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User Guide for SPARSPAK:
Waterloo Sparse Linear Equations
Package

Alan George, Joseph Liu, Esmond Ng

Research Report CS-78-30
(Revised, Jan. 1980)

Department of Computer Science

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ABSTRACT

SPARSPAK USER GUIDE

This document describes the structure and use of SPARSPAK, the Waterloo Sparse Linear Equations Package, which is designed to efficiently solve large sparse systems of linear equations. Computer programs for solving sparse systems of linear equations typically involve fairly complicated data structures and storage management. In many cases the user of such programs simply wants to solve his problem, and should not have to understand how the storage management is done, or how the matrix components are actually stored. One of the attractive features of this package is that it effectively insulates the user from these considerations, while still allowing the package to be used in a variety of ways. Another important feature of the package is the provision of a variety of methods for solving sparse systems, along with convenient means by which the best method for a given problem can be selected.

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SECTION 1

INTRODUCTION AND BASIC STRUCTURE OF SPARSPAK

SPARSPAK offers a collection of methods for solving sparse systems of linear equations

$$A x = b,$$

where A is an N by N nonsingular matrix, and x and b are vectors of length N . We assume the user is aware of the basic issues involved in solving sparse matrix equations, and the basic facts about solving systems of linear equations using Gaussian elimination. For a discussion on the initial design of this package, see [5].

For all the methods provided in SPARSPAK, the user and the package interact to solve the matrix problem through the following basic steps:

- Step 1. The user supplies the nonzero structure of A to the package using a set of subroutines described in Section 2.2.
- Step 2. The package reorders the original problem (finds a permutation P), and allocates storage for the triangular factorization of $PAP' = LU$, as described in Section 2.3 (1).
- Step 3. The user supplies the numerical values for the matrix A to the package, as described in Section 2.4.
- Step 4. The package computes the triangular factors L and U of PAP' , as described in Section 2.5.
- Step 5. The user supplies numerical values for b , as described in Section 2.4. (This step may come before Step 4, and may be intermixed with Step 3.)
- Step 6. The package computes the solution x , using L , U , P and b , as described in Section 2.5.

The different methods provided in SPARSPAK correspond to different algorithms for choosing P (along with appropriate storage methods), and whether or not A is symmetric. When A is symmetric, U is replaced by L' in the above description, and of course only one of L and L' is stored.

1 P' stands for the transpose of the matrix P .

The user chooses a particular method by calling the appropriate subroutines in Steps 2, 3, 4 and 6. The methods are distinguished by a numerical digit i , $1 \leq i \leq 6$, which is the last character of the subroutine names. The subroutines used in Steps 1 and 5 apply to all the methods. The best method to use depends very much on the particular problem, and the context in which it is being solved, so we cannot provide rigid rules as to which method to use. Some guidelines and considerations regarding the choice of method are given in Section 3.

RESTRICTIONS AND ASSUMPTIONS

1. SPARSPAK assumes that the nonzero structure of A is symmetric. If this is not the case, the package will still work, but if A has highly unsymmetric structure, this may lead to some inefficiencies because the matrix will be treated as though its structure is that of $A + A'$. The diagonal elements of A are assumed to be nonzero.
2. SPARSPAK assumes that for any permutation matrix P , Gaussian elimination applied to PAP' without row or column interchanges yields an acceptably accurate factorization LU . In other words, the package assumes that A can be symmetrically permuted without regard for numerical stability. This is true, for example, when A is symmetric and positive definite, or diagonally dominant.

SECTION 2

MODULES OF SPARSPAK AND HOW TO USE THEM

2.1 USER MAINLINE PROGRAM AND AN EXAMPLE

SPARSPAK allocates all its storage from a single one dimensional floating point array⁽²⁾ which for purposes of discussion we will denote by S . In addition, the user must provide its size $MAXS$, which is transmitted to the package via a common block $SPKUSR$, (SPARSPAK USER), which has four variables:

```
COMMON /SPKUSR/ MSGLVL, IERR, MAXS, NEQNS
```

Here $MSGLVL$ is the message level indicator which is used to control the amount of information printed by the package. The second variable $IERR$ is an error code, which the user can examine in his mainline program for possible errors detected by the package. Detailed discussion of the roles of $MSGLVL$ and $IERR$ is provided in Section 7. The variable $NEQNS$ is the number of equations.

The following program illustrates how one might use SPARSPAK. The various subroutines referenced are described in the subsequent parts of this section. The problem that is solved is a 10 by 10 symmetric tridiagonal system $Ax = b$ where the diagonal elements of A are all 4, the superdiagonal and subdiagonal elements are all -1, and the entries in the right hand side vector b are all ones.

² Declared either REAL or DOUBLE PRECISION, depending on the version of SPARSPAK that is available. The examples in this manual assume a single precision version is being used.

```

REAL      S(250), FOUR, ONE
INTEGER   I, IERR, MAXS, MSGVLV, NEQNS
COMMON   /SPKUSR/ MSGLVI, IERR, MAXS, NEQNS
C
      CALL  SPRSPK
      MAXS = 250
C-----
C      INPUT THE MATRIX STRUCTURE.  THE DIAGONAL IS
C      ALWAYS ASSUMED TO BE NONZERO, AND SINCE THE
C      MATRIX IS SYMMETRIC, SPABSPAK ONLY NEEDS TO
C      KNOW THAT THE SUBDIAGONAL ELEMENTS ARE NONZERO.
C-----
      CALL  IJBEGN
      DO 100 I = 2, 10
          CALL INIJ ( I, I-1, S )
100     CONTINUE
      CALL  IJEND ( S )
C-----
C      FIND THE ORDERING AND ALLOCATE STORAGE.
C-----
      CALL  CRDRA1 ( S )
C-----
C      INPUT THE NUMERICAL VALUES FOR A AND B.  SINCE
C      THE MATRIX IS SYMMETRIC, ONLY THE LOWER TRIANGLE
C      AND THE DIAGONAL ARE INPUT.
C-----
      FOUR = 4.0E0
      ONE  = 1.0E0
      DO 200 I = 1, 10
          IF ( I .GT. 1 )
*           CALL INAIJ1 ( I, I-1, (-ONE), S )
          CALL INAIJ1 ( I, I, FOUR, S )
          CALL INBI ( I, ONE, S )
200     CONTINUE
C-----
C      SOLVE THE SYSTEM.
C-----
      CALL  SOLVE1 ( S )
C-----
C      PRINT THE SOLUTION, FOUND IN THE FIRST TEN
C      LOCATIONS OF THE WORKING STORAGE ARRAY S.
C-----
      WRITE (6, 11) (S(I), I = 1, 10)
11     FORMAT ( / 10H SOLUTION / (5F12.5) )
C-----
C      PRINT SOME STATISTICS GATHER BY SPABSPAK.
C-----
      CALL  PSTATS
C
      STOP
      END

```

NOTE: If the SPARSPAK available to you is a double precision version, the REAL declaration in this example should be changed to DOUBLE PRECISION.

The module SPRSPK must be called before any part of the package is used. Its role is to initialize some system parameters (e.g. the logical unit numbers for output files), and to set default values for options (e.g. initializing the timing routine). The routine needs only to be called once in the user program, and the FCRTAN statement is simply

CALL SPRSPK

Note that the only variable in the common block SPKUSR that must be explicitly assigned a value by the user is MAXS.

It is assumed that the subroutines which comprise SPARSPAK have been compiled into a library, and that the user can reference them from his FORTRAN program just as he references the standard FORTRAN library subroutines, such as SIN, COS, etc. Normally, a user will use only a small fraction of the subroutines provided in SPARSPAK.

WARNING

The modules of SPARSPAK communicate with each other through labelled common blocks whose names are SPKUSR, SPKSYS, SPKCON, SPKMAF, SPKDTA, and SPKOPS. Thus, the user must not use labelled common blocks with these names in his program.

If these common block names cause conflicts in your program or at your computer installation, it is possible to have the package distributed with these common blocks having specifically requested labels. These names should be specified when the package is acquired.

2.2 MODULES FOR INPUT OF THE MATRIX STRUCTURE

SPARSPAK has to know the matrix structure before it can determine an appropriate ordering for the system. We now describe the group of routines which provide a variety of ways through which the user can inform the package where the nonzero entries are; that is, those subscripts (i, j) for which the (i, j)-th element of A is nonzero. Before any of these input routines is called, the user must execute an initialization routine called IJBEGN, which tells the package that the structure of a new matrix problem is about to be input:

CALL IJBEGN

a) Input of a nonzero location.

To tell the package that the (i, j)-th element of A is nonzero, the user simply executes the statement

```
CALL INIJ ( I, J, S )
```

where I and J are the subscripts of the nonzero, and S is the working storage array declared by the user for use by the package.

In this example,

```
.  
. .  
. .  
I = 4  
J = 3  
CALL INIJ ( I, J, S )  
. .  
. .
```

the package will record a logical nonzero in the position (4, 3) of the matrix.

b) Input of the structure of a row, or part of a row.

When the structure of a row or part of a row is available, it is more efficient to use the routine INROW. The statement to use is

```
CALL INROW ( I, NIR, IR, S )
```

where I denotes the subscript of the row under consideration, IR is an array containing the column subscripts of some or all of the nonzeros in the I-th row, NIR is the number of subscripts in IR, and S is the user-declared working storage.

For example, in

```
.  
. .  
I = 5  
IR(1) = 2  
IR(2) = 7  
IR(3) = 5  
CALL INRCW ( I, 3, IR, S )  
. .  
.
```

the package is informed of nonzeros in locations (5, 2), (5, 5) and (5, 7) of the matrix. Note that the subscripts in the array IR can be in arbitrary order, and the rows can be input in any order.

c) Input of a submatrix structure.

To provide greater flexibility, the package allows the user to input the structure of a submatrix. The calling statement is

```
CALL INIJIJ ( NIJ, II, JJ, S )
```

where NIJ is the number of input subscript pairs, and II and JJ are the arrays containing the subscripts.

The following example

```
.  
. .  
II(1) = 1  
JJ(1) = 1  
II(2) = 1  
JJ(2) = 3  
II(3) = 2  
JJ(3) = 3  
CALL INIJIJ ( 3, II, JJ, S )  
. .  
.
```

informs the package that there are nonzeros in locations (1, 1), (1, 3) and (2, 3).

d) Input of a full submatrix structure.

The structure of an entire matrix is completely specified if all the full submatrices are given. In applications where they are readily available, the routine INCLQ is useful. Its calling sequence is

```
CALL INCLQ ( NCLQ, CLQ, S )
```

where NCLQ is the size of the submatrix and CLQ is an array containing the subscripts of the submatrix.

Thus, to inform the package that the submatrix corresponding to subscripts 1, 3, 5 and 6 is full, we execute

```
.  
. .  
CLQ(1) = 1  
CLQ(2) = 3  
CLQ(3) = 5  
CLQ(4) = 6  
CALL INCLQ ( 4, CLQ, S )  
. . .
```

The type of structure input routine to use depends on how the user obtains the matrix structure. Anyway, the one or ones that best suit the application can be selected; SPARSPAK allows mixed use of the routines in inputting a matrix structure. The package automatically removes duplications so the user does not have to worry about inputting duplicated subscript pairs.

```
.  
. .  
CLQ(1) = 1  
CLQ(2) = 2  
CLQ(3) = 5  
CALL INCLQ ( 3, CLQ, S )  
IR(1) = 2  
IR(2) = 4  
IR(3) = 5  
CALL INROW ( 4, 3, IR, S )  
CALL INIJ ( 1, 3, S )  
. . .
```

The above code would input the matrix structure

```
[ * * * * ]  
[ * * * * ]  
[ * * * * ]  
[ * * * * ]  
[ * * * * ]
```

into the package.

When all pairs have been input, using one or a combination of the input routines, the user is required to tell SPARSPAK explicitly that structure input is complete by calling the routine IJEND. The statement to use is

```
CALL IJEND ( S )
```

and its purpose is to transform the data from the format used during the recording phase to the standard format used by the later phases. The user does not have to concern himself with this representation or transformation.

IMPORTANT NOTE:

SPARSPAK assumes that the value of NEQNS (the number of equations) is equal to the maximum subscript supplied by the routines which transmit the (i, j) pairs to the package. Thus, it is imperative that the user supplies at least one (i, j) pair for which i or j is equal to NEQNS. The routine IJEND assigns the value of NEQNS found by the package to the corresponding variable in the common block SPKUSR .

Common Errors

The most common cause of error during matrix structure input is insufficient working storage. If we denote the number of offdiagonal nonzeros in the matrix by OFFDA, then the minimum amount of storage necessary to successfully input the structure is given by

$$\text{OFFDA} + 2 * \text{NEQNS} + 1 .$$

Of course sometimes the user does not know the value of OFFDA, and may guess too low. SPARSPAK will still accept and count the (i, j) pairs, even after running out of storage, and the user can obtain an upper bound for OFFDA by calling the module PSTATS, described in Section 7, after all pairs have been input. (The number

reported may be unnecessarily large because duplicate input pairs may not now be detected, and thus may be counted twice by the package.)

For a complete list of errors which may be generated by the structure input modules, see Section 7.3.1.

2.3 MODULES FOR ORDERING AND STORAGE ALLOCATION

With an internal representation of the nonzero structure of the matrix A available, SPARSPAK is ready to reorder the matrix problem. This is initiated by calling an ordering routine, whose name always has the form $ORDRxi$. Here i is a numerical digit between 1 and 6 that signifies the storage method. The character x can take values A or B , which denotes one of two ordering strategies tailored for storage method i .

Executing the statement

```
CALL CREFA1 ( S )
```

will imply the use of storage method 1 and the first ordering algorithm for this method. See Section 3 for a discussion of the various methods provided, and some guidance on which one to use. Section 8 contains a list of ordering strategies provided by the package. The routine $ORDRxi$ not only determines an appropriate ordering for the storage method, it sets up the data structure for the reordered matrix problem. The package is now ready for numerical input.

Common Errors

Just as in the structure input phase, the most common cause of abnormal termination of the $ORDRxi$ module is insufficient working storage. As mentioned above, this module actually performs two functions: ordering, and storage allocation. The ordering step determines the permutation P , and the allocation step sets up the appropriate data structures to store the triangular factors L and U of the permuted matrix PAP' .

In general, the ordering and allocation subroutines require different amounts of storage. Furthermore, their storage requirements are often unpredictable, because the number of data structure pointers, and the number of nonzeros in the factors L and U , are not known until the subroutines have been executed.

Thus, the interface module ORDRxi may terminate in several distinctly different ways:

- a) There was not enough storage to execute the ordering subroutine.
- b) The ordering was successfully obtained, but there was insufficient storage to even initiate execution of the data structure set-up (storage allocation) subroutine.
- c) The data structure set-up subroutine was executed, and the amount of storage required for the data structure pointers etc. was determined, but there was insufficient storage for these pointers.
- d) The data structure was successfully generated, but there is insufficient storage for the actual numerical values, so the next step (input of the numerical values) cannot be executed.
- e) ORDRxi was successfully executed, and there is sufficient storage to proceed to the next step.

If any of the above conditions occurs, the user may execute SAVE, and re-initiate the computation after adjusting his storage declarations (either up or down) and executing RESTRT³). If a) or b) occurs, information is supplied indicating the minimum value of MAXS needed so that c), d) or e) will occur upon re-execution. If c) occurs, the minimum value of MAXS needed for d) and e) is provided.

When c) or d) occurs, after executing SAVE, adjusting our storage declaration, then executing RESTRT, we must again call ORDRxi. However, the interface will detect that the ordering and/or storage allocation have already been performed, and will skip that part of the computation. Note that if a user is simply using SPARSPAK to select a particular method, c) may be an acceptable termination state. (See Example 6 in Section 9.)

2.4 MODULES FOR INPUTTING NUMERICAL VALUES

The modules in this group are similar to those for inputting the matrix structure. They provide a means of transmitting the actual numerical values of the matrix problem to SPARSPAK. Since the data structures for different storage methods are different, the package must have a different matrix input subroutine for each method. For the user's convenience, SPARSPAK uses the same set of

³ See Section 4 for details on how to use SAVE and RESTRT, and Examples 4, 5 and 6 in Section 9.

subroutine names for all the methods, except for the last digit which distinguishes the method, and the parameter lists for all the methods are the same.

IMPORTANT NOTE:

The elements of A and b transmitted to SPARSPAK by these routines are either single or double precision floating point numbers, depending on the version of SPARSPAK being used. The examples in this manual assume a single precision version of the package is being used.

There are three ways of passing the numerical values to SPARSPAK. In all of them, subscripts passed to the package always refer to those of the original given problem. The user need not be concerned about the various permutations to the problem which may have occurred during the ordering step.

a) Input of a single nonzero component.

The subroutine INAIJi is provided for this purpose and its calling sequence is

```
CALL INAIJi ( I, J, VALUE, S )
```

where I and J are the subscripts, and VALUE is the numerical value. The subroutine INAIJi adds the quantity VALUE to the appropriate current value in storage, rather than making an assignment. This is helpful in situations (e.g. in some finite element applications) where the numerical values are obtained in an incremental fashion.

For example, the execution of

```
.  
. .  
. .  
CALL INAIJ2 ( 3, 4, 9.5, S )  
CALL INAIJ2 ( 3, 4, -4.0, S )  
. .  
. .
```

effectively assigns 5.5 to the (3, 4)-th component of A .

b) Input of a row of nonzeros.

The routine INROWi can be used to input the numerical values of a row or part of a row in the matrix. Its calling sequence is similar to that of INROW, described on Section 2.2:

```
CALL INROWi ( I, NIR, IR, VALUES, S )
```

Here the additional parameter VALUES is a floating point array containing the numerical values of the row. Again, the numerical values are added to the current values in storage.

c) Input of a submatrix.

The routine that allows the input of a submatrix is INMATi. Its parameter list corresponds to that of INIJIJ with the additional parameter VALUES that stores the numerical quantities:

```
CALL INMATi ( NIJ, II, JJ, VALUES, S )
```

Again, the numerical values in VALUES are added to those currently held by the package.

Mixed use of the routines INAIJi, INROWi and INMATi is permitted. Thus, the user is free to use whatever routine is most convenient.

The same convenience is provided in the input of numerical values for the right hand side vector b. SPARSPAK includes the routine INEI which inputs an entry of the right hand side vector

```
CALL INEI ( I, VALUE, S )
```

Here I is the subscript and VALUE is the numerical value. Alternatively, the routine INBIBI can be used to input a subvector, and its calling sequence is

```
CALL INBIBI ( NI, II, VALUES, S )
```

where NI is the number of input numerical values, and II and VALUES are vectors containing the subscripts and numerical values respectively. In both routines, incremental calculation of the numerical values is performed.

In some situations where the entire right hand side vector is available, the user can use the routine INRHS which transmits the whole vector to SPARSPAK. It has the form

```
CALL INRHS ( RHS, S )
```

where RHS is the vector containing the numerical values.

In all three routines, the numbers provided are added to those currently held by the package, and the use of the routines can be intermixed. The storage used for the right hand side by SPARSPAK is initialized to zero the first time any of them is executed.

IMPORTANT NOTES:

- a) When the matrix A is symmetric, so that method i, with i odd, is used, SPARSPAK requires that the elements of the lower triangle be provided. Thus, for example, the following statement will cause an error.

```
CALL INAIJ3 ( 3, 5, 1.3, S )
```

- b) The examples which we have given assume that a single precision version of SPARSPAK is being used. If the version is in double precision, the numerical values and numerical variables should be declared as double precision. For example:

```
CALL INAIJ3 ( 5, 3, 1.3D0, S )
```

2.5 MODULES FOR NUMERICAL SOLUTION

The numerical computation of the solution vector is initiated by the FORTRAN statement

```
CALL SOLVEi ( S )
```

where S is the working storage array for SPARSPAK. Again, the last digit i is used to distinguish between solvers for different storage methods.

Internally, the routine SOLVEi consists of both the factorization and forward/backward solution steps. If the factorization has been performed in a previous call to SOLVEi, SPARSPAK will automatically skip the factorization step, and perform the solution step directly. The solution vector is returned in the first NEQNS locations of the

storage vector S . If SOLVEi is called before any right hand side values are input, only the factorization will be performed. The solution returned will be all zeroes. See Examples 3 and 4 in Section 9.

SECTION 3

SOME GUIDELINES ON SELECTING A METHOD

We mentioned in Section 1 that there are six basic methods, distinguished by a numerical digit i satisfying $1 \leq i \leq 6$. These six methods can be viewed as grouped into three odd-even pairs; the only distinction between method i (odd) and method $i+1$ is that method i assumes A is symmetric, and method $i+1$ assumes A is unsymmetric. Thus, we really only provide three essentially distinct methods, with each one having a symmetric and unsymmetric version. Hence, in this section we will largely confine our remarks to methods 1, 3 and 5; comparative remarks about them will also apply to their unsymmetric analogues, methods 2, 4 and 6.

The basic methods are as follows; the remarks comparing them, and the advice provided, should be regarded as at best tentative. Characteristics of sparse matrices vary a great deal.

Method

Basic Strategy and References

1, 2 The objective of these methods is to reorder A so it has a small bandwidth or profile [6]. The well-known reverse Cuthill-McKee algorithm is used. For relatively small problems, say $N \leq 200$, they are probably the best overall methods to use.

3, 4 The objective of these methods is to reduce storage requirements, but the factorization time will usually be substantially higher than any of the other methods. Their storage requirements will usually be substantially less than methods (1, 2) (unless N is very large). The same remark is true about the relative solution times. Thus, these methods are often useful when storage is restricted, and/or when many problems which differ only in the right hand side must be solved (see Section 6).

There are two ordering options provided: ORDRA3 and CRDFE3 (and similarly for the unsymmetric case). The A option is

specifically tailored for 'finite element problems', typical of those arising in structural analysis and the numerical solution of partial differential equations [1]. The B option is effective for less specific problems; and uses a refined quotient tree ordering described in [2].

5, 6 These methods attempt to find orderings which minimize fill-in, and they exploit all zeroes. Their ordering times are almost always greater than those above, but for moderate-to-large problems the reduced factorization times usually are more than compensatory.

Just as for methods (3, 4), there are two ordering options provided. The A option is again specifically designed for finite element problems, and uses a so-called nested dissection ordering [3]. The B option uses the minimum degree algorithm, and is suitable for all sparse problems [4].

To summarize, our tentative advice and guidelines are as follows:

1. For small problems, use method (1, 2).
2. For small to moderate size problems that have to be solved only once, use method (1, 2) if enough storage is available. If not, use method (3, 4). If the problem is quite large, method (5, 6) might be better.
3. For moderate to large problems, use either method (3, 4) or (5, 6). If many problems differing only in the right hand side must be solved, method (3, 4) is probably the best. If the problem is quite large, and many problems having the same structure, but different numerical values, must be solved, then method (5, 6) is probably the best. (See Sections 5 and 6.)

SPARSPAK has been designed so that the ORDRxi modules can be used as aids in selecting a method. The basic strategy, as illustrated in Example 6 in Section 9, is to input the matrix structure, and then run the various ORDRxi modules on it, printing the storage statistics gathered by the package (using PSTATS, described in Section 7) after each ordering module has been executed.

SECTION 4

SAVE AND RESTART FACILITIES

SPARSPAK provides two subroutines called SAVE and RESTRT which allow the user to stop the calculation at some point, save the results on an external sequential file, and then restart the calculation at exactly that point some time later. To save the results of the computation done thus far, the user executes the statement

```
CALL SAVE ( K, S )
```

where K is the FORTRAN logical unit on which the results are to be written, along with other information needed to restart the computation. If execution is then terminated, the state of the computation can be re-established by executing the statement

```
CALL RESTRT ( K, S )
```

Examples 4, 5 and 6 provided in Section 9 illustrate the use of SAVE and RESTRT.

Note that executing SAVE does not destroy any information; the computation can proceed just as if SAVE was not executed.

When errors occur in a module, the routines SAVE and RESTRT are useful in saving the results of previously successfully executed modules (see Section 7.3 and Example 5 in Section 9).

Another potential use of the SAVE and RESTRT modules is to make the working storage array S available to the user in the middle of a sparse matrix computation. After SAVE has been executed, the working storage array S can be used by some other computation.

Finally, the SAVE and RESTRT modules allow the user to segment the computation into several distinct phases, and thereby reduce the amount of program that must be resident in storage at any given time.

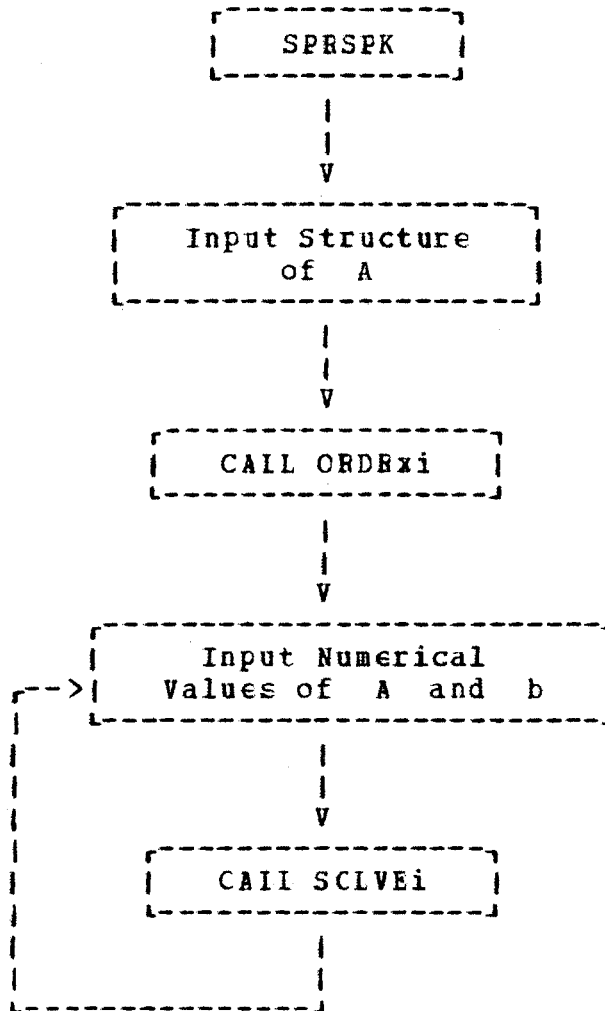
IMPORTANT NOTES:

- a) In the subroutines SAVE and RESTRT, information is either written on or read from the FORTRAN logical unit K using binary_format.
- b) If the subroutines SAVE and RESTRT are used, then before the user executes his program, he must define a file for the FORTRAN logical unit K using the appropriate system control card or command (this depends on the environment in which the program is being executed). Furthermore, this file must be preserved by the user for later access by the RESTRT subroutine.

SECTION 5

SOLVING MANY PROBLEMS HAVING THE SAME STRUCTURE

In certain applications, many problems which have the same sparsity structure, but different numerical values, must be solved. In this case, the structure input, ordering, and data structure set-up needs only to be done once. This situation can be accommodated perfectly well by SPARSPAK. The control sequence is depicted by the following flowchart:



When the numerical input routines (INAIJi, INBI, ..., etc.) are first called after SOLVEi has been called, this is detected by SPARSPAK, and the computer storage used for A and b is initialized to zero.

Note that if such problems must be solved over an extended time period (i.e., in different runs), the user can execute SAVE after executing CRDRxi and thus avoid the input of the structure of A and the execution of OBDRxi in subsequent equation solutions.

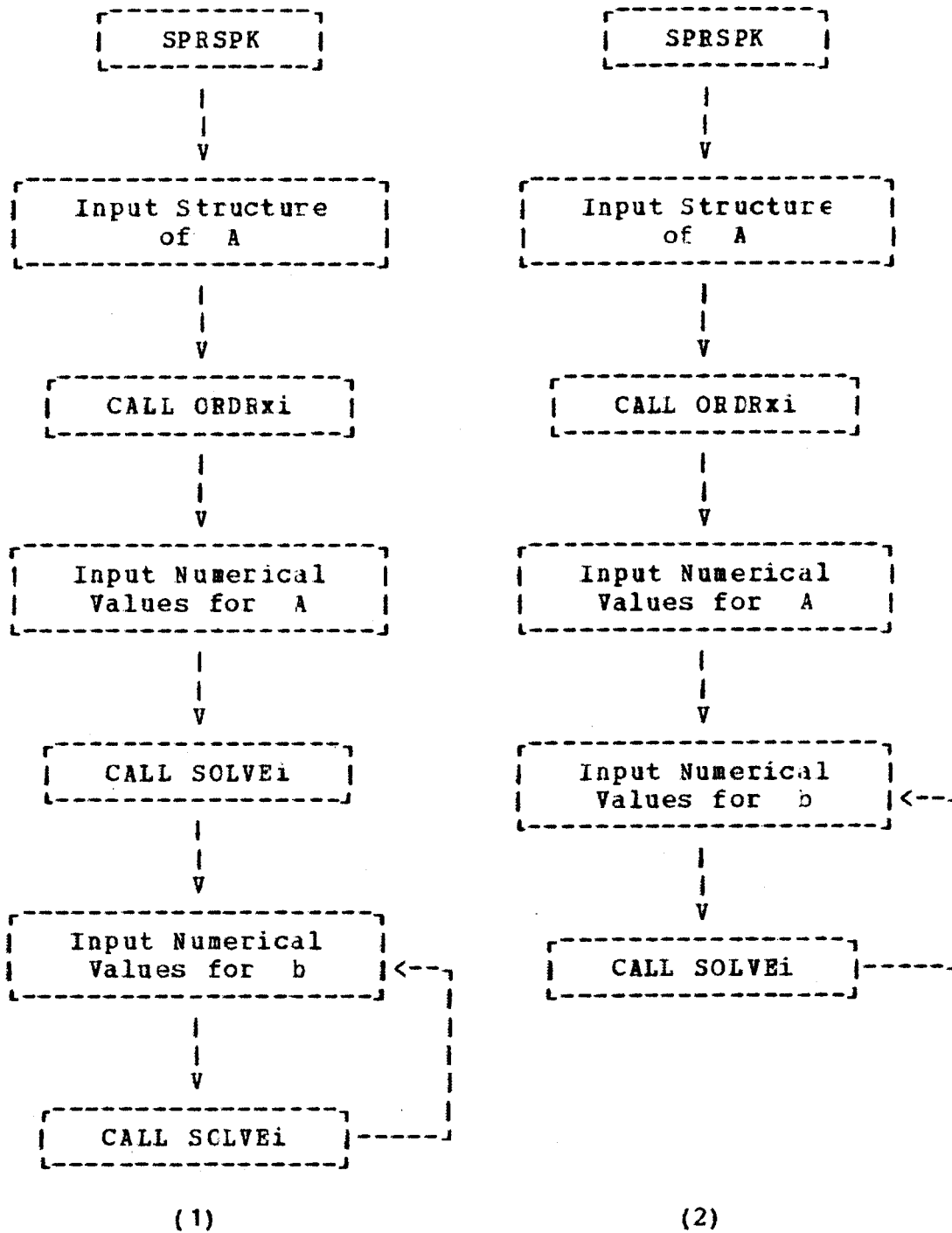
SECTION 6

SOLVING MANY PROBLEMS WHICH DIFFER ONLY IN THEIR RIGHT HAND SIDE

In some applications, numerous problems which differ only in their right hand sides must be solved. In this case, we only want to factor A into LU (or LL') once, and use the factors repeatedly in the calculation of x for each different b. Again, SPARSPAK can handle this situation in a straightforward manner, as illustrated by the flowcharts on the following page.

When SPARSPAK is used as indicated by flowchart (1), the package detects that no right hand side has been provided during the first execution of SOLVEi, and only the factorization is performed. In subsequent calls to SOLVEi, SPARSPAK detects that the factorization has already been performed, and that part of the SOLVEi module is bypassed. In flowchart (2), both factorization and solution are performed during the first call to SOLVEi, with only the solve part performed in subsequent executions of SOLVEi. (See Example 3 in Section 9.)

Note that SAVE can be used after SOLVEi has been executed, if the user wants to save the factorization for use in some future calculation.



SECTION 7

OUTPUT FROM SPARSPAK

As noted earlier in Section 2, the user supplies a one-dimensional floating point array S , from which all array storage is allocated. In particular, the interface allocates the first $NEQNS$ storage locations in S for the solution vector of the linear system. After all the interface modules for a particular method have been successfully executed, the user can retrieve the solution from these $NEQNS$ locations.

In addition to the solution x , SPARSPAK may print other information about the computation, depending upon the value of $MSGVLV$, whether or not errors occur, and whether or not the module $PSTATS$ is called. This section discusses these features of SPARSPAK.

NOTE:

SPARSPAK writes output to two FORTRAN logical output units, whose numbers are given by $IPRNTS$ and $IPRNTE$. The values for these variables are set in the module $SPSEPK$ when the package is installed. Standard output requested by the user is printed on unit $IPRNTS$, while any error messages raised by SPARSPAK are printed on unit $IPRNTE$. In an interactive environment, $IPRNTE$ is usually the user's terminal, while $IPRNTS$ is some other output device on which the output of the (hopefully) successful run is recorded. In a batch oriented environment, $IPRNTS$ and $IPRNTE$ are usually the same. Note that the user and/or the computer installation must ensure that the files associated with $IPRNTS$ and $IPRNTE$ are available to the user's program before execution begins.

7.1 MESSAGE LEVEL (MSGVLV)

The first variable $MSGVLV$ in the common block $SPKUSR$ stands for 'message level', and governs the amount of information printed by the interface modules. Its default value is two, and for this value a relatively small amount of summary information is printed, indicating the initiation of each phase. When $MSGVLV$ is set to one by

the user, only fatal error messages are printed; this option could be useful if SPARSPAK is being used in the 'inner loop' of a large computation, where even summary information would generate excessive output. Increasing the value of MSGLVL (up to 4) provides increasingly detailed information about the computation. Note that the module SPRSPK sets MSGLVL to its default value; if the user wishes MSGLVL to be different from 2, he must reset it after SPRSPK has been called.

In many circumstances, SPARSPAK will be imbedded in still another 'super' package which models phenomena producing sparse matrix problems. Messages printed by SPARSPAK may be useless or even confusing to the ultimate users of the super package, or the super package may wish to field the error conditions and perhaps take some corrective action which makes the error messages erroneous. Thus, all printing by SPARSPAK can be inhibited by setting MSGLVL to zero.

To summarize, we have

<u>MSGLVL</u>	<u>Printed Output</u>
0	No messages
1	Fatal error messages
2	Minimal summary information
3	More detailed information
4	Detailed debugging information

7.2 STATISTICS GATHERING (PSTATS)

SPARSPAK gathers a number of statistics which the user will find useful if he is comparing various methods, or is going to solve numerous similar problems and wants to adjust his working storage to the minimum necessary. The package has a common block called SPKDTA containing variables whose values can be printed by executing the statement

CALL PSTATS

The information printed is:

the number of equations,
the number of off-diagonal nonzeros in the matrix,
the size of the working storage,
the time used to find the ordering,
the time used for data structure set-up,
the time used for the factorization step,
the time used for the triangular solution step,
number of operations required by the factorization step,

number of operations required by the triangular solution,
the storage used by the ordering subroutine,
the storage used by the data structure set-up subroutine,
the storage used by the SOLVEi module.

Since the module PSTATS can be called at any time, some of the above information may not be available, and will not be printed. The word 'operations' here means multiplicative operations (multiplications and divisions). Since most of the arithmetic performed in sparse matrix computation occurs in multiply-add pairs, the number of operations (as defined here) is a useful measure of the amount of arithmetic performed.

The reader is referred to the examples in Section 9 for more discussion about the output from PSTATS.

7.3 ERROR MESSAGES (IERR)

When a fatal error is detected, so that the computation cannot proceed, a positive code is assigned to IERR. The user can simply check the value of IERR to see if the execution of module has been successful. This error flag can be used in conjunction with the save/restart feature described in Section 4 to retain the results of successfully completed parts of the computation, as shown by the program fragment below.

```

.
.
.
CALL ORDRA1 ( S )
IF ( IERR .EQ. 0 ) GO TO 100
CALL SAVE ( 3, S )
STOP
100 CONTINUE
.
.
.

```

The variable IERR is set to the value $10*k+1$, where $0 \leq l \leq 9$ distinguishes the error, and k is determined by the type of module that sets IERR positive:

<u>k</u>	
0	save and restart modules (SAVE, RESTRT)
1	matrix structure input modules (INIJ, INIJIJ, etc.)
2	matrix ordering and allocation modules (ORDRxi)
3	matrix numerical input (INAIJi, ..., etc.)
4	right hand side numerical input (INBI, ..., etc.)
5	factorization and solution modules (SOLVEi)

7.3.1 Save and Restart Routines

IERR

RESTR

- 1 Insufficient storage to restart the computational process. The minimum value of MAXS required is printed in the error message.

7.3.2 Input of the Matrix Structure

IERR

INIJ, INIJIJ, INCLQ

- 11 Insufficient storage was provided in the working storage array. The (i, j) pairs input to INIJ, INIJIJ, and INCLQ will be counted and discarded. Duplicates which are detected will not be counted, but some duplicates may be missed.
- 12 Negative or zero subscript is found.
- 13 Incorrect execution sequence. Probable cause of error: routine IJBEGN was not called before (i, j) pairs input began.

IERR

IJEND

- 16 Insufficient storage to transform matrix structure. The minimum value of MAXS required is printed in the error message.
- 17 Incorrect execution sequence. IJEND was called before new matrix structure has been input.
- 18 NEQNS is zero.

7.3.3 Ordering and Storage Allocation Routines

IERR

- 21 Incorrect execution sequence. Probable cause: subroutine IJEND did not execute successfully.
- 22 Incompatible ordering method. User probably executed part of the ordering subroutine ORDRxi, and then executed SAVE because of insufficient storage. The execution was then restarted, using RESTRT, but ORDRxj was called with $i \neq j$.
- 23 Insufficient storage in working storage array to begin execution
Response: execute SAVE, and call ORDRxi with MAXS at least as large as that indicated in the error message.
- 24 Insufficient storage in working storage array to execute the storage allocation subroutine. The ordering routine has successfully executed.
Response: same as for error 23.
- 25 Working storage array was not large enough. The storage allocation routine was executed, but there was not enough storage to hold the data structure pointers.
Response: same as for error 23.
- 26 Working storage array is large enough for execution of ORDRxi, and it has successfully executed. However, there is not enough storage available for the numerical values, so computation cannot proceed.
Response: execute SAVE, and re-initiate computation after adjusting MAXS to at least the value specified in the error message.

7.3.4 Input of the Numerical Values

IERR INAIJi, INROWi, INMATi

- 31 Incorrect execution sequence. Probable cause:
 unsuccessful execution of the ordering routine
 ORDRxi .

- 32 Incompatible input routine - attempt to use input
 routine INAIJi, INFCWi, or INMATi after using
 ORDRxj , where $i \neq j$. Use the routine specified
 in the error message.

- 33 Attempt to input the (i, j)-th element of matrix
 A for $i < j$. (This error occurs only for
 symmetric matrix methods; i.e., when method is
 odd). Methods for symmetric matrices expect
 elements of the lower triangle to be input.

- 34 Attempt to input an (i, j)-th element of matrix A
 where $i > N, j > N, i < 1, \text{ or } j < 1$.

- 35 Attempt to input a numerical value for the
 (i, j)-th element of matrix A into the data
 structure, but the data structure has no space for
 it. Probable cause: the user has not called INIJ,
 INIJIJ, INCLQ or INROW with all the pairs
 (i, j) for which the (i, j)-th elements of A
 are nonzero. (SPARSEPAK thinks A is sparser than
 it really is.)

IERR INBI, INBIBI, INRHS

- 41 Incorrect execution sequence. Probable cause is
 the unsuccessful execution of ORDRxi .

- 42 Subscript out of range - attempt to input a
 numerical value for the i-th element of b where
 $i > N$ or $i < 1$.

7.3.5 Factorization and Solution

IERR

SOLVEi

- 51 Incorrect execution sequence. Probable cause is unsuccessful execution of the numerical input routines.
- 52 Incompatible ordering and solution routines have been called.
Response: execute SAVE and restart the computation using SCIVEi where i is the value of METHOD specified in the error message.
- 53 Zero pivot or negative square root detected in the factorization routine. Possible causes:
a) incorrect use of the numerical input routines.
b) the matrix may require pivoting in order to preserve numerical stability. In this case the use of SPARSPAK to solve the problem is inappropriate. (See restrictions in Section 1.)

SECTION 8

SUMMARY LISTING OF INTERFACE ROUTINES

SPRSPK	┌			└	
					Initialization of SPARSP
				└	
IJBEGN				┌	
INIJ (I, J, S)					
INROW (I, NIR, IR, S)				└	Structure input
INIJIJ (NIJ, II, JJ, S)					
INCLQ (NCLQ, CLQ, S)					
IJEND (S)				└	
				┌	
ORDRxi (S)				└	Ordering (see next page)
				└	
INAIJi (I, J, VALUE, S)				┌	
INROWi (I, NIR, IR, VALUES, S)				└	Matrix input
INMATi (NIJ, II, JJ, VALUES, S)				└	
				┌	
INBI (I, VALUE, S)				└	
INBIi (NI, II, VALUES, S)				└	Right hand side input
INRHS (RHS, S)				└	
				┌	
SOLVEi (S)				└	Factorization and/or Solution
				└	
PSTATS				┌	
				└	Print statistics
				└	
SAVE (k, S)				┌	
RESTR (k, S)				└	Save and Restart the computation
				└	

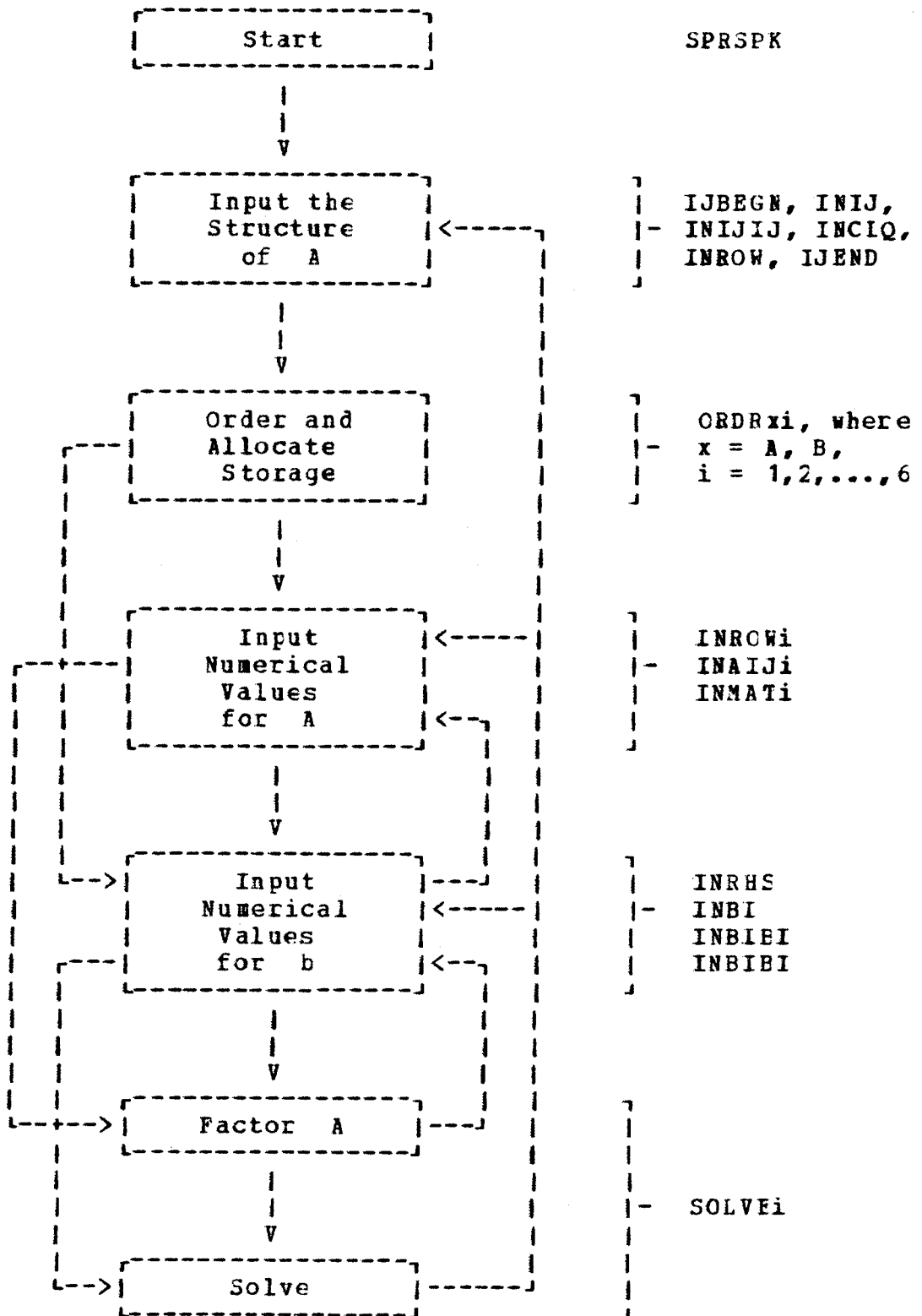
Ordering Choices

ORDRX1

x i

A	1	Reverse Cuthill-McKee ordering [7]; symmetric	A
A	2	Reverse Cuthill-McKee ordering [7]; unsymmetric	A
A	3	One-way Dissection ordering [1]; symmetric	A
A	4	One-way Dissection ordering [1]; unsymmetric	A
B	3	Refined quotient tree ordering [2]; symmetric	A
B	4	Refined quotient tree ordering [2]; unsymmetric	A
A	5	Nested Dissection ordering [3]; symmetric	A
A	6	Nested Dissection ordering [3]; unsymmetric	A
B	5	Minimum Degree ordering [4]; symmetric	A
B	6	Minimum Degree ordering [4]; unsymmetric	A

Sketch of Possible Execution Paths through SPARSPAK Modules



SECTION 9

EXAMPLES

In this section, we provide several programs which illustrate how SPARSPAK can be used. These programs are derived from the one given in Section 2.1.

These examples were run using a standard single precision version of SPARSPAK under the IBM FORTRAN H extended compiler on an IBM 3031 computer. All times reported are in seconds. It should be noted that the results will be different if a different version of SPARSPAK is used.

Example 1

This is an example of the simplest use of SPARSPAK, with each of the modules of method 1 used in sequence. The problem that is solved is a 10 by 10 symmetric system $Ax = b$ where the diagonal elements of A are all 4, and the superdiagonal and subdiagonal elements are all -1. The right hand side vector b is chosen so that the entries of the solution vector x are all ones.

In the program, the nonzero structure of A is input using IJBEGN, INIJ and IJEND. After ORDRA1 is executed, the interface modules INAIJ1 and INBI are used to transmit the numerical values of A and b to the package respectively. The module SOLVE1 is called to do the numerical solution and then PSTATS is called to print out the statistics gathered by the interface during execution. Finally, the error in the computed approximate solution is computed.

Note that the size of the working storage provided was 250, while the maximum amount used by any of the modules was 60, which was the storage requirement for the ORDRA1 and SOLVE1 module. Thus, if the user was going to solve this problem again, he could adjust his storage down to 60.

C
C

M A I N L I N E P R O G R A M

REAL S(250), ERROR, FOUR, ONE, TWO, ZERO
INTEGER I, IERR, IFFNTE, IPRNTE, MAXS, MSGVL, NEQNS
REAL RATIOI, RATIOS, TIME
COMMON /SPKUSR/ MSGVL, IERR, MAXS, NEQNS
COMMON /SPKSYS/ IPRNTE, IERNTS, RATIOS, RATICI, TIME

C

```
CALL SPRSPK
MAXS = 250
CALL IJBEGN
DO 100 I = 2, 10
    CALL INIJ ( I, I-1, S )
100 CONTINUE
CALL IJEND ( S )
CALL ORDRA1 ( S )
ZERO = 0.0E0
ONE = 1.0E0
TWO = 2.0E0
FOUR = 4.0E0
DO 200 I = 1, 10
    IF ( I .GT. 1 ) CALL INAIJ1 ( I, I-1, -ONE, S )
    CALL INAIJ1 ( I, I, FOUR, S )
    CALL INBI ( I, TWO, S )
200 CONTINUE
CALL INBI ( 1, ONE, S )
CALL INBI ( 10, ONE, S )
CALL SOLVE1 ( S )
CALL PSTATS
ERROR = ZERO
DO 300 I = 1, 10
    ERROR = AMAX1 ( ERROR, ABS ( S(I)-ONE ) )
300 CONTINUE
WRITE (IPRNTE, 11) ERROR
11 FORMAT ( / 15H MAXIMUM ERROR , 1PE15.3 )
STOP
END
```

***** UNIVERSITY OF WATERLOO
 ***** SPARSE MATRIX PACKAGE
 ***** (S P A R S P A K)
 ***** RELEASE 2
 ***** (C) JANUARY 1979
 ***** STANDARD VERSION
 ***** SINGLE PRECISION
 ***** LAST UPDATE JANUARY 1980

OUTPUT UNIT FOR ERROR MESSAGE 6
 OUTPUT UNIT FOR STATISTICS 6

IJBEGN- BEGIN STRUCTURE INPUT...

INIJ- INPUT OF ADJACENCY PAIRS

IJEND- END OF STRUCTURE INPUT

CRDRA1- RCM ORDERING

INAIJ1- INPUT OF MATRIX COMPONENTS

INBI- INPUT OF RIGHT HAND SIDE

SOLVE1- ENVELOPE SOLVE

PSTATS- STATISTICS

NUMBER OF EQUATIONS	10
OFF-DIAGONAL NONZEROS	18
SIZE OF WORKING STORE (MAXS)	250
TIME FOR ORDERING	0.003
STORAGE FOR ORDERING	60.
TIME FOR ALLOCATION	0.0
STORAGE FOR ALLOCATION	60.
STORAGE FOR SOLUTION	60.
TIME FOR FACTORIZATION	0.0
TIME FOR SOLVING	0.003
OPERATIONS IN FACTORIZATION	18.
OPERATIONS IN SOLVING	38.

MAXIMUM ERROR 1.013E-06

Example 2

This is the same as Example 1, except that the matrix A is unsymmetric. The diagonal elements of A are all 4, the superdiagonal elements are all 1, and the subdiagonal elements are all -1. The right hand side vector b is chosen so that the entries of the solution vector x are all ones.

```
C          M A I N L I N E   P R O G R A M
C
INTEGER      I, IERR, IPRNTE, IPRNTS, MAXS, MSGLVL, NEQNS
REAL         S(250), ERROR, FOUR, ONE, ZERO
REAL         RATICL, RATICS, TIME
COMMON /SPKUSR/ MSGLVL, IERR, MAXS, NEQNS
COMMON /SPKSYS/ IPRNTE, IERNTS, RATIOS, RATIOI, TIME
C
      CALL SPRSPK
      MAXS = 250
      CALL IJBEGN
      DO 100 I = 2, 10
          CALL INIJ ( I, I-1, S )
100     CONTINUE
      CALL IJEND ( S )
      CALL ORDRA2 ( S )
      ZERO = 0.0E0
      ONE = 1.0E0
      FOUR = 4.0E0
      DO 200 I = 1, 10
          IF ( I .GT. 1 ) CALL INAIJ2 ( I, I-1, -ONE, S )
          IF ( I .LT. 10 ) CALL INAIJ2 ( I, I+1, ONE, S )
          CALL INAIJ2 ( I, I, FOUR, S )
          CALL INBI ( I, FOUR, S )
200     CONTINUE
      CALL INBI ( 1, ONE, S )
      CALL INBI ( 10, -ONE, S )
      CALL SOLVE2 ( S )
      CALL PSTATS
      ERROR = ZERO
      DO 300 I = 1, 10
          ERROR = AMAX1 ( ERROR, ABS ( S(I)-ONE ) )
300     CONTINUE
      WRITE (IPRNTS, 11) ERROR
11     FORMAT ( / 15H MAXIMUM ERROR , 1PE15.3 )
      STOP
END
```



```

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```

```

          OUTPUT UNIT FOR ERROR MESSAGE      6
          OUTPUT UNIT FOR STATISTICS        6

```

IJBEGN- BEGIN STRUCTURE INPUT...

INIJ- INPUT OF ADJACENCY PAIRS

IJEND- END OF STRUCTURE INPUT

ORDRA2- RCM ORDERING

INAIJ2- INPUT OF MATRIX COMPONENTS

INBI- INPUT OF RIGHT HAND SIDE

SOLVE2- ENVELOPE SOLVE

PSTATS- STATISTICS

```

          NUMBER OF EQUATIONS                10
          OFF-DIAGONAL NONZEROS             18
          SIZE OF WORKING STORE (MAXS)      250
          TIME FOR ORDERING                  0.003
          STORAGE FOR ORDERING              60.
          TIME FOR ALLOCATION                 0.0
          STORAGE FOR ALLOCATION             60.
          STORAGE FOR SOLUTION              69.
          TIME FOR FACTORIZATION             0.003
          TIME FOR SOLVING                   0.0
          OPERATIONS IN FACTORIZATION       18.
          OPERATIONS IN SOLVING            28.

```

MAXIMUM ERROR 0.0

Example 3

This is similar to Example 1, except that method 3 is used (with the A ordering option), and two problems differing only in their right hand sides are solved. After solving the problem whose solution vector contains all ones, a new right hand side is input which corresponds to a different problem whose solution vector contains all twos. When the module SOLVE3 is called a second time, the interface detects that the factorization has already been done, and only the triangular solution is performed.

C
C

M A I N L I N E P R O G R A M

INTEGER I, IERR, IPRNTE, IPRNTS, MAXS, MSGIVL, NEQNS
REAL S(250), ERROR, FOUR, ONE, TWO, ZEFC
REAL RATICL, RATICS, TIME
COMMON /SPKUSR/ MSGLVL, IERR, MAXS, NEQNS
COMMON /SPKSYS/ IPRNTE, IPRNTS, RATIOS, RATIOI, TIME

C

```
CALL SPRSPK
MAXS = 250
CALL IJBEGN
DO 100 I = 2, 10
    CALL INIJ ( I, I-1, S )
100 CONTINUE
CALL IJEND ( S )
CALL ORDRA3 ( S )
ZERO = 0.0E0
ONE = 1.0E0
TWO = 2.0E0
FOUR = 4.0E0
DO 200 I = 1, 10
    IF ( I .GT. 1 ) CALL INAIJ3 ( I, I-1, -ONE, S )
    CALL INAIJ3 ( I, I, FOUR, S )
    CALL INBI ( I, TWO, S )
200 CONTINUE
CALL INBI ( 1, ONE, S )
CALL INBI ( 10, ONE, S )
CALL SOLVE3 ( S )
CALL PSTATS
ERROR = ZERO
DO 300 I = 1, 10
    ERROR = AMAX1 ( ERROR, ABS ( S(I)-ONE ) )
300 CONTINUE
WRITE (IPRNTS, 11) ERROR
11 FORMAT ( / 15H MAXIMUM ERROR , 1PE15.3 )

C
DO 400 I = 1, 10
    CALL INBI ( I, FOUR, S )
400 CONTINUE
CALL INBI ( 1, TWO, S )
CALL INBI ( 10, TWO, S )
CALL SOLVE3 ( S )
CALL PSTATS
ERROR = ZERO
DO 500 I = 1, 10
    ERROR = AMAX1 ( ERROR, ABS ( S(I)-TWO ) )
500 CONTINUE
WRITE (IPRNTS, 11) ERROR
STOP
END
```

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OUTPUT UNIT FOR ERRCR MESSAGE 6
 OUTPUT UNIT FOR STATISTICS 6

IJBEGN- BEGIN STRUCTURE INPUT...

INIJ- INPUT OF ADJACENCY PAIRS

IJEND- END OF STRUCTURE INPUT

ORDRA3- ONE WAY DISSECTION ORDERING

INAIJ3- INPUT OF MATRIX COMPONENTS

INBI- INPUT OF RIGHT HAND SIDE

SOLVE3- IMPLICIT BLOCK SOLVE

PSTATS- STATISTICS

NUMBER OF EQUATIONS	10
OFF-DIAGONAL NONZEROS	18
SIZE OF WORKING STORE (MAXS)	250
TIME FOR ORDERING	0.007
STORAGE FOR ORDERING	91.
TIME FOR ALLOCATION	0.0
STORAGE FOR ALLOCATION	94.
STORAGE FOR SOLUTION	94.
TIME FOR FACTORIZATION	0.003
TIME FOR SOLVING	0.003
OPERATIONS IN FACTORIZATION	18.
OPERATIONS IN SOLVING	38.

MAXIMUM ERROR 1.013E-06

INBI- INPUT OF RIGHT HAND SIDE

SOLVE3- IMPLICIT BLOCK SOLVE

FACTORIZATION ALREADY DONE.

PSTATS- STATISTICS

NUMBER OF EQUATIONS	10
OFF-DIAGONAL NONZEROS	18
SIZE OF WORKING STORE (MAXS)	250
TIME FOR ORDERING	0.007
STORAGE FOR ORDERING	91.

TIME FOR ALLOCATION	0.0
STORAGE FOR ALLOCATION	94.
STORAGE FOR SOLUTION	94.
TIME FOR FACTORIZATION	0.003
TIME FOR SOLVING	0.003
OPERATIONS IN FACTORIZATION	18.
OPERATIONS IN SOLVING	38.

MAXIMUM ERROR 1.907E-06

Example 4

This example illustrates the use of the save/restart feature of SPARSPAK. After the factorization is computed, SAVE is executed, which writes the current state of the computation on FORTRAN logical unit 3. In a second program the module RESTRT is executed to read the information from unit 3, and the computation resumes at the point at which SAVE was invoked.

```
C           M A I N L I N E       P R O G R A M
C
INTEGER      I, IERR, MAXS, MSGVLV, NEQNS
REAL         S(250), ERROR, FOUR, ONE
COMMON /SPKUSR/ MSGVLV, IERR, MAXS, NEQNS
C
      CALL SPRSPK
      MAXS = 250
      CALL IJBEGN
      DO 100 I = 2, 10
          CALL INIJ ( I, I-1, S )
100    CONTINUE
      CALL IJEND ( S )
      CALL ORDRA1 ( S )
      ONE = 1.0E0
      FOUR = 4.0E0
      DO 200 I = 1, 10
          IF ( I .GT. 1 )      CALL INAIJ1 ( I, I-1, -ONE, S )
          CALL INAIJ1 ( I, I, FOUR, S )
200    CONTINUE
      CALL SOLVE1 ( S )
      CALL PSTATS
      CALL SAVE ( 3, S )
      STOP
END
```

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OUTPUT UNIT FOR ERROR MESSAGE 6
 OUTPUT UNIT FOR STATISTICS 6

IJBEGN- BEGIN STRUCTURE INPUT...

INIJ- INPUT OF ADJACENCY PAIRS

IJEND- END OF STRUCTURE INPUT

CRDRA1- RCM ORDERING

INAIJ1- INPUT OF MATRIX COMPONENTS

SOLVE1- ENVELOPE SOLVE
 NO RIGHT HAND SIDE PROVIDED,
 SOLUTION WILL BE ALL ZEROS.

PSTATS- STATISTICS
 NUMBER OF EQUATIONS 10
 OFF-DIAGONAL NONZEROS 18
 SIZE OF WORKING STORE (MAXS) 250
 TIME FOR ORDERING 0.003
 STORAGE FOR ORDERING 60.
 TIME FOR ALLOCATION 0.0
 STORAGE FOR ALLOCATION 60.
 STORAGE FOR SOLUTION 60.
 TIME FOR FACTORIZATION 0.0
 TIME FOR SOLVING 0.0
 OPERATIONS IN FACTORIZATION 18.
 OPERATIONS IN SOLVING 0.

SAVE- STORAGE VECTOR SAVED

```

C           M A I N L I N E   P R O G R A M
C
INTEGER      I, IERR, IPRNIE, IPRNTS, MAXS, MSGGLVL, NEQNS
REAL         S(250), ERROR, ONE, TWO, ZERO
REAL         RATIOI, RATICS, TIME
COMMON /SPKUSR/  MSGGLVL, IERR, MAXS, NEQNS
COMMON /SPKSYS/  IPRNTE, IPRNTS, RATIOS, RATIOI, TIME
C
      CALL SPRSPK
      MAXS = 250
      CALL RESTART ( 3, S )
      ZERO = 0.0E0
      ONE = 1.0E0
      TWO = 2.0E0
      DO 100 I = 1, 10
          CALL INBI ( I, TWO, S )
100      CONTINUE
          CALL INBI ( 1, ONE, S )
          CALL INBI ( 10, ONE, S )
          CALL SOLVE1 ( S )
          CALL PSTATS
          ERROR = ZERO
          DO 200 I = 1, 10
              ERROR = AMAX1 ( ERROR, ABS ( S(I)-ONE ) )
200      CONTINUE
          WRITE (IPRNTS, 11)  ERROR
          FORMAT ( / 15H MAXIMUM ERROR , 1PE15.3 )
          STOP
      END

```


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OUTPUT UNIT FOR ERROR MESSAGE 6
 OUTPUT UNIT FOR STATISTICS 6

RESTR- RESTART SYSTEM

INBI- INPUT OF RIGHT HAND SIDE

SOLVE1- ENVELOPE SOLVE
 FACTORIZATION ALREADY DONE.

PSTATS- STATISTICS

NUMBER OF EQUATIONS	10
OFF-DIAGONAL NONZEROS	18
SIZE OF WORKING STORE (MAXS)	250
TIME FOR ORDERING	0.003
STORAGE FOR ORDERING	60.
TIME FOR ALLOCATICN	0.0
STORAGE FOR ALLOCATION	60.
STORAGE FOR SCLUTION	60.
TIME FOR FACTORIZATION	0.0
TIME FOR SOLVING	0.0
OPERATIONS IN FACTORIZATION	18.
OPERATIONS IN SOLVING	38.

MAXIMUM ERROR 1.013E-06

Example 5

This example consists of four runs of essentially the same program, illustrating how the SAVE and RESTART modules can be used to avoid repeating successfully completed computations when the execution cannot proceed further because of lack of working storage. In the first run, MAXS was too small to accommodate the structure, and a message was printed indicating that MAXS must be at least 999 in order to input the structure. A second run with MAXS = 999 was executed, and the structure was successfully input; however, the ORDRA5 module could not execute because MAXS was less than 1400. The module SAVE was then executed and the run terminated.

The third run had MAXS = 2500, and the ordering and storage allocation were successfully performed. However, ORDRA5 terminated with an error because it detected that too little storage was available for the numerical computation (SOLVE5), so SAVE was again executed. Finally, the last run was executed with MAXS set to 2509 (the maximum value, printed in the third run), and the solution to the problem was obtained.

NOTE:

The following examples were run using a single precision version of SPARSPAK. The working storage required will therefore be different if a different version of SPARSPAK is used.

C
C

M A I N L I N E P R O G R A M

INTEGER I, IERR, IPRNTE, IPRNTE, MAXS, MSGLVL, NEQNS
REAL S(900), ERROR, FOUR, ONE, TWO, ZEEC
REAL RATICL, RATICS, TIME
COMMON /SPKUSR/ MSGLVL, IERR, MAXS, NEQNS
COMMON /SPKSYS/ IPRNTE, IPRNTE, RATIOS, RATIOI, TIME

C

```
CALL SPRSPK
MAXS = 900
CALL IJBEGN
DO 100 I = 2, 200
    CALL INIJ ( I, I-1, S )
100 CONTINUE
    CALL IJEND ( S )
    IF ( IERR .EQ. 0 ) GO TO 200
    CALL PSTATS
    STOP
200 CALL ORDRA5 ( S )
    IF ( IERR .EQ. 0 ) GO TO 300
    CALL SAVE ( 3, S )
    CALL PSTATS
    STOP
300 ZERO = 0.0E0
    ONE = 1.0E0
    TWO = 2.0E0
    FOUR = 4.0E0
    DO 400 I = 1, 200
        IF ( I .GT. 1 ) CALL INAIJ5 ( I, I-1, -ONE, S )
        CALL INBI ( I, TWO, S )
400 CONTINUE
    CALL INBI ( 1, ONE, S )
    CALL INBI ( 200, ONE, S )
    CALL SOLVE5 ( S )
    CALL PSTATS
    ERROR = ZERO
    DO 500 I = 1, 200
        ERROR = AMAX1 ( ERROR, ABS ( S(I)-ONE ) )
500 CONTINUE
    WRITE ( IPRNTE, 11 ) ERROR
    11 FORMAT ( / 15H MAXIMUM ERROR , 1PE15.3 )
    STOP
END
```

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OUTPUT UNIT FOR ERROR MESSAGE 6
 OUTPUT UNIT FOR STATISTICS 6

IJBEGN- BEGIN STRUCTURE INPUT...

INIJ- INPUT OF ADJACENCY PAIRS

IJEND- END OF STRUCTURE INPUT

IJEND - ERROR NUMBER 16
 TOO LITTLE STORAGE,
 MAXS MUST AT LEAST BE 999

PSTATS- STATISTICS
 NUMBER OF EQUATIONS 200
 OFF-DIAGONAL NONZEROS 398
 SIZE OF WORKING STORE (MAXS) 900

C
C

MAINLINE PROGRAM

INTEGER I, IERR, IPRNTE, IPRNTS, MAXS, MSGIVI, NEQNS
REAL S(999), ERROR, FOUR, ONE, TWO, ZERC
REAL RATICL, RATICS, TIME
COMMON /SPKUSR/ MSGLVL, IERR, MAXS, NEQNS
COMMON /SPKSYS/ IPRNTE, IPRNTS, RATIOS, RATIOI, TIME

C

```
CALL SPRSPK
MAXS = 999
CALL IJBEGN
DO 100 I = 2, 200
    CALL INIJ ( I, I-1, S )
100 CONTINUE
CALL IJEND ( S )
IF ( IERR .EQ. 0 ) GO TO 200
CALL PSTATS
STOP
200 CALL ORDRA5 ( S )
IF ( IERR .EQ. 0 ) GO TO 300
CALL SAVE ( 3, S )
CALL PSTATS
STOP
300 ZERO = 0.0E0
ONE = 1.0E0
TWO = 2.0E0
FOUR = 4.0E0
DO 400 I = 1, 200
    IF ( I .GT. 1 ) CALL INAIJ5 ( I, I-1, -ONE, S )
    CALL INBI ( I, TWO, S )
400 CONTINUE
CALL INBI ( 1, ONE, S )
CALL INBI ( 200, ONE, S )
CALL SOLVE5 ( S )
CALL PSTATS
ERROR = ZERC
DO 500 I = 1, 200
    ERROR = AMAX1 ( ERROR, ABS ( S(I)-ONE ) )
500 CONTINUE
WRITE ( IPRNTS, 11 ) ERROR
11 FORMAT ( / 15H MAXIMUM ERROR , 1PE15.3 )
STOP
END
```

```

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```

```

          OUTPUT UNIT FOR ERROR MESSAGE      6
          OUTPUT UNIT FOR STATISTICS        6

```

IJBEGN- BEGIN STRUCTURE INPUT...

INIJ- INPUT OF ADJACENCY PAIRS

IJEND- END OF STRUCTURE INPUT

ORDRA5- NESTED DISSECTION ORDERING

```

ORDRXI  (X=A, B. I=1,2,3,4,5,6)
        - ERROR NUMBER 23

```

```

INSUFFICIENT STORAGE FOR ORDERING.
MAXS MUST BE AT LEAST          1400

```

SAVE- STORAGE VECTOR SAVED

```

PSTATS- STATISTICS
          NUMBER OF EQUATIONS          200
          OFF-DIAGONAL NONZEROS       398
          SIZE OF WORKING STORE (MAXS) 999

```

C
C

MAINLINE PROGRAM

INTEGER I, IERR, IPRNTE, IPRNTS, MAXS, MSGLVL, NEQNS
REAL S(2500), ERRCR, FOUR, ONE, TWO, ZERO
REAL RATIOI, RATICS, TIME
COMMON /SPKUSR/ MSGLVL, IERR, MAXS, NEQNS
COMMON /SPKSYS/ IPRNTE, IPRNTS, RATIOS, RATICL, TIME

C

```
CALL SPRSPK
MAXS = 2500
CALL RESTRT ( 3, S )
CALL ORDRAS ( S )
IF ( IERR .EQ. 0 ) GO TO 100
CALL SAVE ( 3, S )
CALL PSTATS
STOP
100 ZERO = 0.0E0
ONE = 1.0E0
TWO = 2.0E0
FOUR = 4.0E0
DO 200 I = 1, 200
IF ( I .GT. 1 ) CALL INAIJ5 ( I, I-1, -ONE, S )
CALL INAIJ5 ( I, I, FOUR, S )
CALL INBI ( I, TWO, S )
200 CONTINUE
CALL INBI ( 1, ONE, S )
CALL INBI ( 200, ONE, S )
CALL SOLVE5 ( S )
CALL PSTATS
ERROR = ZERO
DO 300 I = 1, 200
ERROR = AMAX1 ( ERROR, ABS ( S(I)-ONE ) )
300 CONTINUE
WRITE (IPRNTS, 11) ERRCR
11 FORMAT ( / 15H MAXIMUM ERROR , 1PE15.3 )
STOP
END
```

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OUTPUT UNIT FOR ERROR MESSAGE 6
 OUTPUT UNIT FOR STATISTICS 6

RESTRT- RESTART SYSTEM

ORDRA5- NESTED DISSECTION ORDERING

ORDRXI (X=A, B. I=1,2,3,4,5,6)
 - ERROR NUMBER 26

INSUFFICIENT STORAGE
 FOR SOLVEI (I=1,2,3,4,5,6)
 MAXS MUST BE AT LEAST 2509

SAVE- STORAGE VECTOR SAVED

PSTATS- STATISTICS

NUMBER OF EQUATIONS	200
OFF-DIAGONAL NONZEROS	398
SIZE OF WORKING STORE (MAXS)	2500
TIME FOR ORDERING	0.083
STORAGE FOR ORDERING	1400.
TIME FOR ALLOCATION	0.030
STORAGE FOR ALLOCATION	2324.
STORAGE FOR SOLUTION	2509.

C
C

MAINLINE PROGRAM

INTEGER I, IERR, IPRNTE, IPRNTS, MAXS, MSGLVL, NEQNS
REAL S(2509), ERROR, FOUR, ONE, TWO, ZERO
REAL RATIO1, RATICS, TIME
COMMON /SPKUSR/ MSGLVL, IERR, MAXS, NEQNS
COMMON /SPKSYS/ IPRNTE, IPRNTS, RATIOS, RATIO1, TIME

C

```
CALL SPRSPK
MAXS = 2509
CALL RESTRT ( 3, S )
CALL ORDRA5 ( S )
IF ( IERR .EQ. 0 ) GO TO 100
CALL SAVE ( 3, S )
CALL PSTATS
STOP
100 ZERO = 0.0E0
ONE = 1.0E0
TWO = 2.0E0
FOUR = 4.0E0
DO 200 I = 1, 200
  IF ( I .GT. 1 ) CALL INAIJ5 ( I, I-1, -ONE, S )
  CALL INAIJ5 ( I, I, FOUR, S )
  CALL INBI ( I, TWO, S )
200 CONTINUE
CALL INBI ( 1, ONE, S )
CALL INBI ( 200, ONE, S )
CALL SOLVE5 ( S )
CALL PSTATS
ERROR = ZERO
DO 300 I = 1, 200
  ERROR = AMAX1 ( ERROR, ABS ( S(I)-ONE ) )
300 CONTINUE
WRITE (IPRNTS, 11) ERROR
11 FORMAT ( / 15H MAXIMUM ERROR , 1PE15.3 )
STOP
END
```

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OUTPUT UNIT FOR ERROR MESSAGE 6
 OUTPUT UNIT FOR STATISTICS 6

RESTRT- RESTART SYSTEM

ORDRA5- NESTED DISSECTION ORDERING

INAIJ5- INPUT OF MATRIX COMPONENTS

INBI- INPUT OF RIGHT HAND SIDE

SOLVE5- GENERAL SPARSE SOLVE

PSTATS- STATISTICS

NUMBER OF EQUATIONS	200
OFF-DIAGONAL NONZEROS	398
SIZE OF WORKING STORE (MAXS)	2509
TIME FOR ORDERING	0.083
STORAGE FOR ORDERING	1400.
TIME FOR ALLOCATION	0.030
STORAGE FOR ALLOCATION	2324.
STORAGE FOR SOLUTION	2509.
TIME FOR FACTORIZATION	0.043
TIME FOR SOLVING	0.017
OPERATIONS IN FACTORIZATION	953.
OPERATIONS IN SOLVING	1168.

MAXIMUM ERROR 1.550E-06

Example 6

This is a program to illustrate how one might use SPARSPAK to choose a method. The matrix is 300 by 300, it has nonzeros on the diagonal, the first column and the last row. The structure of the matrix is input using IJBEGN, INIJ and IJEND, and then saved on FORTRAN unit 3. The modules ORDRA1, ORDRA3 and ORDRA5 are then executed, each one followed by a call to PSTATS to obtain the storage information. Note that RESTRT is called after execution of ORDRA1 and ORDRA3, to restore the package to the state that existed immediately after the structure inputting routines were executed. Note also that SAVE could have been used after each ordering module (with different output unit numbers). After one of the methods was chosen, RESTRT (with the appropriate unit number) could be used to initiate the computation, avoiding re-executing the ordering module corresponding to the method chosen.

```
C           M A I N L I N E   P R O G R A M
C
C           INTEGER      I, IERR, MAXS, MSGGLVL, NEQNS
C           REAL         S(7500)
C           COMMON /SPKUSR/ MSGGLVL, IERR, MAXS, NEQNS
C
C           CALL SPKSPK
C           MAXS = 7500
C           CALL IJBEGN
C           DO 100 I = 1, 300
C               CALL INIJ ( I, 1, S )
C               CALL INIJ ( 300, I, S )
100          CONTINUE
C           CALL IJEND ( S )
C           CALL SAVE ( 3, S )
C           CALL ORDRA1 ( S )
C           CALL PSTATS
C           CALL RESTRT ( 3, S )
C           CALL ORDRA3 ( S )
C           CALL PSTATS
C           CALL RESTRT ( 3, S )
C           CALL ORDRA5 ( S )
C           CALL PSTATS
C           STOP
END
```

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OUTPUT UNIT FOR ERROR MESSAGE 6
 OUTPUT UNIT FOR STATISTICS 6

IJBEGN- BEGIN STRUCTURE INPUT...

INIJ- INPUT OF ADJACENCY PAIRS

IJEND- END OF STRUCTURE INPUT

SAVE- STORAGE VECTOR SAVED

ORDRA1- RCM ORDERING

PSTATS- STATISTICS

NUMBER OF EQUATIONS	300
OFF-DIAGONAL NONZEROS	1194
SIZE OF WORKING STORE (MAXS)	7500
TIME FOR ORDERING	0.087
STORAGE FOR ORDERING	2396.
TIME FOR ALLOCATION	0.013
STORAGE FOR ALLOCATION	2396.
STORAGE FOR SOLUTION	2098.

RESTR1- RESTART SYSTEM

ORDRA3- ONE WAY DISSECTION ORDERING

PSTATS- STATISTICS

NUMBER OF EQUATIONS	300
OFF-DIAGONAL NONZEROS	1194
SIZE OF WORKING STORE (MAXS)	7500
TIME FOR ORDERING	0.137
STORAGE FOR ORDERING	3297.
TIME FOR ALLOCATION	0.060
STORAGE FOR ALLOCATION	3300.
STORAGE FOR SOLUTION	3002.

RESTR2- RESTART SYSTEM

ORDRA5- NESTED DISSECTION ORDERING

PSTATS- STATISTICS

NUMBER OF EQUATIONS	300
OFF-DIAGONAL NONZEROS	1194

SIZE OF WORKING STORE (MAXS)	7500
TIME FOR ORDERING	0.123
STORAGE FOR ORDERING	2696.
TIME FOR ALLOCATION	0.043
STORAGE FOR ALLOCATION	3599.
STORAGE FOR SOLUTION	3301.

Appendix A

IMPLEMENTATION OVERVIEW

In this section, we describe briefly the use of labelled common blocks in the internal implementation of SPARSPAK and the various methods of communication between modules.

A.1 USER/MODULE COMMUNICATION

As described in previous sections of this user guide, the user supplies a one-dimensional floating point array S, from which all array storage is allocated. In particular, the interface allocates the first NEQNS storage locations in S for the solution vector of the linear system of equations. After all the interface modules for a particular method have been successfully executed, the user can retrieve the solution from these NEQNS locations.

There is one labelled common block that the user must provide, having four variables:

```
COMMON /SPKUSR/ MSGlvl, IERR, MAXS, NEQNS
```

The variable MAXS is the declared size of the one-dimensional floating point array S and it must be set by user at the beginning of his program. For each module in the interface that allocates storage (e.g. INIJ, IJFND, ORDERxi), MAXS is used to make sure that there is enough storage to carry out the particular phase.

A.2 MODULE/MODULE COMMUNICATION

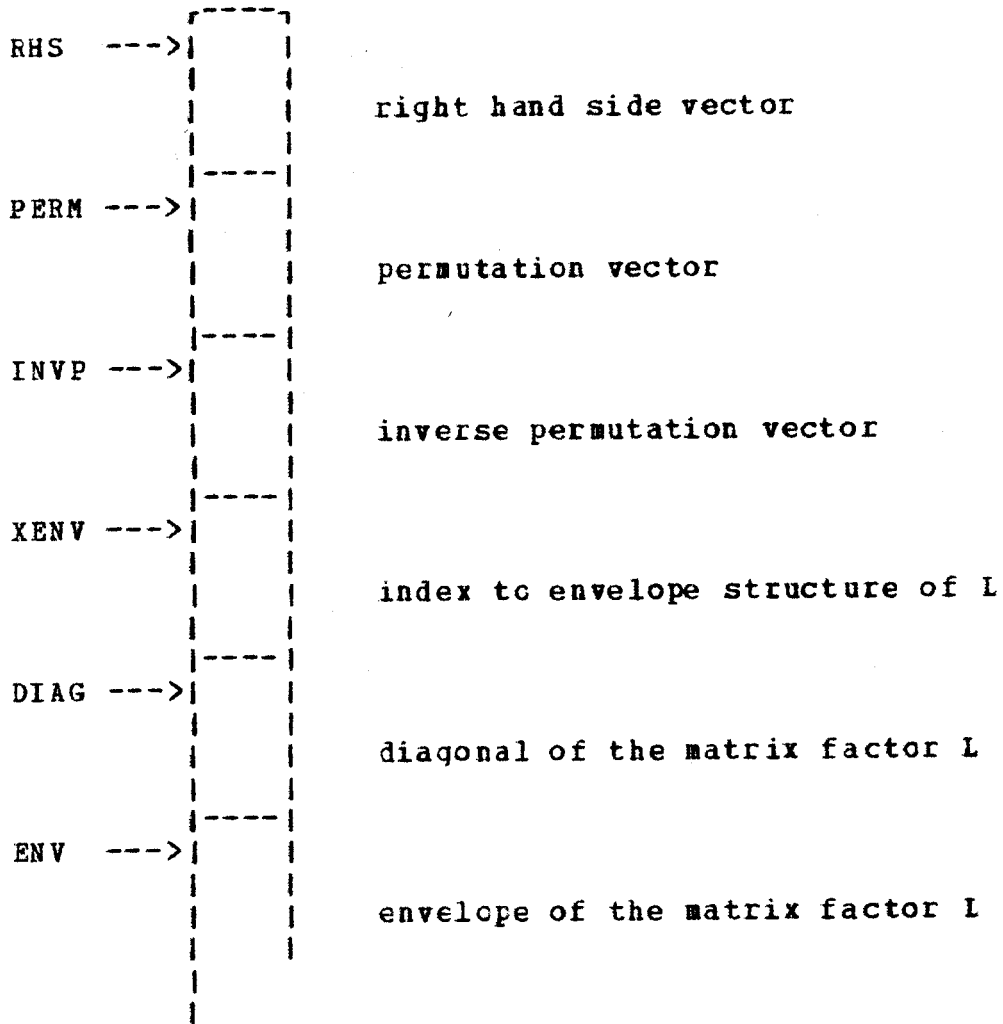
There are several labelled common blocks used for communication among modules within the interface. Two important ones are the control block and the storage map block:

```
COMMON /SPKCON/ STAGE, MXUSED, MXREQD, NEQNS,  
               NEDGES, METHCD, {and other  
               methcd-related control variables}
```

```
COMMON /SPKMAP/ PERM, INVP, RHS,
```

{and other method-related data
structure pointers}

The control block has fourteen integer variables and contains control information about the specific problem being solved. There are fifteen variables in the storage map block, which keep the locations (origins in S) of the various arrays used in the particular storage scheme. These storage schemes differ in complexity across the methods, so the same storage map block must be used in the corresponding routines ORDRxi, INAIJi, INROWi, INMATi, and SOLVEi. An example is given below.



Storage allocation for the symmetric envelope method
(CRIRA1)

A.3 SAVE/RESTART IMPLEMENTATION

The SAVE routine saves the control information in the control block, the storage pointers in the storage map block, as well as the storage vector S. In this way, the state of the computation can be re-established by executing the module RESTRT, which restores the control block and the storage map block, and the storage vector S.

The variable MXUSED in the control block is used to avoid saving irrelevant data from S. After the successful completion of each phase, MXUSED is set to the maximum number of storage locations in S used thus far. It is then only necessary to save the first MXUSED locations of S whenever the routine SAVE is called.

Some operation systems allow a program to change the space it occupies in main storage during execution. Thus, in some installations the user of SPARSPAK may be able to dynamically increase or decrease the size of the working storage S. He can determine what the value of MAXS should be by declaring the labelled common block SPKCON in his mainline program, and examining the value of MXREQD. At the end of each successfully executed phase of the computation, MXREQD is set to the minimum value of MAXS required to successfully execute the next phase of the computation.

It is often the case that when this dynamic growing of program space is provided, the effect is to increase the space allocated to the unlabelled COMMON, which is usually assigned the highest memory locations in the user's program area. In such a circumstance, the array S in the user's program would have to be declared in blank common.

A.4 METHOD CHECKING

As we discussed in the introduction, using a particular 'method' means calling the appropriate interface routines ORDRxi, INAIJi, INROWi, INMATi, and SOLVEi, where the last character is a numerical digit denoting the method. These ordering, input, and solve modules cannot be mixed since they in general involve different data structures. In order to ensure that these modules are not inadvertently mixed by the user, ORDRxi sets the variable METHOD in the control block SPKCON equal to $(10*i + k)$, where k is an integer that distinguishes orderings A and B. This variable is checked by subsequently executed input and solve modules.

A.5 STAGE(SEQUENCE) CHECKING

Another control variable that deserves comment is STAGE. As its name implies, it is used to keep track of the current step or stage of the execution. This variable is particularly important in connection with SAVE and RESTRT modules. In restarting the system using the RESTRT routine, the variable STAGE in the control block SPKCON is restored, and it indicates the last successfully completed stage or phase before the routine SAVE was called. In this way, the execution can be restarted without repeating already successfully completed steps.

Another function of this variable is to enforce the correct execution sequence of the various interface routines. Before the actual execution of each interface routine, the variable STAGE is used to check that all previous interface modules have been successfully completed. This avoids producing erroneous results due to improper processing sequence, or accidental omission of steps.

The content of the variable STAGE is only changed after a phase has been successfully executed. When an error occurs during the execution of the phase, the variable STAGE remains unchanged. This prevents the execution of all the subsequent phases, even if they are invoked by the user. The variable STAGE is also used by the modules to determine whether some initialization is necessary in a module, or whether part of the module has already successfully executed during a previous call to it.

A.6 STORAGE ALLOCATION OF INTEGER AND FLOATING POINT ARRAYS

The ANSI FORTRAN standard specifies that the number of bits used to represent integers and floating point numbers are the same. However, some vendors provide the user with the option of specifying 'short' integers, either explicitly in the declarations such as 'INTEGER*2', or via a parameter to the FORTRAN processor which automatically represents all integers using fewer bits than used for floating point numbers. Since a significant portion of the storage used in sparse matrix computations involves integer data for pointers, subscripts etc., it is desirable to try to exploit these 'short' integer features whenever it makes sense to do so.

SPARSPAK contains parameters RATIOS and RATIOI, set in the module SPRSPK(*), which specify the ratios of the number of bits used for floating point numbers to the number used for 'short' and 'long' integers. For example, in a double precision IBM version of the package which exploits 'short' integers, RATIOS is 4 and RATIOI is 2. Let $U(x)$ be the smallest integer m such that $m \geq x$. The package then uses $RATIOS \cdot RATIOI$ to allocate only $U(p/RATIOS)$ ($U(p/RATIOI)$) elements of S for 'short' ('long') integer arrays of length p .

SPARSPAK assumes that the declaration of S that the user makes in his program is of the same type as that used for floating point computation. We also make the reasonable assumption that $RATIOS \geq 1$ and $RATIOI \geq 1$.

A.7 STATISTICS GATHERING

SPARSPAK contains a labelled common block called SPKDTA which appears below. These variables are used to provide the output described in Section 7.2.

```
COMMON /SPKDTA/  ORDTIM, ALOCTM, FCTIME, SLVTIM, FCTCES  
                SILVCS, ORBSTF, ALOSTR, SLVSTR, OVERHD
```

In order to supply timing information, SPARSPAK assumes the existence of a real function DTIME which returns the processor execution time that has elapsed since DTIME was last referenced. Thus, the DTIME function is also installation dependent.

* Thus SPRSPK is an installation dependent subroutine.

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