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# Generating Functions for the Digital Sum and Other Digit Counting Sequences 

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#### Abstract

A numeration system associates a unique string, $\Xi(n)$, with each positive integer $n$, where each string is over the same finite alphabet. Various digit counting statistics of $\Xi(n)$ are of interest with respect to a numeration system. An example is the digital sum, which is the sum of the digits in the number. We present a unifying framework for deriving identities for the generating functions of such statistics in many of the more popular numeration systems.


## 1 Introduction

Numeration systems provide a rich source of integer sequences. There are many interesting digit counting statistics that arise from the various numeration systems. A typical example, the digital sum, is explained below.

Given a number $n$ represented in binary, $n=\left(b_{d} b_{d-1} \cdots b_{1} b_{0}\right)_{2}$, the (binary) digital sum of $n$, denoted $s_{2}(n)$ is $b_{0}+b_{1}+\cdots+b_{d}$. The digital sum goes under several other names including sideways sum, sideways addition, population count, and Hamming weight. It is denoted $\nu n$ in Knuth [5], but we use the notation used by Allouche and Shallit [1] (but due to earlier researchers; e.g., Bésineau [2] and Coquet and Toffin [3]). There are two natural ways in which we might extend the idea to $k$-ary numbers, either by summing digits, or by counting non-zeros. We use the notation $s_{k}(n)$ (again following [1]) for the sum-of-digits function and $c_{k}(n)$ for the counting non-zeros function.

According to OEIS A053735, the ordinary generating function of $s_{k}(n)$ has the beautiful expression given in Theorem 1 below. For $k=2$ this generating function may be found in Knuth [5]; see exercise 7.1.3.41. The purpose of this short paper is to provide a proof of this and some other related generating functions that are in the OEIS - as part of a more generalized setting. Table 1 contains some of the sequences to which our results apply.

Theorem 1. For all $k \geq 2$,

$$
\begin{aligned}
\sum_{n \geq 0} s_{k}(n) z^{n} & =\frac{1}{1-z} \sum_{m \geq 0} \frac{z^{k^{m}}+2 z^{2 k^{m}}+\cdots+(k-1) z^{(k-1) k^{m}}}{1+z^{k^{m}}+z^{2 k^{m}}+\cdots+z^{(k-1) k^{m}}} \\
& =\frac{1}{1-z} \sum_{m \geq 0} \frac{z^{k^{m}}-k z^{k^{m+1}}+(k-1) z^{(k+1) k^{m}}}{\left(1-z^{k^{m}}\right)\left(1-z^{k^{m+1}}\right)}
\end{aligned}
$$

## 2 Numeration as a sequence of columns

Imagine a table comprised of infinite columns of numbers, $C_{0}, C_{1}, C_{2}, \ldots$ The numbers in each column are indexed starting at 0 and the numbers found in all of the columns all come from the same finite set. In a numeration system each row of the table is distinct. For example, in the binary number system, $C_{j}$ consists of the periodic repetition of $2^{j} 0 \mathrm{~s}$ followed by $2^{j} 1 \mathrm{~s}$.

What is the generating function function for the row sums of those columns? Suppose that the generating function for the $m$-th column is $C_{m}(z)$. Then

$$
\begin{equation*}
A(z)=\sum_{m \geq 0} C_{m}(z) \tag{1}
\end{equation*}
$$

is the generating function for the row sums. That is, $\left\langle z^{n}\right\rangle A(z)$ is the sum of the numbers in the $n$-th row, where $\left\langle z^{n}\right\rangle$ means "coefficient of $z^{n}$ ". The generating function (1) will exist so long as there are constants $c_{n}$ such that $\left\langle z^{n}\right\rangle C_{m}(z)=0$ for all $m \geq c_{n}$. In many numeration systems the $m$-th column can be described as a infinite string of the form $\mathbf{s}_{m} \mathbf{t}_{m}^{\infty}$, where $\mathbf{s}_{m}$ and $\mathbf{t}_{m}$ are strings with $\mathbf{t}_{m} \neq \varepsilon$ and $\mathbf{t}_{m}^{\infty}$ denotes the infinite string $\mathbf{t}_{m}^{\infty}=\mathbf{t}_{m} \mathbf{t}_{m} \mathbf{t}_{m} \cdots$. Let $|\mathbf{s}|$ be the length of the string s. If $S_{m}(z)$ and $T_{m}(z)$ are the generating functions (which are actually polynomials) of $\mathbf{s}$ and $\mathbf{t}$, respectively. Then

$$
\begin{equation*}
C_{m}(z)=S_{m}(z)+\frac{z^{\left|\mathbf{s}_{m}\right|} T_{m}(z)}{1-z^{\left|\mathbf{t}_{m}\right|}} \tag{2}
\end{equation*}
$$

Often $\mathbf{s}_{m}$ and $\mathbf{t}_{m}$ will have a special form that allows for further simplification of $C_{m}(z)$
In this paper the most general from that we use is shown below. Here the $m$-th column depends on integers $b_{m}, a_{m}$ and $u_{m}$, and the sequence of numbers $\alpha_{0}, \alpha_{1}, \alpha_{2}, \ldots$ The following string is called the column pattern:

$$
\begin{equation*}
\mathbf{s}_{m} \mathbf{t}_{m}^{\infty}=0^{b_{m}}\left(\alpha_{0}^{a_{m}} \alpha_{1}^{a_{m}} \cdots \alpha_{\left(u_{m}-1\right)}^{a_{m}}\right)^{\infty} \tag{3}
\end{equation*}
$$

For example, the column pattern for the binary digital sum is $\left(0^{2^{m}} 1^{2^{m}}\right)^{\infty}$; here $b_{m}=0$, $a_{m}=2^{m}, \alpha_{0}=1, \alpha_{1}=1$, and $u_{m}=2$. The generating function of $\alpha_{0}^{a_{m}} \alpha_{1}^{a_{m}} \cdots \alpha_{\left(u_{m}-1\right)}^{a_{m}}$ is

$$
A_{m}(z)=\left(\alpha_{0}+\alpha_{1} z^{a_{m}}+\alpha_{1} z^{2 a_{m}}+\cdots+\alpha_{u_{m}-1} z^{\left(u_{m}-1\right) a_{m}}\right) \frac{1-z^{a_{m}}}{1-z}
$$

Thus, by (2), the generating function for (3) is

$$
\frac{z^{b_{m}} A_{m}(z)}{1-z^{u_{m} a_{m}}}=\frac{z^{b_{m}}}{1-z} \cdot \frac{\alpha_{0}+\alpha_{1} z^{a_{m}}+\alpha_{2} z^{2 a_{m}}+\cdots+\alpha_{u_{m}-1} z^{\left(u_{m}-1\right) a_{m}}}{1+z^{a_{m}}+z^{2 a_{m}}+\cdots+z^{\left(u_{m}-1\right) a_{m}}} .
$$

Summing over $m \geq 0$ we obtain

$$
\begin{align*}
A(z) & =\frac{1}{1-z} \sum_{m \geq 0} \frac{z^{b_{m}}\left(1-z^{a_{m}}\right)}{1-z^{u_{m} a_{m}}}\left(\alpha_{0}+\alpha_{1} z^{a_{m}}+\alpha_{2} z^{2 a_{m}}+\cdots+\alpha_{u_{m}-1} z^{\left(u_{m}-1\right) a_{m}}\right)  \tag{4}\\
& =\frac{1}{1-z} \sum_{m \geq 0} z^{b_{m}} \frac{\alpha_{0}+\alpha_{1} z^{a_{m}}+\alpha_{2} z^{a_{m}}+\cdots+\alpha_{u_{m}-1} z^{\left(u_{m}-1\right) a_{m}}}{1+z^{a_{m}}+z^{2 a_{m}}+\cdots+z^{\left(u_{m}-1\right) a_{m}}} . \tag{5}
\end{align*}
$$

In the sections to follow we apply this generating function to various numeration systems, starting with a new addition to the OEIS.

## 3 The Balanced Ternary System

In the balanced ternary system each natural number $n$ is expressed as a sum of distinct signed powers of 3 . For example $5=9-3-1=3^{3}-3^{1}-3^{0}$. The digital sum is OEIS A065363. Following Knuth [4] we use $\overline{1}$ to denote -1 . It is outside the scope of this paper, but it is not difficult to show that the pattern of the $m$-th column is

$$
0^{3^{m}-3^{m-1}-\cdots-3^{0}}\left(1^{3^{m}} \overline{1^{m}} 0^{3^{m}}\right)^{\infty} .
$$

Since $3^{m}-3^{m-1}-\cdots-3^{0}=\left(3^{m}+1\right) / 2$, the generating function for the sum of the digits of the balanced ternary representation of $n$ is

$$
A(z)=\frac{1}{1-z} \sum_{m \geq 0} z^{\left(3^{m}+1\right) / 2} \frac{1-z^{3^{m}}}{1+z^{3^{m}}+z^{2 \cdot 3^{m}}}=\frac{1}{1-z} \sum_{m \geq 0} z^{\left(3^{m}+1\right) / 2} \frac{\left(1-z^{3^{m}}\right)^{2}}{1-z^{3^{m+1}}} .
$$

## 4 The $k$-ary Numeration System and Morphisms

The column pattern for $k$-ary numbers is

$$
\begin{equation*}
\left(\alpha_{0}^{k^{m}} \alpha_{1}^{k^{m}} \ldots \alpha_{k-1}^{k^{m}}\right)^{\infty} . \tag{6}
\end{equation*}
$$

Here the row sum of the $n$-th row, where

$$
n=\sum_{k \geq 0} b_{k} k^{m}, \quad \text { is } \quad \sum_{k \geq 0} \alpha_{b_{k}} k^{m} .
$$

That is, each digit $b$ is "weighted" by $\alpha_{b}$. For the digital sum, $\alpha_{b}=b$.
The special form of (6) implies that there is a "morphism" that underlies the construction; for the digital sum it is $j \rightarrow j, j+1, j+2, \ldots, j+k-1$. For example, when $k=3$ we get the sequence $s_{3}(0), s_{3}(1), s_{3}(2), \ldots$ as the limit (i.e., fixed-point) of a morphism noted by Robert G. Wilson in A053735, which gives us successively

$$
0 \rightarrow 012 \rightarrow 012123234 \rightarrow 012123234123234345234345456 \rightarrow \cdots
$$

This limit will exist for any morphism of the form

$$
\begin{equation*}
j \rightarrow j+\alpha_{0}, j+\alpha_{1}, \ldots, j+\alpha_{k-1} \tag{7}
\end{equation*}
$$

so long as $\alpha_{0}=0$.
Theorem 2. The row sums of the column pattern (6) are generated by the morphism (7) so long as $\alpha_{0}=0$.

Proof. The column pattern (6) is invariant under the following two-step operation: (a) Take column $m$ and replace each entry in the column by $k$ identical entries, calling the new column $C_{m+1}^{\prime}$. (b) Form a new column $C_{0}^{\prime}$ with the pattern $(01 \cdots(k-1))^{\infty}$. The invariance is that $C_{m}^{\prime}=C_{m}$ for $m=0,1,2, \ldots$.

A row sum $j$ under operations (a) and (b) becomes the $k$ row sums $j+\alpha_{0}, j+\alpha_{1}, \ldots, j+$ $\alpha_{k-1}$. This is the morphism (7).

With the pattern (6) equation (5) gives us the theorem below.
Theorem 3. If $k$ is an integer with $k \geq 2$ and $\alpha_{0}=0$, then the generating function of the limit of the morphism (7) is

$$
\begin{equation*}
A(z)=\frac{1}{1-z} \sum_{m \geq 0} \frac{\alpha_{1} z^{k^{m}}+\alpha_{2} z^{2 k^{m}}+\cdots+\alpha_{k-1} z^{(k-1) z^{m}}}{1+z^{k^{m}}+z^{2 k^{m}}+\cdots+z^{(k-1) k^{m}}} \tag{8}
\end{equation*}
$$

Note that Theorem 1 is the special case where $\alpha_{i}=i$ for $i=0,1, \ldots, k-1$. The second equality in Theorem 1 follows from the fact that $z+2 z+\cdots+(k-1) z^{k-1}=$ $\left(z-k z^{k}+(k-1) z^{k+1}\right) /(1-z)^{2}$. We now return to the non-zero count function, $c_{k}(n)$, which can be expressed without the inner sums used in (8).
Corollary 4. The generating function of $c_{k}(n)$ is

$$
\begin{equation*}
C_{k}(z)=\frac{1}{1-z} \sum_{m \geq 1} \frac{z^{k^{m-1}}-z^{k^{m}}}{1-z^{k^{m}}} \tag{9}
\end{equation*}
$$

Proof. Here the morphism is $j \rightarrow j, j+1, \ldots, j+1$ and so Theorem 3 gives us the first equality below.

$$
\begin{align*}
C_{k}(z) & =\frac{1}{1-z} \sum_{m \geq 0} \frac{z^{k^{m}}+z^{2 k^{m}}+\cdots+z^{(k-1) k^{m}}}{1+z^{k^{m}}+z^{2 k^{m}}+\cdots+z^{(k-1) k^{m}}}  \tag{10}\\
& =\frac{1}{1-z} \sum_{m \geq 0} \frac{\left(1-z^{k^{m+1}}\right) /\left(1-z^{k^{m}}\right)-\left(1-z^{k^{m}}\right) /\left(1-z^{k^{m}}\right)}{\left(1-z^{k^{m+1}}\right) /\left(1-z^{k^{m}}\right)}
\end{align*}
$$

Cancelling common denominators and simplifying gives (9).

Other morphisms would give counts of the number of times individual digits occur in the obvious way. For example, $j \rightarrow j, j, j+1, j$ is the morphism for the number of 2 's that occur in the 4 -ary expansion of $n$ (here the pattern is $\left.\left(0^{4^{m}} 0^{4^{m}} 1^{4^{m}} 0^{4^{m}}\right)^{\infty}\right)$.

Theorem 5. Let $d$ be an integer with $0<d<k$. The generating function for the number of digits equal to $d$ in the $k$-ary expansion of $n$ is

$$
\begin{equation*}
\frac{1}{1-z} \sum_{m \geq 0} \frac{z^{d k^{m}}}{1+z^{k^{m}}+z^{2 k^{m}}+\cdots+z^{(k-1) k^{m}}}=\frac{1}{1-z} \sum_{m \geq 0} \frac{z^{d k^{m}}\left(1-z^{k^{m}}\right)}{1-z^{k^{m+1}}} . \tag{11}
\end{equation*}
$$

The generating function for the number of 0 digits in the $k$-ary expansion of $n$ is

$$
\begin{equation*}
\frac{1}{1-z} \sum_{m \geq 0} \frac{z^{k^{m+1}}}{1+z^{k^{m}}+z^{2 k^{m}}+\cdots+z^{(k-1) k^{m}}}=\frac{1}{1-z} \sum_{m \geq 0} \frac{z^{k^{m+1}}\left(1-z^{k^{m}}\right)}{1-z^{k^{m+1}}} . \tag{12}
\end{equation*}
$$

Proof. Equation (11) follows from Theorem 3 with the morphism $j \rightarrow j, \ldots, j, j+1, j, \ldots, j$ where the $j+1$ occurs in position $d$, counting from 0 . To prove (12) we use the generating function

$$
T(z)=\frac{1}{1-z} \sum_{m \geq 0} z^{k^{m}}
$$

for $1+\left\lfloor\log _{k} n\right\rfloor$, which is the number of $k$-ary digits in $n$. Adding (10) and (12) we clearly obtain $T(z)$.

A second way of finishing the proof is to note that the column pattern

$$
0^{k^{m}}\left(1^{k^{m}} 2^{k^{m}} \cdots(k-1)^{k^{m}} 0^{k^{m}}\right)^{\infty}
$$

also describes the $k$-ary listing of numbers. The useful aspect of expressing it this way is that the leading 0s are correspond to the initial $0^{k^{m}}$ above. Thus the pattern for counting (non-leading) 0s is

$$
0^{k^{m}}\left(0^{k^{m}} 0^{k^{m}} \cdots 0^{k^{m}} 1^{k^{m}}\right)^{\infty}
$$

According to (5) the numerator inside the sum of the generating function is $z^{b_{m}} z^{\left(u_{m}-1\right) a_{m}}=$ $z^{k^{m}} z^{(k-1) k^{m}}=z^{k^{m+1}}$, as desired.

### 4.1 Digit counts in specific positions

Let $C(n, r, d)$ be the number of 1 bits in the binary representation of $n$ that are in positions that are congruent to $r \bmod d$. As usual, the "positions" are indexed starting at 0 on the right. For example, $888=(1101111000)_{2}$, so $C(888,0,3)=3, C(888,1,3)=1$ and $C(888,2,3)=2$.

Theorem 6. For all integers $d \geq 0$ and integers $r$ with $0 \leq r<d$,

$$
\begin{equation*}
\sum_{n \geq 0} C(n, r, d) z^{n}=\sum_{m \geq 0} \frac{z^{2^{r+d m}}}{1+z^{2 r+d m}} \tag{13}
\end{equation*}
$$

| OEIS | description | comment | pattern |
| :---: | :---: | :---: | :---: |
| A023416 | 0s in binary | same as A080791 | $0^{2^{m}}\left(0^{2^{m}} 1^{2^{m}}\right)^{\infty}$ |
| A000120 | 1s in binary | digital sum, base 2 | $\left(0^{2^{m}} 1^{2^{m}}\right)^{\infty}$ |
| A077267 | 0s in base 3 | same as A081602 | $0^{3^{m}}\left(0^{3^{m}} 0^{3^{m}} 1^{3^{m}}\right)^{\infty}$ |
| A062756 | 1 s in base 3 |  | $\left(0^{3^{m}} 1^{3^{m}} 0^{3^{m}}\right)^{\infty}$ |
| A081603 | 2 s in base 3 |  | $\left(0^{3^{m}} 0^{3^{m}} 1^{3^{m}}\right)^{\infty}$ |
| A160380 | 0s in base 4 |  | $0^{4^{m}}\left(0^{4^{m}} 0^{4^{m}} 0^{4^{m}} 1^{4^{m}}\right)^{\infty}$ |
| A160381 | 1 s in base 4 |  | $\left(0^{4^{m}} 1^{4^{m}} 0^{4^{m}} 0^{4^{m}}\right)^{\infty}$ |
| A160382 | 2 s in base 4 |  | $\left(0^{4^{m}} 0^{4^{m}} 1^{4^{m}} 0^{4^{m}}\right)^{\infty}$ |
| A160383 | 3 s in base 4 |  | $\left(0^{4^{m}} 0^{4^{m}} 0^{4^{m}} 1^{4^{m}}\right)^{\infty}$ |
| A160384 | non-0 base 3 |  | $\left(0^{3^{m}} 1^{3^{m}} 1^{3^{m}}\right)^{\infty}$ |
| A160385 | non-0 base 4 |  | $\left(0^{4^{m}} 1^{4^{m}} 1^{4^{m}} 1^{4^{m}}\right)^{\infty}$ |
| A053735 | digital, base 3 |  | $\left(0^{3^{m}} 1^{3^{m}} 2^{3^{m}}\right)^{\infty}$ |
| A053737 | digital, base 4 |  | $\left(0^{4^{m}} 1^{4^{m}} 2^{4^{m}} 3^{4^{m}}\right)^{\infty}$ |
| A034968 | digital, factorial base | see also $\underline{\text { A139365 }}$ |  |
| A065363 | digital, balanced base 3 |  | $0^{\left(3^{m}+1\right) / 2}\left(1^{3^{m}} \overline{1}^{3^{m}} 0^{3^{m}}\right)^{\infty}$ |
| A139351 | 1 s or 3 s in base 4 | 1s in even positions in binary | $\left(0^{4^{m}} 1^{4^{m}} 0^{4^{m}} 1^{4^{m}}\right)^{\infty}$ |
| A139352 | 2 s or 3 s in base 4 | 1 s in odd positions in binary | $\left(0^{4^{m}} 0^{4^{m}} 1^{4^{m}} 1^{4^{m}}\right)^{\infty}$ |

Table 1: Relevant sequences in OEIS.

Proof. Let $B(r, d)=\left\{s \in \mathbb{Z}: 1 \leq s<2^{d}\right.$ and $\left\lfloor s / 2^{r}\right\rfloor$ is odd $\}$; in other words, the $d$ bit binary numbers with the $r$-th bit equal to 1 . Let $\bar{B}(r, d)=\left\{0,1, \ldots, 2^{d}-1\right\} \backslash B(r, d)$. For example, $B(1,3)=\{2,3,6,7\}$ and $B(1,3)=\{0,1,4,5\}$.

Now consider the number $n$ written both in binary and in base $2^{d}$. Note that, in the binary representation of $n$, the number of 1 bits in positions that are congruent to $r \bmod d$ is the same as the number of digits from the set $B(r, d)$ in the $2^{d}$-ary representation of $n$. Thus we may apply Theorem 3 to get the generating function

$$
\sum_{n \geq 0} C(n, r, d) z^{n}=\sum_{m \geq 0} \frac{\sum_{s \in B(r, d)} z^{s 2^{d m}}}{\sum_{0 \leq s<2^{d}} z^{s 2^{d m}}}
$$

Note that the numerator above can be written

$$
\sum_{s \in B(r, d)} z^{s 2^{d m}}=z^{2^{r} 2^{d m}} \sum_{s \in \bar{B}(r, d)} z^{s 2^{d m}}
$$

and the denominator as

$$
\sum_{0 \leq s<d} z^{s 2^{d m}}=\sum_{s \in B(r, d)} z^{s 2^{d m}}+\sum_{s \in \bar{B}(r, d)} z^{s 2^{d m}}=\left(1+z^{2^{r} 2^{d m}}\right) \sum_{s \in \bar{B}(r, d)} z^{s 2^{d m}} .
$$

Canceling the common sum gives the right hand side of (13).
The following corollary allows us to give a generating function for A139351 and A139352.

Corollary 7. The generating function for the number of 1's in even positions in the binary expansion of $n$, and the corresponding generating function for the number of 1's in odd positions, are given below.

$$
\frac{1}{1-z} \sum_{m \geq 0} \frac{z^{4^{m}}}{1+z^{4^{m}}}, \quad \frac{1}{1-z} \sum_{m \geq 0} \frac{z^{2 \cdot 4^{m}}}{1+z^{2 \cdot 4^{m}}}
$$

## 5 Multi-Radix Numeration Systems

In this section we consider numbers written in the multi-radix base $k_{0} \times k_{1} \times k_{2} \times \cdots$. If each $k_{i}=k$ then we get the $k$-ary numeration system considered in the previous section. It will prove useful to adopt the following notation: (a) $k_{j}^{\prime}=k_{j}-1$, (b) $\bar{k}_{j}=k_{0} k_{1} \cdots k_{j-1}$, with the usual convention for the empty product, $\bar{k}_{0}=1$. Then the column pattern is

$$
\left(\alpha_{0}^{\bar{k}_{m}} \alpha_{1}^{\bar{k}_{m}} \cdots \alpha_{k_{m}^{\prime}}^{\bar{k}_{m}}\right)^{\infty} .
$$

Theorem 8. The generating function for the digital sum of the number $n$ written in the multi-radix base $k_{0} \times k_{1} \times k_{2} \times \cdots$ is

$$
\frac{1}{1-z} \sum_{m \geq 0} \frac{z^{\bar{k}_{m}}+2 z^{2 \bar{k}_{m}}+\cdots+k_{m}^{\prime} z^{z_{m}^{\prime} \cdot \bar{k}_{m}}}{1+z^{\bar{k}_{m}}+z^{2 \bar{k}_{m}}+\cdots+z^{k_{m}^{\prime} \cdot \bar{k}_{m}}} .
$$

Proof. The generating function is (5) with $b_{m}=0, \alpha_{i}=i, u_{m}=k_{m}^{\prime}$, and $a_{m}=\bar{k}_{m}$.
In the factorial base, $k_{j}=j+1$, so that $\bar{k}_{j}=j$ !. For example, $99=3 \cdot 4!+0 \cdot 3!+2 \cdot 2!+1 \cdot 1$ !. We now obtain a generating function for A034968 in the following corollary.

Corollary 9. The generating function for the digital sum of the number $n$ written in the factorial base is

$$
\frac{1}{1-z} \sum_{m \geq 1} \frac{z^{m!}+2 z^{2 m!}+\cdots+m z^{m \cdot m!}}{1+z^{m!}+z^{2 m!}+\cdots+z^{m \cdot m!}}
$$

Proof. This follows directly from the previous theorem. Note that the numerator is zero when $m=0$, so that the summation starts at 1 .

## 6 Final Remarks

It is also possible to approach the derivation of the generating functions used here using "divide-and-conquer" recurrence relations. See Stephan [8] for examples of this approach in the $k=2$ case. For example the recurrence relation corresponding to the morphism (7) is $a(0)=0$ and $a(k m+i)=\alpha_{i}+a(m)$ for integer indices $i$ with $0 \leq i<k$. These recurrence relations are very useful for actually computing the sequences.

As we have shown here, finding a generating function for the sum of the digits is straightforward when dealing with a simple radix, or a mixed radix system where each positional
multiplier is a multiple of the previous one. When this does not hold, the problem is much more difficult.

The simplest example of such a system is the Zeckendorf [9] or "base Fibonacci" representation (A014417, digital sum in A007895). Attempting the same sort of one digit at a time approach, the low order digit is the infinite Fibonacci word, A003849. Since this sequence includes arbitrarily long repeated segments, but is not periodic, it does not have a rational generating function.

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