Steps in Query Processing

- 1. Translation
 - check SQL syntax
 - check existence of relations and attributes
 - replace views by their definitions
 - generate internal query representation
- 2. Optimization
 - consider alternative **plans** for processing the query
 - select an efficient plan
- 3. Processing
 - execute the plan
- 4. Data Delivery

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Example

```
select DeptNo, Deptname, count(*), sum(Salary)
from Employee, Department
where WorkDept = DeptNo and DeptNo like 'D%'
group by DeptNo, Deptname
having sum(Salary) > 1000000
```

An Execution Plan

- 1. Scan the Employee table, select all tuples for which WorkDept starts with 'D', call the result R_1 .
- 2. Join R_1 and Department, eliminate attributes other than DeptNo, Deptname, and Salary. Call the result R_2 . This may involve:
 - sorting R_1 on WorkDept
 - sorting Department on Deptno
 - joining the two sorted relations to produce R_2
- 3. Group the tuples of R_2 . Call the result R_3 . This may involve:
 - sorting R_2 on DeptNo and Deptname
 - group tuples with identical values of DeptNo and Deptname
 - count tuples in each group, and add their Salaries
- 4. Scan R_3 , select all tuples with **sum**(Salary) > 1000000

Pictorial Access Plan



Pipelined Plans and Iterators

- In a pipelined plan, each tuples stream from one operator to another.
- Pipelining allows for parallel execution of operators, and avoids unnecessary *materialization* of intermediate results. (Sometimes materialization may be necessary...)
- Iterators are a common model for plan operators:
 - every operator is an iterator
 - an iterator provides the following interface: Open, GetNext, and Close
 - each iterator implements its interface, using calls to the interface functions of its child (or children)

DB2 Access Plan



DB2 Access Plan with Index



Some Basic Query Processing Operations

- Data Access and Filtering
 - Index scans
 - Table scans
- Projection
- Joining
 - nested loop join
 - hash join
 - sort-merge join
 - and others ...
- Sorting
- Grouping and Duplicate Elimination
 - by sorting
 - by hashing

Joining Relations

```
select DeptName, LastName
from Department, Employee
```

```
where DeptNo = WorkDept
```

```
Conceptually, a nested-loop join works like this:
```

```
foreach tuple d in Department do
  foreach tuple e in Employee do
    if d.DeptNo = e.WorkDept then
        output d,e
    end
end
```

Block Nested Loop Join

select DeptName, LastName
from Department, Employee

where DeptNo = WorkDept

Process outer relation a chunk at a time

```
foreach chunk C of Department
  foreach tuple e in Employee do
    foreach tuple d in C
        if d.DeptNo = e.WorkDept then
            output d,e
        end
    end
end
```

Other Techniques for Join

• If there is an index on the WorkDept attribute of the Employee relation, an **index join** can be used:

```
foreach tuple d in Department do
    use the index to find Employee tuples where d.DeptNo = Wor
    for each such tuple e
    output d,e
end
```

- Examples of other join techniques:
 - Sort-Merge Join: sort the tuples of Employee on WorkDept and the tuples of Department of DeptNo, then merge the sorted relations.
 - Hash Join: assign each tuple of Employee and of Department to a "bucket" by applying a hash function to its WorkDept (DeptNo) value. Within each bucket, look for Employee/Department tuple pairs for which WorkDept = DeptNo.

Hash Join Example















External Merge Sort: Run Formation



External Merge Sort: Run Formation (cont'd)



External Merge Sort: Run Formation (cont'd)



External Merge Sort: Merging Runs



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Summary

- A plan describes how a query is executed, including:
 - the sequence of basic operations (select, project, join, sort, etc.)
 used to process the query
 - how each operation will be implemented, e.g., which join method will be used, which indices will be used to perform a selection.