

Z. Ésik: A fixed point theorem for non-monotonic functions

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The Knaster-Tarski fixed point theorem

Theorem (Knaster–Tarski) Suppose that L is a complete lattice and $f : L \rightarrow L$ is monotonic. Then f has extremal fixed points. Moreover, the set of all fixed points of f is a complete lattice by itself.

Applications in CS

- automata and languages
- programming logics and verification
- semantics
- algorithms, etc

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Fixed points of non-monotonic functions

Fixed points of non-monotonic functions in CS

- Boolean automata and language equations (Brzozowski-Leiss 1980)
- logic programming (Przymusiński 1989, Rondogiannis-Wadge 2005)
- semantics of parallelism (Chen 2003)

Logic programming

- Przymusiński's well founded semantics
 - abstract fixed point theory by Denecker, Fitting, Gelder, Truszczyński...
- Rondogiannis-Wadge semantics
 - abstract fixed point theory: this work

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Stratified complete lattices

κ a fixed limit ordinal

Definition Suppose that $L = (L, \leq)$ is a complete lattice. We call L a **stratified complete lattice** if it is equipped with preorderings \sqsubseteq_α , $\alpha < \kappa$ such that:

- $\forall x, y \forall \alpha, \beta (x \sqsubseteq_\alpha y \wedge \beta < \alpha \Rightarrow x =_\beta y)$
- $\forall x, y ((\forall \alpha x =_\alpha y) \Rightarrow x = y)$

Example

$L = (V^Z, \leq) : V : F_0 < F_1 < \dots < F_\alpha < \dots < 0 < \dots < T_\alpha < \dots < T_1 < T_0 (\alpha < \kappa)$

$f \sqsubseteq_\alpha g$ iff

1. $\forall z \forall \beta < \alpha ((f(z) = F_\beta \Leftrightarrow g(z) = F_\beta) \wedge (f(z) = T_\beta \Leftrightarrow g(z) = T_\beta))$
2. $\forall z ((f(z) = T_\alpha \Rightarrow g(z) = T_\alpha) \wedge (g(z) = F_\alpha \Rightarrow f(z) = F_\alpha))$

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Axioms

$$\text{A1. } \forall x \forall \alpha \exists y (x =_{\alpha} y \wedge \forall z (x \sqsubseteq_{\alpha} z \Rightarrow y \leq z))$$

Fact: y is uniquely determined by x and α . Notation $x|_{\alpha}$

$$\text{A2. } \forall (x_i)_{i \in I (\neq \emptyset)} \forall y \forall \alpha (x_i =_{\alpha} y, i \in I \Rightarrow \bigvee_{i \in I} x_i =_{\alpha} y)$$

$$\text{A3. } \forall x, y \forall \alpha < \kappa (x \leq y \Rightarrow x|_{\alpha} \leq y|_{\alpha})$$

$$\text{A4. } \forall x, y \forall \alpha < \kappa ((x \leq y \wedge (\forall \beta < \alpha x =_{\beta} y)) \Rightarrow x \sqsubseteq_{\alpha} y)$$

Example When $f \in V^Z$, $\alpha < \kappa$:

$$f|_{\alpha}(z) = \begin{cases} f(z) & \text{if } f(z) \in \{F_{\beta}, T_{\beta} : \beta < \alpha\} \\ F_{\alpha+1} & \text{otherwise} \end{cases}$$

Stratified complete lattices satisfying A1–A4 will be called **models**.

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Lattice theorem

Definition Let L be a model and $x, y \in L$. We define $x \sqsubseteq y$ iff $x = y$ or $\exists \alpha x \sqsubset_{\alpha} y$.

Theorem (with P. Rondogiannis) If L is a model, then (L, \sqsubseteq) is a complete lattice with the same extremal elements.

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Weakly monotonic functions

Definition Suppose that L is a model and $f : L \rightarrow L$. We call f α -**monotonic** for some $\alpha < \kappa$ if

$$\forall x, y (x \sqsubseteq_{\alpha} y \Rightarrow f(x) \sqsubseteq_{\alpha} f(y))$$

Moreover, we call f **weakly monotonic** if it is α -monotonic for all $\alpha < \kappa$.

Example Let $\kappa = \Omega$ and define $\wedge : V^2 \rightarrow V$ as the binary minimum (infimum) operation and $\neg : V \rightarrow V$ by

$$\neg(x) = \begin{cases} T_{\alpha+1} & \text{if } x = F_{\alpha} \\ F_{\alpha+1} & \text{if } x = T_{\alpha} \\ 0 & \text{if } x = 0 \end{cases}$$

P is a propositional logic program over Z i.e., P is a countable set of instructions

$$z \leftarrow p_1; \dots; p_m; \neg q_1; \dots; \neg q_n$$

where $z, p_i, q_j \in Z$. Let $\Phi_P : V^Z \rightarrow V^Z$ be defined by:

$$\Phi(f)(z) = \bigvee_{z \leftarrow p_1; \dots; p_m; \neg q_1; \dots; \neg q_n \in P} f(p_1) \wedge \dots \wedge f(p_m) \wedge \neg f(q_1) \wedge \dots \wedge f(q_n)$$

Then Φ_P is weakly monotonic.

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Fixed point theorem

Theorem (with P. Rondogiannis) Suppose that L is a model and $f : L \rightarrow L$ is weakly monotonic. Then f has a least fixed point and a greatest fixed point w.r.t. the ordering \sqsubseteq .

Theorem (ZE) Suppose that L is a model and $f : L \rightarrow L$ is weakly monotonic. Then the fixed points of f form a complete lattice w.r.t. the ordering \sqsubseteq .

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Fixed point theorem

Example

$$P: \quad p \leftarrow ; \quad s \leftarrow \neg q ; \quad t \leftarrow \neg t$$

$$p = T_0$$

$$s = \neg q$$

$$t = \neg t$$

Minimal model: $(p, T_0), (q, F_0), (s, T_1), (t, 0)$

p is 'more true' than s since there is a rule that says p is true, while s is true only because q is false by default

Another model: $(p, T_0), (q, T_0), (s, F_1), (t, 0)$

Well-founded model: $(p, T), (q, F), (s, T), (t, 0)$

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Inverse limit models

Inverse system of **projections**

$$\begin{aligned}h_{\beta}^{\alpha} : L_{\alpha} &\rightarrow L_{\beta}, & \beta < \alpha < \kappa \\h_{\gamma}^{\beta} \circ h_{\beta}^{\alpha} &= h_{\gamma}^{\alpha}, & \gamma < \beta < \alpha\end{aligned}$$

Inverse limit complete lattice:

$$L_{\infty} = \{(x_{\alpha})_{\alpha < \kappa} \in \prod_{\alpha < \kappa} L_{\alpha} : \forall \beta < \alpha \ h_{\beta}^{\alpha}(x_{\alpha}) = x_{\beta}\}$$

equipped with the pointwise ordering

Definition Let $x = (x_{\gamma})_{\gamma < \kappa}$ and $y = (y_{\gamma})_{\gamma < \kappa}$ in L_{∞} and $\alpha < \kappa$. We define $x \sqsubseteq_{\alpha} y$ iff $x_{\alpha} \leq y_{\alpha}$ and $x_{\beta} = y_{\beta}$ for all $\beta < \alpha$.

Proposition (ZE) L_{∞} is a stratified complete lattice. It is a model iff each h_{β}^{α} is locally completely additive:

$$\forall X \subseteq L_{\alpha} \forall y \in L_{\beta} \ (h_{\beta}^{\alpha}(X) = \{y\}) \Rightarrow h_{\beta}^{\alpha}(\bigvee X) = y)$$

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Representation theorem

Theorem (ZE) Every model is isomorphic to the model determined by the limit of an inverse system over complete lattices with locally completely additive projections.

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Future work

Comparison with well-founded fixed points

Applications

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