

# IM 2005, Nice, France

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## Introduction

#### System Design

**System Prototype** 

**Admission Control** 

**Load Estimation** 

Introduction

System Design

Admission Control

Load Estimation

System Prototype

Evaluation

Wrap Up

**Evaluation** 

#### Wrap Up





# Introduction

### **QoS in the Internet**

- long debated, many research proposals, limited deployment
- concerns: end-to-end scope, technical complexity, scalability
- technical tools
  - admission control and traffic regulation
  - differentiated packet scheduling

#### Introduction

System Design

**Our Goals** 

Load Estimation

Admission Control

- System Prototype
- Evaluation
- Wrap Up

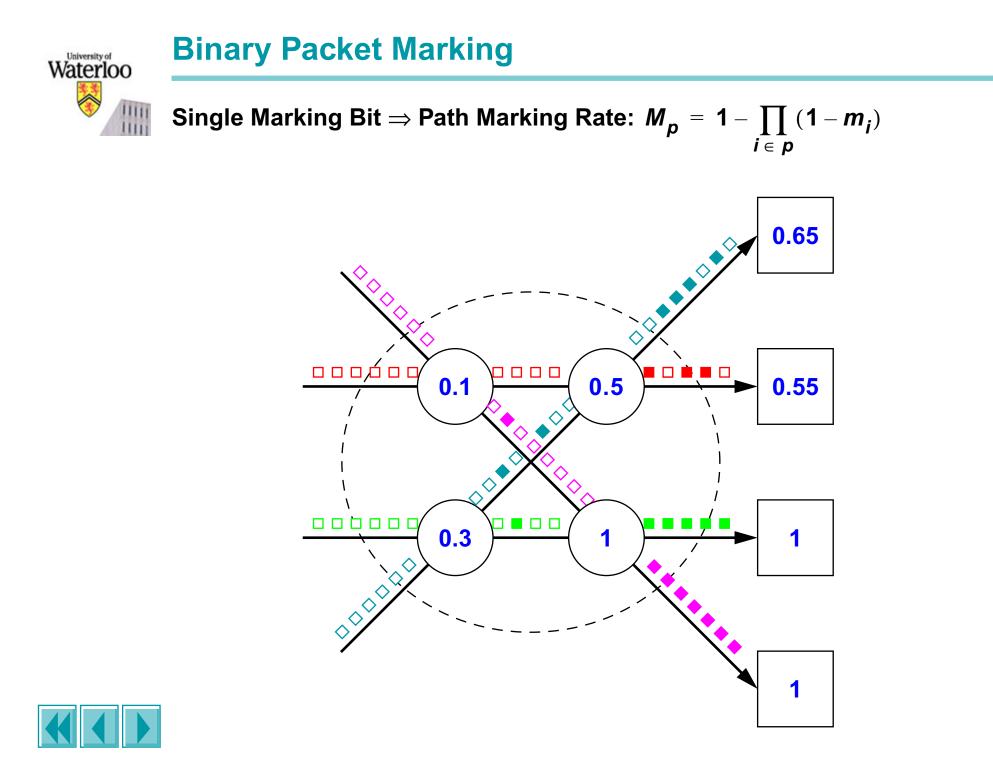
- admission control and load-based routing for network domain
- ultra-scalability: low complexity and little functionality in core
- modularization of control functions

#### **Basic Mechanism: Binary Packet Marking**

- Kelly et al. results for strictly concave utility curves
- TCP, TFRC and ECN
- other admission control proposals, e.g. RMD

#### **Key Aspect**

support admission AND load-based routing



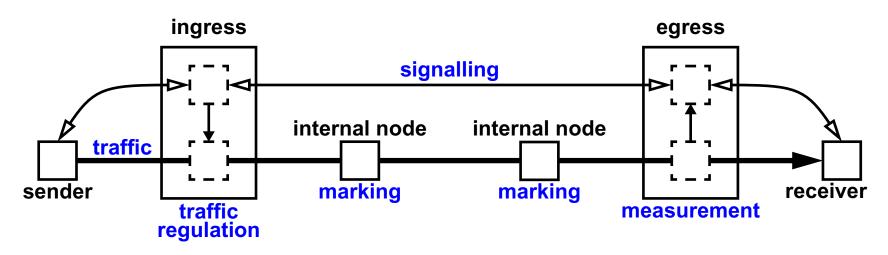
# **System Design** Waterloo **Network Domain Overview** load estimation neighbour domain neighbour domain edge gateway core node request signalling transmission path $\rightarrow$ load reporting





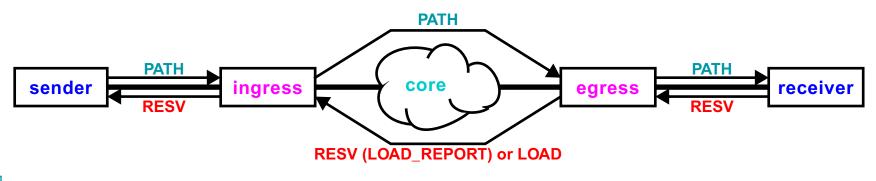
# **System Design - Details**

Combination of Proactive & Reactive Resource Allocationproactive resource allocation based on network load feedback



inherent feedback delay between egress and ingress

#### Signalling Design: RSVP Extension



purely local extension (LOAD\_REPORT object & LOAD message)



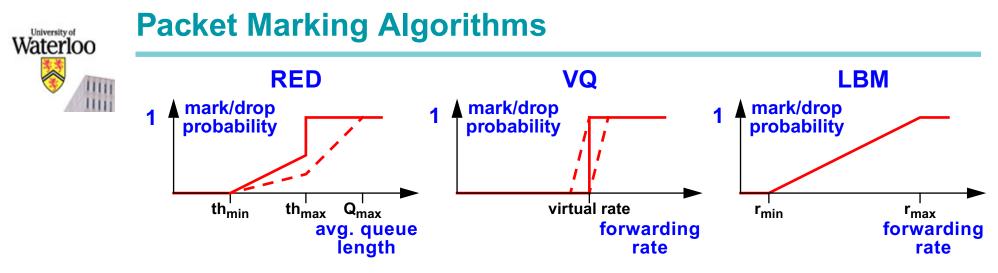
# **Admission Control**

- **Admission Control**
- decide about connection acceptance
- inherent per-domain concept
  - reliability of end systems
- feedback between edge gateways
- suitable for inelastic flows
- vs. Flow Control
- throttle sending rate according to load signal
- feedback from receiver
- suitable for elastic flows
- $\Rightarrow$  Support for Both Needed

#### **Feedback Delay**

- traffic load  $\rightarrow$  measurement  $\rightarrow$  marking/transmission  $\rightarrow$  measurement
- admission control at ingress vs. egress  $\rightarrow$  no difference
  - measuring always at egress
  - traffic control always at ingress
- safety margin in resource utilization
- vs. arrival of service requests?





#### **Random Early Detection (RED) & Variants**

- queue-based feedback
- ineligible packets  $\rightarrow$  random drop (ok)
- meaning of path marking rate for inelastic flows (?)

#### Virtual Queue (VQ) & Variants

- hybrid feedback, time-scale dependent
- ineligible packets  $\rightarrow$  bursty dropping (?)
- inelastic flows  $\rightarrow$  binary path marking rate

#### Load Based Marking (LBM)

- rate-based feedback
- ineligible packets  $\rightarrow$  continuous random dropping (?)
- path marking rate is product of local load values
- use relative load of link or node (!)





# **Load Estimation**

Goals

- provide load information to constraint-based routing
- decouple centralized routing from resource allocation system

#### **Capacity-oblivious Load Estimation**

- service class capacity adaptation by independent allocation system
- wireless or overlay links with varying capacity
- complex notion of load, e.g. combination of processing and link load
- heterogeneous notion of load at different nodes

#### **Hybrid Load Estimation**

- marking-based load estimation: multiplicative error propagation
- usage  $\rightarrow$  capacity estimation: only additions, less sensitive to errors

#### **Packet Marking**

- see Binary Packet Marking  $\rightarrow$  system of equations  $\rightarrow$  individual load
  - need continuous marking function
  - usually over-specified  $\rightarrow$  use last N load reports
- real-world engineering challenge: three 'signals' in two bits





# **System Prototype**

#### Lab Prototype

- FreeBSD, Linux, Solaris
- signalling  $\rightarrow$  KOM RSVP engine (user-level daemon)
- packet handling (internal & edge) → FreeBSD/ALTQ (kernel modules)
- traffic generation & measurement

#### Simulation

- ns-2
- RSVP-based signalling fully shared code basis with prototype
- packet handling partially shared code basis with prototype
- traffic generation & measurement mostly separate code
- load estimation simulation only

#### **Packet Marking**

- using ECN bits
- RED, VQ, AVQ, LBM
- threshold-based marking (TBM)  $\rightarrow$  simplification of VQ
  - forwarding rate > threshold  $\rightarrow$  mark or drop



# **Evaluation**



#### **Admission Control**

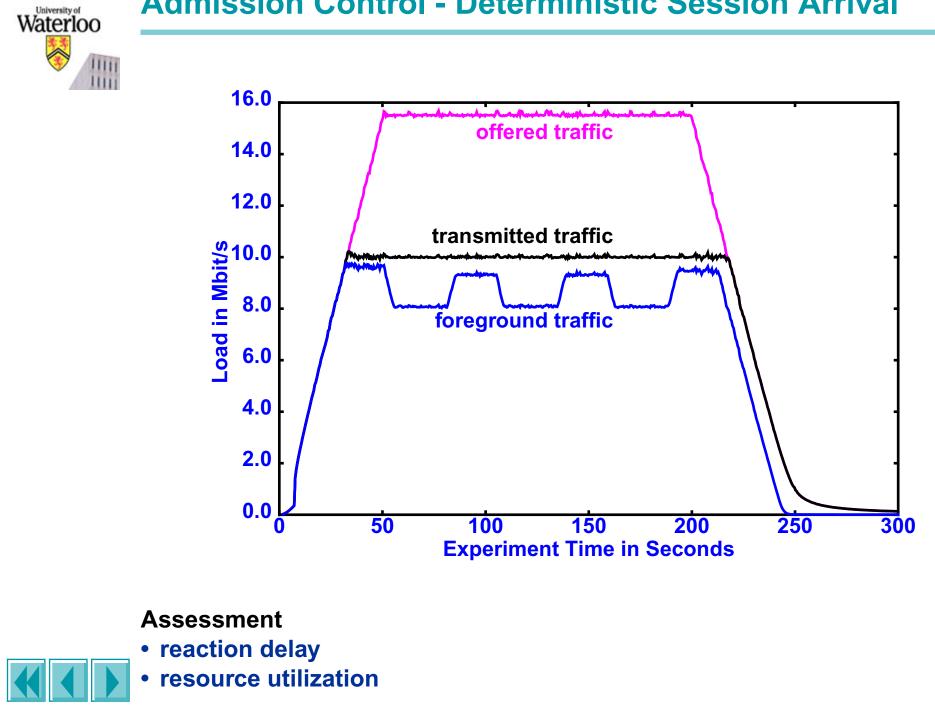
- lab experiment with real prototype
  - reconciliation/calibration for larger simulation experiments
- load system and test whether admission control works
- assessment: reaction delay
- assessment: traffic discrimination

#### **Load Estimation**

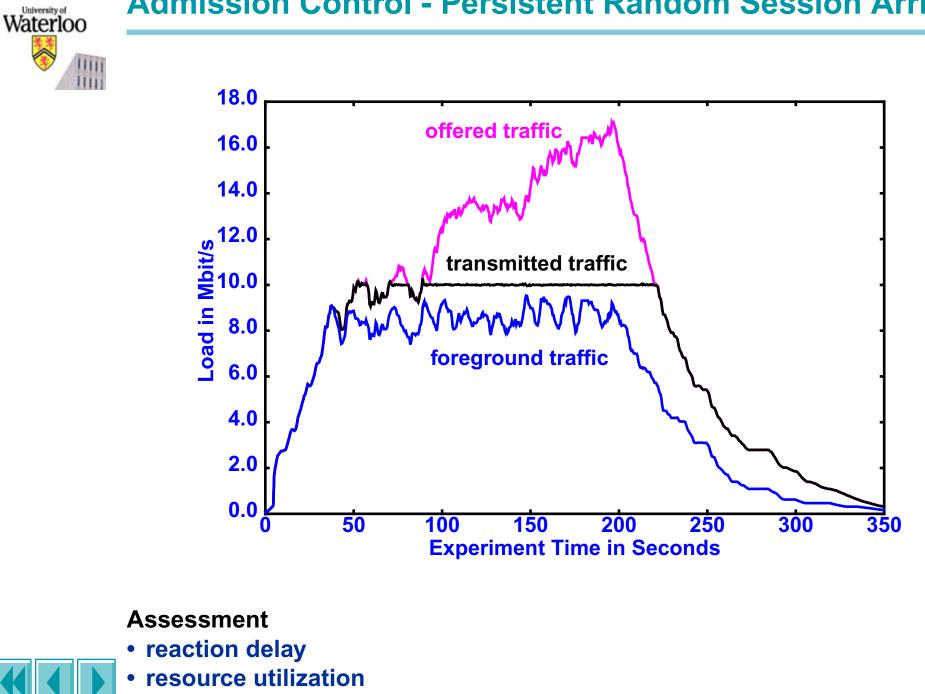
- simulation experiments
- load system and compare measured local load with estimated load
- assessment: reporting delay
- assessment: estimation precision



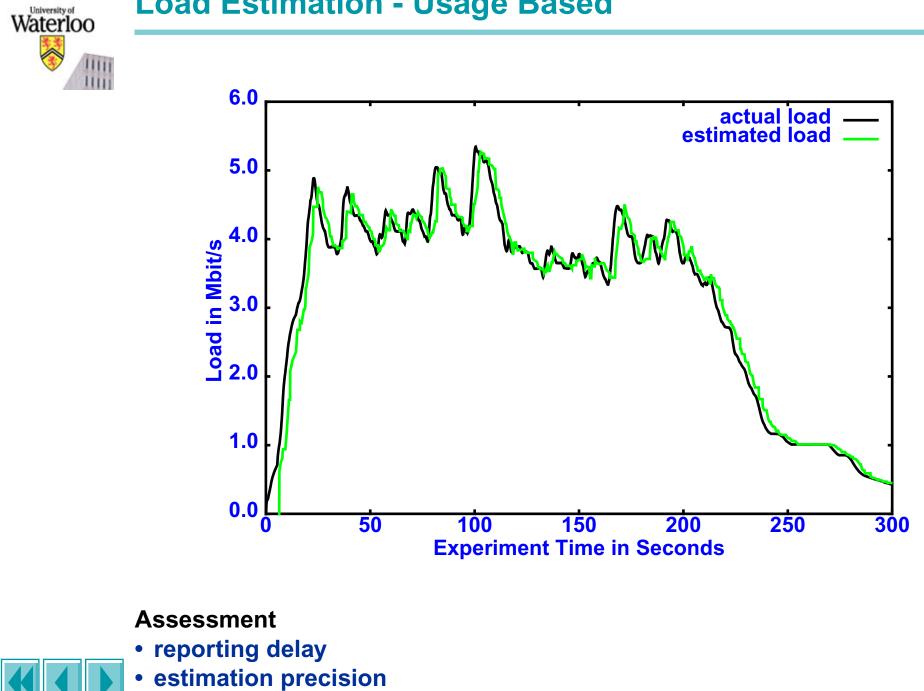
# **Admission Control - Deterministic Session Arrival**



# **Admission Control - Persistent Random Session Arrival**

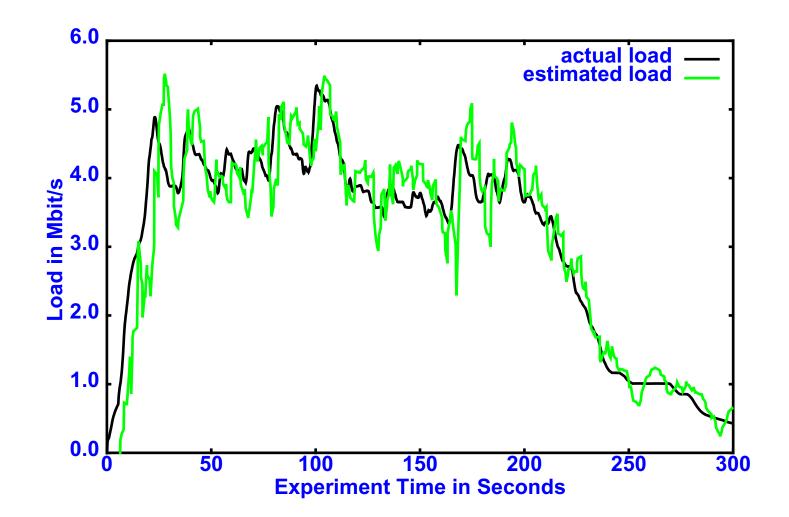


# **Load Estimation - Usage Based**



# **Load Estimation - Load Based**

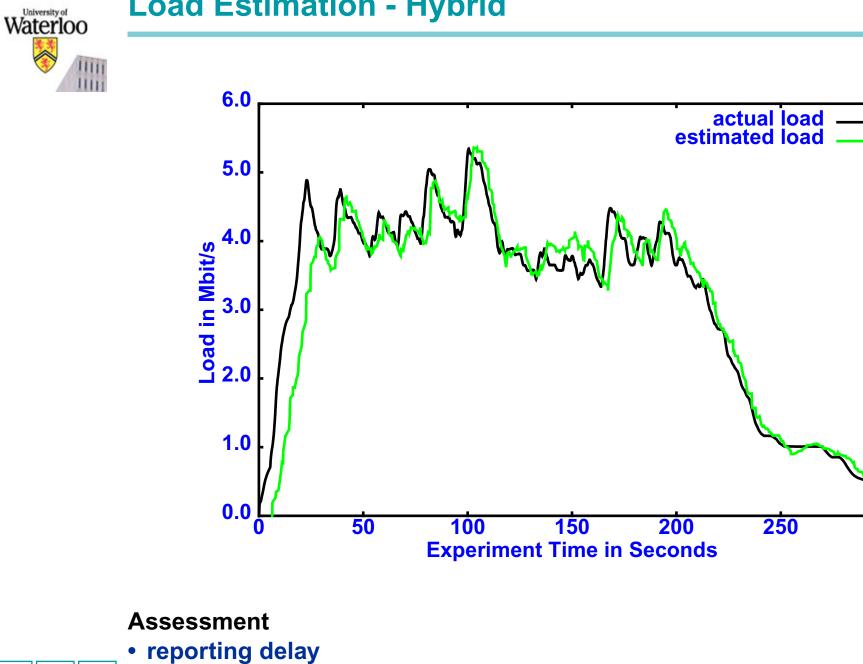




# Assessmentreporting delayestimation precision



# **Load Estimation - Hybrid**





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# Wrap Up



Integrated Load Control System

- admission control and load estimation
- packet marking and evaluation at network edge
- system design, architecture, and prototype (simulation and lab)
- experimental validation

#### **Admission Control**

- feedback signalling & admission control  $\rightarrow$  RSVP extensions
  - fundamentals about admission control & feedback delay
- various packet marking algorithms

#### Load Estimation

- early, speculative work ightarrow basic proof of concept
- enabling technology

#### Remarks

- system design is orthogonal to DiffServ
- deployment path: RED/ECN & DiffServ for dedicated service class
  - per-node/per-path deployment possible

