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The Design of a Fast Energy-Efficient Stack

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Once Upon a Time....



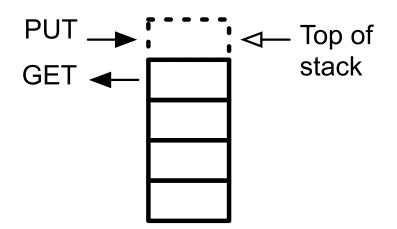
Stacks and Clockless Circuits

- Stack = Last-In First-Out (LIFO) Data Structure
 - Pushdown automaton
 - Like FIFOs, an often used data structure
- Why clockless circuits?
 - Clockless circuit aka asynchronous circuit
 - Speed, Energy Efficiency, ...
 - "There is some meat in there"



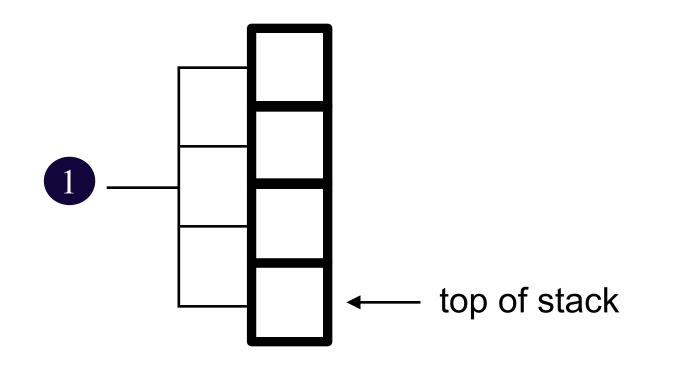
Stacks

- Have many applications
 - function call/return
 - expression evaluation
- Software implementation
 - linked lists, arrays
- Hardware implementation
 - SRAMs, shift registers
- Can we find a better implementation?
 - "better" = less delay
 - "better" = less energy
- Delay = # gate delays per PUT or GET
- Energy = # data moves per PUT or GET



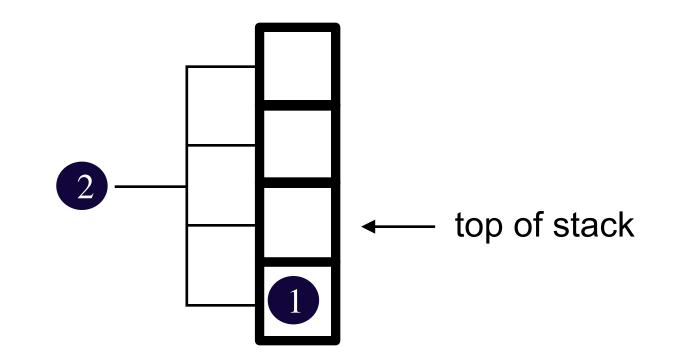


• PUT(1)



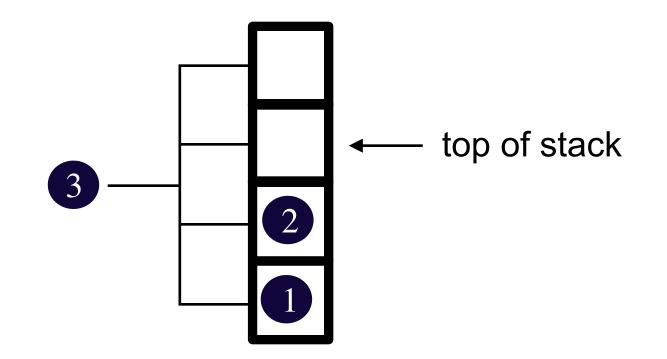






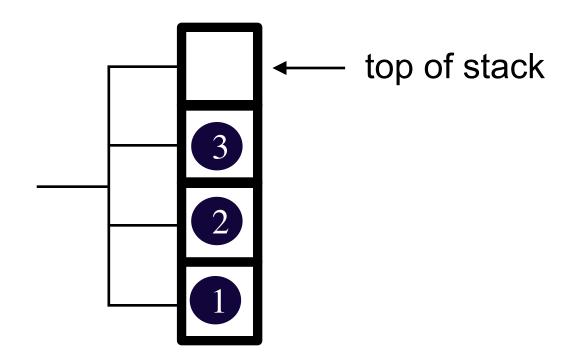


• PUT(3)





• GET





A Pointer Stack

- Has only one move per PUT or GET, but ...
- Has large fan-out and fan-in
- Has large delay
 - SRAM access time > 4 ns, in TSMC 180nm
- Consumes much energy per PUT or GET
 - energy per SRAM access > 15 pJ in TSMC 180nm
- A stack doesn't need random access

A Simple Linear Stack

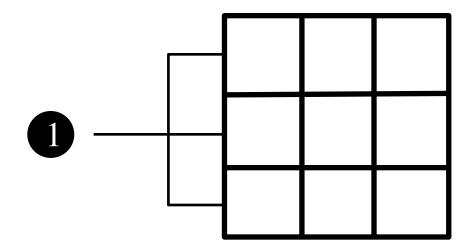
• Shift register, synchronized to clock

• Ripple shift register, clockless

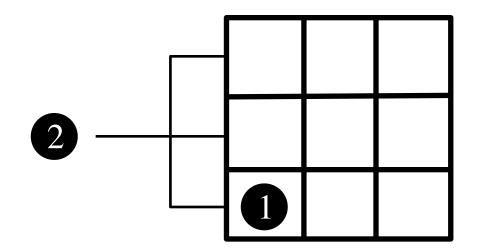
Simple Linear Stack

- Has small fan-in and fan-out
 - Enables short cycle time, but ...
- Has many moves per PUT or GET
 - Consumes much energy per PUT or GET
- Can we combine good properties of pointer and linear stack?
- Main goals:
 - Keep number of moves per PUT or GET small
 - Keep fan-in and fan-out small

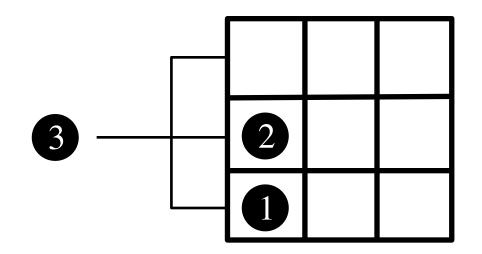
• PUT(1)



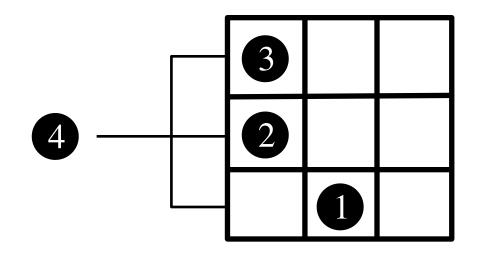




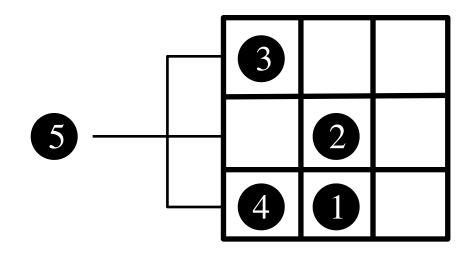
• PUT(3)



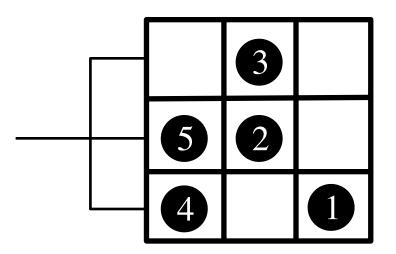
• PUT(4)



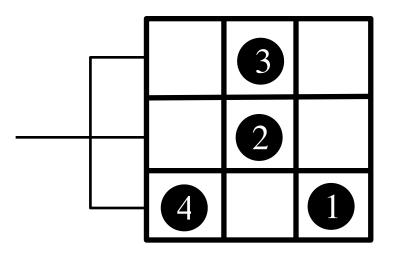
• PUT(5)



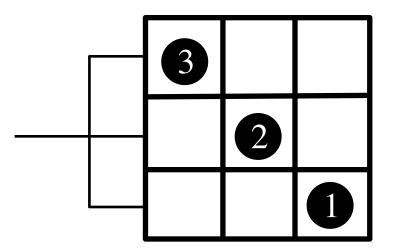
• GET



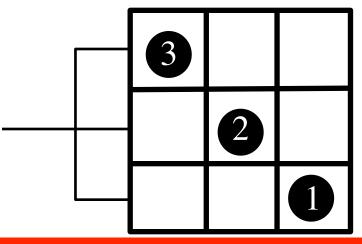
• GET

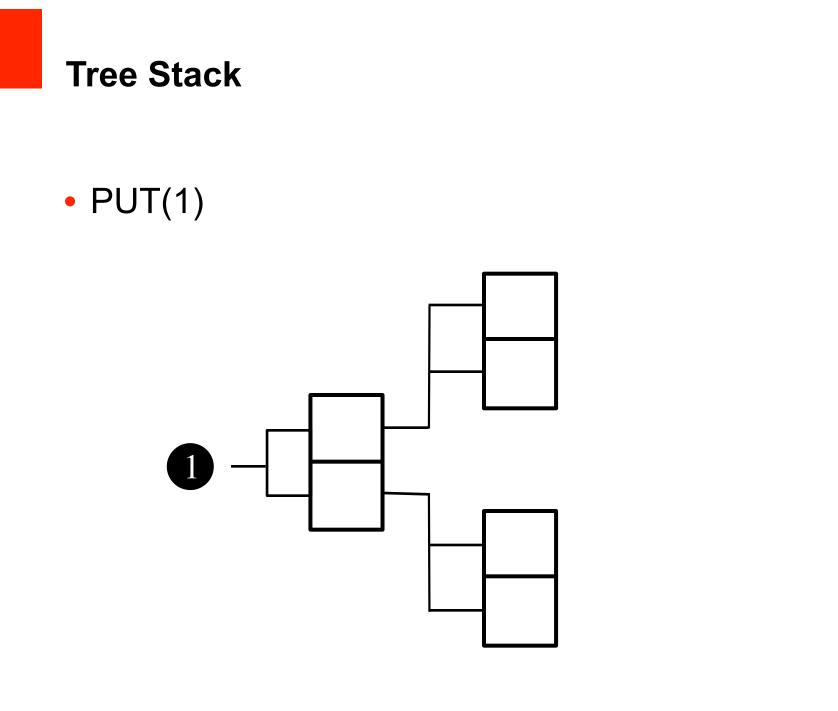


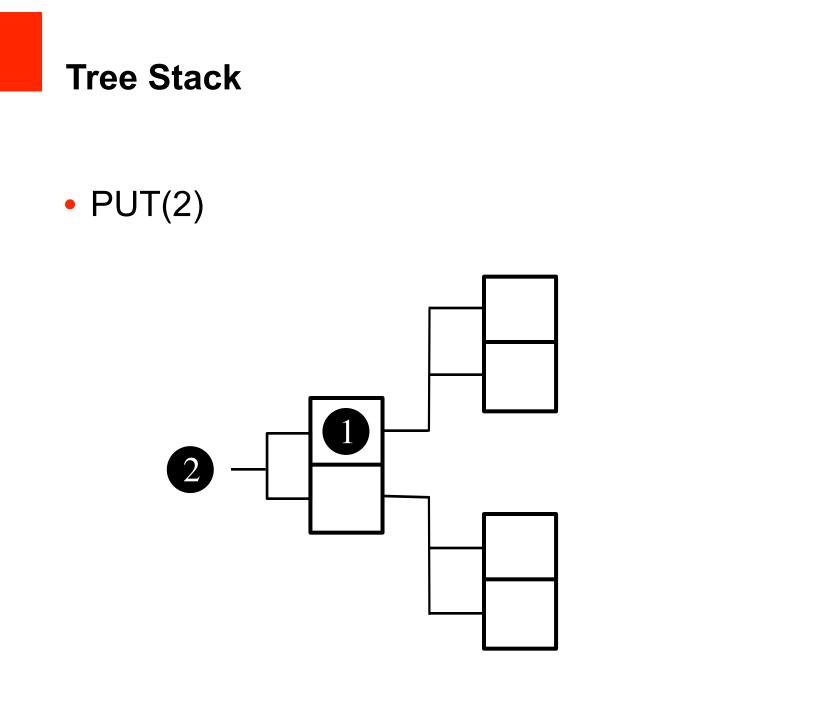
• GET



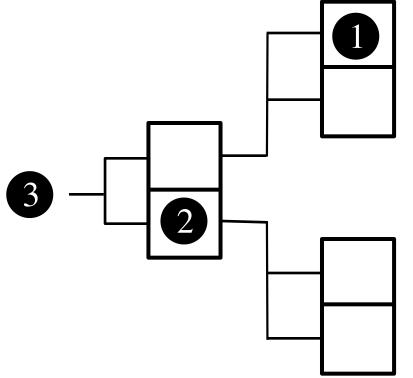
- Has small fan-in and fan-out
- Has small cycle time
- Other attractive properties
 - Moves occur only between neighboring cells
 - Only a fraction of all stack items move per PUT or GET
 - a PUT propagates further only when column-stack has 2 items
 - a GET propagates further only when column-stack has 1 item







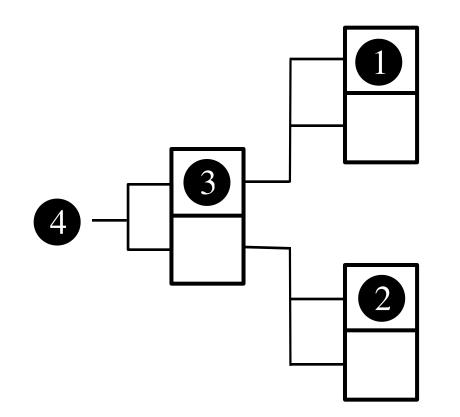




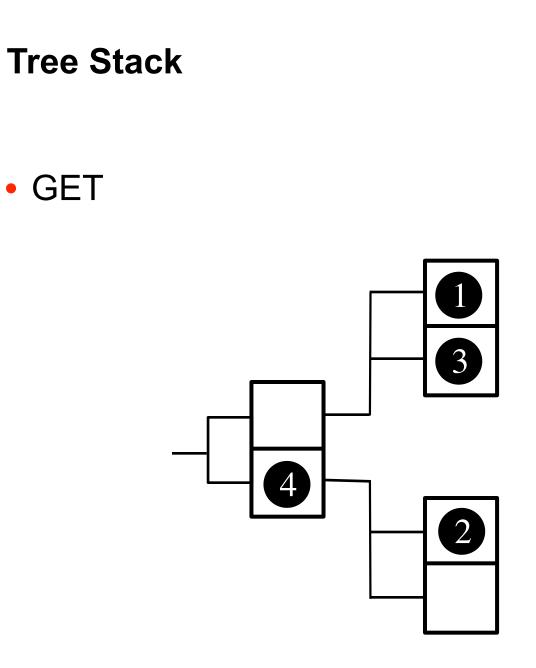




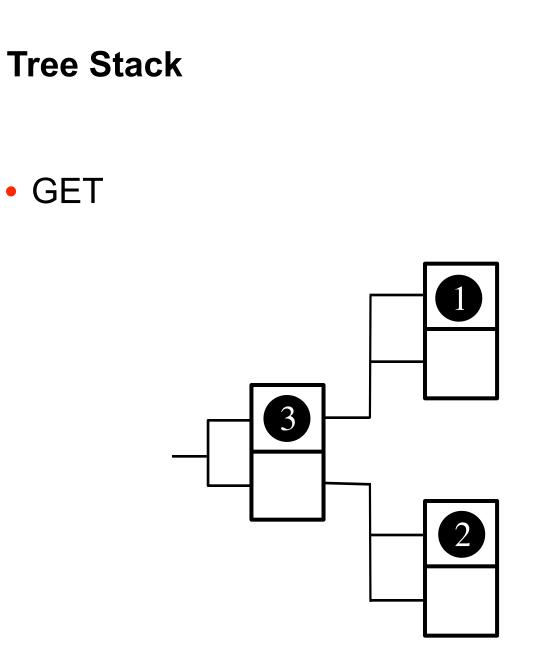
• PUT(4)

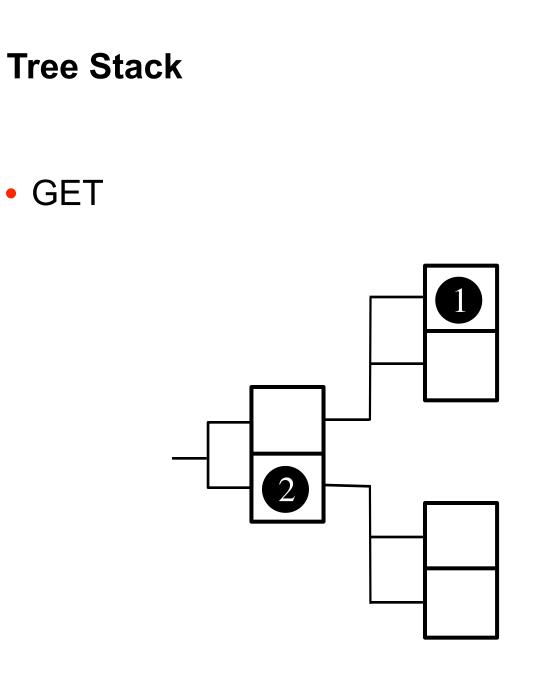






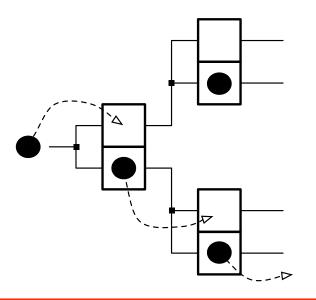






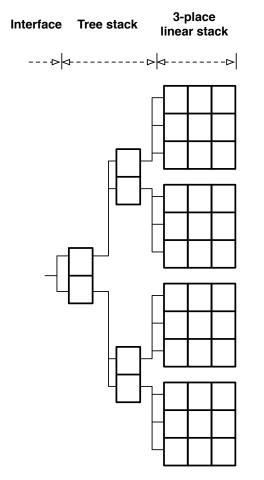
Tree Stack

- Has small fan-in and fan-out
- Has small cycle time
- Other attractive properties
 - Moves occur only between neighboring cells
 - Only a fraction of all stack items move per PUT or GET



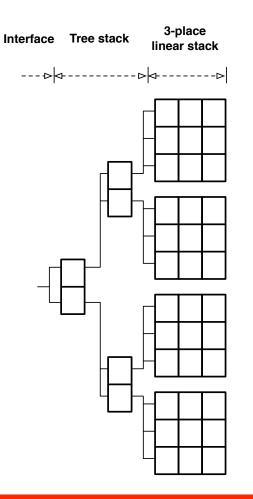
Hybrid Stack

Combination of pointer, linear, and tree stack



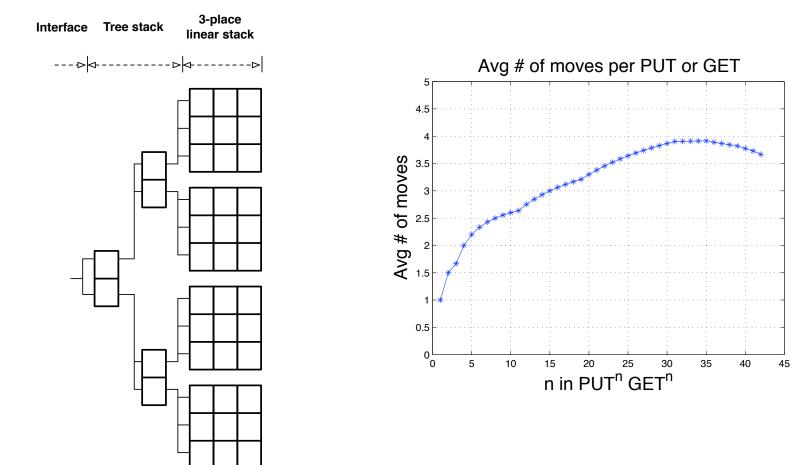
Properties of Hybrid Stack: Response Time

Constant response time



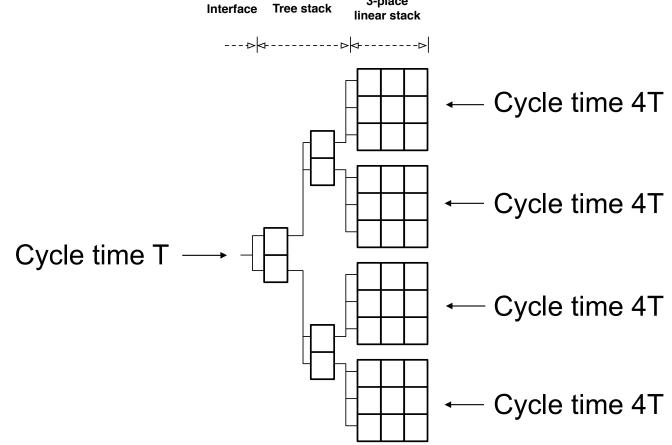
Properties of Hybrid Stack: number of moves

Average number of data moves per PUT or GET

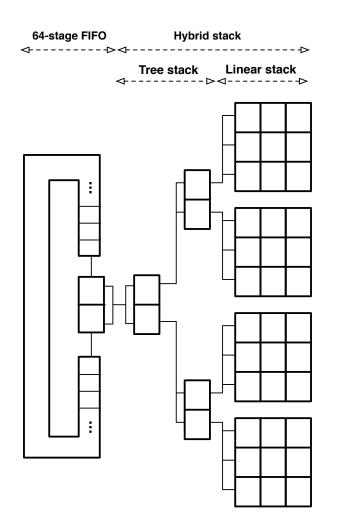


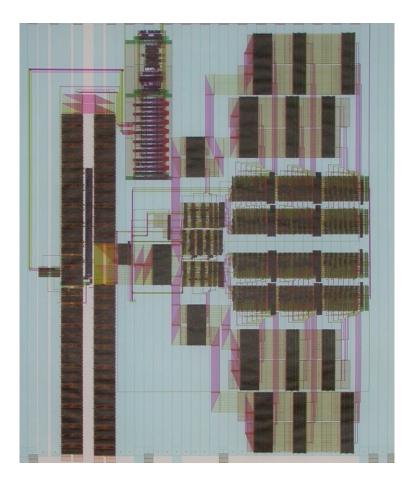
Properties of Hybrid Stack: Cycle Times

 Cycle times at leaf nodes can be 4X larger than cycle time at root node



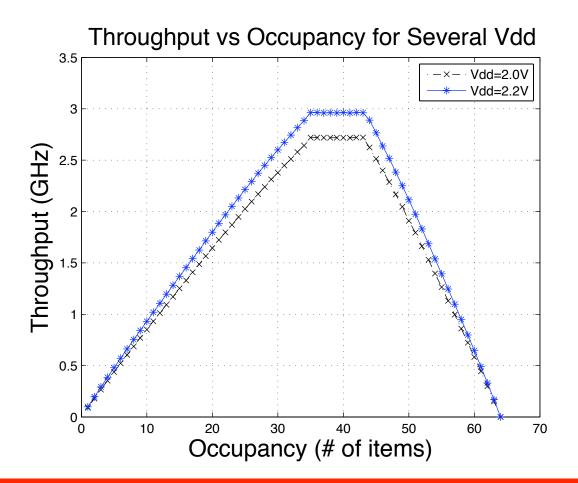
Stack Implementation





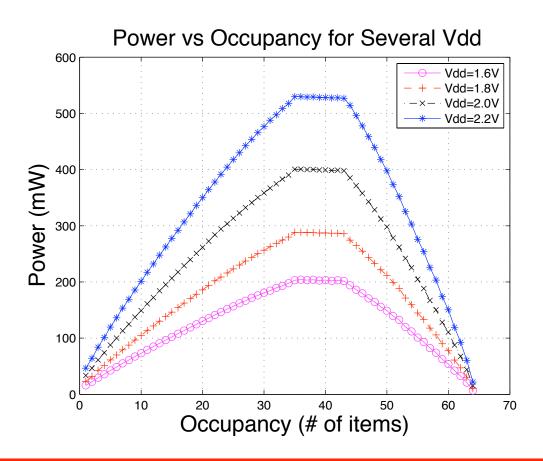
Throughput vs Occupancy

• 2.75GHz at 2.0V



Power Consumption vs Occupancy

Power consumption is proportional to number of moves



Concluding Remarks

- A stack implementation
 - Cycle time of 6 gate delays, (368ps at 2V or 2.72GHz in 180nm)
 - Cycle time is independent of # items or data path width
 - Average number of moves per PUT or GET grows (very) slowly
- Dynamic energy per move is small (6-8pJ)
 - and can be made smaller by using more efficient latches
- Exhibits much concurrency
- A problem for you
 - How does avg number of moves depend on input sequences of PUTs and GETs?
 - What is lower bound for avg number of moves for a stack implementation given bounded fan-out?

More Reading

- "An Evaluation of Asynchronous Stacks," IEEE Design and Test, Sept/Oct 2011, pp 52--61.
- A Fast and Energy-Efficient Stack, "Proc of 10th IEEE Int'I Symp. Advanced Research in Asynchronous Circuits and Systems (ASYNC)," IEEE CS Press, 2004, pp. 7--16.

