

Fréchet distance between polygonal curves

[Computing the Fréchet distance between two polygonal curves](#)

H Alt, M Godau - International Journal of Computational Geometry & ..., 1995 - World Scientific

As a measure for the resemblance of **curves** in arbitrary dimensions we consider the so-called **Fréchet-distance**, which is compatible with parametrizations of the **curves**. For **polygonal** chains P and Q consisting of p and q edges an algorithm of runtime $O(pq \log(\dots))$

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From: http://scholar.google.ca/scholar?q=computing+the+Fréchet+distance+between+two+polygonal+curves&btnG=&hl=en&as_sdt=0%2C5

Figures here, notes follow.

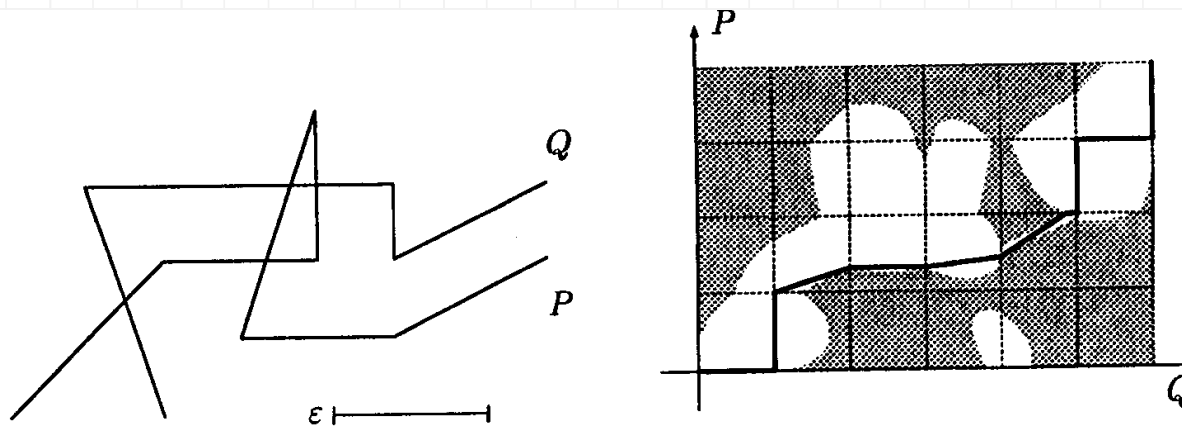


Fig. 3. Diagram for polygonal chains P, Q and the given ϵ

Frechet distance between polygonal curves

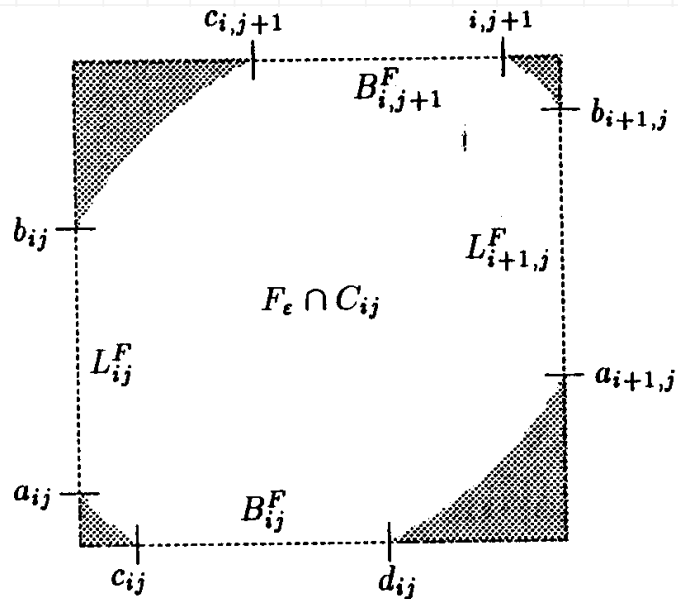


Fig. 4. Intervals of free space on the boundary of a cell.

[Efficient algorithms for geometric optimization](#)

PK Agarwal, M Sharir - ACM Computing Surveys (CSUR), 1998 - dl.acm.org

Abstract We review the recent progress in the design of **efficient algorithms** for various problems in **geometric optimization**. We present several techniques used to attack these problems, such as parametric searching, **geometric** alternatives to parametric searching, ...

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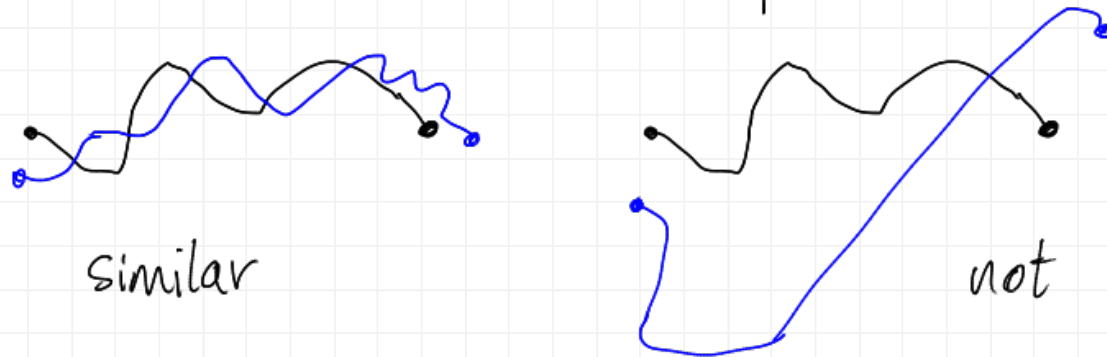
From: http://scholar.google.ca/scholar?q=Efficient+Algorithms+for+Geometric+Optimization&btnG=&hl=en&as_sdt=0%2C5

Frechet distance between polygonal curves

Distances — to measure similarity

e.g. edit distance in strings (can be formulated as distance in a graph, "reconfiguration")

Distance between curves in the plane



similar

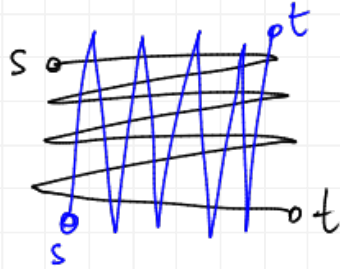
not

Hausdorff distance between A and B .

- for each point $a \in A$ take min distance to B
 - for each point $b \in B$ take min distance to A
- take max of these (need sup, inf in general)

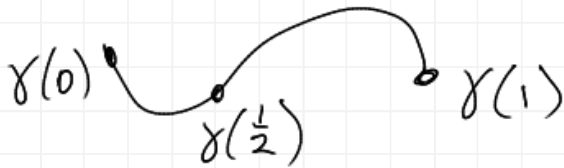
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Hausdorff distance does not capture similarity



only takes into account the set of points, not their order on the curve

curve $\gamma: [0, 1] \rightarrow \mathbb{R}^2$

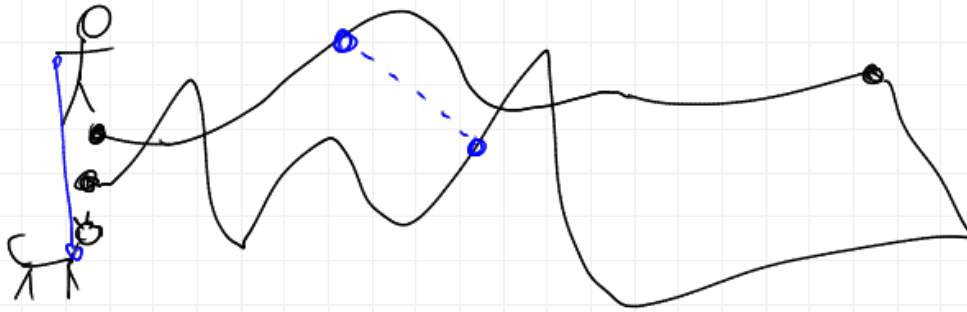


same curve may have different parameterizations

polygonal curve
composed of n
line segments



Frechet distance between polygonal curves



person moves on curve A } never backwards
 dog on curve B

minimum length of leash. = Frechet distance
 more formally (Frechet 1906)

$$\delta_F(\alpha, \beta) = \inf_{\text{reparameterizations } \alpha', \beta'} \max_{t \in [0, 1]} \|\alpha'(t) - \beta'(t)\|$$

how long the leash must be

Frechet distance between polygonal curves

Alt, Godau 1995

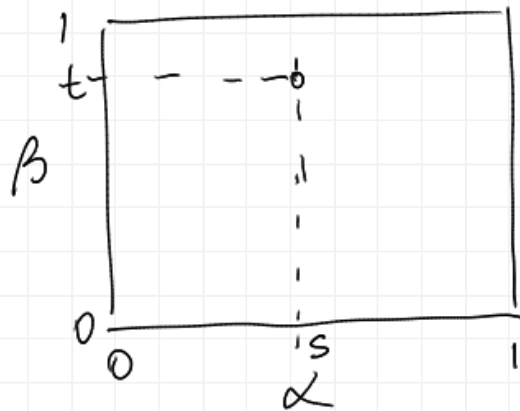
alg. to compute Frechet distance between 2 polygonal curves with a, b vertices

- decision version $O(ab)$

- computing $O(ab \log ab)$

} in real RAM.

Main Idea

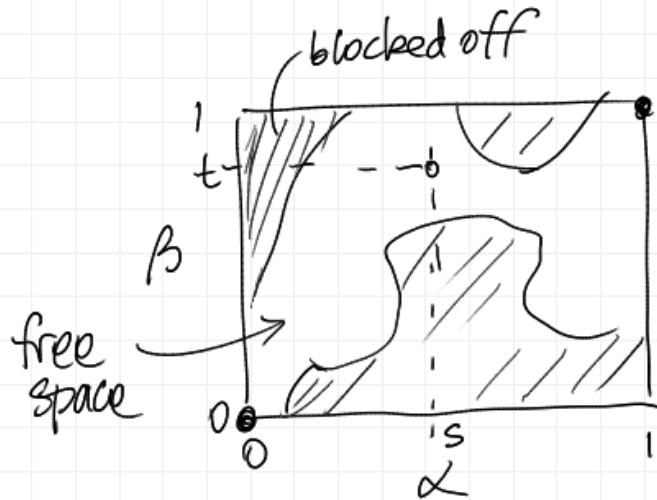


choosing position on α , on β gives a pt. in rectangle

$$d(\alpha(s), \beta(t))$$

= least length for those positions.

Frechet distance between polygonal curves



$$\text{leash} \leq \epsilon$$

For decision problem, block off
pts in rectangle where leash
is $> \epsilon$.

ϵ -diagram

See Fig. 3.

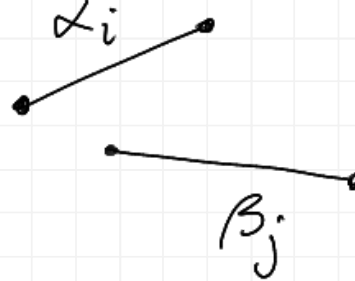
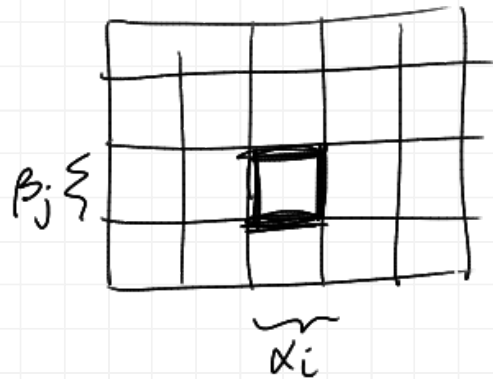
Lemma Frechet distance is $\leq \epsilon$ iff there is
a monotone path from $(0,0)$ to $(1,1)$ in the free
space of the ϵ -diagram -

monotone = x coord never decreases
 y - - - - -

Question: what does the length of the path correspond to?

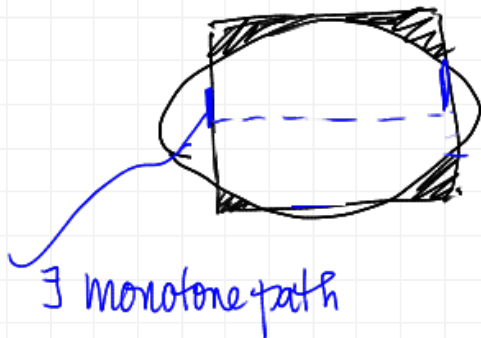
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ϵ -diagram has $a \times b$ rectangles - each corresponds to one segment of α versus one segment of β .



Lemma. Inside one rectangle the free space is an ellipse (Fig.4)

Along each side, the free space is an interval



If we know the subintervals on left & bottom reachable via monotone path from $(0,0)$ then we can compute same for right/top in constant time

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This solves decision problem in $O(ab)$ time.

To actually compute the Frechet distance:

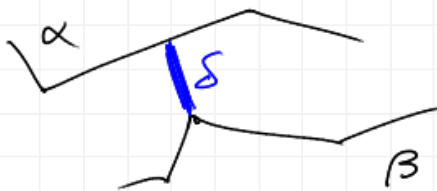
We need $\varepsilon > d(\alpha(0), \beta(0))$ initial & final
 $\varepsilon > d(\alpha(1), \beta(1))$ leash lengths.

Starting from this ε_0 , imagine increasing ε .

Combinatorial changes at critical values of ε

1. a new passage opens between neighbouring cells
2. a new horizontal / vertical passage opens.

1 corresponds to



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Idea:

- determine all critical values (there are $O(a^2b + b^2a)$ of them)
- sort them
- do binary search for ϵ (use decision alg.)

↑
 ───────────┬──────────┐ all ϵ values

test this ϵ and recurse on appropriate side

Time $O(\underbrace{(a^2b + b^2a)}_{\text{sort}} \log(ab) + \underbrace{ab}_{\text{decision}} \underbrace{\log(ab)}_{\text{binary search}})$

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For improved run-time of $O(ab \log(ab))$
use Megiddo's parametric search.

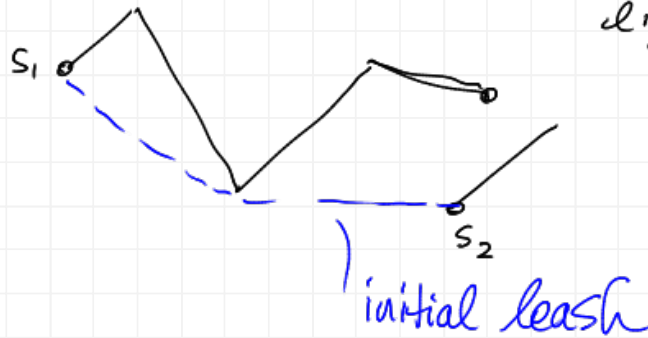
(+ Cole's variant using Ajtai-Komlos-Szemerédi sorting network — high constant and impractical).

See survey paper by Agarwal & Sharir.

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Further work

- can get run time below $O(n^2 \log n)$ - at least for $(1+\epsilon)$ -approx.
- Frechet distance inside a polygon - the leash must remain inside



e.g. the curves themselves might act as boundaries

- Frechet distance between 2 surfaces (up one dimension).