

Vigenère Cipher

Caesar Ciphers (and Affine transformation ciphers) are too weak and suffer from frequency analysis attacks (or more simply brute force attacks).

For Example: Encrypt MILLENNIUM using the keyword YTWOK

M	I	L	L	E	N	N	I	U	M	S	Y	T	W	O	K
12	8	11	11	4	13	13	8	20	12	18	24	19	22	14	10
+ 24 19 22 14 10 24 19 22 14 10 $\text{mod } 26$															
10	1	7	25	14	11	6	4	8	22	16					
K	B	H	Z	O	L	G	E	I	W	Q					

Idea: Use each letter in the keyword to encrypt the plaintext.

Also combine with a full affine transformation cipher. For example

A	P	P	L	E	Keys: $\left[\begin{matrix} k_1 & k_2 \\ a_1 & b_1 & a_2 & b_2 \end{matrix} \right] \quad C \equiv a_i P + b_i \text{ mod } 26$	$C \equiv a_i P + b_i \text{ mod } 26$	
0	15	15	11	4		$C \equiv 2P + 1 \text{ mod } 26$	$C \equiv 3P + 3 \text{ mod } 26$
k_1, k_2, k_1, k_2, k_1						$C \equiv 3(0) + 1 \equiv 1 \text{ mod } 26$	$C \equiv 3(15) + 3 \equiv 22 \text{ mod } 26$
1	22	20	10	13		$C \equiv 3(15) + 1 \equiv 20 \text{ mod } 26$	$C \equiv 3(11) + 3 \equiv 10 \text{ mod } 26$
B	W	U	K	N	$C \equiv 2(4) + 1 \equiv 9 \text{ mod } 26$		

We will call the second encryption an Affine Vigenère Cipher.

Example: Encrypt the message

PINOCCHIO IS A BOY

using a Vigenere cipher with the key word TOY. For reference, you may use

	A	B	C	D	E	F	G	H	I	J	K	L	M
P	00	01	02	03	04	05	06	07	08	09	10	11	12
C	19												
T													
C	14												
O													
C	24												
Y													
	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
P	13	14	15	16	17	18	19	20	21	22	23	24	25
C													
C													
C													

TOY
19 14 24

if you'll
those w/
you can encrypt/
decrypt
just as
(potentially).

P I N O C C H I O I S A B O Y

15 8 13 14 2 2 7 8 14 8 18 0 1 14 24

+ 19 14 24 19 14 24 19 14 24 19 14 24 19 14 24 mod 26

8 22 11 7 16 0 0 22 12 1 6 24 20 2 22

I W L H Q A A W M B G Y U C W

Now decipher the message

ZSNIS RMCGL OKTB

using the same encryption scheme above

Z S N I S R M C G L O K T B

25 18 13 8 18 17 12 2 6 11 14 10 19 1

- 19 14 24 19 14 24 19 14 24 19 14 24 19 14 24 mod 26

6 4 15 15 4 19 19 14 8 18 0 12 0 13

G E P P E T T O I S A M A N

Frequency Analysis:

What is the most common English letter?

E	E	T	A	O	I	N	S	... see wikipedia
12.7%	12.7%	9.1%	8.2%	7.5%				

So in an encrypted text, if M is the most likely letter, it is likely that $E \mapsto M$.

This would help with breaking Caesar and Affine Transformation ciphers.

This alone will not help with a Vigenère cipher. First attack ~ 1863 named after Friedrich Kasiski.

Kasiski's Test: Look for ~~repeated~~ repeated strings of length 3 or more.

Once you have a guess of the key length, attempt to decipher OR verify your guess using the index of coincidence (see Friday's lecture).

IS IMPORTANT TO BE USEFUL FOR
ENCIPHERING MESSAGES.

Consequently, we made the correct guess. If we had tried this transformation, and instead of plaintext, it produced garbled text, we would have tried another likely transformation based on the frequency count of letters in the ciphertext.

Example 8.5. Suppose we know that an affine transformation of the form $C \equiv aP + b \pmod{26}$, $0 \leq C \leq 25$, has been used for encryption. For instance, suppose that we wish to cryptanalyze the encrypted message

USLEL JUTCC YRTPS URKLT YGGFV
EELYUS LRYXD JURTU ULVCU URJRK
QLLQL YXSRV LBRYZ CYREK LVEXB
RYZDG HRGUS LJLLM LYPDJ LJ TJU
FALGU PTGVT JULYU SLDAL TJRWU
SLJFE OLP

The first thing to do is to count the occurrences of each letter; this count is displayed in Table 8.7.

Letter	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
Number of Occurrences	2	2	4	4	5	3	6	1	0	10	3	22	1	0	1	4	2	12	7	8	16	5	1	3	10	2

Table 8.7 *The number of occurrences of letters in a ciphertext.*

With this information, we guess that

Example 8.8. Suppose that the ciphertext produced by encrypting plaintext using a Vigenère cipher is

QW	HID	DNZEM	WTLMT	BKTTIT	EMWLZ
WW	VCEV	HLTBS	TUDLG	WNUJE	WEUL
EX	WQO	SLNZ	NLHYQ	ALWEH	VOQWD
VQ	TBW	ILURY	STIJW	CLHW	RNSIH
MN	UDI	YFAVD	ELAGB	LSNZA	NSMIF
GN	ZEM	WALWL	CXEFA	BYJTS	SNXLH
YH	LK	UCLOZ	ZAJHI	HWSM	

We describe the steps we use to break this message. We first use the Kasiski test, looking for repeated triples of letters in the ciphertext. We list our finding in a table.

Triple Starting positions Differences in starting positions

BMW	9, 21, 129	12, 108, 120
ZEM	8, 128	120
ZAN	59, 119	60
NZE	7, 127	120
NZA	58, 118	60
NAV	59, 140	97

$g_{\text{d}}(12, 108, 120, 60, 87, 66)$
 $\boxed{= 3}$