Computing the Smith form with multipliers of a nonsingular integer matrix

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Smith normal form

Given

- a nonsingular integer matrix $A \in \mathbb{Z}^{n \times n}$,

Determine

- the Smith normal form $S = \mathrm{diag}(s_1, s_2, \ldots, s_n) \in \mathbb{Z}^{n \times n}$.
- $s_1 \,|s_2\, | \ldots |\, s_n$. (invariant factors)

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- S obtained using unimodular row and column operations.

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Uses?

- Solving integer linear equations
- Classifying finite abelian groups:

  SNF of generators $\rightarrow \mathbb{Z}_{s_1} \oplus \cdots \oplus \mathbb{Z}_{s_n}$
- Useful for solving polynomial systems invariant under finite abelian groups via Gröbner bases [Faugère, Svartz, 2013]
- Rational invariants and rewrite rules for systems invariant under finite abelian group [Hubert, L., 2016]
- Outer Adjoint formula [Storjohann, Birmpilis, L., Storjohann]
- ...

Example

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Suppose we want to solve

$$x \overbrace{\begin{bmatrix} -13 & 10 & -20 & 27 \\ 27 & 30 & 15 & 30 \\ 0 & 15 & 15 & 6 \\ -21 & 0 & -15 & 9 \end{bmatrix}}^{A} = \overbrace{\begin{bmatrix} 277 & 50 & 290 & -132 \end{bmatrix}}^{b}.$$

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$$(xU) \begin{matrix} \\ \\ \\ \\ \end{matrix} \overset{S}{\begin{bmatrix} 1 & & & \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{bmatrix}} = \overset{(bV)}{\begin{bmatrix} -9106 & 1701 & 6885 & 18795 \end{bmatrix}}.$$

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End up with

$$\bar{x} \overset{S}{\begin{bmatrix} 1 & & & \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{bmatrix}} = \overset{\bar{b}}{\begin{bmatrix} -9106 & 1701 & 6885 & 18795 \end{bmatrix}}.$$

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$$A$$

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Exponential growth of intermediate integers

Modern History

Citation	Time complexity	U, V	Type
Kannan and Bachem (1979)	$poly(n, \log \|\|A\|\|)$	✓	Det
Iliopoulos (1989)	$n^5 (\log \|\|A\|\|)^2$	✓	Det
Hafner and McCurley (1991)	$n^5 (\log \|\|A\|\|)^2$		Det
Storjohann (1996, 2000)	$n^{\omega+1} \log \|\|A\|\|$	✓	Det
Eberly, Giesbrecht and Villard (2000)	$n^{2+\omega/2} \log \|\|A\|\|$		MC
Kaltofen and Villard (2004)	$n^{2.695591} \log \|\|A\|\|$		MC
Birmpilis, Labahn, Storjohann (2021)	$n^{\omega} \log \|\|A\|\|$	✓	LV

- ω exponent of matrix multiplication
- Complexity is given without the extra $\log n$ and $\log \log \|\|A\|\|$ factors.
- Det = deterministic, MC = Monte Carlo or LV = Las Vegas.

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Randomization:

- Our algorithm is of type Las Vegas.
- Return the correct output with probability at least $1/2$ or fail.

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Example

$$\overbrace{\begin{bmatrix} \frac{3}{7} & -\frac{5}{21} & \frac{4}{21} & -\frac{13}{21} \\ \frac{101}{35} & -\frac{12}{7} & \frac{11}{7} & -\frac{419}{105} \\ -\frac{69}{35} & \frac{122}{105} & -\frac{106}{105} & \frac{19}{7} \\ -\frac{16}{7} & \frac{29}{21} & -\frac{26}{21} & \frac{67}{21} \end{bmatrix}}^{A^{-1}} \overbrace{\begin{bmatrix} 5 \\ 5 \\ 3 \\ 2 \end{bmatrix}}^{b} = \begin{bmatrix} \frac{2}{7} \\ \frac{272}{105} \\ -\frac{173}{105} \\ -\frac{13}{7} \end{bmatrix}$$

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For two vectors, the probability of success is at least $1/2$.

- "On computing the determinant and Smith form of an integer matrix." [Eberly, Giesbrecht, Villard (2000)]

Fast linear system solving

Any rational vector $v \in \mathbb{Q}^{n\times 1}$ with denominator $s \in \mathbb{Z}_{>0}$ has an integral $q \in \mathbb{Z}^{n\times 1}$ and a fractional part $r \in (\mathbb{Z}/s)^{n\times 1}$ such that

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$$A^{-1}b = \begin{matrix} v \\ \left[\begin{array}{c} \frac{87798811184766974 07}{11711} \\ \frac{36103271414459480 05}{23422} \\ \frac{54168639766491175 43}{11711} \\ \frac{13839883865944116065}{23422} \end{array}\right] \end{matrix}$$

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- Deterministic method for system solving based on high-order lifting [Birmpilis, Labahn, Storjohann (ISSAC, 2019)]
- Deterministic variant of integrality certification [Birmpilis, Labahn Storjohann (ISSAC, 2020)]

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$$\begin{bmatrix} s_n/s_n & & & \\ & s_n/s_{n-1} & & \\ & & \ddots & \\ & & & s_n/s_{n-r+1} \end{bmatrix} \bmod s_n.$$

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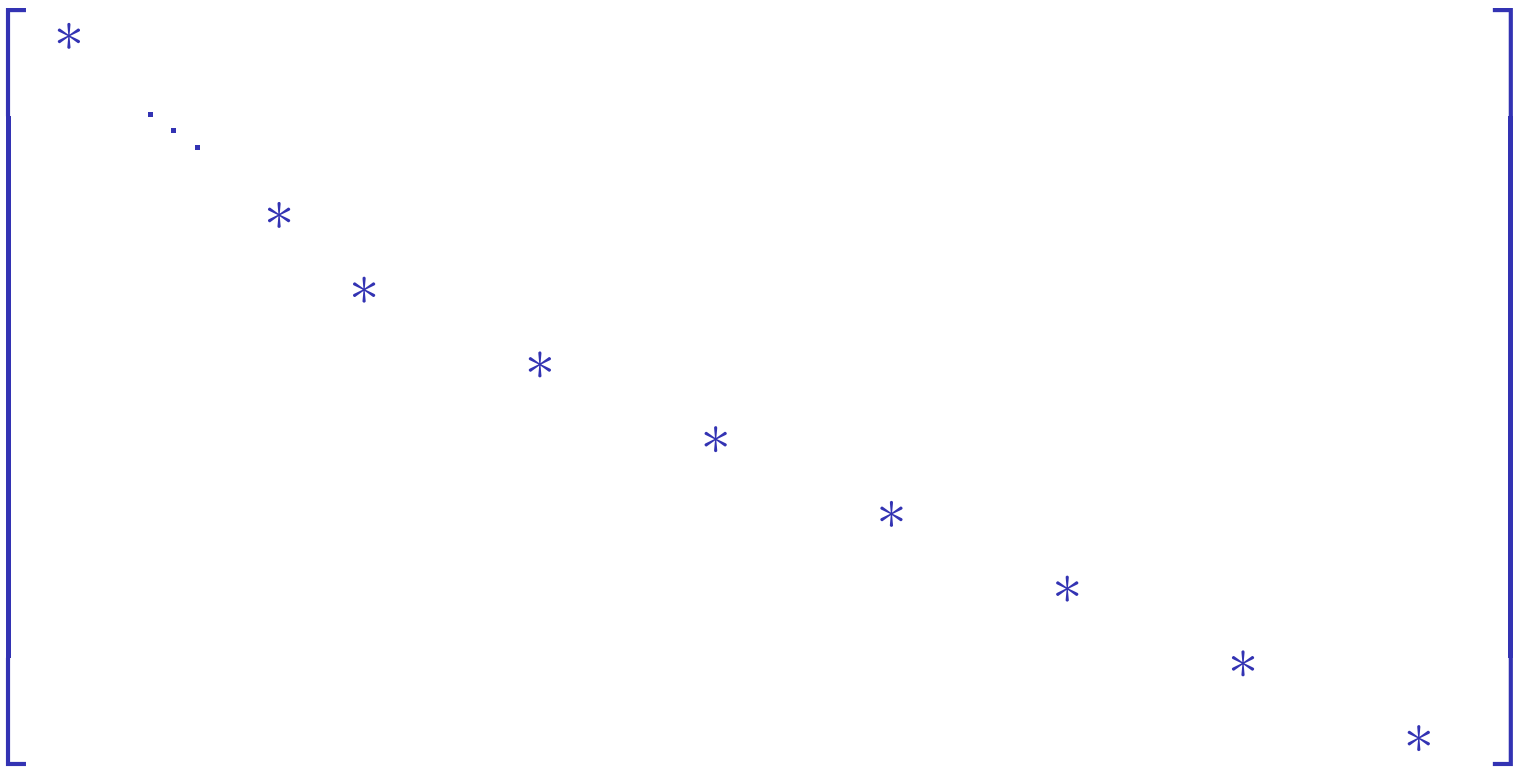
- The length of s_n can be n times the length of entries in A.

Instead we use dimension $\times$ precision $\leq$ invariant

Computing the invariant factors of A:



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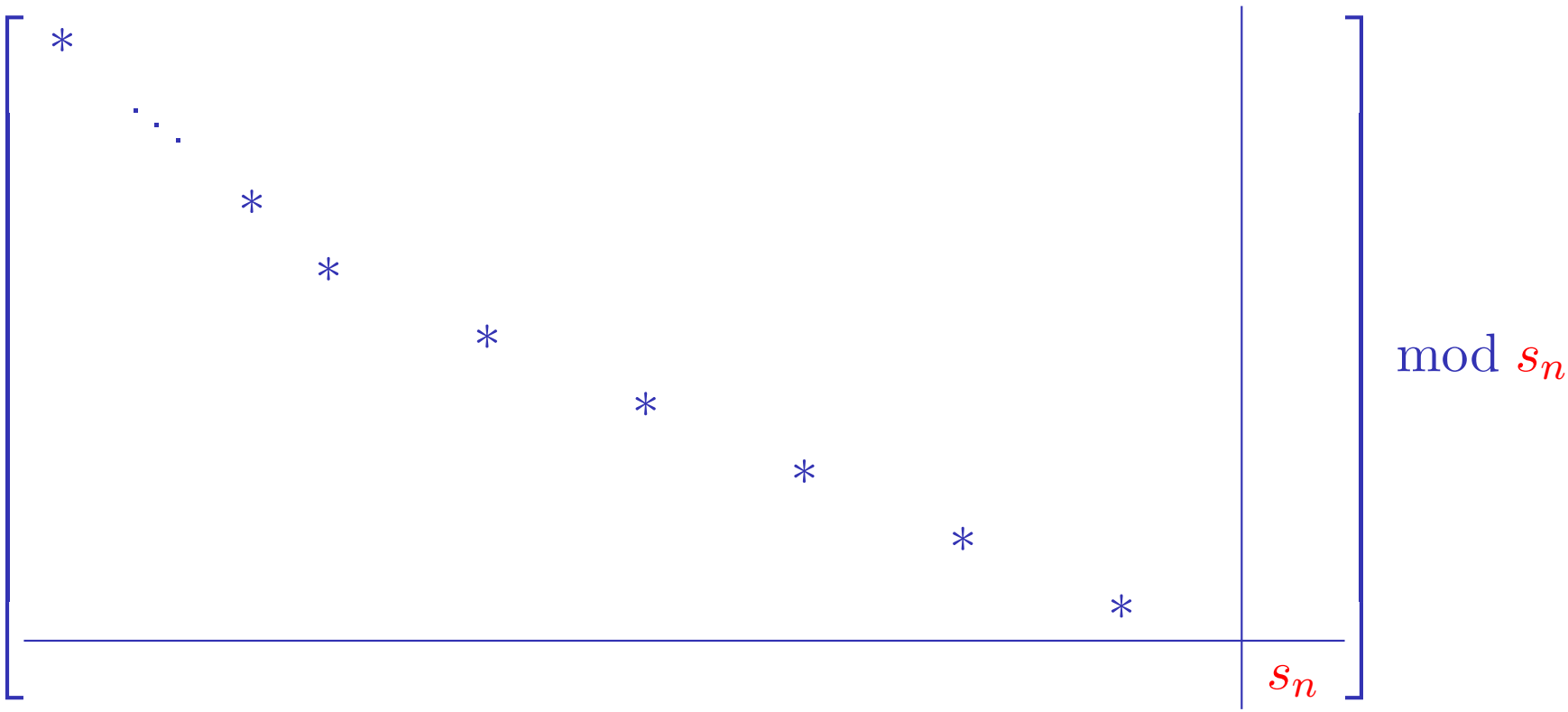

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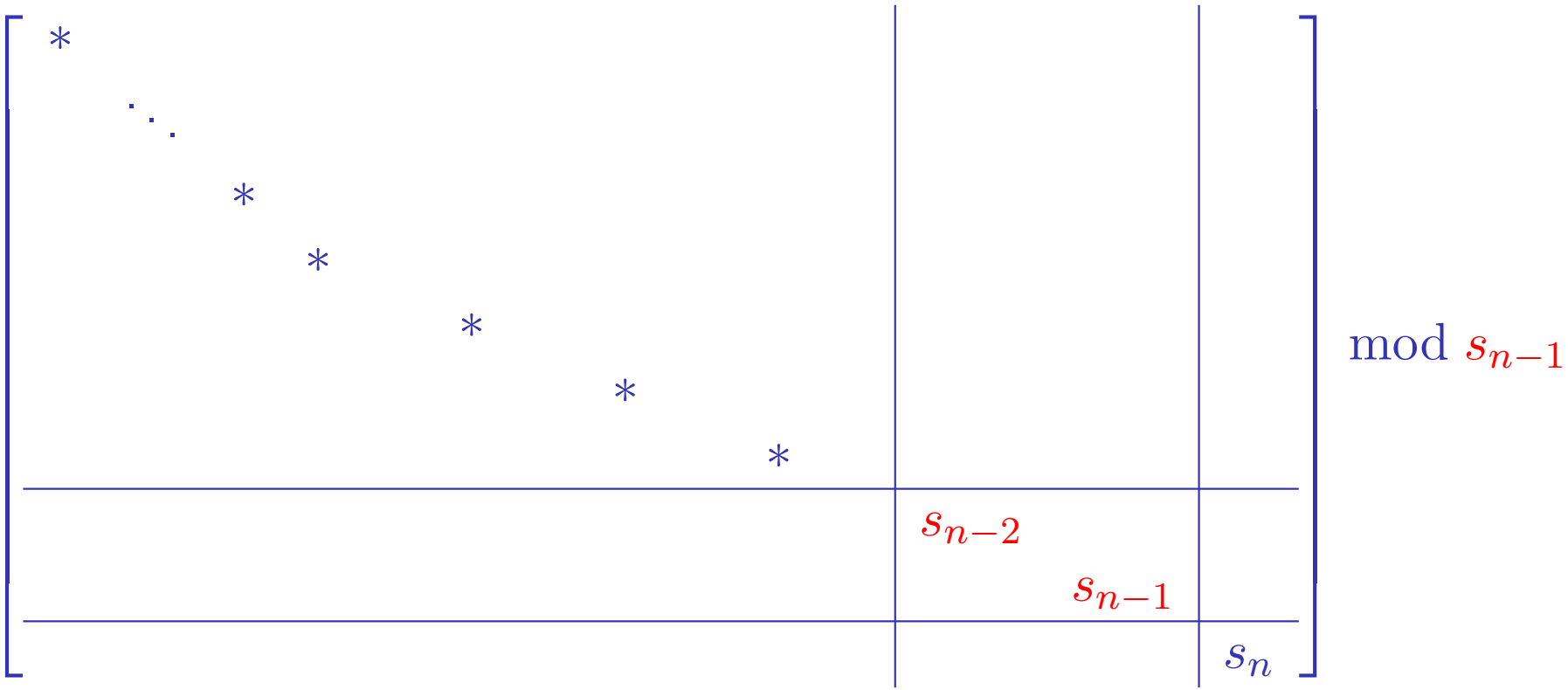

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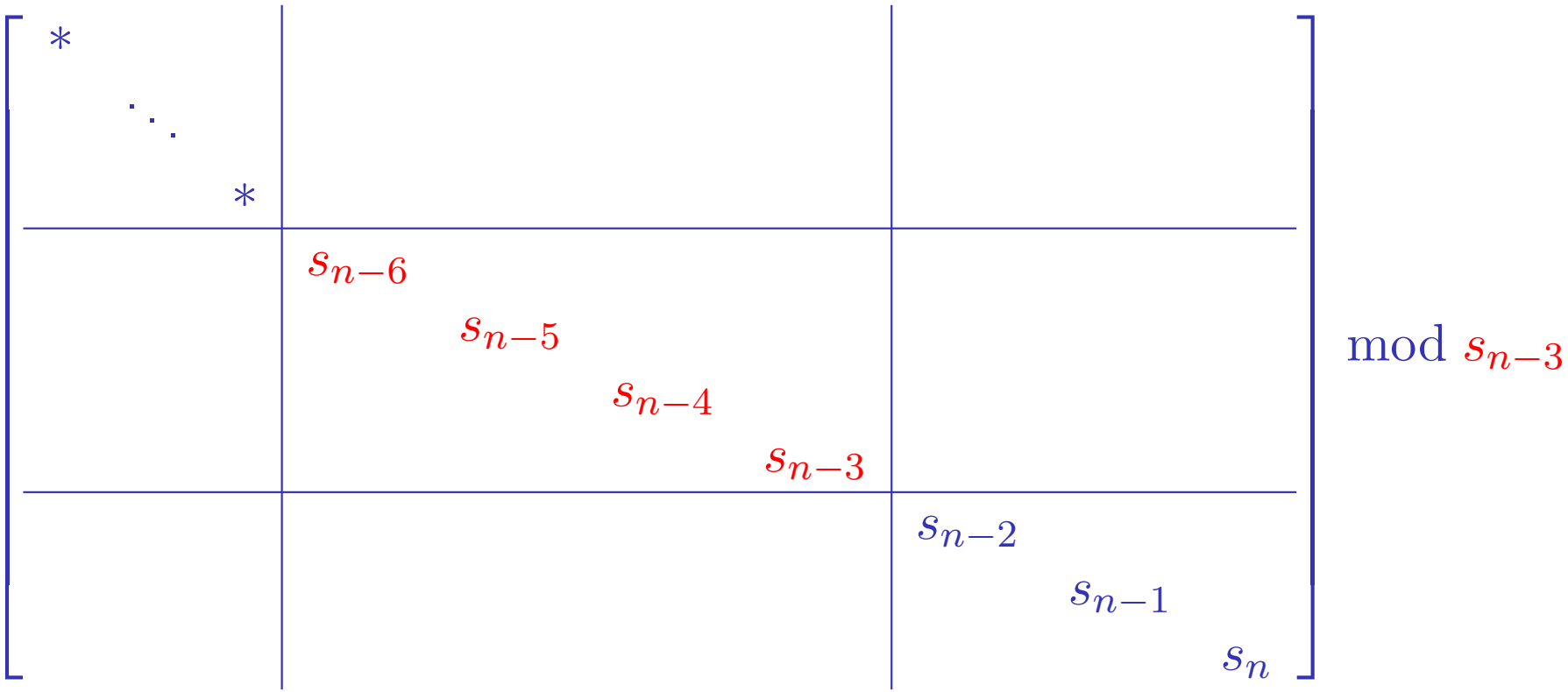

Example

$$B_0 := \operatorname{diag}(A, I_n) =$$

$$\begin{bmatrix} -13 & 10 & -20 & 27 & & & & \\ 27 & 30 & 15 & 30 & & & & \\ 0 & 15 & 15 & 6 & & & & \\ -21 & 0 & -15 & 9 & & & & \\ & & & & 1 & & & \\ & & & & & 1 & & \\ & & & & & & 1 & \\ & & & & & & & 1 \end{bmatrix}$$

$S = \operatorname{diag}(*, *, *, *)$ and $\det B_0 = \det A$

Example

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$S = \operatorname{diag}(*, *, *, 105)$ and $\det B_0 = \det A$

Example

$$B_1 = \begin{bmatrix} -13 & 10 & -20 & 27 & & & & 0 \\ 27 & 30 & 15 & 30 & & & & 43 \cdot 105 \\ 0 & 15 & 15 & 6 & & & & 9 \cdot 105 \\ -21 & 0 & -15 & 9 & & & & -15 \cdot 105 \\ & & & & 1 & & & 0 \\ & & & & & 1 & & 0 \\ & & & & & & 1 & 0 \\ 0 & 41 & 17 & 67 & & & & 40 \cdot 105 \end{bmatrix}$$

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$S = \operatorname{diag}(*, 3, 15, 105)$ and $\det B_1 = \det A/105$

Example

$$B_2 = \left[\begin{array}{cccc|ccc} -13 & 10 & -20 & 27 & -2\cdot 3 & 12\cdot 15 & 0 \\ 27 & 30 & 15 & 30 & 24\cdot 3 & 46\cdot 15 & 43 \\ 0 & 15 & 15 & 6 & 7\cdot 3 & 16\cdot 15 & 9 \\ -21 & 0 & -15 & 9 & -9\cdot 3 & -5\cdot 15 & -15 \\ \hline & & & & 1 \quad 0 & 0 & 0 \\ 0 & 1 & 2 & 0 & 1\cdot 3 & 1\cdot 15 & 0 \\ 0 & 8 & 0 & 0 & 0 & 7\cdot 15 & 0 \\ 0 & 41 & 17 & 67 & 28\cdot 3 & 59\cdot 15 & 40 \end{array}\right]$$

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Example

$$B_2 :=$$

$$\left[\begin{array}{cccc|ccc} -13 & 10 & -20 & 27 & -2 & 12 & 0 \\ 27 & 30 & 15 & 30 & 24 & 46 & 43 \\ 0 & 15 & 15 & 6 & 7 & 16 & 9 \\ -21 & 0 & -15 & 9 & -9 & -5 & -15 \\ \hline & & & & 1 \; 0 & 0 & 0 \\ 0 & 1 & 2 & 0 & 1 & 1 & 0 \\ 0 & 8 & 0 & 0 & 0 & 7 & 0 \\ 0 & 41 & 17 & 67 & 28 & 59 & 40 \end{array}\right]$$

$S = \operatorname{diag}(*, 3, 15, 105)$ and $\det B_2 = \det A/(105 \cdot 15 \cdot 3)$

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$S = \operatorname{diag}(1, 3, 15, 105)$ and $\det B_3 = \det A/(105 \cdot 15 \cdot 3 \cdot 1)$

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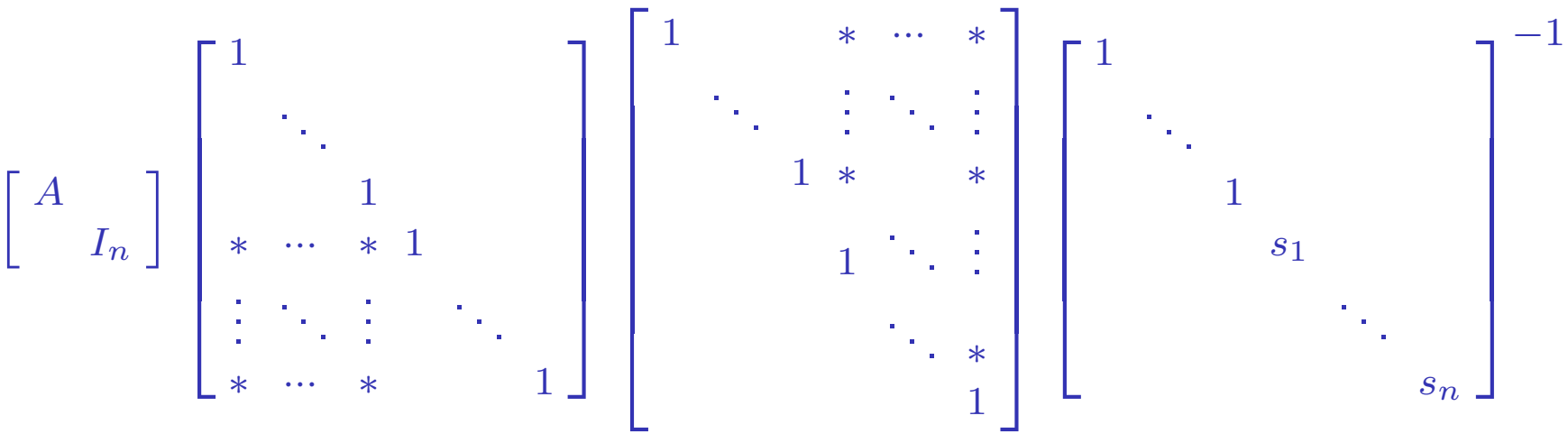
$S = \operatorname{diag}(1, 3, 15, 105)$ and $\det B_3 = -1$

Basically constructing a Smith massager



Basically constructing a Smith massager

- Augment A with I_n.

$$\begin{bmatrix} A & \\ & I_n \end{bmatrix} \begin{bmatrix} 1 & & & & & \\ & \ddots & & & & \\ & & 1 & & & \\ * & \cdots & * & 1 & & \\ \vdots & \ddots & \vdots & & \ddots & \\ * & \cdots & * & & & 1 \end{bmatrix} \begin{bmatrix} 1 & & & * & \cdots & * \\ & \ddots & & \vdots & \ddots & \vdots \\ & & 1 & * & & * \\ & & & 1 & \ddots & \vdots \\ & & & & \ddots & * \\ & & & & & 1 \end{bmatrix} \begin{bmatrix} 1 & & & & & \\ & \ddots & & & & \\ & & 1 & & & \\ & & & s_1 & & \\ & & & & \ddots & \\ & & & & & s_n \end{bmatrix}^{-1}$$

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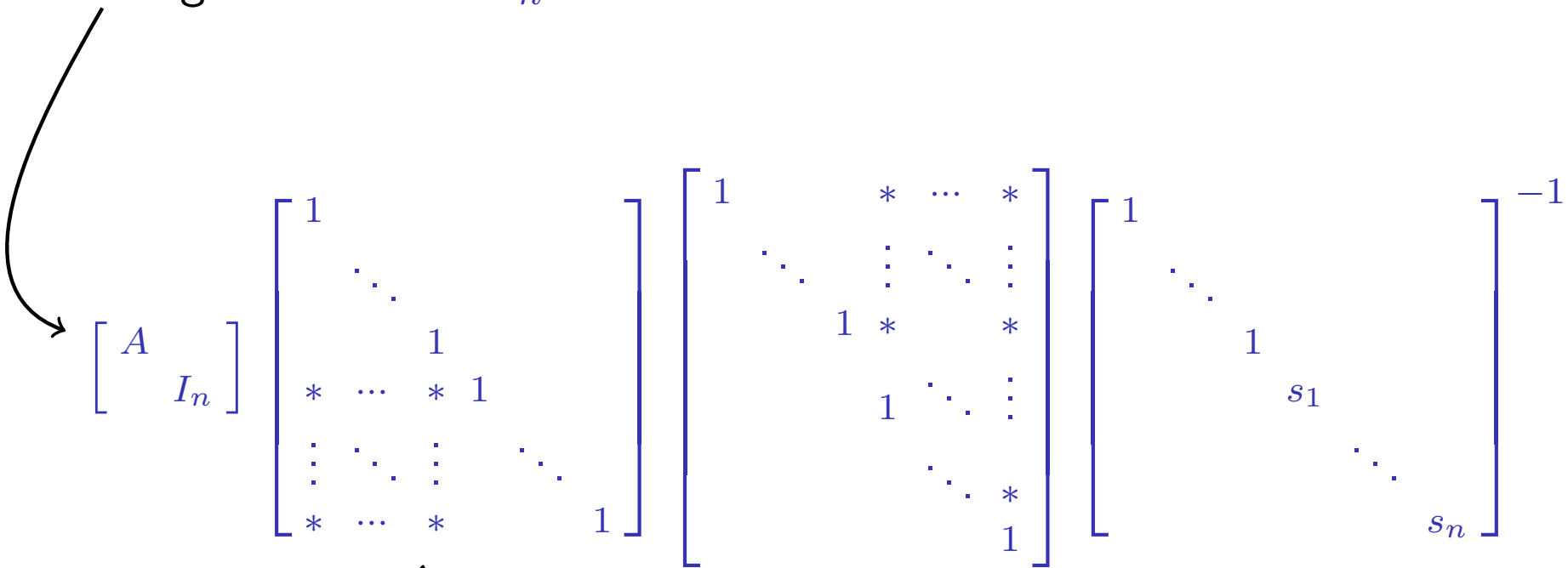

Conditioner matrix $\begin{bmatrix} I_n & \\ C & I_n \end{bmatrix}$.

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Massager $\begin{bmatrix} I_n & M \\ & T \end{bmatrix}$.

- AMS^{-1} is integral.
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Conditioner matrix $\begin{bmatrix} I_n & \\ C & I_n \end{bmatrix}$.

- If the product is unimodular, then S is the Smith form of A.

Part 2: What about the multiplier matrices?

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e.g. $\overset{A}{\begin{bmatrix} -13 & 10 & -20 & 27 \\ 27 & 30 & 15 & 30 \\ 0 & 15 & 15 & 6 \\ -21 & 0 & -15 & 9 \end{bmatrix}}$ with Smith form $\overset{S}{\begin{bmatrix} 1 & & & \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{bmatrix}}$.

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How do we get the unimodular matrices? That is, U, V such that

$$A \overset{V}{\begin{bmatrix} -26 & 5 & 15 & 55 \\ -62 & 0 & 22 & 137 \\ -20 & 2 & 17 & 41 \\ -53 & 2 & 25 & 115 \end{bmatrix}} = \overset{U}{\begin{bmatrix} -1313 & -17 & 24 & 28 \\ -4452 & 75 & 138 & 92 \\ -1548 & 14 & 49 & 32 \\ 369 & -39 & -23 & -7 \end{bmatrix}} S.$$

Seems to be a harder problem

Consider matrix $A = \begin{bmatrix} a & -b \\ b & a \end{bmatrix}$ with gcd$(a, b) = 1$.

Suppose $\begin{array}{rcl} au + bv & = & 1 \\ a(-b) + b(a) & = & 0 \end{array}$.

Then $\begin{bmatrix} a & -b \\ b & a \end{bmatrix} \begin{bmatrix} u & b \\ -v & a \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ bu - av & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & a^2 + b^2 \end{bmatrix}$

Main tool : Smith Massager (again)

Recall

- From before Smith massager satisfied:

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b. $(M'M - I) \equiv 0 \,\mathbf{c}\mathrm{mod}\, S$

Better view of Smith massager

Better view of Smith massager

Definition

Let $A \in \mathbb{Z}^{n\times n}$ be nonsingular with Smith normal form $S \in \mathbb{Z}^{n\times n}$. A matrix $M \in \mathbb{Z}^{n\times n}$ is a *Smith massager* for A if

a. it satisfies that

$$AM \equiv 0 \pmod{S},$$

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- A Smith massager can be always reduced column modulo S.
- The length of the average entry in M is in $(\log \|A\|)^{1+o(1)}$.

Example

a. $$\underset{A}{\begin{bmatrix} -13 & 10 & -20 & 27 \\ 27 & 30 & 15 & 30 \\ 0 & 15 & 15 & 6 \\ -21 & 0 & -15 & 9 \end{bmatrix}} \underset{M}{\begin{bmatrix} 0 & 2 & 0 & 55 \\ 0 & 0 & 7 & 32 \\ 0 & 2 & 2 & 41 \\ 0 & 2 & 10 & 10 \end{bmatrix}} = 0 \,\mathbf{c}\mathrm{mod} \underset{S}{\begin{bmatrix} 1 & & & \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{bmatrix}}$$

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b. $$\overset{M'}{\begin{bmatrix} 13 & 5 & 35 & 84 \\ 43 & 80 & 45 & 100 \\ 57 & 2 & 11 & 97 \\ 5 & 86 & 69 & 76 \end{bmatrix}} \overset{M}{\begin{bmatrix} 0 & 2 & 0 & 55 \\ 0 & 0 & 7 & 32 \\ 0 & 2 & 2 & 41 \\ 0 & 2 & 10 & 10 \end{bmatrix}} = \overset{I}{\begin{bmatrix} 1 & & & \\ & 1 & & \\ & & 1 & \\ & & & 1 \end{bmatrix}} \mathbf{c}\mathrm{mod} \overset{S}{\begin{bmatrix} 1 & & & \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{bmatrix}}$$

M looks like a relaxed version of V

For Smith multipliers, we are looking for $V \in \mathbb{Z}^{n \times n}$ that satisfies

1a. AVS^{-1} is integral, and

1b. V is unimodular, namely, there exists a matrix $V' \in \mathbb{Z}^{n \times n}$ such that $V'V = I_n$.

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Question

How do we go from M to V?

We arrive at V by perturbation

Let $\overset{M}{\begin{bmatrix} 0 & 2 & 0 & 55 \\ 0 & 0 & 7 & 32 \\ 0 & 2 & 2 & 41 \\ 0 & 2 & 10 & 10 \end{bmatrix}}$ be a relaxed version of $\overset{V}{\begin{bmatrix} -26 & 5 & 15 & 55 \\ -62 & 0 & 22 & 137 \\ -20 & 2 & 17 & 41 \\ -53 & 2 & 25 & 115 \end{bmatrix}}$.

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Because of the addition of $\mathbf{c} \bmod S$,

- we might expect V to satisfy $V = M + QS$ for some Q.

This implies that

- we just need a suitable perturbation for M scaled by S.

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Idea

A random perturbation should work just fine!

- *"On computing the determinant and Smith form of an integer matrix." [Eberly, Giesbrecht, Villard (2000)]*

Example: Going from M to V

$$M = \begin{bmatrix} 0 & 2 & 0 & 55 \\ 0 & 0 & 7 & 32 \\ 0 & 2 & 2 & 41 \\ 0 & 2 & 10 & 10 \end{bmatrix}$$

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- Perturb M by a random $R \in \mathbb{Z}^{n \times n}$ times Smith form S.

Example: Going from M to V

$$\underset{M}{\begin{bmatrix} 0 & 2 & 0 & 55 \\ 0 & 0 & 7 & 32 \\ 0 & 2 & 2 & 41 \\ 0 & 2 & 10 & 10 \end{bmatrix}} + \underset{R}{\begin{bmatrix} 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}} \cdot \underset{S}{\begin{bmatrix} 1 & & & \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{bmatrix}}$$

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$$M + RS$$
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- $M + RS$ won't be unimodular. But with high probability,
  a. it will be nonsingular, and
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Example: Going from M to V

$$\overset{M+RS}{\begin{bmatrix} 1 & 5 & 15 & 55 \\ 0 & 0 & 22 & 137 \\ 1 & 2 & 17 & 41 \\ 0 & 2 & 25 & 115 \end{bmatrix}} \cdot \overset{H}{\begin{bmatrix} 3849 & & & \\ 256 & 1 & & \\ 485 & & 1 & \\ 1664 & & & 1 \end{bmatrix}}^{-1}$$

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We can obtain U, V from S, M in a Las Vegas fashion in time

$$(n^\omega \log \|A\|)^{1+o(1)}.$$

Conclusion

1. A Las Vegas algorithm which computes the Smith form S and a Smith massager M for a nonsingular matrix A in time

$$O(n^\omega \, \mathsf{B}(\log \|A\| + \log n)(\log n)^2).$$

$B(d) =$ cost of integer gcds, $M(d) =$ cost of integer multiplication

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3. Application: we can compute an outer product adjoint formula $(\bar{V}, S, \bar{U})$ for A (without further randomization) in extra time

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Further details found in:

Series of papers on fast integer matrix arithmetic:

- S. Birmpilis, G. Labahn, A. Storjohann, *Deterministic reduction of integer nonsingular linear system solving to matrix multiplication,* Proc. of ISSAC'19, July 15-18, Beijing, China, (2019), 58-65.

- S. Birmpilis, G. Labahn, A. Storjohann, *A Las Vegas Algorithm for Computing the Smith Form of a Nonsingular Integer Matrix*, Proc. of ISSAC'20, July 21-23, Kalamata, Greece, (2020).

- S. Birmpilis, G. Labahn, A. Storjohann, *A fast algorithm for computing the Smith normal form with multipliers for a nonsingular integer matrix*. Submitted to Journal of Symbolic Computation

- S. Birmpilis, G. Labahn, A. Storjohann, *A softly cubic algorithm for computing the Hermite normal form of a nonsingular integer matrix*. In preparation.