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If  $u$  and  $v$  have the same set of subwords of length  $m$ , they can be intertwined in such a way that no new subword of length  $m$  appears.

São Paulo, April 30, 1981

Dear Jacques

Just for a change I am answering your two letters promptly. The card you sent me is really beautiful, thanks a lot.

① I agree with your remarks on Dominique's editorial job with one exception: I also do not like the word definition, but as you noted it is a matter of taste. Besides, I think that uniformity throughout the book must <sup>have</sup> been the reason for his suppression of all our definitions. Anyway... The constructiveness thing I put in Note A.

② Exercise 2.2. The notation  $f \equiv 0 [3m]$  is great. I propose to put it like this. ~~Not in Note A.~~

a)  $f \equiv 0 [3m] \Rightarrow \text{~~some condition~~ } |f| \geq km$

b)  ~~$f \equiv 0 [3m]$~~   
 Let  $f \in A^{km}$ . Then  $f$  has a factorization  $f = f_1 f_2 \dots f_k$  with  $\text{alph}(f_1) = \text{alph}(f_2) = \dots = \text{alph}(f_k) = A$

c)  ~~$f \equiv 0 [3m]$~~   
 Let  $f, g \in A^*$ , with  $A = \{a, b\}$ .  
 (i)  $af \equiv bg [3m] \Rightarrow af \equiv bg \equiv 0 [3m]$

Let  $A$  be a  $k$ -letter alphabet. We note  $f \equiv 0 [3m]$   
 if  $S(m, f) = A^{\leq m}$ . Show that

(a)  $f \equiv 0 [3m] \iff fg \equiv gf \equiv 0 [3m]$  for all  $g \in A^*$ .

(b) If  $k=2$  and  $A=\{a, b\}$  then for any  $f, g \in A^*$

(b.1)  $af \equiv bg \equiv 0 [3m] \implies af \equiv bg \equiv 0 [3m]$

(b.2)  $abf \equiv 0 [3m] \iff f \equiv 0 [3m-1]$

(c)  $f \equiv 0 [3m] \implies |f| = km$

(d) Let  $f \in A^{km}$ . Then

$f \equiv 0 [3m] \iff f$  has a factorization  
 $f = f_1 f_2 \dots f_k$  with  
 $\text{alph}(f_1) = \text{alph}(f_2) = \dots = \text{alph}(f_k) = A$ .

③ About Exercise 2.3: I am getting excited  
 with your ideas. I always thought that  
 enumerating the classes is ~~at~~ out of reach but  
 I might have changed my mind. I have  
 the impression that this thing is connected with  
 the Schutzenberger-Simon theorem which I call  
 the reconstruction theory. I confess that  
 my ideas are fuzzy and the whole thing might  
 be nonsense. Anyway if you keep interested in



UNDECIDABILITY  
AND THE  
TROPICAL SEMIRING  
A LA  
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$F_0, F_1, \dots$

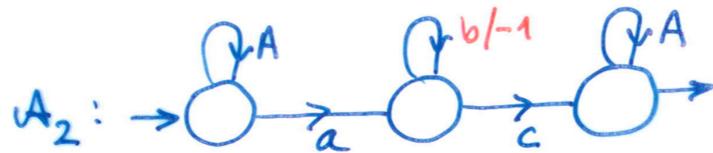
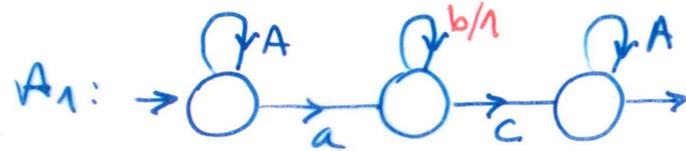
$$wF_0 = \begin{cases} 0 & \text{if } w \in U \\ -1 & \text{otherwise} \end{cases}$$

will allow to define  $wF_i$  on words in  $\bar{U}$  provided  $wF_i \leq 0$  (arbitrarily)

$$F_1 = \|A_1\| \odot \|A_2\| \quad sF_1 = \begin{cases} w_{MIN} - w_{MAX} & \text{if } s \in I \\ \leq 0 & \text{if } s \notin I \end{cases}$$

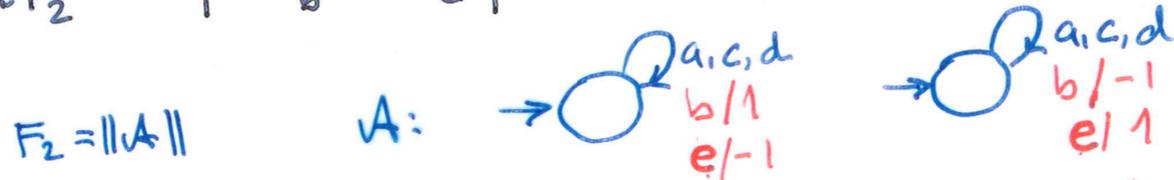
$$w = s(de^+d)^{-1}$$

For  $s \in \bar{U}$   $sF_1 = 0$  iff  $s \in C$



deterministic automaton assures  $wF_1 = 0$  on uninteresting words. Belongs to both  $A_1$  &  $A_2$

$$sF_2 = - | |s|_b - |s|_e | \quad sF_2 = \min \{ |s|_b - |s|_e, -|s|_b + |s|_e \}$$



For  $s \in C$   $sF_2 = 0$  iff  $s \in S$

For  $p$  composite use:  $F' = F_0 \odot F_1 \odot F_2 \odot F_3 \odot F_4 \odot F_5$

with  $sF_3 = 0$  iff  $z = p$   
 $sF_4 = 0$  iff  $x > 1$   
 $sF_5 = 0$  iff  $y > 1$

YET ANOTHER PROOF

OF HASHIGUCHI'S

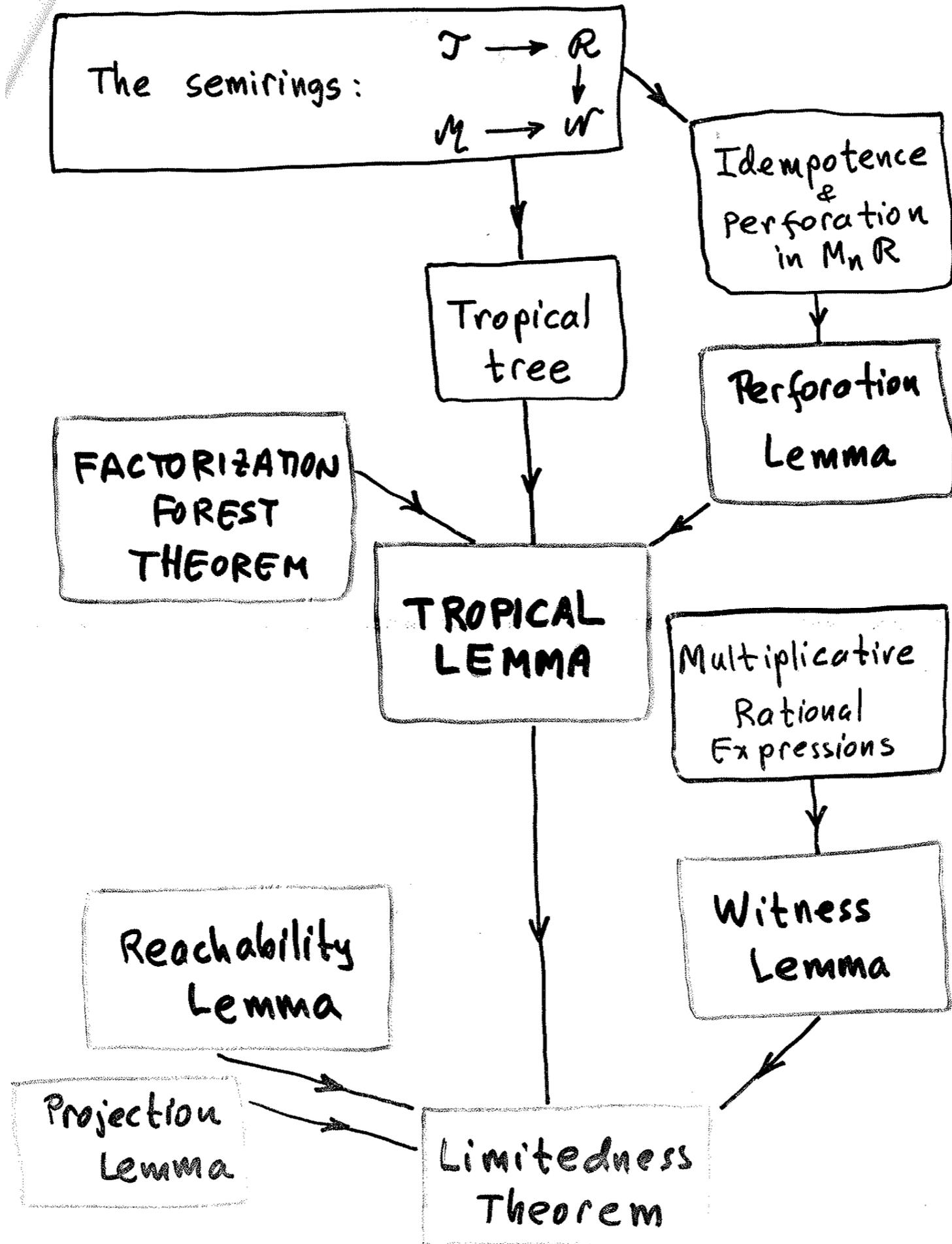
LIMITEDNESS THEOREMS

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# An overview



# LIMITED SUBSETS OF A FREE MONOID

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## 1. INTRODUCTION

While attending the 7<sup>th</sup> SWAT Conference in 1966, J.A.Brzozowski formulated the following problem [3]:

"Is it decidable whether for a given regular set  $R$ ,  $R^* = (\lambda \cup R)^m$  for some integer  $m \geq 1$ ?"

Even though some effort has been made to clarify the above question, it has been answered only for a few particular subclasses of regular sets [10,11,12,17]. The main objective of the present paper is to give a positive answer to Brzozowski's problem. The algorithm we present

iff there exists an integer  $m \geq 1$ , such that  $A^* = (1 \cup A)^m$ . Since

$$A^* = 1 \cup A \cup A^2 \cup \dots \cup A^n \cup \dots, \quad (1)$$

while

$$(1 \cup A)^m = 1 \cup A \cup A^2 \cup \dots \cup A^m, \quad (2)$$

it follows that  $A$  is limited iff the infinite union in (1) can be replaced by the finite union in (2), for some  $m$ . Equivalently,  $A$  is limited iff any concatenation of a finite number of elements of  $A$  can also be obtained by the concatenation of at most  $m$  elements of  $A$ .

Let us consider the set  $M = \mathbb{N} \cup \infty$ , where  $\mathbb{N}$  is the set of all integers  $n \geq 0$ . We extend to  $M$  addition, multiplication and order of  $\mathbb{N}$  in the usual way ( $0 \cdot \infty = \infty \cdot 0 = 0$ ). In the sequel we consider  $M$  as a semiring [6, p. 122] with operations

$$a \oplus b = \min\{a, b\} \text{ and } a \otimes b = a + b.$$

We will also need the semiring  $N$  with  $N = \{0, 1, \infty\}$  and with operations

$$a \oplus b = \min\{a, b\} \text{ and } a \otimes b = \begin{cases} a + b & \text{if } a + b \in N \\ 1 & \text{otherwise.} \end{cases}$$

Notice that  $N$  is a homomorphic image of  $M$ , and that

#### 4. LOCALLY FINITE SEMIGROUPS

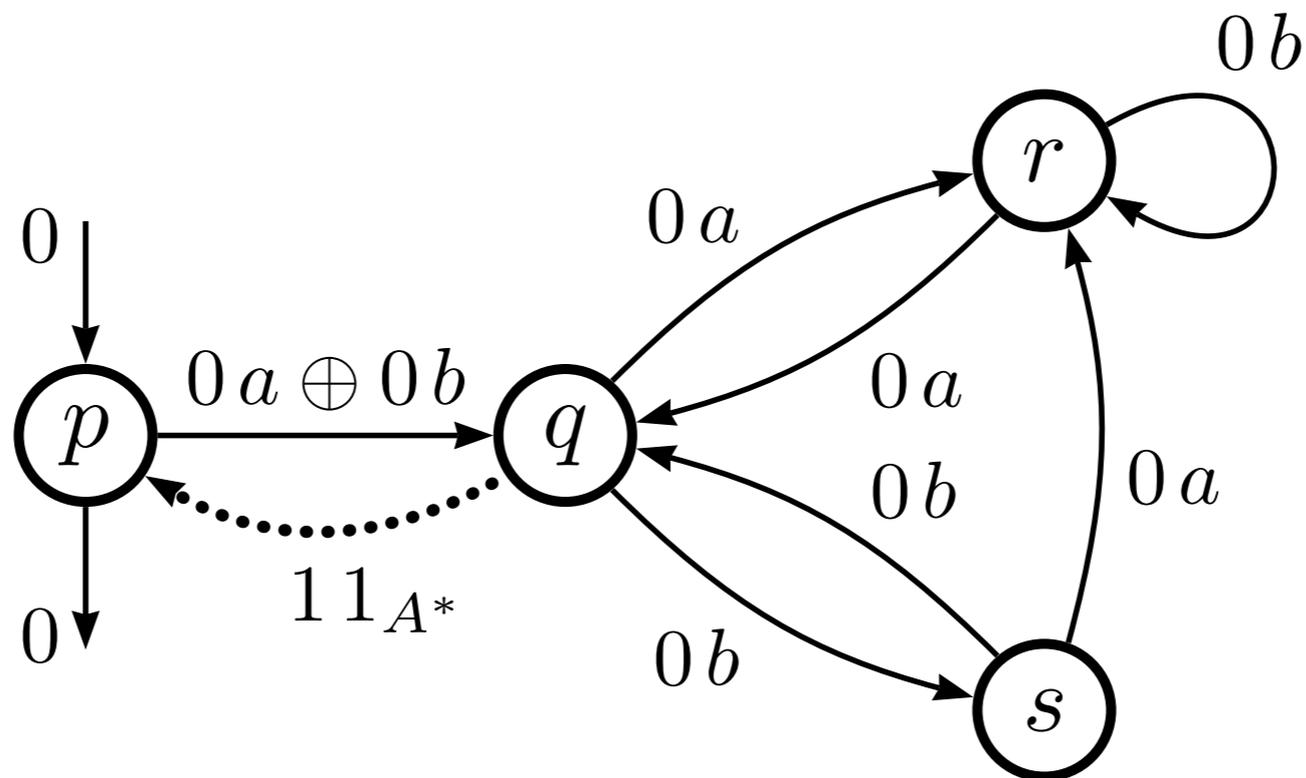
A semigroup  $S$  is *locally finite* if every finitely generated subsemigroup of  $S$  is finite. A semigroup  $S$  is *torsion* if every element in  $S$  generates a finite subsemigroup of  $S$ .

The next result is the key fact in the proofs of Theorems A and B.

THEOREM C. Every torsion subsemigroup of  $M_n M$  is locally finite.

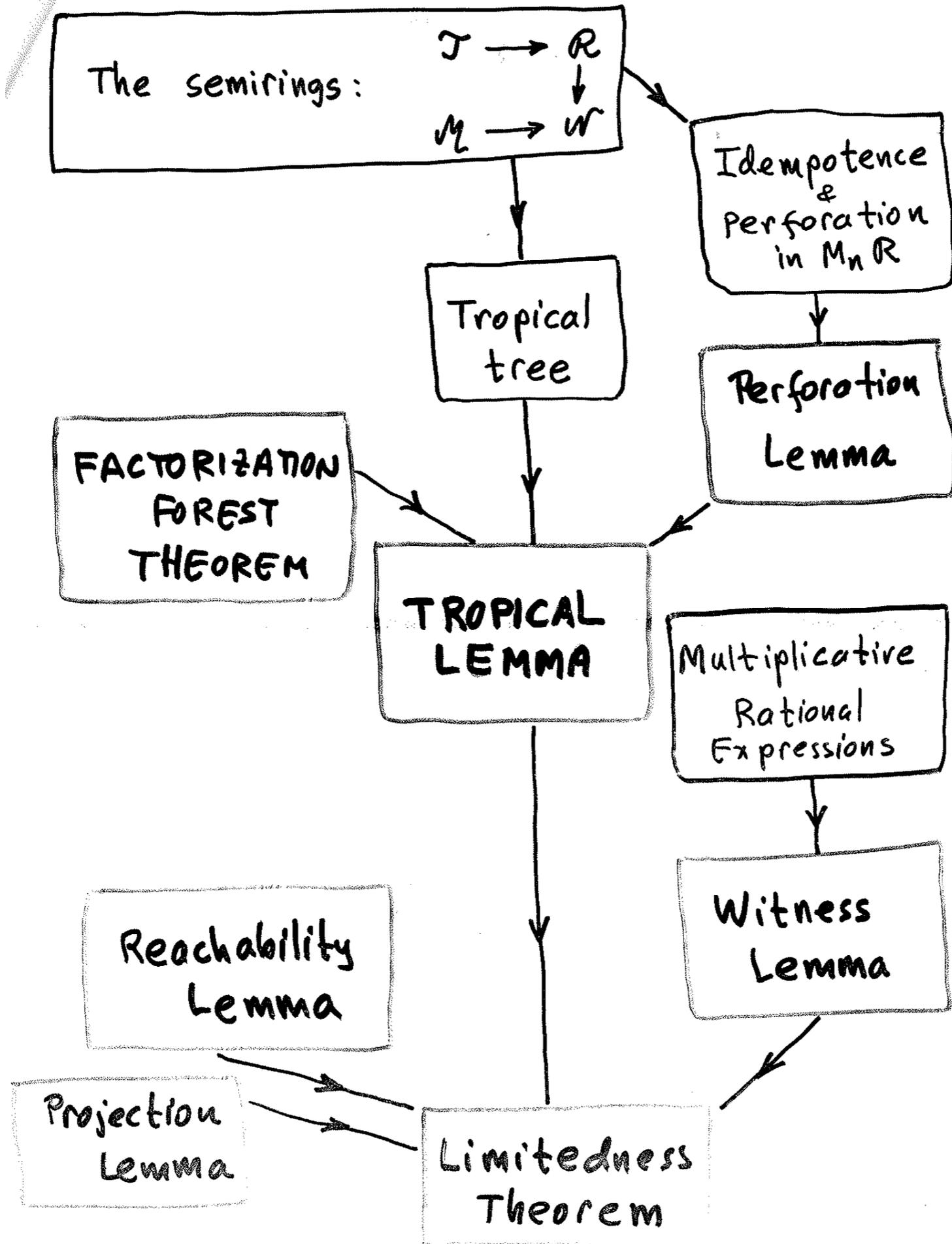
This result is interesting on its own right. Indeed, Thue's Theorem mentioned in

$$S_3 = (a + b)(bb + ab^*a + bab^*a)^*$$



$\mathcal{M}$ -automaton  $S_3$

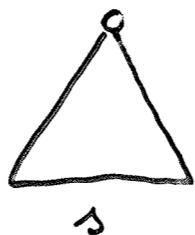
# An overview



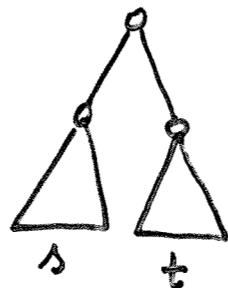
# Factorization forests

Let  $f: A^+ \rightarrow S$  be a morphism, with  $S$  finite

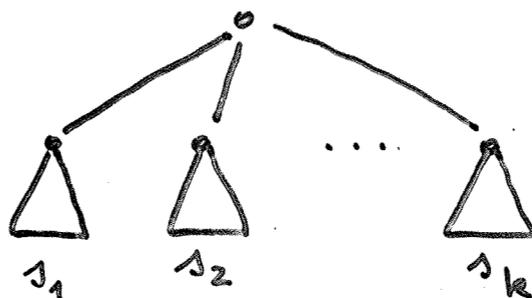
For every  $s \in A^+$  we construct a tree with frond  $s$  using 2 rules:



Rule 1

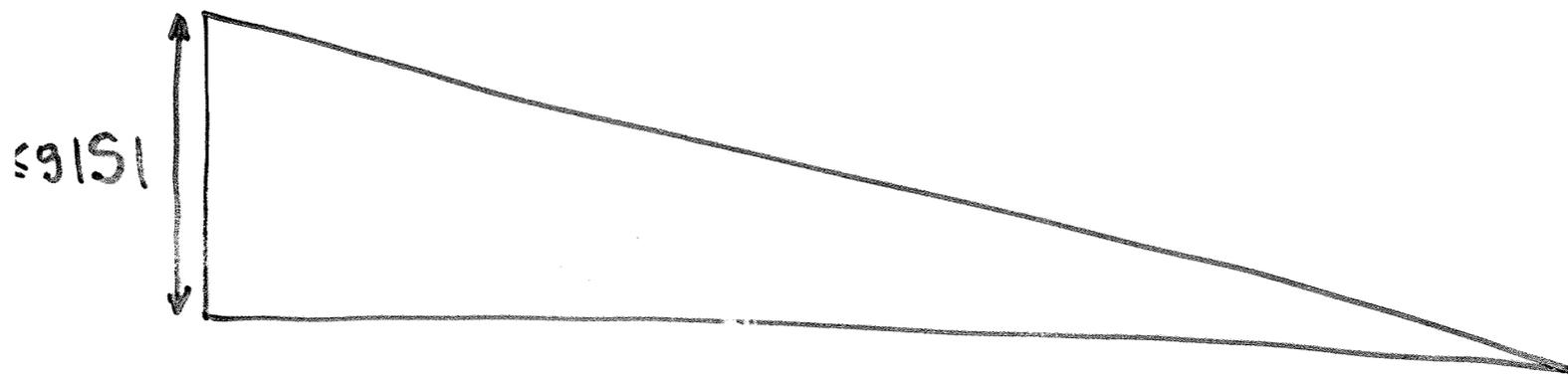


Rule 2

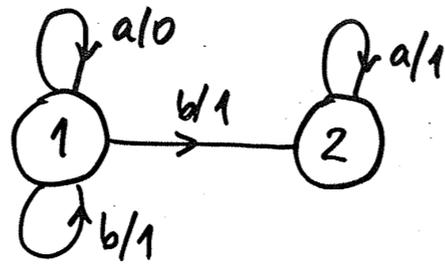


$$s_1 f = s_2 f = \dots = s_k f = e = e^2 \in S$$

**THEOREM** Every morphism  $f: A^+ \rightarrow S$ ,  $S$  finite admits a Ramseyan factorization forest of height at most  $g|S|$



# An example



$$A^+ \xrightarrow{\psi} M_n \mathcal{J}$$

$$\downarrow \quad \quad \downarrow \psi$$

$$B^+ \xrightarrow{f} M_n \mathcal{R}$$

$$A^+ \psi = \{a, b\}$$

$$A = \{a, b\} \quad af = \begin{bmatrix} 0 & \infty \\ \infty & 1 \end{bmatrix} \quad bf = \begin{bmatrix} 1 & 1 \\ \infty & \infty \end{bmatrix}$$



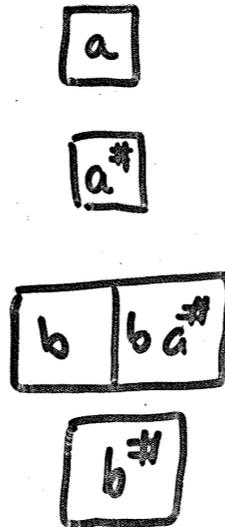
$$B = \{a, b, a^\#, b^\#\} \quad a^\#f = \begin{bmatrix} 0 & \infty \\ \infty & w \end{bmatrix} \quad b^\#f = \begin{bmatrix} w & w \\ \infty & \infty \end{bmatrix}$$

$$S = B^+f = \{a, b, a^\#, b^\#, ba^\#\}f \quad ba^\#f = \begin{bmatrix} 1 & w \\ \infty & \infty \end{bmatrix}$$

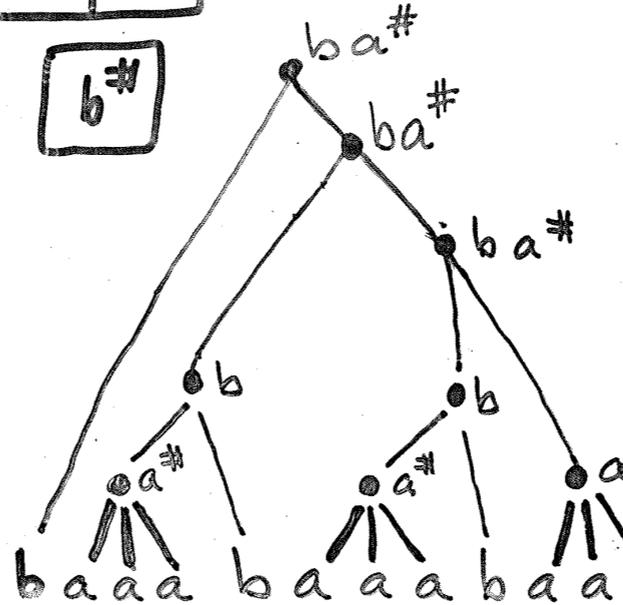
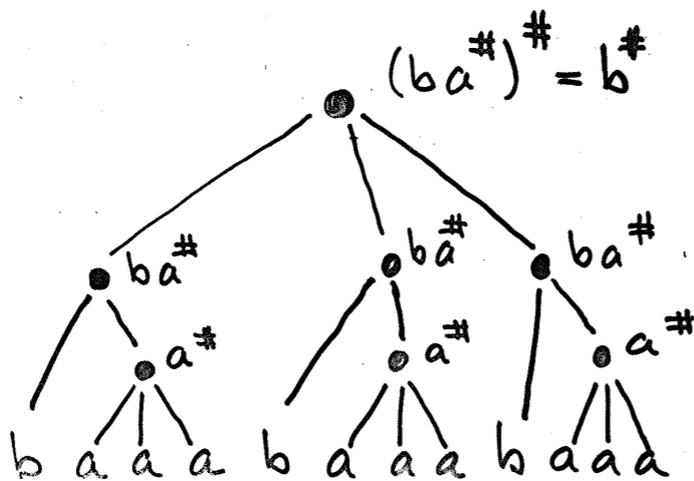
S: a: identity  
b: zero

	b	a <sup>#</sup>	ba <sup>#</sup>
b	b	ba <sup>#</sup>	ba <sup>#</sup>
a <sup>#</sup>	b	a <sup>#</sup>	ba <sup>#</sup>
ba <sup>#</sup>	b	ba <sup>#</sup>	ba <sup>#</sup>

⊙-structure



⊙: discontinuous vertices  
first components of labels not shown





**O desbravador  
da computação brasileira**