

A Computational Toolkit for Engagement and Scalable Assessment in a Large Logic Course

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Abstract—Large required courses in theoretical computer science face two related challenges: helping students engage with abstract material and supporting reliable student assessment at scale. This paper describes Logicat, a lightweight computational toolkit developed for CS 245, *Logic and Computation*, at the University of Waterloo. The course is required for undergraduate computer science students and serves a large annual cohort.

The main pedagogical objective is to help students concretize the ideas they encounter in lectures and assignments. Handwritten formulas and proof steps do not give students immediate correctness feedback. This can slow their development of confidence in formal reasoning and makes assessment harder to apply consistently at scale. Logicat addresses this by allowing students to manipulate formulas, transformations, clauses, valuations, and proof steps as computational objects in Racket, building directly on their Scheme/Racket experience from the first-year curriculum. The system is lighter than a general proof assistant such as Coq or Lean and uses notation aligned with the course. It exposes composable functions students can invoke individually or use to program their own automations.

The paper presents the design rationale, system organization, and planned course integration of Logicat as a practical model for using computational tools to support engagement, conceptual concreteness, and more consistent assessment in large formal-reasoning courses.

Index Terms—computer science education, learning technologies, logic education, student engagement, scalable assessment, large-enrollment courses

I. INTRODUCTION

Large required courses in theoretical computer science have to make abstract material usable for students while also supporting assessment at scale. In a course on logic, both problems are underscored by the same feature: correctness depends on small formal details. Students often discover such errors only after graded work is returned or, worse, not at all because the errors are missed by the graders ploughing through a large volume of handwritten work.

This paper describes Logicat, a Racket toolkit developed for CS 245, *Logic and Computation*, at the University of

Waterloo [1]. CS 245 is required for undergraduate computer science students and develops mathematical reasoning skills through propositional logic, first-order logic, formal proof, computability, and program verification [2]. The recommended text is Lu’s *Mathematical Logic for Computer Science*, used mainly for definitions, notation, and formal deduction, with additional material developed through lectures and course notes [2], [3].

The aim of Logicat is to make course objects computationally concrete. A formula can be entered, parsed, inspected, transformed, treated under valuations, converted to normal forms, used as a clause in a resolution procedure or as part of a formal deduction proof. These are the same objects students see in lectures and assignments, but now with operations they can apply and test.

The use of Racket is part of the design. Waterloo computer science students have experience with Scheme/Racket from their first-year curriculum, and Racket is well suited to symbolic data and functional composition [4], [5]. Logicat therefore presents logic through a programming medium that students already know. The tool is provided as a set of functions that students can call, combine, and inspect.

The system also addresses a grading problem. In large courses, teaching assistants must repeatedly judge whether symbolic steps are well formed, whether transformations preserve equivalence, whether clauses have been resolved correctly, and whether a proof line follows from the cited rule. This is tedious and error-prone. Logicat can check these conditions automatically. The instructor and teaching assistants can then spend more attention on explanation, strategy, and conceptual understanding.

Logicat was made available on an optional basis during a recent offering of the course and is being prepared for broader integration. This paper gives a design and system account. It focuses on the parts of the toolkit that most directly support the instructional goals: logical formula transformation, reasoning under valuations, resolution procedures and formal deduction proofs. Other capabilities, such as simplification of normal forms, are treated as supporting infrastructure rather than as the main emphasis.

II. COURSE CONTEXT AND INSTRUCTIONAL PROBLEM

Waterloo's CS 245 introduces logic as a tool for representation, reasoning, and computation. The official course description emphasizes mathematical reasoning, propositional and first-order logic, the distinction between syntax and semantics, formal proof, and the limits of computation [2]. The course outline includes logical connectives, truth tables, well-formed formulas, semantics, logical equivalences, normal forms, Boolean algebra, formal deduction, soundness and completeness, resolution, the Davis–Putnam procedure, first-order logic, Peano arithmetic, undecidability, and Hoare-style program verification [2].

This range of topics asks students to move among several views of the same formal material: A formula has a syntactic structure. It has a truth value under a valuation or interpretation. It can be transformed by equivalence laws. It can be converted to a normal form and used as input to an algorithm. It can appear as a premise or conclusion in a formal proof. Students learn the topics separately, but the course expects them to see how the views fit together.

Handwritten work makes this integration difficult. A student may write several lines after an early invalid transformation, or may build a formal proof whose third line is not actually justified by the first two. The feedback may arrive days later, if at all. At that point the student has lost the connection between the error and the reasoning decision that produced it. A small formal error becomes a delayed and often opaque comment in the margin.

The same issue appears in assessment. A large course requires a large grading team. Even with clear rubrics, judgments are hard to apply uniformly. One grader may treat a variation as harmless where another sees an ambiguity. One may give generous partial credit for a proof with an invalid early step, while another may treat the later lines more strictly. More importantly, grading handwritten logic under time pressure is error-prone. These aspects are not failures of goodwill — they are consequences of requiring a team to grade dense logic work by hand.

Logicat was designed to reduce the delay and variability of such evaluation. Students still write proofs and explanations, but they also have a way to test the formal objects that their explanations refer to. In a large class, this can change the students' learning process and graders' workflow and accuracy. Some errors can be found while students are still thinking, and graders can check results more easily, or even have them pre-screened.

III. PEDAGOGICAL AIMS

A. Concretizing Formal Objects

Students first meet logical formulas as strings of notation. To reason with them reliably, they must see more structure. A formula has an inductive structure. A rule application has a scope. A resolution step has two parent clauses and a resolvent. A proof step has cited premises, a rule, a substitution, and a conclusion. Logicat makes such structure accessible through data objects and operations.

This matters because many errors in beginning logic are structural. A student may apply De Morgan's law to wrong subformulas, distribute over the wrong connectives, mis-parenthesize, cite a rule whose requirements have not been met, or resolve two clauses on a symbol that does not occur with opposite signs. Immediate checks on such steps help students separate two questions: whether they chose a good overall strategy, and whether a particular operation was legitimate.

B. Building on Racket Experience

Logicat is a Racket library rather than a standalone application. This choice keeps the tool close to the programming experience students already have. A student can enter a formula as a string, inspect its internal representation, call a transformation function, or write a small program that performs a sequence of desired steps.

The design also lets the course use programming as a way to understand logic. A formula can be traversed and transformed. A resolution procedure can be read as a program over clauses. A proof checker can be connected to the same rule-instantiation discipline students are expected to follow by hand. The point is not to replace formal reasoning by programming, but to let students use a familiar computational environment to work with the formal objects of the course.

C. Supporting Assessment at Scale

Logicat can provide local correctness checks before grades are assigned. It can check whether formulas are well formed, whether transformations are valid, whether two propositional formulas agree under valuations, whether a resolution step is valid, or whether a formal deduction step follows from its justification. These checks leave the quality of a student's explanation for human judgment, while removing uncertainty about many formal details.

This gives the grading process a better starting point and overall reliability. Students can submit work that has already passed some mechanical checks, and graders can focus on the remaining questions: the choice of method, the clarity of explanation, and the student's understanding of the reasoning.

IV. SYSTEM ORGANIZATION

Logicat represents logical formulas as structured Racket data. Atomic propositions are represented by symbols; compound formulas are represented by expressions whose first component is an operator and whose remaining components are arguments. For example, the formula

$$A \rightarrow (B \wedge C)$$

has the internal representation

$$(\text{impl } A \text{ (and } B \ C)).$$

The representation is simple enough for students to inspect directly and regular enough to support parsing, formatting, traversal, transformation, evaluation, and proof checking.

The system accepts input as text strings or as Racket data structures. Students may write formulas using Unicode logical symbols, ASCII alternatives, or LaTeX syntax. This flexibility is useful because logic texts and software use a wide range of notation [3], [6], [7]. Logicat normalizes inputs to a common internal form, while output routines can print formulas in a logical notation, Boolean-algebra notation, or LaTeX. This allows students to copy inputs and outputs into assignments and notes.

The supporting operations include formula input and output, structural inspection, truth valuations, truth tables, normal-form conversion, application of specific equivalence rules, and automatic simplification routines. Propositional logic is supported most fully, while formal deduction for first-order logic is supported as needed for the course.

Additionally, an interactive logical formula editor allows students to apply transformations to whole formulas and subexpressions as needed.

These features are useful in assignments, but they are mainly infrastructure for the parts of the system where students work with course reasoning: local transformations, resolution, and formal deduction. The rest of the paper concentrates on those parts.

V. REPRESENTATIONS USED BY THE REASONING TOOLS

The choice of data representation is kept visible because it helps students connect the course topics. A formula entered in course notation is not stored as a string. It is stored as a structured expression whose operator and arguments can be extracted. This is a basic implementation point, but it also has a pedagogical value: it keeps students focused on formula structure.

The representation also uses n -ary forms for associative connectives such as conjunction and disjunction. This avoids forcing students to choose an artificial binary parenthesization when the course treats a formula such as $A \wedge B \wedge C$ as a conjunction of three parts (after associativity is established). At the same time, operations are available to re-associate a formula when a binary rule or a particular proof format requires it. Thus the system can respect the informal classroom convention while still exposing the formal structure.

For resolution and the Davis–Putnam procedure in propositional logic, Logicat uses a separate term representation for conjunctions and disjunctions of literals. A term records, for each propositional symbol, whether the symbol appears positively, negatively, or not at all. This is a compact way to represent clauses and implicants. It also makes resolution easier to explain computationally: two clauses resolve on a symbol when that symbol appears positively in one clause and negatively in the other, and the resolvent is obtained by combining the remaining literals, unless a tautology is produced.

This representation is useful in teaching because it gives students a concrete view of clauses where the set notation can be confusing. A clause can be displayed as a formula, but it can also be displayed compactly as a row of signs relative

to a fixed symbol order. The latter form is close to the way an algorithm processes the data and makes algorithm traces intelligible. Students can therefore compare the mathematical rule with the data structure that implements it.

Formal deduction uses another family of objects. Entailments, rules, substitutions, justifications, steps, and proofs are represented explicitly. These objects correspond closely to the annotations students write in a proof: the line number, the statement proved on that line, the rule used, the cited previous lines, and the instantiation of the rule. By making these parts explicit, Logicat can check validity without requiring any prose explanation.

These representation choices support the same instructional pattern. Students begin with familiar notation, but the tool enforces the structure that notation abbreviates. Once that structure is visible, the student can ask more precise questions: which subformula is selected, which clauses are being resolved, which substitution instantiates the rule, or which earlier lines justify the present one. Those questions are where students often make errors.

VI. FORMULA TRANSFORMATIONS AND FEEDBACK

Formula transformation is the first place where Logicat gives students immediate feedback on course-level reasoning. The system provides transformation rules for a range of properties proven in the course: implication elimination, equivalence elimination, double-negation removal, associativity, commutativity, distribution, and De Morgan transformations. Some rules act at the top level of a formula; others operate recursively, for example to eliminate non-basic connectives, push negations down to atoms, or convert formulas to DNF or CNF.

The local nature of rule application is important. In handwritten work, students may write a transformed formula without recording exactly which subformula changed. Logicat makes the selected subformula explicit. A student can select a part of the formula, apply a rule to that part, and see the new whole formula. This gives a concrete counterpart to the usual classroom instruction that equivalence transformations must be applied to well-defined subexpressions.

Logicat also includes a small formula editor. The editor maintains a current formula and a current selection. Students can move the selection within the formula tree, grow it across adjacent arguments of an associative connective, and apply a rule to the selected part.

For propositional logic, truth tables and normal forms remain available as checks. A student can transform a formula and then test whether the result is equivalent to the original. The tool can also show valuations under which two formulas differ. These checks are simple, but they often reveal the kind of error that would otherwise be found only in high-quality grading.

This use of the tool also helps move students toward planning. If local steps are checked as they are made, students can spend less effort wondering whether a copied connective or a local equivalence would be valid, and more effort deciding

```

;; Apply the Davis-Putnam procedure on the formulas in lf-list,
;; eliminating proposition symbols in the order given by sym-list.
(define (dpp atom-list lf-list)
  (define (make-clause lf) (make-term-with-lf 'or atom-list lf))
  (define ((has-pos? ix) term) (equal? tlit-pos (term-ref term ix)))
  (define ((has-neg? ix) term) (equal? tlit-neg (term-ref term ix)))

  (let* ([cnf-list (map lf-cnf lf-list)]
        [conjuncts (map lf-conjuncts cnf-list)]
        [uconjuncts (apply lset-union #:equal? lf=? conjuncts)]
        [clauses (filter term? (map make-clause uconjuncts))])
    (do ([i 0 (+ 1 i)])
        [(= i (length atom-list))]
        (let* ([atom (list-ref atom-list i)]
              [pos-atom-clauses (filter (has-pos? i) clauses)]
              [neg-atom-clauses (filter (has-neg? i) clauses)]
              [atom-clauses (append pos-atom-clauses neg-atom-clauses)]
              [resolvents '()])
          (for ([pos pos-atom-clauses])
            (for ([neg neg-atom-clauses])
              (let ([r (resolve-terms i pos neg)])
                (when (term? r) ;; Exclude #t terms.
                  (set! resolvents (lset-union1 r resolvents)) ) ) )
              (set! clauses (lset-minus clauses atom-clauses))
              (set! clauses (lset-union clauses resolvents ) ) )
            ;; At this point clauses is either empty
            ;; or has only the empty clause.
            (if (null? clauses) clauses (list (lf-or))) ) )
  )

```

Fig. 1. Programming example: the Davis-Putnam procedure

what form they are trying to reach. In a normal-form exercise, for example, the question becomes not only “what is the next legal rewrite?” but “which legal rewrite moves the formula toward the intended shape?”

VII. PROGRAMMABILITY

Students can write their own Racket programs to perform sequences of steps. This can be useful when they wish to put formulas into some particular form, or for more ambitious purposes.

They can also inspect example programs that embody the algorithms taught in the course. One such example is the Davis-Putnam procedure for resolution theorem proving. We use this as an illustration.

The `dpp` procedure converts formulas to CNF, extracts clauses, partitions clauses by positive and negative occurrences of a symbol, resolves pairs, and adds the resolvents back to the clause set. Figure 1 shows the structure of the implementation. Details such as tracing and error handling are omitted, as are the definitions of the subordinate functions. An important point is that the code uses the same exposed operations that are available elsewhere in the library, *e.g.* `lf-cnf` which converts a logical formula to conjunctive normal form.

The example is included because it shows the role of composability. CNF conversion, conjunct extraction, term

construction, term inspection, set operations on lists, and resolution are separate pieces. The Davis-Putnam procedure is obtained by composing them. This makes the algorithm available at two levels. A student can run it as a command to check an example, or can read and modify the program to understand how resolution is being applied.

When run in verbose mode (the code for which is not shown), the procedure displays formulas, clauses, selected symbols, resolving pairs, resolvents, and the remaining clause set. This trace is useful in a course setting because it connects the formal resolution rule to the changing state of the algorithm. The trace also helps students see why an implementation must make choices that are easy to omit in a paper proof, such as the order in which symbols are eliminated and the treatment of tautological resolvents.

The same code can also become an assignment object. Students can be asked to run the procedure on a small inconsistent set, explain the empty-clause result, change the symbol order, or compare the trace with a manual resolution derivation. These are modest programming tasks, but they require students to connect the course rule to an executable procedure.

Successful proof: substitutions inferred

```
(fd-check-proof (mk-fd-proof ' (
  [1 "A->B, B->C, A |- A->B"
    (mem)      ]
  [2 "A->B, B->C, A |- A"      (mem)      ]
  [3 "A->B, B->C, A |- B"      (impl- 1 2)]
  [4 "A->B, B->C, A |- B->C" (mem)      ]
  [5 "A->B, B->C, A |- C"      (impl- 4 3)]
  [6 "A->B, B->C |- A->C"      (impl+ 5) ]
)))
```

Output:

```
Checking step 1: Subst found
  A = A->B, Sigma = {B->C,A}. OK
Checking step 2: Subst found
  A = A, Sigma = {A->B,B->C}. OK
Checking step 3: Subst found
  A = A, B = B,
  Sigma = {A->B,B->C,A}. OK
Checking step 4: Subst found
  A = B->C, Sigma = {A->B,A}. OK
Checking step 5: Subst found
  A = B, B = C,
  Sigma = {A->B,B->C,A}. OK
Checking step 6: Subst found
  B = C, A = A,
  Sigma = {A->B,B->C}. OK
Proof OK!
```

Failing proof: explicit substitutions

```
(fd-check-proof (mk-fd-proof ' (
  [S1 "A->B, B->C, A |- A->B" (mem)
    "A=A->B, Sigma={B->C,A}" ]
  [S2 "A->B, B->C, A |- B" (mem)
    "A=A, Sigma={A->B,B->C}" ]
  [S3 "A->B, B->C, A |- B" (impl- S1 S2)
    "A=A, B=B, Sigma={A->B,B->C,A}" ]
  [S4 "A->B, B->C, A |- B->C" (mem)
    "A=B->C, Sigma={A->B,A}" ]
  [S5 "A->B, B->C, A |- C" (impl- S4 S3)
    "A=B, B=C, Sigma={A->B,B->C,A}" ]
  [S6 "A->B, B->C |- A->C" (impl+ S5)
    "A=A, B=C, Sigma={A->B,B->C}" ]
)))
```

Output:

```
Checking step S1: Subst given. OK
Checking step S2: Subst given. Failed
Proof fails at step S2.
Proof failed.
```

Fig. 2. Formal deduction examples of correct and incorrect proofs

VIII. FORMAL DEDUCTION CHECKING

Formal deduction is the part of the course where delayed feedback can be most damaging. Students often write incorrect proofs with plausible overall shape, but where an early line does not follow from its cited rule. Several later lines may then depend on a statement that has not been established. Delayed feedback means that students do not get sufficient practice at discriminating between valid and invalid operations. One of the principal motivations in the development of Logicat was to help students learn to navigate this challenge.

Logicat represents the elements of a formal deduction proof explicitly: entailments, rules, substitutions, justifications, steps, and proofs. An entailment has the form

$$P_1, \dots, P_n \vdash C,$$

where the premises may be formulas or symbols representing sets of formulas, and the conclusion is a formula. A rule specifies required entailments and a resulting entailment. A proof step contains a claimed entailment, a justification naming a rule and previous steps, and a substitution showing how the rule is instantiated.

The checker verifies proof steps in sequence. It checks that the cited rule exists, that the cited previous steps exist, and that the rule's requirements and result match the cited entailments

and the current claim after substitution and canonicalization. If the substitution is not supplied, the system infers it if possible. The result is a check on whether a proof step has been justified.

Figure 2 gives examples of correct and incorrect proofs of the transitivity of implication in the course's formal-deduction notation. The first proof uses the theorem `mem` for membership in the premise set, implication elimination `impl-` ($\rightarrow -$) and implication introduction `impl+` ($\rightarrow +$). The student supplies the proof steps but not the substitutions; Logicat finds the rule instantiations.

The checker returns `#t`. In verbose mode it also reports the substitutions that it inferred for the rule applications, as shown. This output is useful pedagogically because it shows more than success or failure. Students can verify how a schematic inference rule is matched to a concrete line of proof. The substitutions make explicit what is often left implicit, and often wrong, in handwritten work. The set of surrounding premises, represented here by `Sigma`, is part of the rule instantiation rather than a background assumption. Note the ASCII forms " $\rightarrow$ " and " $\vdash$ " in the example are accepted input forms for the course notations " $\rightarrow$ " and " $\vdash$ ".

The system also allows substitutions to be supplied explicitly. This is useful for debugging, for tests, and for assignments in which the instantiation itself is part of what the student is expected to provide.

The second “proof” has explicit substitutions, but line S2 incorrectly claims membership of B from premises that do not include B . The checker returns $\#f$ and reports the first failing line.

This example illustrates the assessment problem at the scale of the course. A short proof of six lines is easy to inspect after the error is pointed out. In assignments, proofs may be twenty or thirty lines, with later lines depending on earlier ones. A graduate teaching assistant can find such errors, but doing so reliably across hundreds of submissions is slow and uneven. The checker gives a uniform answer to the local question: does this line follow from the cited rule and previous steps?

It also changes the student’s task. Without checking, students often produce a sequence of lines and hope that the derivation will be accepted as a proof. With checking, they can find and correct their mistakes while working. The work shifts toward proof planning: identifying useful intermediate claims, choosing when to introduce or eliminate a connective, and arranging the proof so that it follows from justified steps. The checker does not choose the plan, but it prevents unsupported lines from silently becoming part of the plan.

This differs in scope from using a full proof assistant. Systems such as Coq and Lean support machine-checked mathematics and verified software at a much larger scale [8], [9]. They also bring their own syntax, libraries, tactic languages, and proof cultures — and consequently would take considerable course time to introduce. Logicat is designed for the notation and inference rules of the specific course. It gives students feedback on the formal steps used in the course without making the operation of an advanced proof assistant itself a major topic.

IX. USING CHECKED WORK IN LEARNING ACTIVITIES

The examples above suggest several ways to use Logicat in assignments without changing the character of the course. The tool can be used privately by students while they work, as a source of transcripts or program fragments in submitted work, or as a checking aid for graders. These uses can be combined gradually. A first assignment might ask students only to use the tool to confirm a transformation. A later assignment might ask them to submit a short Racket expression that builds and checks part of a proof.

One useful pattern is a guided transformation exercise. Students are given a formula and a target form, such as a formula using only $\neg$, $\wedge$, and $\vee$, or a formula in conjunctive normal form. They must record the transformations they chose, but they can use Logicat to check each local application. The written part of the answer can then ask why those transformations were chosen. This keeps the emphasis on reasoning while reducing time spent on uninformative copying errors.

A second pattern is a resolution exercise. Students can first perform a manual resolution derivation and then compare it with a Davis–Putnam trace. If the two differ, the question becomes diagnostic: did the manual proof use a different elimination order, did it omit a resolvent, or did it introduce

a clause that does not follow? The tool does not replace the manual derivation; it gives students another representation of the same process. This is particularly helpful because resolution is both a proof rule and an algorithmic procedure.

A third pattern is proof repair. Students can be given a nearly correct formal deduction proof containing one or two invalid lines. Their task is to run the checker, identify where the proof first fails, explain why the cited rule does not justify the line, and repair the proof. This activity is close to what teaching assistants do when grading. It also helps students see that a proof is not just a sequence of plausible statements. Each line has a local contract: a rule, cited previous steps, and an instantiation.

A fourth pattern uses the programmability of the library. Students can write small Racket functions that generate families of examples, apply a normalization procedure, or test a conjectured equivalence over a list of formulas. Such exercises use the students’ existing programming background to deepen their contact with the formal material. This is consistent with evidence that active engagement improves learning in STEM courses [10], [11], and with previous work using interactive computational tools in theoretical computer science topics [12].

These activities have a common form. The student is not asked simply to produce a final formula or a final proof. The student is asked to work with intermediate objects: selected subformulas, clauses, resolvents, substitutions, and checked proof lines. Those intermediate objects are where much of the learning happens. They also give the course staff better evidence about what the student understood and where an error entered the reasoning.

For a large required course, this matters. Some students will quickly see how to extend the library, while others will only use the examples provided with the assignments. Both levels of use are useful. The basic commands give immediate feedback on local correctness. The programming interface gives stronger students a way to explore the logic more deeply, while students who are less enthusiastic about the course still benefit from more engagement.

X. COURSE INTEGRATION AND ASSESSMENT WORKFLOW

Logicat was introduced on an optional basis in a recent course offering and is being prepared for broader integration. The initial use was deliberately conservative. Students could use the system to explore examples and check work, but success in the course did not depend on adopting a new software workflow before the documentation and assignment design had stabilized.

The planned integration is staged. Early use should be light: students enter formulas, run simple checks, and compare the output with what they have done by hand. Later assignments can require more structured artifacts, such as a transcript of transformations, a checked proof fragment, or a short Racket expression that constructs and verifies part of a solution. This progression allows students to begin with direct commands and move gradually toward more programmatic use.

This staged approach is important in a required course. Some students will quickly use the library as a programming environment and write their own helpers. Others will use only the commands needed for an assignment. Both levels of use are acceptable. The basic commands give immediate feedback on local correctness, while the programming interface gives stronger or more curious students a way to explore the formal material more deeply.

The submitted work can still include handwritten or typeset explanations. The difference is that some formal details can be checked before submission. A student building a formal deduction proof, for example, can verify that each line follows from its cited rule and previous lines while the proof is being written. If a step fails, the student receives feedback at the point where the reasoning decision was made, rather than days later in a grading comment.

The assessment workflow can use the same artifacts at several levels. First, students may use Logicat privately as a checking aid. Second, an assignment may ask for selected output or code showing how part of a solution was checked. Third, course staff may re-run selected checks when grading. In each case, the tool supplies a consistent answer to local formal questions, while teaching assistants remain responsible for evaluating explanation, strategy, and understanding.

This changes the role of grading without removing human judgment. The teaching assistants need not be the first source of feedback on every syntax error, invalid transformation, or unjustified proof line. Nor need they rely only on visual inspection of a long proof that may contain an early unsupported step. Their effort can be directed toward the questions that require judgment: whether the student chose an appropriate method, explained the reasoning clearly, and reached the required result in a principled way.

XI. LIMITATIONS AND FUTURE WORK

Logicat is still under development. Its mature components support propositional formula manipulation, local transformations, normal forms, simplification, resolution, Davis–Putnam-style consistency checking, and formal deduction checking in the style used by the course. First-order formal deduction and resolution proving has recently been added, as well as program verification as used in Waterloo’s CS 245.

The system also requires careful assignment design. Running commands is not evidence of understanding by itself. Assignments must ask for explanations, intermediate choices, and interpretation of results, with Logicat used to check and explore the formal objects involved. This is especially important for formal deduction: a correct transcript can show that a proof is justified, but students must still explain why the chosen intermediate claims are useful.

A further limitation is that the current paper does not report measured learning gains. The next stage is full integration into the course, followed by evaluation. Possible measures include student engagement, patterns of tool use, reduction in common local errors, grading consistency across teaching assistants, and student attitudes toward formal reasoning. Active-learning

studies in STEM education provide a useful background for this evaluation, but course-specific evidence will be needed [10], [11]. Work on interactive tools for theoretical computer science, such as JFLAP, also provides a point of comparison for studying how computational artifacts affect student work in formal subjects [12].

Future development will focus on three practical issues. The first is making assignment templates that use the tool without turning assignments into software tutorials. The second is improving diagnostics so that students receive feedback that is useful without giving away the proof strategy. The third is improving support for the first-order material in the course, including richer syntax and proof patterns.

XII. CONCLUSION

Logicat was developed to help students in a large required theoretical computer science course engage more actively. Logical formulas, proofs and related objects are represented as data structures in Racket. This builds on students’ first-year Scheme/Racket experience and gives them immediate checks on many of the formal details that otherwise receive delayed feedback.

The strongest uses of the system are where students see reasoning as a sequence of checkable operations: selecting and transforming a subformula, following a resolution trace, or building a formal deduction proof whose lines are justified. These activities help move student work from producing strings of symbols toward planning and checking a derivation.

The same properties support assessment at scale. Many routine formal judgments can be made accurately and consistently by the tool, leaving teaching assistants to evaluate explanation, strategy, and understanding. Full course integration and evaluation remain future work, but the design provides a practical way to connect programming experience with formal reasoning in a large logic course.

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