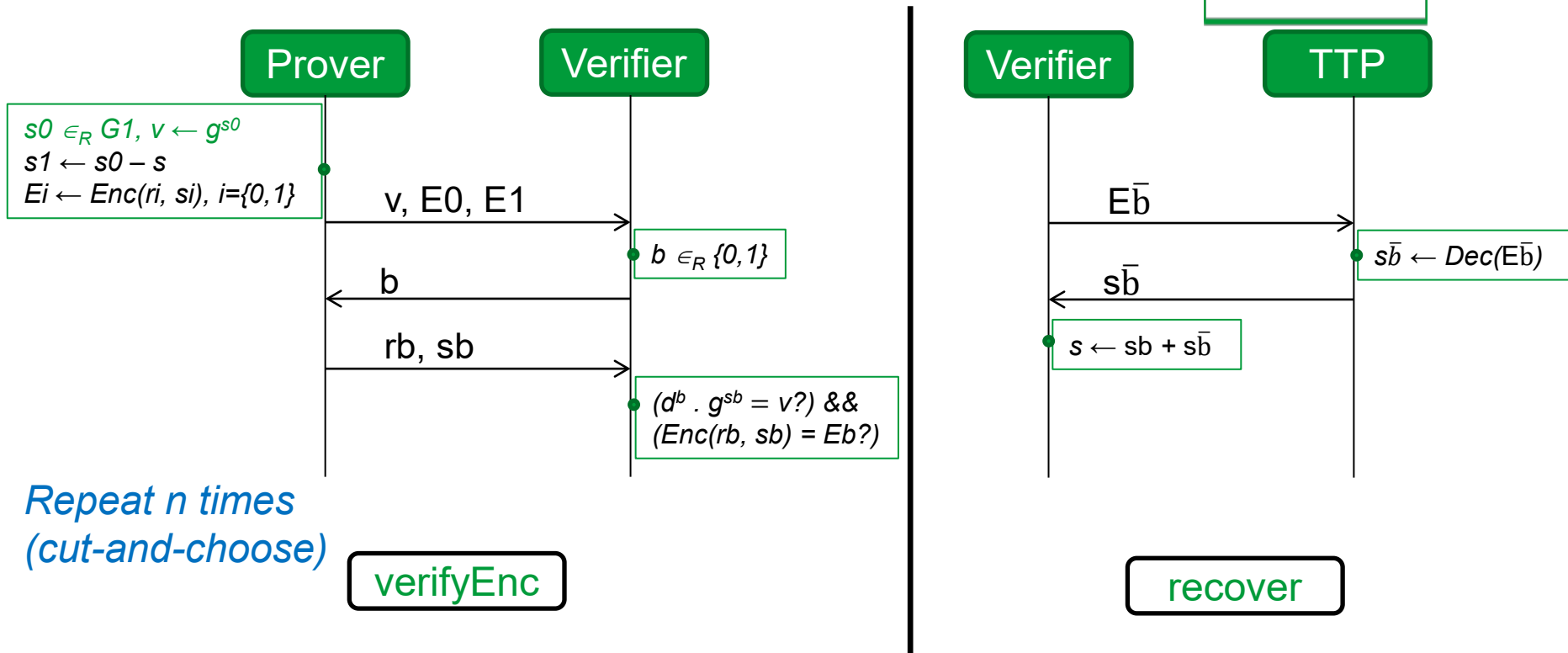
- 
- [ASW97] [“Optimistic Protocols for Fair Exchange”](#), ACM CCS '97
  - [ASW98] [“Asynchronous Protocols for Optimistic Fair Exchange”](#), IEEE S&P '98
  - [ASW00] [“Optimistic Fair Exchange of Digital Signatures”](#), JSAC 18(4): 593-610 (2000)

# Optimistic Fair Exchange: the aftermath

- Someone has to run the Third Party
  - Wants to monetize *every* transaction!

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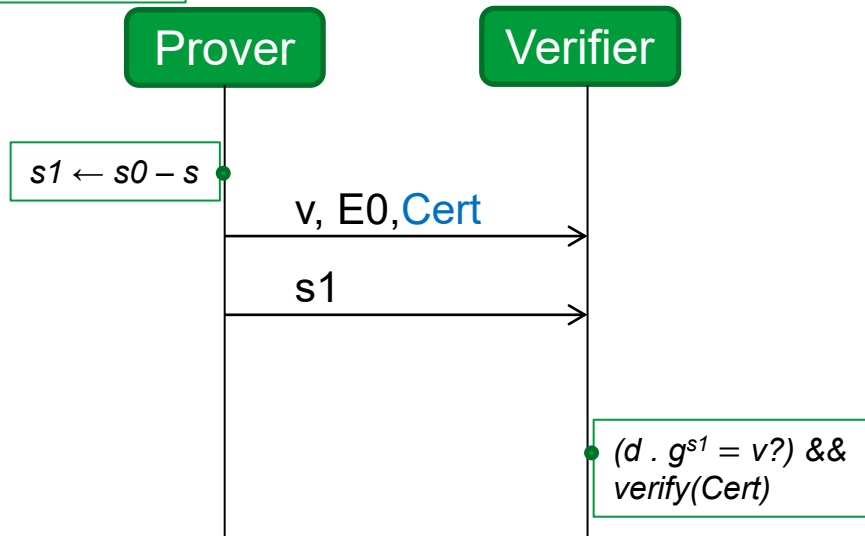


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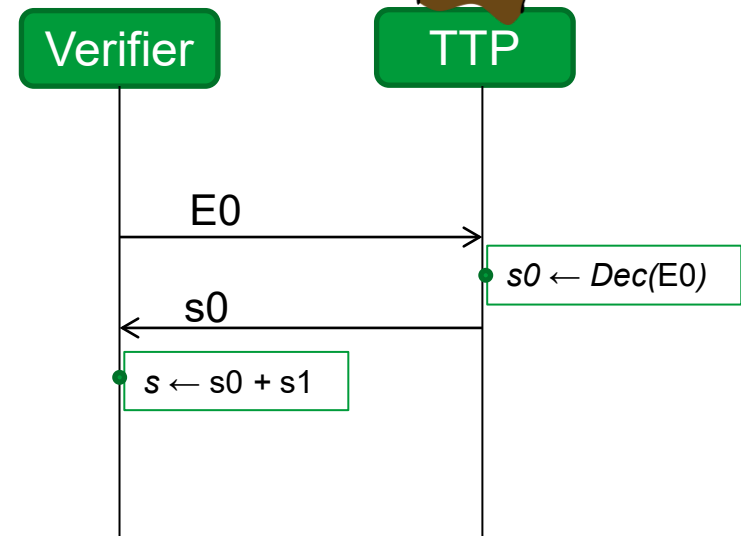
$s_0 \in_R G_1, v \leftarrow g^{s_0}$   
 $E_0 \leftarrow \text{Enc}(r_0, s_0)$   
 $\text{Cert} \leftarrow \text{Sig}_{\text{TTP}}(v, E_0)$

Setting: secret =  $s \in G_1$ , desc  $d = g^s$  (in  $G_2$ )



*Repeat n times  
(cut-and-choose)*

verifyEnc



recover

**Pre-paid coupons bought from the TTP to be used for every optimistic transaction!**

# Optimistic Fair Exchange: the aftermath

- Someone has to run the Third Party
  - Wants to monetize *every* transaction!
- Two decades on, current status:
  - Reputation systems
  - In-line TTP (e.g., E-bay escrow service)





# Continuing “impact” in research circles!

The image displays three overlapping screenshots of Google Scholar search results, illustrating the impact of research in the field of fair exchange. Each screenshot shows search results for a specific query, with key elements highlighted by green circles.

- Left Screenshot:** Search query: "fair exchange". Results: About 14,400 results (circled). Top article: "Optimistic protocols" by N Asokan, M Schunter, et al. (2010), cited by 530.
- Middle Screenshot:** Search query: optimistic "fair exchange". Results: About 3,550 results (circled). Top article: "Optimistic fair exchange" by M Schunter (2000), cited by 44.
- Right Screenshot:** Search query: optimistic "fair exchange". Results: About 105 results (0.06 sec) (circled). Top article: "Ambiguous optimistic fair exchange: Definition and construction" by Q Huang, G Yang, DS Wong, W Susilo (2015), cited by 105. A filter for "Any time" is set to "Since 2015" (circled).

The rightmost screenshot also shows a second article: "Optimistic fair exchange in the enhanced chosen-key model" by Y Wang, MH Au, W Susilo (2015), cited by 105. A third article is partially visible: "How to protect privacy in Optimistic Fair Exchange of digital contracts" by Q Huang, DS Wong, W Susilo (2015), cited by 105.

Autumn 2015

# Continuing “impact” in research circles!

The image displays three overlapping screenshots of Google Scholar search results. The leftmost screenshot shows the search for "fair exchange" with approximately 23,700 results. The middle screenshot shows the search for "optimistic fair exchange" with approximately 5,490 results. The rightmost screenshot shows the search for "optimistic fair exchange" with approximately 154 results, with the "Since 2022" filter highlighted. Each screenshot shows the search query, the number of results, and a list of articles with their titles and authors.

**Search 1: "fair exchange"**  
About 23,700 results  
Articles  
Any time  
Since 2022  
Since 2021  
Since 2018  
Custom range...  
Sort by relevance  
Sort by date  
Any type  
Review articles  
 include patents  
 include citations  
 Create alert

**Search 2: optimistic "fair exchange"**  
About 5,490 results (0.05 seconds)  
Articles  
Any time  
Since 2022  
Since 2021  
Since 2018  
Custom range...  
Sort by relevance  
Sort by date  
Any type  
Review articles  
 include patents  
 include citations  
 Create alert

**Search 3: optimistic "fair exchange"**  
About 154 results (0.14 seconds)  
Articles  
Any time  
**Since 2022**  
Since 2021  
Since 2018  
Custom range...  
Sort by relevance  
Sort by date  
Any type  
Review articles  
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Nov 2022

# Optimistic Fair Exchange: the aftermath

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  - Wants to monetize *every* transaction!
- Two decades on, current status:
  - Reputation systems
  - In-line TTP (e.g., E-bay escrow service)
- Impact in academia vs. real world impact
- Biggest impact of SEMPER?

<http://logging.apache.org/log4j/2.x/>



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# Optimistic Fair Exchange: lessons learned

- Don't just guess security requirements; Ask stakeholders
- Desiderata for deployment and research can be different
  - “the more (independent) parties you require for your scheme, the less likely it will be deployed”
- Capturing researcher interest → (Tech transfer) Impact
- “90-10 rule” applies to deploying security
  - “Good enough beats perfect”


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**Secure Device Pairing**

How can the process of pairing two devices be made easy to use without compromising security or adding to cost?

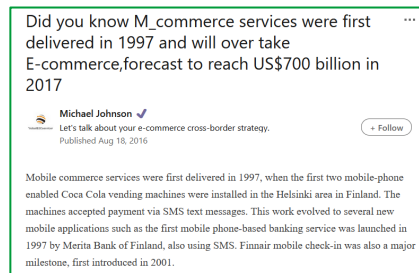
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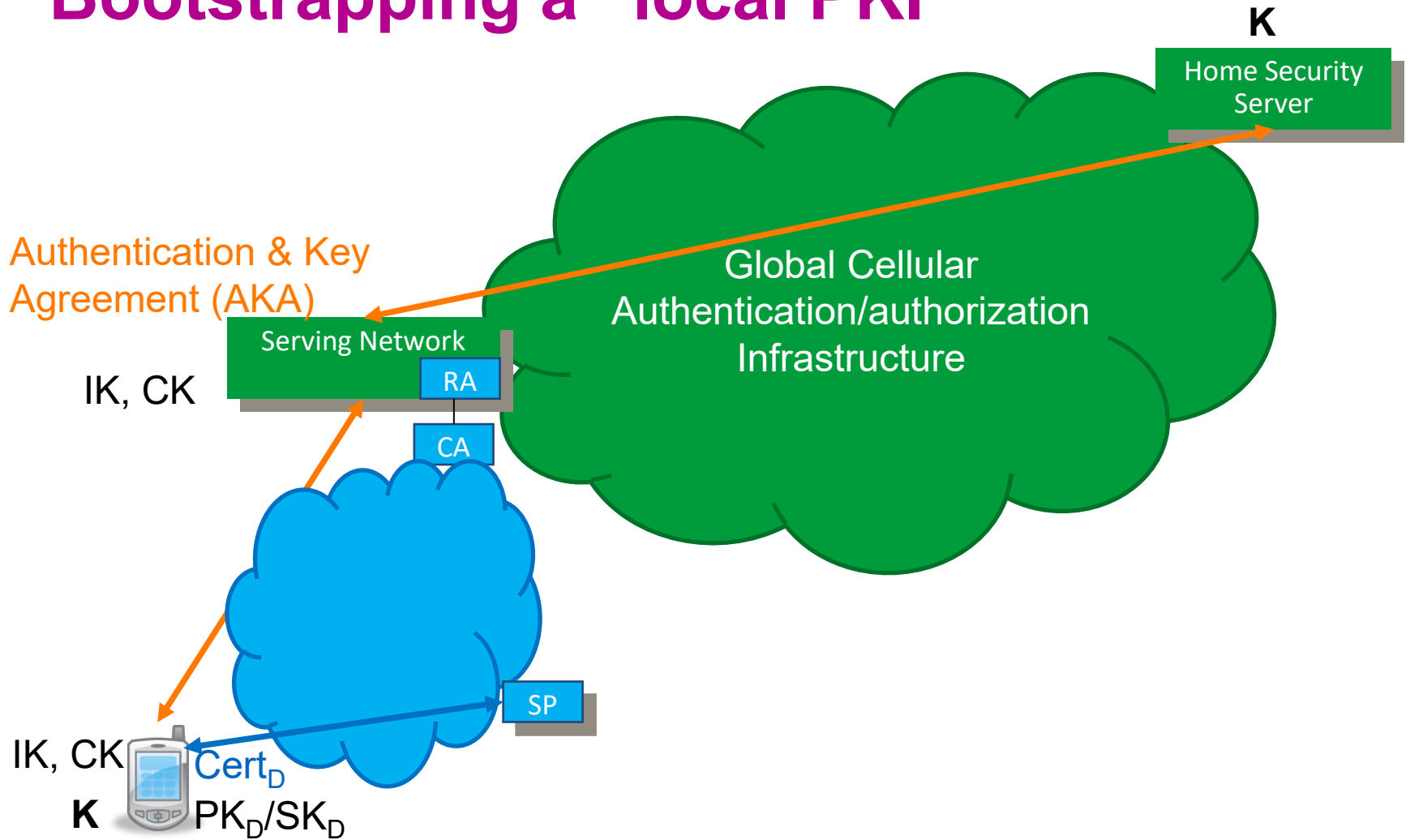
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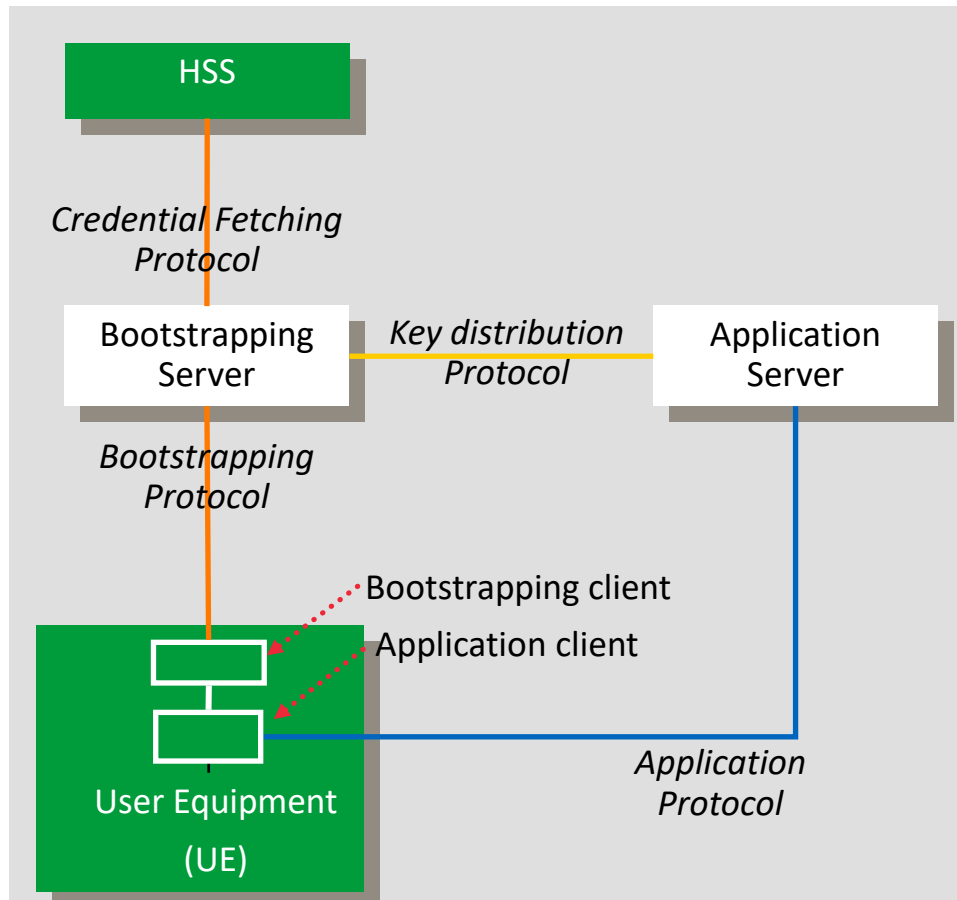


# Bootstrapping a “local PKI”





# 3GPP “Generic Authentication Architecture”



## Two-layer architecture

- Generic Bootstrapping Architecture (GBA)
- Specialized Application Servers
  - E.g., for “subscriber certificates”

[HLGNA08] “[Cellular Authentication for Mobile and Internet Services](#)”, Wiley, 2008  
Relevant 3GPP documents: E.g., [\[33.919\]](#), [\[33.220\]](#)

# GAA: the aftermath

- Standardized in 3GPP
    - Variants: GBA and GBA\_U (implemented in the smartcard, UICC)
    - GBA productized for some services
      - none of which has taken off (e.g., Mobile TV), at least not yet!
    - Revived in 5G “[Authentication and key management for applications](#)” (AKMA)
  - Today’s solutions:
    - Bootstrapping: Facebook, Google, ...
      - Some mobile carriers even deployed PKI-enabled SIM cards
    - Roaming: iPass, Shibboleth, ...
  - Variants of the idea had more success
    - E.g., EAP SIM
-

# GAA: lessons learned

- (Standardization) Politics can suffocate a good idea
- (Tech transfer) Impact  $\nrightarrow$  Capturing researcher interest
- “90-10 rule” applies to deploying security


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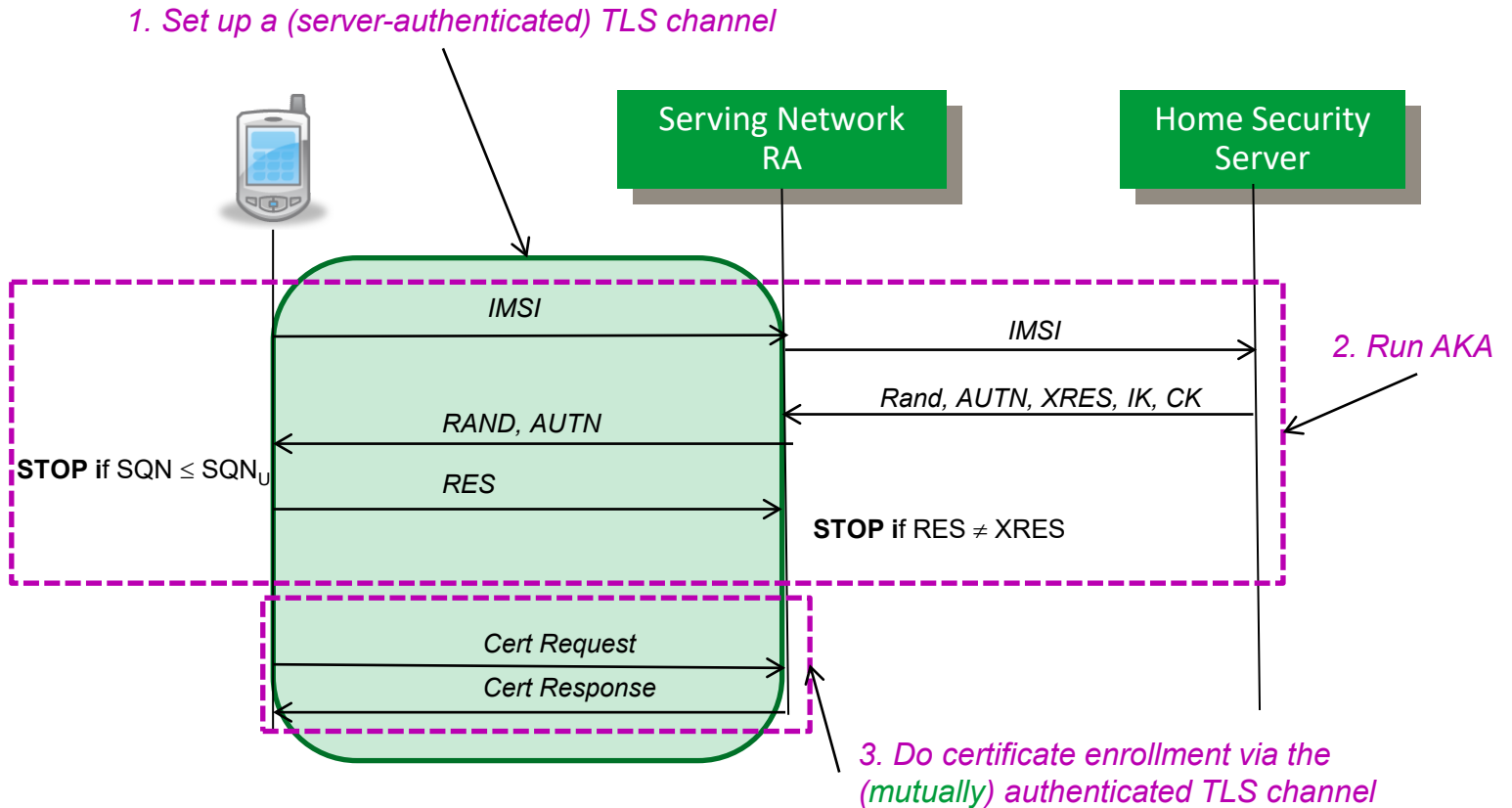
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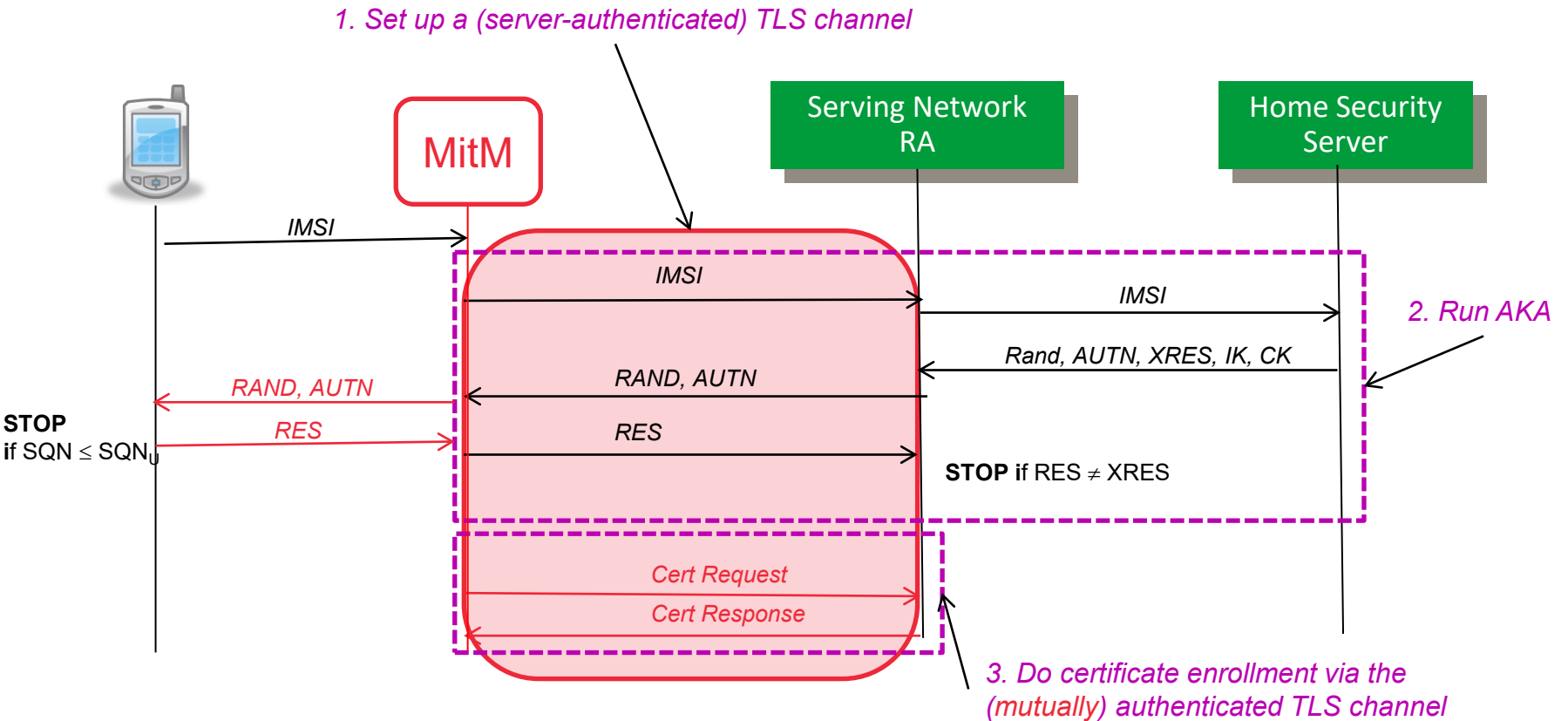
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# Bootstrapping certificate enrollment



# Bootstrapping certificate enrollment



Channel binding: Use of **cryptographic binding** to compose two authenticated channels

[ANN03] "[Man-in-the-middle in Tunnelled Authentication Protocols](#)", Security Protocols, 2003

# Channel binding: the aftermath

- Fiery reception at Security Protocols workshop!
  - “But you are using the worst rackets in industry as a justification for what you’re doing. There are all sorts of people just generating garbage protocols, a couple of which you have already mentioned here. We’re trying to reverse their work, whereas you’re trying to advocate we use all these garbage protocols.”
  - For an entertaining read, see [transcript of discussion during my talk](#) at SPW '03!
- Impact in IETF
  - Closing down of *ipsra* working group; channel binding in IKEv2
  - Continued attention: e.g., [RFC 6813](#)

□ [Man-in-the-middle in tunnelled authentication protocols](#)  
N Asokan, V Niemi, K Nyberg  
International Workshop on Security Protocols, 28-41

345

2003



# Channel Binding: lessons learned

- Negative results are useful for security practitioners
- Standardization can make a good idea see light of day
- (Tech transfer) Impact  $\rightarrow$  Capturing researcher interest


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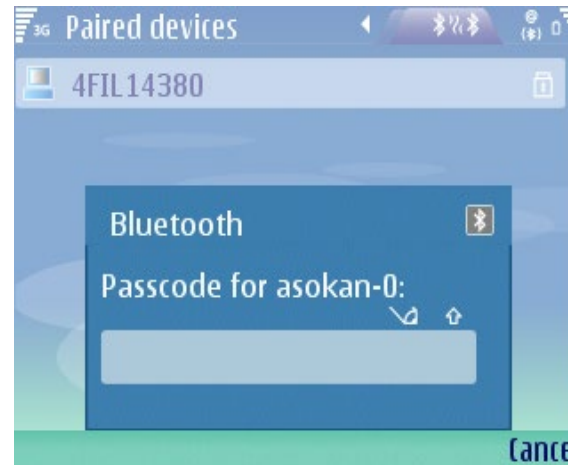
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# Secure Device Pairing: ca. 2005



**Cracking the Bluetooth PIN\***

Yaniv Shaked and Avishai Wool

*School of Electrical Engineering Systems,  
Tel Aviv University, Ramat Aviv 69978, ISRAEL*  
shakedy@eng.tau.ac.il, yaashia@acm.org

**Abstract** This paper describes the implementation of an attack on the Bluetooth security mechanism. Specifically, we de-

new primitives to be risky, because new cryptography is less tested and may contain hidden flaws. Furthermore, Bluetooth is designed for short-range communication (nominal range of about 10m). This short-range is

**Security Weaknesses in Bluetooth**

Markus Jakobsson and Susanne Wetzel

Lucent Technologies - Bell Labs  
Information Sciences Research Center  
Murray Hill, NJ 07974  
USA  
{markusj,sgwetzel}@research.bell-labs.com

**Abstract.** We point to three types of potential vulnerabilities in the Bluetooth standard, version 1.0B. The first vulnerability opens up the system to an attack in which an adversary under certain circumstances is able to determine the key exchanged by two victim devices, making

# Naïve usability measures damage security

http://www.helsinki-hs.net/news.asp?id=20030930IE16

## HELSINGIN SANOMAT INTERNATIONAL EDITION

TODAY

THIS WEEK

WEBORTAGE

THIS IS

Consumer - Tuesday 30.9.2003

### **Pictures taken with mobile phone showed up on neighbour's TV**

► Default password must be changed when starting to use Bluetooth-equipped devices; read the manual!

elsewhere as well. It is, therefore, absolutely essential that the password is changed immediately when the device is first installed."

"This is clearly printed in the user's manual", Rosenberg points out. How often have we heard *that* before?

"Once the digital receiver's password has been changed, the new password also has to be entered in the transmitting device, in this


# Naïve security erodes usability

## Pairing

To create a connection using Bluetooth wireless technology, you must exchange Bluetooth passcodes with the device you are connecting to for the first time for reasons of security. This operation is called pairing. The Bluetooth passcode is a 1- to 16-character numeric code, which you must enter in both devices. You only need this passcode once.


## SIM access mode

In SIM access mode, if the car kit finds a compatible mobile phone that supports the Bluetooth SIM access profile standard, the car kit shows a randomly chosen, 16-character numeric code on the display, which you must enter on the compatible mobile phone to be paired with the car kit. Note that you must be prepared to do this quickly within 30 seconds. Follow the instructions on the display of your mobile phone.

If pairing is successful, Paired with, followed by the name of your mobile phone is displayed. Then Create connection is displayed. Press  to establish the Bluetooth wireless connection.



## Note

When pairing a mobile phone in SIM access mode, a 16-character numeric passcode is generated in the car kit. You can delete this passcode if desired: within 3 seconds, press  to delete the Bluetooth passcode. Then enter an arbitrary 16-character numeric code into the car kit using the Navi wheel number editor.

## Car kits

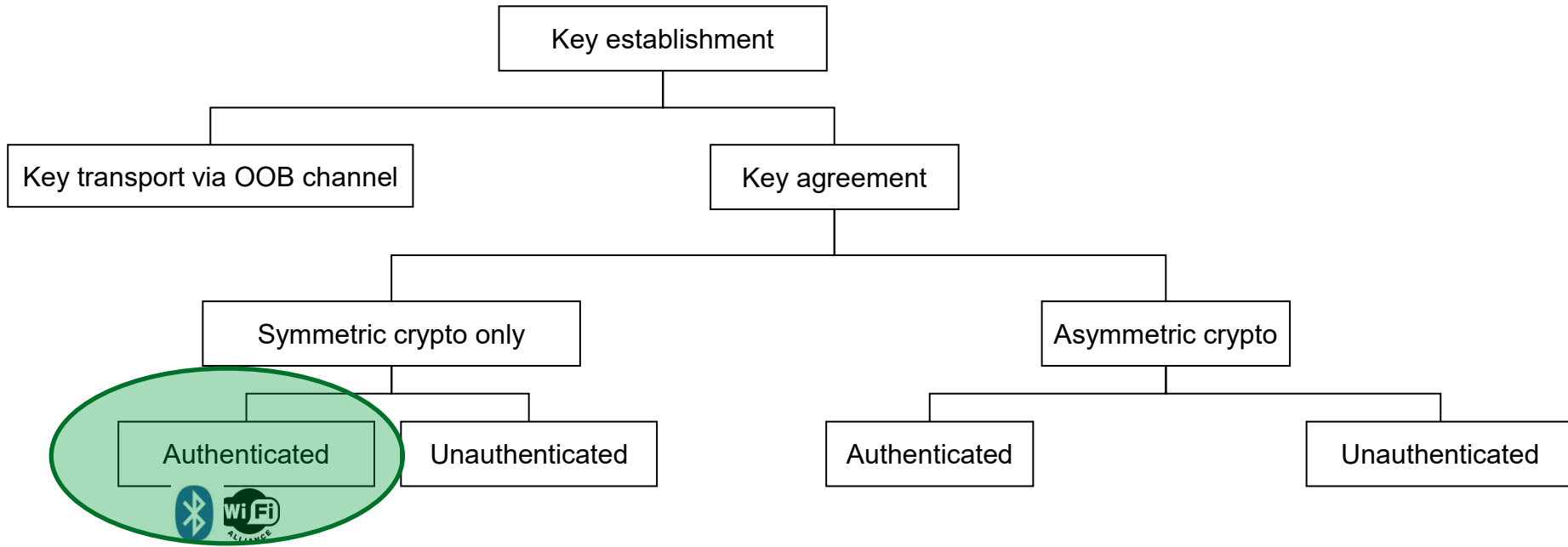
- Allow hands-free phone usage in cars
- Retrieve/use session keys from phone SIM
- require higher level of security

➤ users must enter 16-character passcodes

More secure = Harder to use?

**Cost:**  
Calls to Customer

# Key establishment for secure pairing ~2005

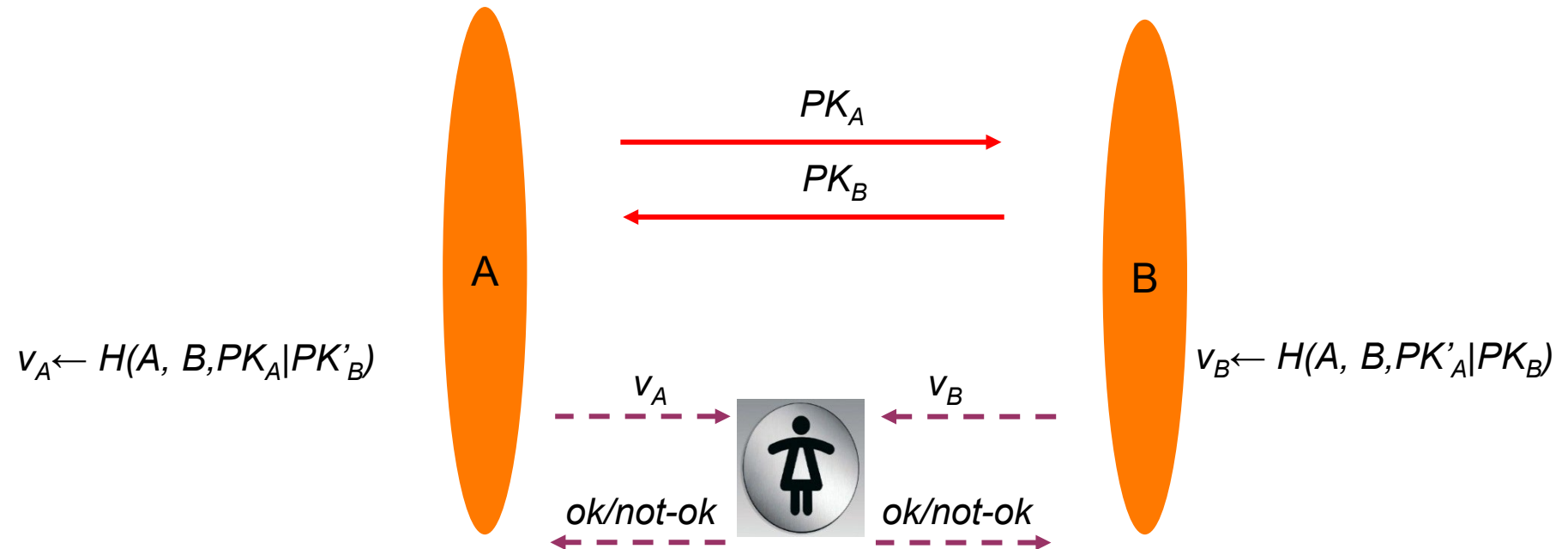


*Short keys vulnerable to passive attackers*

*Secure against passive attackers*



# Authentication by comparing short strings



$v_A$  and  $v_B$  are short strings (e.g., 4 digits),

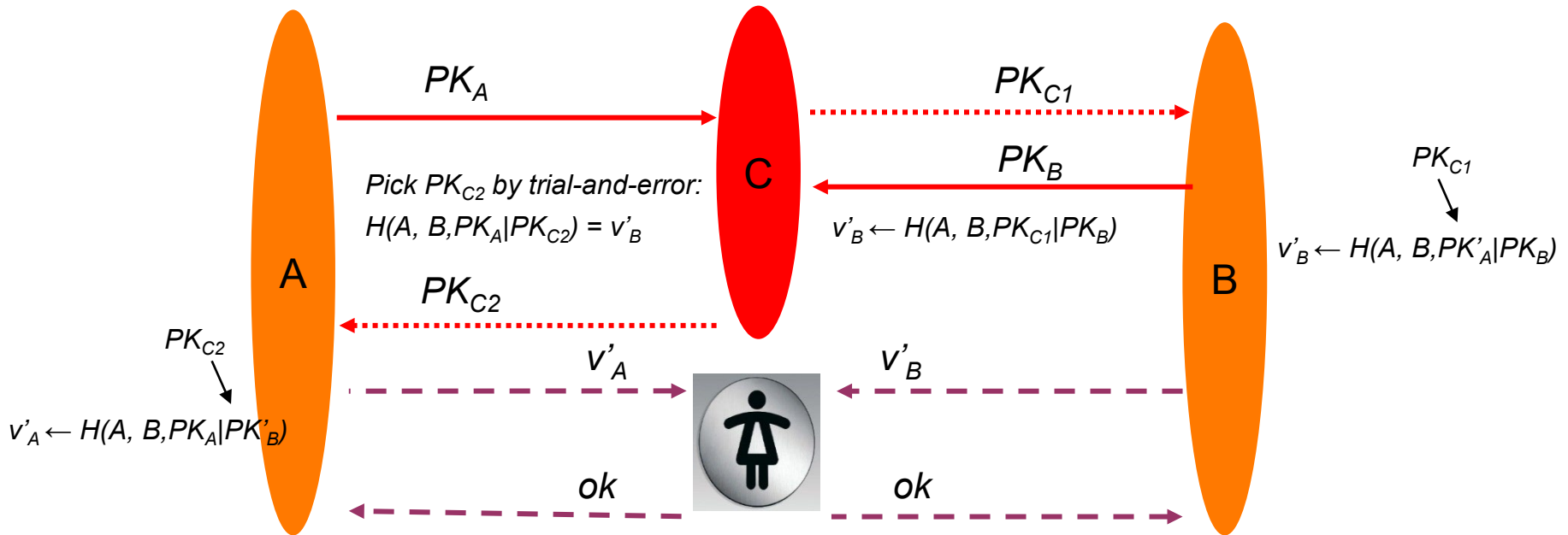
User approves acceptance if  $v_A$  and  $v_B$  match

A man-in-the-middle can easily defeat this protocol

---



# MitM in comparing short strings



Guess a value  $SK_{C2}/PK_{C2}$  until  $H(A, B, PK_A | PK_{C2}) = v'_B$

If  $v'_B$  is n digits, attacker needs at most  $10^n$  guesses; Each guess costs one hash calculation

A typical modern PC can calculate 100000 MACs in 1 second

# Authentication by comparing short strings

Choose long random  $R_A$

Calculate commitment

$$h_A \leftarrow h(A, R_A)$$



key agreement: exchange  $PK_A, PK_B$



Choose long random  $R_B$

Verify commitment

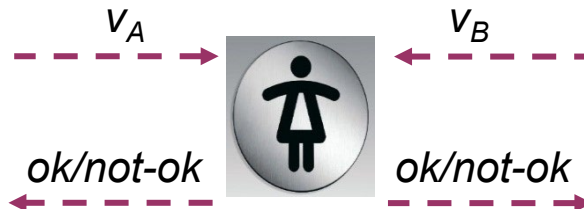
$$h'_A \stackrel{?}{=} h(A, R'_A)$$

Abort on mismatch

$$v_B \leftarrow H(A, B, PK'_A | PK_B, R'_A, R_B)$$



$$v_A \leftarrow H(A, B, PK_A | PK'_B, R_A, R'_B)$$



User approves acceptance if  $v_A$  and  $v_B$  match

$2^{-l}$  (“unconditional”) security against man-in-the-middle ( $l$  is the length of  $v_A$  and  $v_B$ )

$h()$  is a hiding commitment; in practice SHA-256

# Key establishment for secure pairing ~2008

	Unauthenticated Diffie-Hellman	Authenticated Diffie-Hellman		
		short-string comparison	short PIN	Out-of-band channel
WiFi Protected Setup	“Push-button”		√	NFC
Bluetooth 2.1	“Just-works”	√	√	NFC
Wireless USB		√		USB Cable

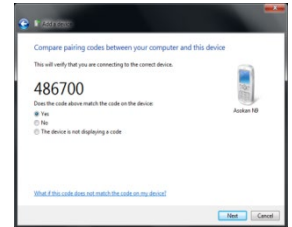
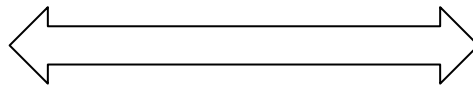
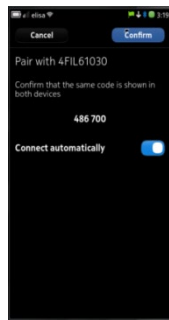
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[AN10] [“Security associations for wireless devices”](#) (Overview, book chapter)

[SVA09] [“Standards for security associations in personal networks: a comparative analysis”](#) IJSN 4(1/2):87-100 (survey of standards)

# Secure Pairing: the aftermath

- Widely deployed (Bluetooth SSP, WiFi Protected Setup)
- **Improving usability/security → fundamental protocol changes**



[UKA07] [“Usability Analysis of Secure Pairing Methods”](#), USEC '07  
[SEKA06] [“Secure device pairing based on a visual channel”](#), IEEE S&P '06

# Secure Device Pairing: lessons learned

- Address pain points - builds credibility with stakeholders
- Don't just guess security requirements; Ask stakeholders
- Desiderata for deployment and research can be different
- Standardization can make a good idea see light of day



- How to choose the "right" problems?
  - Don't just guess security requirements: Ask stakeholders
  - Consider: for deployment and research can be different
  - "60-10 rule" applies to deploying security
- How to identify "good" results?
  - Negative results are useful for security practitioners
  - Capturing researcher interest  $\neq$  (Tech transfer) Impact
  - (Tech transfer) Impact  $\neq$  Capturing researcher interest
- How to find paths to deployment?
  - Address pain points - builds credibility with stakeholders
  - (Standardization) Politics can suffocate a good idea
  - Standardization can make a good idea see light of day

# The remaining examples

- On-board Credentials
  - How to make hardware TEEs safely accessible to developers?  
(Deployments in Nokia devices, but quietly!)


# Five examples

- Optimistic Fair Exchange
- Generic Authentication Architecture
- Channel Binding in Protocol Composition
- Secure Device Pairing
- On-board Credentials

**Fair Exchange**

How can two mutually distrustful parties exchange digital "Items" on the Internet?

Existing solutions:



Gradual Exchange protocols      Trusted Third Party protocols

**Generic Authentication Architecture**

Can we bootstrap a general-purpose global-scale authentication and authorization infrastructure from the existing cellular security infrastructure?

- Need was evident:
  - "Global PKIs will not happen"
- Ad-hoc bootstrapping already in use
  - e.g., Coke vending machine accepting payments via SMS, 1997
- Idea: Bootstrap short-lived certificates from "local PKIs"

**Channel Binding in protocol composition**

Composing two secure authentication protocols carelessly can lead to a man-in-the-middle vulnerability

- Protocol composition can ease deployment
- Examples:
  - Server auth. using TLS + user auth. with password
  - Authentication for VPN access using legacy credentials
  - Bootstrapping a "local PKI"

**Secure Device Pairing**

How can the process of pairing two devices be made easy to use without compromising security or adding to cost?

**On-board Credentials**

Can we safely open up widely deployed secure hardware on mobile devices for use by app developers?

# On-board Credentials

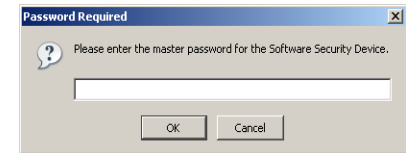
Can we safely open up widely deployed secure hardware on mobile devices for use by app developers?



# Authentication on the Internet

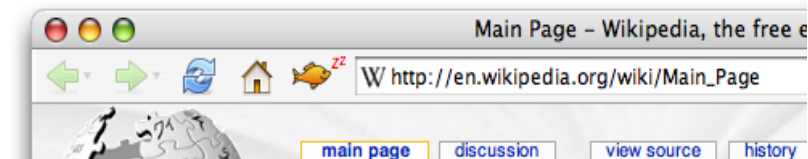
Username/password rules the Internet

- Cheap, easy-to-deploy, portable
- Annoying, vulnerable (phishing, dictionary attacks, password-stealing trojans...)



Attempts to improve usability and security

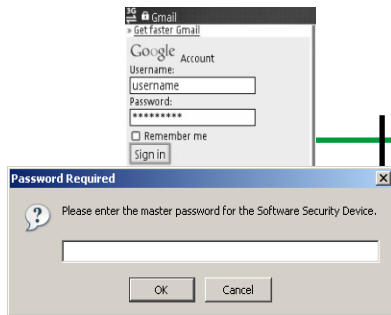
- Password-managers
- Single Sign-On
- Better protocols



# Hardware tokens

## Deployed for specific-services

- More secure, sometimes more intuitive
- More expensive, usually no trusted path to user,
- Single-purpose or issuer-controlled



SW-only credentials



HW credentials

# Trusted hardware is widely deployed

- Trusted Execution Environments on smartphones have been available for years
  - Introduced for manufacturer and operator needs
  - Not accessible for app developers



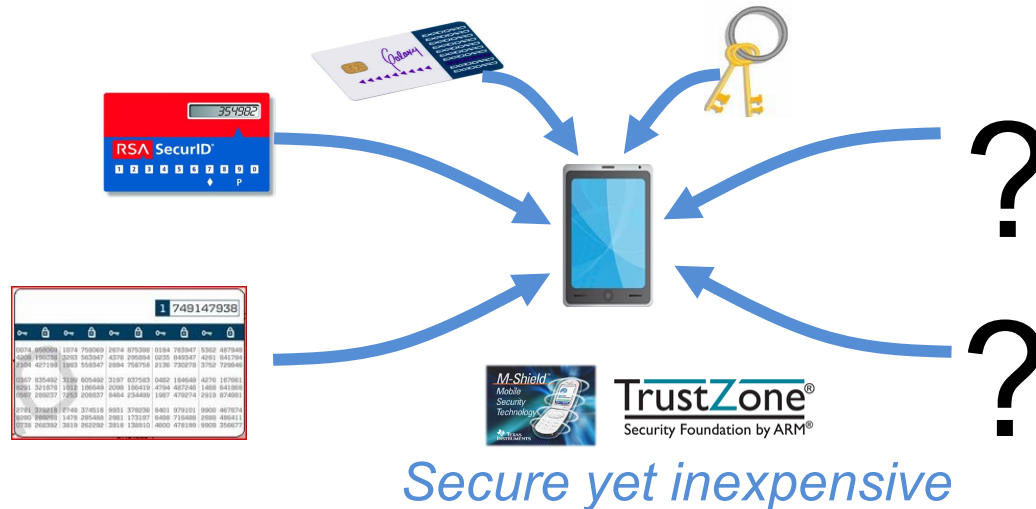
# On-board Credentials

**ObC: the aftermath**

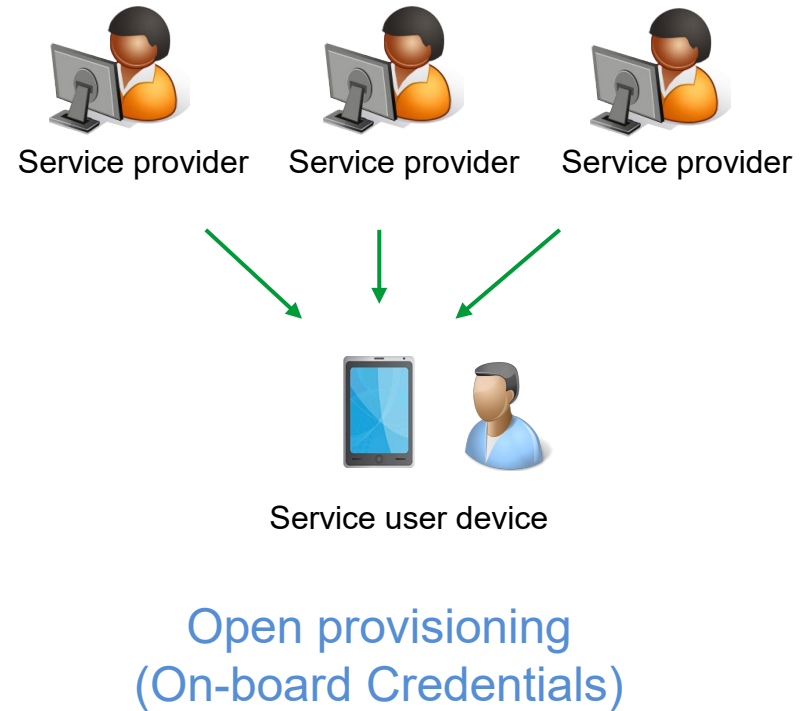
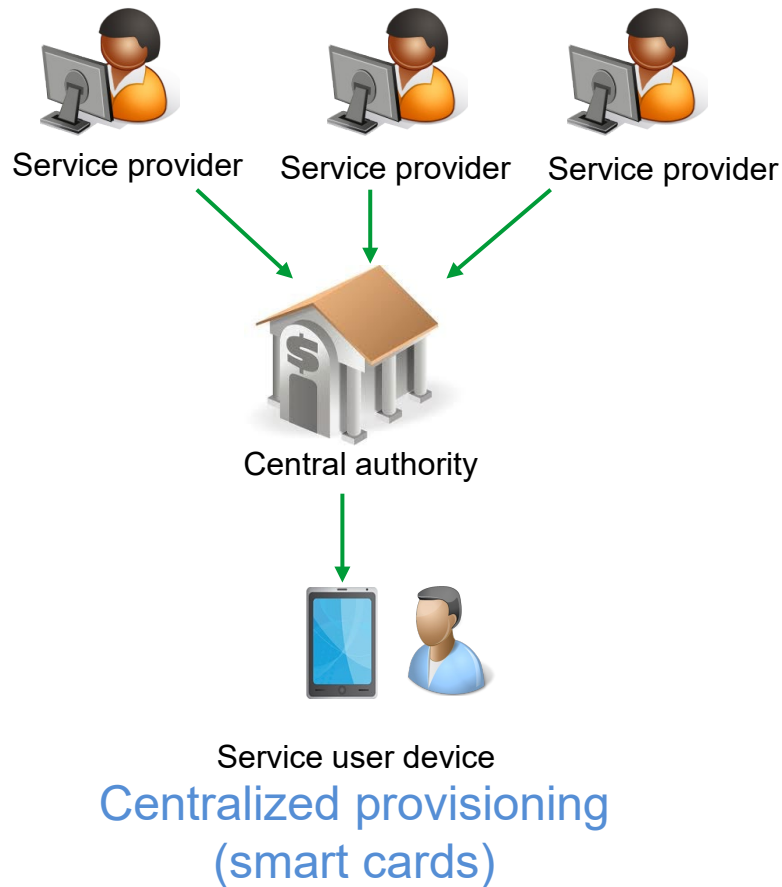
- Initial prototypes ca. 2008
  - RSA SecurID, SoftSIM
- (Sliently) deployed in recent Lumia devices
  - Used for, e.g. [MirrorLink](#), attestation, LIRR ticketing trial
- Stumbling blocks:
  - "who takes liability?" "avoid stepping on toes"
- Related standardization
  - Global Platform device committee
  - Open provisioning is elusive

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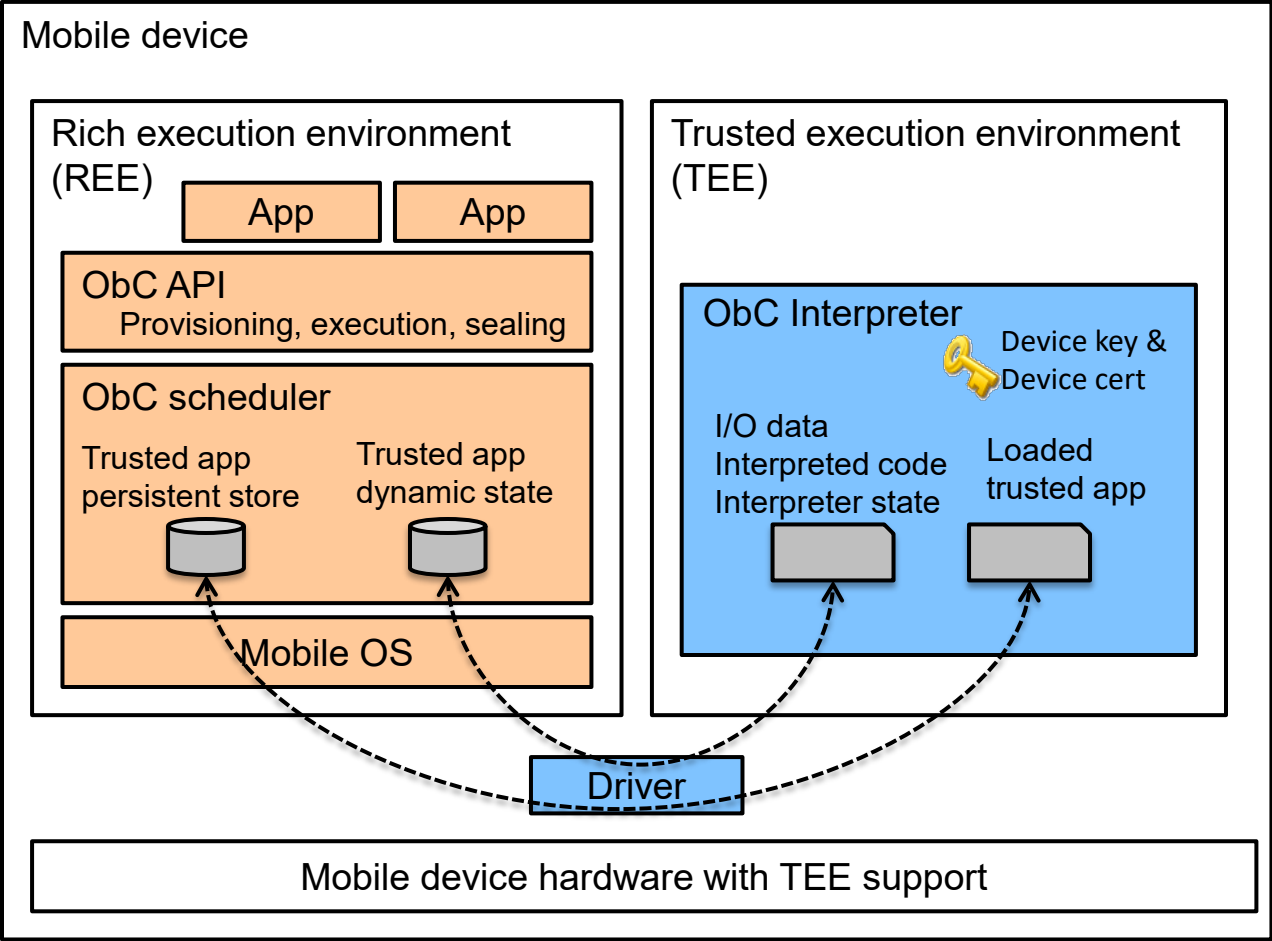
An **open** credential platform that leverages existing mobile TEEs



# Centralized vs. open provisioning

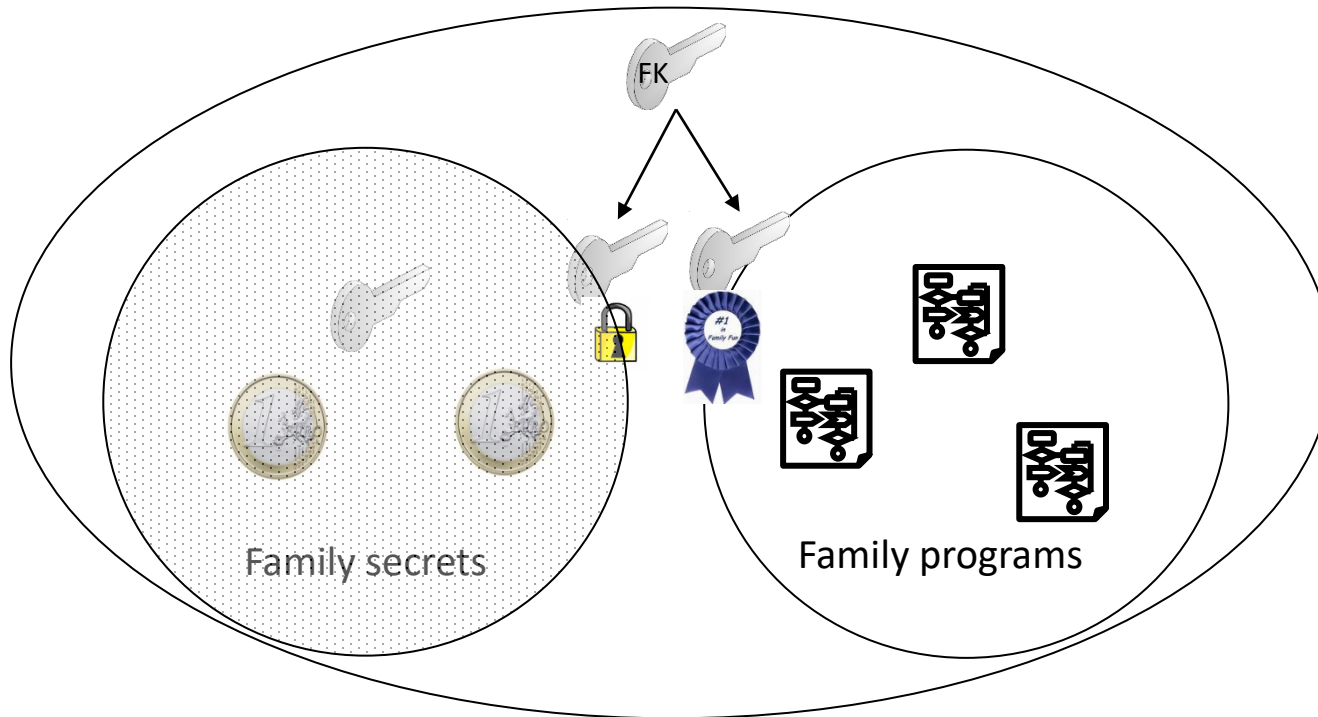


# On-board Credentials (ObC) architecture



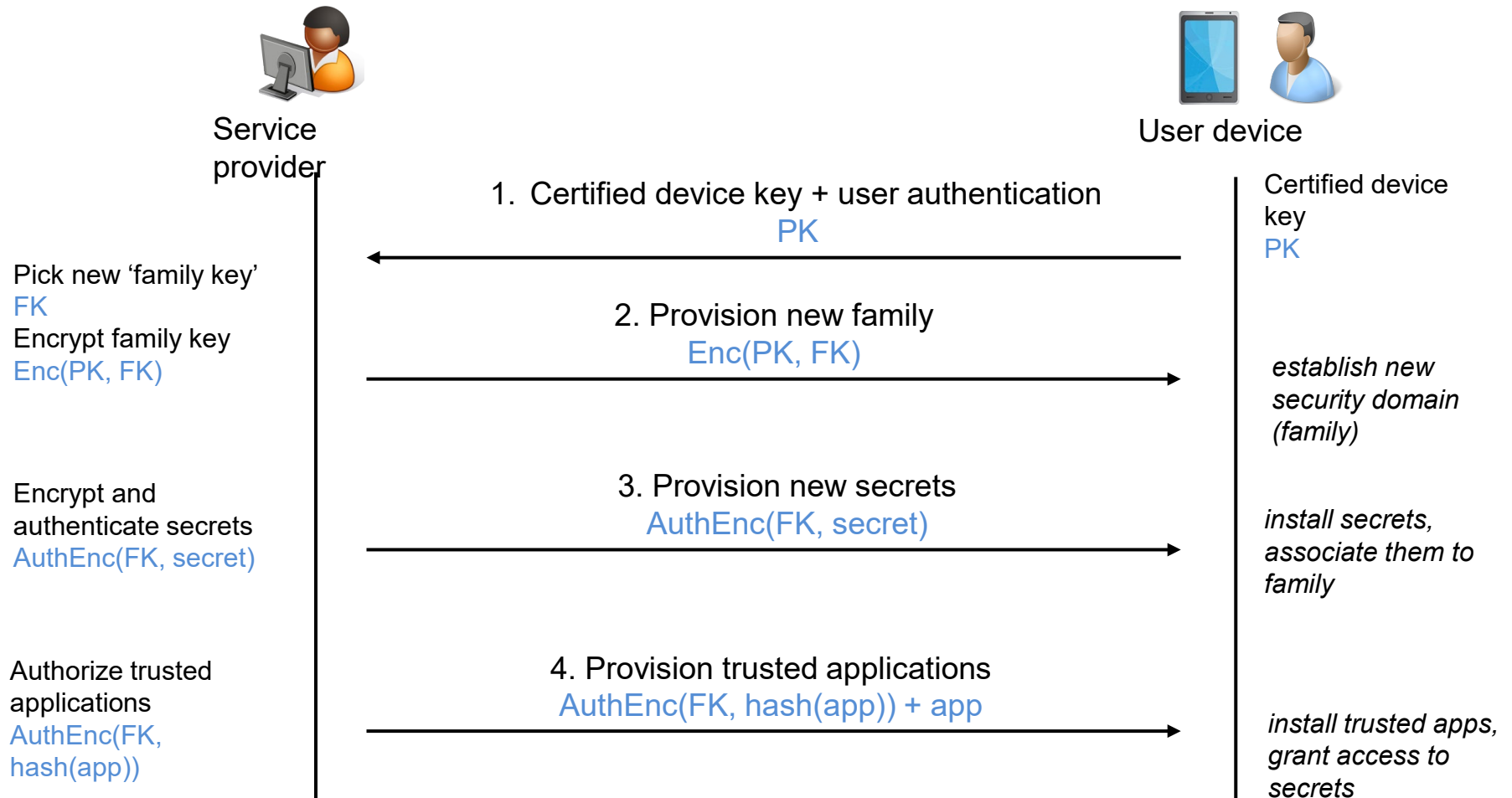
# ObC Provisioning (1/2)

Basic Idea: the notion of a **family** of credential secrets and credential programs endorsed to use them



Principle of same-origin policy

# Open provisioning model



[KEAR09] ["On-board Credentials with Open Provisioning"](#). ASIACCS 2009.

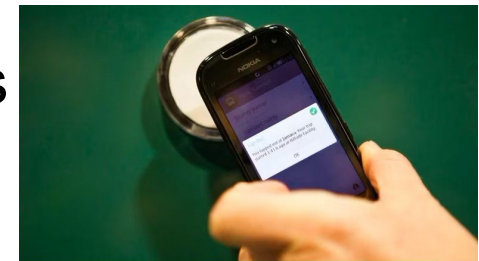
Ekberg. [Securing Software Architectures for Trusted Processor Environments](#). Dissertation, Aalto University 2013.

Kostiainen. [On-board Credentials: An Open Credential Platform for Mobile Devices](#). Dissertation, Aalto University 2012.

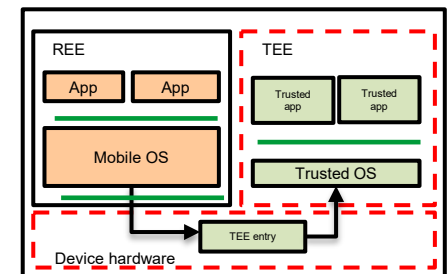


# ObC: the aftermath

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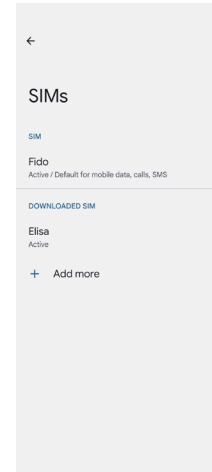
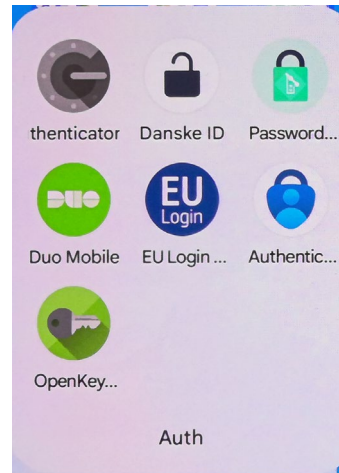
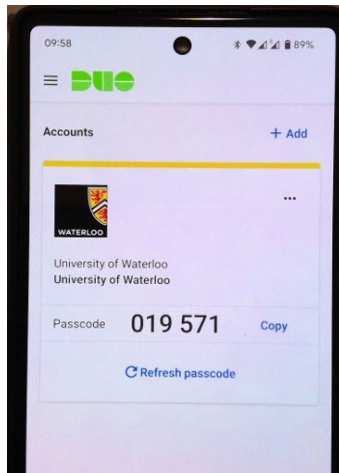
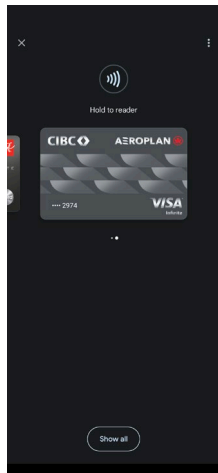
<https://www.newsday.com/long-island/transportation/lirr-tests-smartphone-payment-system-u04362>



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[GP12] [“A New Model: The Consumer-Centric Model and How It Applies to the Mobile Ecosystem”](#)

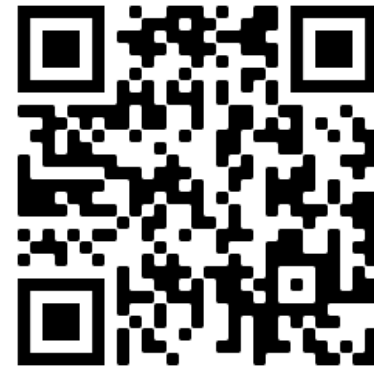
# “On-board Credentials” on my phone



# ObC: Lessons Learned

- Address pain points - builds credibility with stakeholders
- Politics can suffocate a good idea
- Standardization can make a good idea see light of day
- (Tech transfer) Impact → Capturing researcher interest

# Lessons Learned



<https://asokan.org/asokan/research>

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    - Don’t just guess security requirements; Ask stakeholders
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