# CS 886: Multiagent Systems Introduction to Social Choice 

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## Outline

(9) Introduction

- Motivation
- Formal Model
(2) Two Alternatives: A Special Case
(3) Three or More Alternatives
- Case 1: Agents Specify Top Preference
- Case 2: Agents Specify Complete Preferences

4 Properties for Voting Protocols

- Properties
- Arrow's Theorem
(5) Summary


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## What Is Social Choice Theory

- Study of decision problems in which a group has to make the decision
- The decision affects all members of the group
- Their opinions should count!
- Applications
- Political elections
- Other elections
- Allocations problems (e.g. allocation of money among agents, alocation of goods, tasks, resources....)


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## CS Applications of Social Choice

- Multiagent Planning
- Computerized Elections
- Accepting a joint project
- Rating Web articles
- Rating CD's, movies,...


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## Formal Model

- Set of agents $N=\{1,2, \cdots, n\}$
- Set of outcomes $O$
- Set of strict total orders on $O, L$
- Social choice function: $f: L^{n} \rightarrow O$
- Social welfare function: $f: L^{n} \rightarrow L^{-}$where $L^{-}$is the set of weak total orders on $O$


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## Assumptions

- Agents have preferences over alternatives
- Agents can rank order outcomes
- Voters are sincere
- They truthfully tell the center their preferences
- Outcome is enforced on all agents


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Assume that there are only two alternatives, $x$ and $y$. We can represent the family of preferences by

$$
\left(\alpha_{1}, \ldots, \alpha_{n}\right) \in \mathbb{R}^{n}
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where $\alpha_{i}$ is 1,0 , or -1 according to whether agent $i$ preferes $x$ to $y$, is ambivalent between them, or prefers $y$ to $x$.

Definition (Paretian)
A social choice function is paretian if it respects unanimity of strict preferences on the part of the agents.

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## Majority Voting

$$
f\left(\alpha_{1}, \ldots, \alpha_{n}\right)=\operatorname{sign} \sum_{i} \alpha_{i}
$$

$f(\alpha)=1$ if and only if more agents prefer $x$ to $y$ and -1 if and onyl if more agents prefer $y$ to $x$. Clearly majority voting is paretian.

## Additional Properties

- Symmetric among agents
- Neutral between alternatives
- Positively responsive

> Theorem (May's Theorem)
> A social choice function $f$ is a majority voting rule if and only if it is symmetric among agents, neutral between alternatives, and positively responsive.

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## Plurality Voting

The rules of plurality voting are probably familiar to you (e.g. the Canadian election system)

- One name is ticked on a ballot
- One round of voting
- One candidate is chosen
- Candidate with the most votes

> Is this a "good" voting system?

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## Plurality Example

- 3 candidates
- Lib, NDP, C
- 21 voters with the following preferences
- 10 C>NDP>Lib
- 6 NDP>Lib>C
- 5 Lib>NDP>C
- Result: C 10, NDP 6, Lib 5

The Conservative candidate wins, but a majority of voters (11)
prefer all other parties more than the Conservatives.

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## What Can We Do?

Majority system works well when there are two alternatives, but has problems when there are more alternatives.

Proposal: Organize a series of votes between 2 alternatives at
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## Agendas

- 3 alternatives $\{A, B, C\}$
- Agenda: $\langle A, B, C\rangle$

where X is the outcome of majority vote between A and B , and Y is the outcome of majority vote between X and C .


## Agenda Paradox: Power of the Agenda Setter

 3 types of agents: $A>C>B$ (35\%), $B>A>C$ (33\%), $C>B>A$ (32\%).3 different agendas:


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## Pareto Dominated Winner Paradox

## 4 alternatives and 3 agents

- $X>Y>B>A$
- $A>X>Y>B$
- $B>A>X>Y$


BUT Everyone prefers X to Y

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## Maybe the problem is with the ballots

Now have agents reveal their entire preference ordering.
Condorcet proposed the following

- Compare each pair of alternatives
- Declare "A" is socially preferred to "B" if more voters strictly prefer A to B

Condorcet Principle: If one alternative is preferred to all other candidates, then it should be selected.

## Definition (Condorcet Winner)

An outcome $a \in O$ is a Condorcet Winner if $\forall O^{\prime} \in O$,
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## Condorcet Example

- 3 candidates
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## Result: NDP win since 11/21 prefer them to the Conservatives and 16/21 prefer them to the Liberals.

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## There Are Other Problems With Condorcet Winners

- 3 candidates: Liberal, NDP, Conservative
- 3 voters with preferences
- Liberal > NDP>Conservative
- NDP>Conservative>Liberal
- Conservative>Liberal>NDP

Result: Condorcet winners do not always exist.

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- Conservative>Liberal>NDP

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## Borda Count

- Each ballot is a list of ordered alternatives
- On each ballot, compute the rank of each alternative
- Rank order alternatives based on decreasing sum of their ranks

$$
\begin{aligned}
& A>B>C \\
& A>C>B \\
& C>A>B
\end{aligned} \quad \Rightarrow \quad B: 8, \begin{aligned}
& A: 4 \\
& C: 6
\end{aligned}
$$

$$
\therefore \quad \theta \quad \equiv \text { 三 }
$$

## Borda Count

## - The Borda Count is simple

- There is always a Borda winner
- BUT the Borda winner is not always the Condorcet winner

> 3 voters: 2 with preferences $B>A>C>D$ and one with $A>C>D>B$ Borda scores: A:5, B:6, C:8, D:11
> Therefore $A$ wins, but $B$ is the Condorcet winner.

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## Other Borda Count Issues: Inverted-Order Paradox

## Agents

- $\mathrm{X}>\mathrm{C}>\mathrm{B}>\mathrm{A}$
- $A>X>C>B$
- $B>A>X>C$
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- $\mathrm{X}>\mathrm{C}>\mathrm{B}>\mathrm{A}$

Borda Scores

- X:13, A:18, B:19, C:20

Remove X

- C:13, B:14, A:15


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## Vulnerability to Irrelevant Alternatives

3 types of agents

- $X>Z>Y(35 \%)$
- $Y>X>Z(33 \%)$
- $Z>Y>X$ (32\%)

The Borda winner is $X$.
Remove alternative Z. Then the Borda winner is Y.

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## Other Scoring Rules

- Copeland
- Do pairwise comparisons of outcomes.
- Assign 1 point if an outcome wins, 0 if it loses, $\frac{1}{2}$ if it ties
- Winner is the outcome with the highest summed score
- Kemeny
- Given outcomes $a$ and $b$, ranking $r$ and vote $v$, define $\delta_{a, b}(r, v)=1$ if $r$ and $v$ agree on relative ranking of $a$ and $b$
- Kemeny ranking $r^{\prime}$ maximises $\sum_{v} \sum_{a, b} \delta_{a, b}{ }^{\prime}(r, v)$


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## Properties for Voting Protocols

## Property (Universality)

A voting protocol should work with any set of preferences.

## Property (Transitivity)

A votina protocol should produce an ordered list of alternatives (social welfare function).

Property (Pareto efficiency)
If all agents prefer $X$ to $Y$, then in the outcome $X$ should be prefered to $Y$. That is, SWF $f$ is pareto efficient if for any


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## More Properties

## Property (Independence of Irrelevant Alternatives (IIA))

Comparison of two alternatives depends only on their standings among agents' preferences, and not on the ranking of other alternatives. That is, SWF $f$ is IIA if for any $o_{1}, o_{2} \in O$

Property (No Dictators)
A SWF $f$ has no dictator if $\neg \exists i \forall O_{1}, O_{2} \in O, O_{1}>_{i} O_{2} \Rightarrow O_{1}>_{f} O_{2}$

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## Arrow's Theorem

## Theorem (Arrow's Theorem)

If there are 3 or more alternatives and a finite number of agents, then there is no SWF which satisfies the 5 desired properties.

## Is There Anything That Can Be Done?

Can we relax the properties?

- No dictator?
- Fundamental for a voting protocol
- Paretian?
- Also pretty fundamental
- Transitivity?
- Maybe you only need to know the top ranked alternative?
- Stronger form of Arrow's theorem says that you are still in trouble
- IIA?
- Universality
- Some hope here (1 dimensional preferences, spacial preferences...)


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## Take-home Message

- Despair?
- No ideal voting method
- That would be boring!
- A group of more complex that an individual
- Weigh the pro's and cons of each system and understand the setting they will be used in
- Do not believe anyone who says they have the best voting system out there!


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## For Further Reading I

A. Mas-Colell, M. Whinston, and J. Green.

Microeconomic Theory.
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